## CS5220 Project 2: Shallow Water Equations

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#### 1 Introduction

In this project, we worked with the a finite volume solver for the shallow water equations. These PDEs model the movement of waves in scenarios where the waves are long compared to the depth of the water. We were provided with a moderately performant single threaded version of the code, and made an effort to improve performance by parallelizing the algorithm. For our implementation, we chose to use MPI as our parallelization paradigm. Our algorithm divides the grid up into a series of subdomains, and assigns one subdomain to each processor. Processors are responsible for calculating the behavior of the water in their portion of the domain. They communicate information about the cells on the edge of their domain with the processors of neighboring domains so that information propagates properly across domain boundaries.

The aims for this project are: 1) using MPI as parallelization implementation for the single-core code; 2) conducting weak and strong scaling studies on the shallow water problem to test the robustness of our implementation; 3) profiling and tuning the codes for higher-level optimization based on computational resources.

### 2 The Algorithm

For the purposes of this algorithm, we are considering two consecutive calls to central2d\_step to be "one time step," in order to alleviate any confusion with the shifting grids in Jiang and Tadmor.

To address the problem of parallelizing the method of Jiang and Tadmor, we use a domain decomposition method like the one introduced for Conway's Game of Life. For the problem, we assume that the domain  $\Omega = (0, L)^2$  is divided into square cells, with a resolution of M cells in each direction. Because the problem is periodic, this amounts to M "points" at which the height h, x-velocity u, and y-velocity v are known. To parallelize the code, we divide the domain into an  $N_x \times N_y$  grid of subdomains, where each subdomain is owned by a process. By this construction, we are using a total of  $N_x N_y$  processors.

As a part of this construction, each subdomain is responsible for knowing the value of  $m_x \times m_y$  points at all times where  $m_x = M/N_x$  and  $m_y = M/N_y$ . However, due to the way information flows through a hyperbolic system of equations, we also require each processor to store a "halo" of ghost cells with width  $m_g$ . Each processor is not tasked with knowing the value of the ghost cells, but rather retrieves the values of the ghost cells from its neighbors. The width of the halo is related to how many time steps each processor can perform independently before information must be exchanged – called  $m_t$  – by the relation  $m_g = 4m_t$ . Independent of the number of processors (assuming M is fixed), the we will call the total number of time steps is  $m_t N_t$ .

We must also keep track of the maximum speed of information across all of the processors, so that each processor takes the same time step. In the code, we are currently sharing this information at every time step, but this turns out not to be a large burden for the sizes of problems we are considering.

In total, the algorithm can be summarized by Alg. 1. Inside of the main loop, we first share the ghost cells. This goes in the order of pass-left and receive-right; pass-right and receive-left; pass-up and receive-down; pass-down and receive-up. In the left/right passing of ghost cells, we pass a block of size  $m_g m_y$ , whereas for up/down we pass a block of size  $m_g (m_x + 2m_g)$ , as we need to communicate the corners of the halo along with the sides in this step.

#### Algorithm 1 Parallel Jiang-Tadmor Main Loop

```
for i \leftarrow 1 to N_t do

Share ghost cells via 4 calls of MPI_Sendrecv

for j \leftarrow 1 to m_t do

Get local time step using speed

Get global time step via MPI_Allreduce

Perform time step

end for

end for
```

#### 2.1 Model Speed Up

The general process:

- 1. Used MPI
- 2. Transfer Boundary information to left and right neighbors, the top and bottom (transferring the corners to top and bottom)
- 3. Transfer boundary information every  $\tau$  steps
- 4. Transfer time step information every  $\tau$  steps as well (This is what is being used in the model anyway)

The model

- 1. Four different contributions to batches of time steps
- 2. Give predictions for how strong scaling, weak scaling, and optimal time batching look like.

### 3 Scaling and Profiling Results

#### 3.1 Scaling Studies

We performed both strong and weak scaling studies to analyze the performance of our algorithm. For the strong scaling study, we used the dam break scenario, and fixed the resolution of the grid cells at 1000 cells in both the X and Y dimensions. We then varied the number of processors used to compute the simulation up until a fixed simulation time. The size of the sub-domains varied with the number of processors in order to keep the problem size constant. We compare the wall clock time required to solve a fixed problem as the amount of compute resources available varies. For the weak scaling study, we again used the dam break scenario, but varied the resolution of the grid cells with the number of processors used. The resolution of the grid was varied such that each sub-domain was always 300 cells on each side. The simulation was run until a fixed amount of time had been simulated. In order to account for the different CFL conditions with different grid resolutions, we recorded the total number of time steps processed, and divided the wall clock time by this number to determine the average amount of time required to compute a single tick of the simulation. We compare the wall clock time required to compute a single simulation tick as both the size of the problem and the computational resources available scale. Additionally, for both types of scaling tests we analyzed the impact of varying the frequency of communication between sub-domains. These tests explored the tradeoff between frequency of suffering the penalty of communication and amount of data communicated and duplicated computation performed.

The data collected from the strong scaling test is shown in Table 1. The data in this table is the wall clock time required for the computation. We see the processing time decrease as the number of processors increases. We have also presented this data in graphical form in Figure 1.

For the weak scaling study, we collected data on the amount of work accomplished per time. To get this number, we multiplied the number of cells in the simulation by the number of time steps computed, and divided by the wall clock time. This data is presented in Table 2. As expected, the size of the problem we can compute in a given amount of time increases as the number of processors increase. We present a graph of this data in Figre 2.

sub-domain grid	nproc		simulation ticks per communication						
		1	2	3	4	5	6	7	8
1x1	1	138.95	140.02	136.78	130.53	119.90	130.60	160.26	171.99
2x2	4	48.23	48.17	48.62	29.87	47.93	51.60	30.35	51.15
3x3	9	30.36	30.65	32.47	34.00	30.08	35.37	32.00	39.20
4x4	16	9.50	10.78	12.11	13.88	14.88	17.04	16.48	20.62
5x5	25	3.76	3.72	3.89	4.41	4.63	6.23	6.52	8.85
6x6	36	2.69	2.62	2.65	2.76	2.81	3.08	2.99	3.67
7x7	49	2.22	2.07	2.13	2.17	2.20	2.39	2.30	2.70
8x8	64	1.92	1.72	1.73	1.83	1.83	2.05	1.88	2.32
9x9	81	1.62	1.46	1.41	1.51	1.53	1.78	1.56	1.89

Table 1: Strong scaling study data

### Compute Time vs Number of Processors

for different numbers of steps between communication

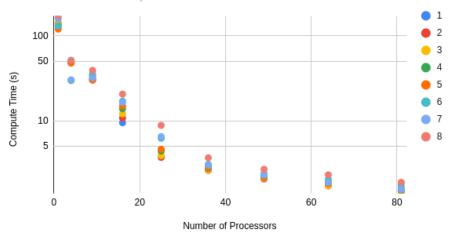


Figure 1: Strong scaling study performance

sub-domain grid	nproc		simulation ticks per communication						
		1	2	3	4	5	6	7	8
1x1	1	5.33e4	3.47e4	3.57e4	4.87e4	3.88e4	4.19e4	2.99e4	2.59e4
2x2	4	6.22e4	4.04e4	3.93e4	4.20e4	3.59e4	3.53e4	2.92e4	6.26e4
3x3	9	3.96e4	4.12e4	3.72e4	3.64e4	3.67e4	3.55e4	3.27e4	3.88e4
4x4	16	5.45e4	5.21e4	5.04e4	4.62e4	4.92e4	4.84e4	4.04e4	4.37e4
5x5	25	6.79e4	6.44e4	6.10e4	5.89e4	5.48e4	5.51e4	5.00e4	5.15e4
6x6	36	8.03e4	7.62e4	7.32e4	7.13e4	6.82e4	6.23e4	5.98e4	5.85e4
7x7	49	9.38e4	8.82e4	8.49e4	8.17e4	7.67e4	7.35e4	6.97e4	6.67e4
8x8	64	1.02e5	9.59e4	9.38e4	8.99e4	8.63e4	8.32e4	7.82e4	7.45e4
9x9	81	1.16e5	1.11e5	1.05e5	9.98e4	9.69e4	9.43e4	8.77e4	8.39e4

Table 2: Weak scaling study data

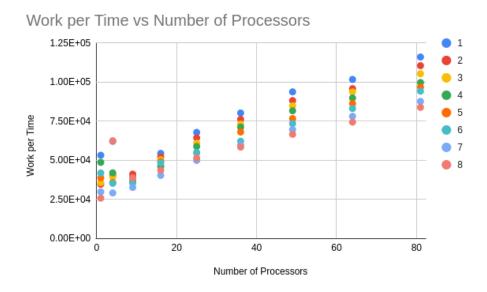


Figure 2: Weak scaling study performance

#### 3.2 Profiling

We used Tau package on Comet as the profiler tool to study the performance of our parallelization and higher-level optimization of the original codes. We used compiler-based instrumentation method, which records timing for both the MPI routines and user-defined functions in the codes. By changing the compiler from mpic to tau\_cc.sh and adding -tau\_options=-optCompInst to the argument for compiling, running the program returns profiling results for each processor and we can look them up by pprof command. The profiling result shows the proportion of time spent, exclusive time, the total inclusive time, number of calls, number of subroutines as well as the average inclusive time per call for each of the functions in the program. As the parallel computing on each individual processor, except for the rank 0 core where we expect to see different behavior due to all the subdomain data gathering to output solution after certain time steps for a rationale check, mostly the mean times spent on all the cores are compared straightforward.

We first did a direct profiling comparison between a 200 × 200 grid and a 1000 × 1000 grid, where in both cases 9 processors (blocking into 3 × 3) were called based on Comet terminal environment. Since the two are computed using the same number of cores, the main difference is the size of subdomains divided into each core, which also indicates blocking size affects the performance of codes. For a small grid problem, the MPI\_Init() used 35.5% of the total profiling time; the major of the code central2d\_xrun consists of three parts, central2d\_step, central2d\_periodic and a MPI routine MPI\_Allreduce(), taking up 53.0%, 7.0% and 0.0%, respectively. central2d\_periodic is the modified code for inter-core communication of ghost cell; MPI\_Allreduce() is used for synchronizing each subdomain at the same physical time. Numerical PDE computing as well as core communication does not take up fully. When using same amount of computational resources, the situation of the large scale becomes different, where central2d\_step, central2d\_periodic and MPI\_Allreduce() take up 95.1%, 3.3% and 0.0% of the total time; MPI\_Init() is only 0.1%. This indicates the bottleneck in a large scale problem running on 9 processors is positioned on how to make the PDE computing itself faster. The most time is spent on the actual PDE numerical computing, whereas the cost of core communication is pretty low. The profiling tables are attached at the end of the report.

#### 3.3 Tuning with Blocking

So a direct way to speed up the computation in the large-scale problem is to introduce more resources, i.e. to decompose into more blocks and assign to more processors. We used sbatch command on both Graphite and Comet to submit the task in order to take advantage of the huge computational resources on these two large-scale server. A code profiling was also run by Tau on Comet utilizing 81 cores (blocking into  $9 \times 9$ ), for comparison with the one using only 9 cores. Under this circumstance, the total time spent on running the whole program with profile instrumentation is greatly reduced, from 15min to 2min. When looking into the computing as well as communication, central2d\_step, central2d\_periodic and MPI\_Allreduce() use up 69.3%, 6.0% and 15.0%, respectively. It demonstrates that when deploying more cores to compute, the proportion of computing itself goes down, while communication and synchronization among each cores goes up. Compared with the small-scale, where the blocking size in each core is more comparable, the PDE computing is in the same scale but  $9 \times 9$  consumes more average inclusive time per call since it still uses nearly twice as large the block size; same as the communication where  $9 \times 9$  is around twice average per call since there are more data to exchange in the same amount of neighboring. The same trends can be found when compared with  $3 \times 3$  on the same large-scale. However, more decomposition gives more distributed tasks to be done, meaning when the communication and synchronization is needed, it will pay more. Luckily, our parallelization optimizes the way of communication so that only the ghost cells have to be updated from the left, right, top and bottom neighboring blocks. Nevertheless, the cost of synchronization, as of MPI\_Allreduce(), grows tremendously from  $22 \text{usec/call in } 3 \times 3$  to  $18823 \text{usec/call in } 9 \times 9$ . The blocking strategy is constraint to how many computational resources are available. There is also trade-off between the individual computing and synchronization. Block size also plays a key role in utilizing the cache. In our test, blocking works prominently well in large-scale study.

#### 3.4 Tuning with Time-step Batching

After making the parallelized code speedup by having more individual processors, another way we tried on is to advance the time-step. In the original code, the iteration is based on a two-time-step staggered grids.

During each iteration, the time is forwards two step to an updated status. In our primary parallelization, this means after two time-step, the periodic condition or individual core communication is needed in the area of the ghost cells. As a way to possible speedup, we can move forwards from the two time-step to more, so that the numbers of communication among each cores can be reduced, herein called time-step batching. We did a profile on the large-scale problem with  $9 \times 9$  cores and a four time-step batching. The numbers of call of central2d\_periodic is reduced half, as expected half less of communication in four time-step batching; the inclusive total time is also decreased from 9164msec to 7437msec. However, since more time-step is forwards, which is proportional to the increase of numbers of ghost cells, as indicated that the average inclusive time per call is increased from 7562usec/call to 12437usec/call. At the same, due to the demand of time-step synchronization, dt has to be communicated in every time-step. So this time-step batching does not help with the timing of textttMPI\_Allreduce(). The time-step batching plays the trade-off between the numbers and the amount of required communication. However, we did not see improvement in the implementation of large-scale study.

#### 4 Conclusion

What we would add for next time:

- 1. More careful cache performance
- 2. Tuning number of ghost cells to block size
- 3. Different sized domains/initial conditions
- 4. Deal with the fact that not all processors are the same
- 5. Think about the tradeoff between a conservative time step and communicating every step (how uneven are the time steps really?)
- 6. Communication between processors at different nodes is different than that of the same nodes?

## A Profiling table: small scale on 9 cores

FUNCTION SUMMARY (mean):

%Time	Exclusive msec	Inclusive total msec	#Call	#Subrs	Inclusive usec/call	Name
100.0	6	16,162	1	1	16162524	.TAU application
100.0	0.954	16,156	1	3	16156274	main
63.8	0.596	10,313	1	119.444	10313260	run_sim
60.0	0.0372	9,703	50	50	194062	central2d_run
60.0	0.538	9,703	50	1250	194061	central2d_xrun
53.0	7	8,559	500	37833.3	17118	central2d_step
35.5	5,737	5,737	1	0	5737190	<pre>MPI_Init()</pre>
26.7	49	4,322	500	218000	8644	central2d_predict
26.0	4,175	4,205	213500	50014.3	20	limited_deriv1
26.0	68	4,200	500	309001	8401	central2d_correct
25.9	4,150	4,179	213500	49986.7	20	limited_derivk
7.0	1	1,135	250	7000	4542	central2d_periodic
7.0	1,131	1,131	1000	0	1131	<pre>MPI_Sendrecv()</pre>
3.0	0.0558	482	51	96.3333	9452	gather_sol
2.5	0.037	398	45.3333	45.3333	8787	send_full_u
2.5	398	398	45.3333	0	8787	MPI_Send()
0.6	104	104	1	0	104869	<pre>MPI_Finalize()</pre>
0.5	80	81	51	9614.89	1597	copy_u
0.5	0.0621	73	45.3333	90.6667	1631	recv_full_u
0.4	60	60	5.66667	0	10762	solution_check
0.4	43	58	100001	100001	1	limdiff [THROTTLED]
0.2	12	28	36833.3	36833.3	1	shallow2d_flux

0.2	25	25	5.66667	0	4555	viz_frame
0.2	25	25	1	0	25248	<pre>MPI_Barrier()</pre>
0.1	20	20	100001	0	0	<pre>central2d_correct_sd [THROTTLED]</pre>
0.1	16	16	36833.3	0	0	shallow2dv_flux
0.1	15	15	100001	0	0	xmin2s [THROTTLED]
0.0	7	7	0.111111	0	66771	viz_close
0.0	5	5	250	0	22	<pre>MPI_Allreduce()</pre>
0.0	3	5	1	13333.3	5336	lua_init_sim
0.0	2	2	6000	0	0	copy_subgrid
0.0	2	2	1	0	2708	viz_open
0.0	0.133	2	250	250	10	shallow2d_speed
0.0	2	2	250	0	9	shallow2dv_speed
0.0	2	2	45.3333	0	47	MPI_Recv()
0.0	1	1	11837	0	0	central2d_offset
0.0	1	1	11111.2	0	0	central2d_offset [THROTTLED]
0.0	0.0189	0.0201	1	2	20	central2d_init
0.0	0.00733	0.00733	1	0	7	copy_basic_info
0.0	0.00122	0.00122	2	0	1	central2d_free
0.0	0.000778	0.000778	1	0	1	<pre>MPI_Comm_size()</pre>
0.0	0.000444	0.000444	1	0	0	<pre>MPI_Comm_rank()</pre>

# B Profiling table: large scale on 9 cores

FUNCTION	SUMMARY	(mean)	):
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%Time	Exclusive msec	Inclusive total msec		#Subrs	Inclusive usec/call	Name
100.0	5	15:34.549	1	 1	934549844	.TAU application
100.0	0.461	15:34.543			934543980	main
99.8	0.763	15:32.687	1	119.444	932687449	run_sim
98.4	0.0854	15:19.270	50	50	18385400	central2d_run
98.4	4	15:19.269	50	6060	18385398	central2d_xrun
95.1	2,382	14:48.441			366519	central2d_step
47.7	1,531	7:25.437 7:21.766	2424	4.93526E+06	183761	central2d_predict
47.3	7:21.737	7:21.766	4.91345E+06	49979	90	limited_derivk
47.1	2,236	7:20.259	2424	4.99163E+06	181625	central2d_correct
47.1	7:20.090	7:20.119	4.91345E+06	50022	90	limited_deriv1
3.3	9	30,591	1212	33936	25240	central2d_periodic
3.3	30,494	30,494				MPI_Sendrecv()
1.1	0.131	10,712	51	96.3333	210043	gather_sol
0.9	0.102	8,610	45.3333	45.3333	189947	send_full_u
0.9	8,610	8,610	45.3333	0	189945	MPI_Send()
0.2	1,977	1,977	51	0		copy_u
0.2	14	1,877	45.3333	90.6667	41415	recv_full_u
0.2	1,508	1,508		0	266283	solution_check
0.1	1,302	1,302	1	0	1302127	MPI_Init()
0.1	616	616	5.66667	0	108842	viz_frame
0.1	553	553	1	0	553943	<pre>MPI_Finalize()</pre>
0.1	503	503	1	0	503194	<pre>MPI_Barrier()</pre>
0.0	34	361	100001	100001	4	shallow2d_flux [THROTTLED]
0.0	327	327	100001	0	3	shallow2dv_flux [THROTTLED]
0.0	1	205	1212	1212	169	shallow2d_speed
0.0	204	204	1212	0	168	shallow2dv_speed
0.0	109	109	45.3333	0	2419	MPI_Recv()
0.0	87	87	29088	0	3	copy_subgrid
0.0	55	70	1	100001	70665	lua_init_sim
0.0	43	58	100001	100001	1	limdiff [THROTTLED]
0.0	42	42	100001	0	0	central2d_correct_sd [THROTTI
0.0	27	27	1212	0	22	<pre>MPI_Allreduce()</pre>
0.0	15	15	100001	0	0	xmin2s [THROTTLED]
0.0	14	14	100001	0	0	<pre>central2d_offset [THROTTLED]</pre>
0.0	4	4	0.111111	0	36135	viz_close
0.0	0.639	0.639	2	0	320	central2d_free
0.0	0.227	0.227	1	0	227	viz_open
0.0	0.0163	0.0171	1	2	17	central2d_init

0.0	0.00878	0.00878	1	0	9 copy_basic_info
0.0	0.000556	0.000556	1	0	1 MPI_Comm_size()
0.0	0.000222	0.000222	1	0	0 MPI Comm rank()

## C Profiling table: large scale on 81 cores

FUNCTION SUMMARY (mean):

%Time	msec	total msec		#Subrs	usec/call	Name
100.0	 6	2:32.563				.TAU application
100.0	255	2:32.556			152556601	
97.2	0.538	2:28.236		109.272	148236114	run_sim
90.3	0.0432	2:17.766		50	2755333	central2d_run
90.3	2	2:17.766	50	6060		central2d_xrun
69.3	140	1:45.760	2424	104849	43631	central2d_step
35.0	392	53,411	2424	1.70326E+06	22034	central2d_predict
34.3	52,316	52,345	1.68145E+06	49967	31	limited_derivk
34.2	52,222	-	1.68145E+06			limited_deriv1
34.2	511	52,112		1.75963E+06	21498	central2d_correct
15.0	22,813	22,813				MPI_Allreduce()
6.3	0.0596	9,664		101.37		gather_sol
6.2	0.0427	9,431				send_full_u
6.2	9,431	9,431		0	187251	MPI_Send()
6.0	7	9,164		33936	7562	central2d_periodic
6.0	9,139	9,139	4848	0	1885	MPI_Sendrecv()
1.7	2,568	2,568	1	0	2568150	<pre>MPI_Init()</pre>
1.0	1,496	1,496	1	0	1496864	MPI_Finalize()
0.4	549	549	1	0	549138	<pre>MPI_Barrier()</pre>
0.2	0.0568	229	50.3704	100.741	4554	recv_full_u
0.1	225	226	51	769.988	4433	copy_u
0.1	167	167	0.62963	0	266338	solution_check
0.1	34	96	100001	100001	1	shallow2d_flux [THROTTLED]
0.0	69	69	0.62963	0	109638	viz_frame
0.0	62	62	100001	0	1	<pre>shallow2dv_flux [THROTTLED]</pre>
0.0	43	58	100001	100001	1	limdiff [THROTTLED]
0.0	0.486	25	1212	1212	21	shallow2d_speed
0.0	25	25	1212	0	21	shallow2dv_speed
0.0	22	22	100001	0	0	central2d_correct_sd [THROTT
0.0	17	17	29088	0	1	copy_subgrid
0.0	15	15	100001	0	0	xmin2s [THROTTLED]
0.0	9	14	1	37037	14530	lua_init_sim
0.0	6	6	50.3704	0	125	MPI_Recv()
0.0	5	5	36572.4	0	0	central2d_offset
0.0	3	3	0.0123457	0	321839	viz_close
0.0	0.182	0.182	1234.58	0	0	central2d_offset [THROTTLED]
0.0	0.118	0.118	2	0	59	central2d_free
0.0	0.0138	0.0145	1	2	15	central2d_init
0.0	0.0119	0.0119	1	0	12	viz_open
0.0	0.00809	0.00809	1	0	8	copy_basic_info
0.0	0.000481	0.000481	1	0	0	MPI_Comm_size()
0.0	0.000272	0.000272	1	0	0	<pre>MPI_Comm_rank()</pre>

## D Profiling table: large scale on 81 cores with 4 time-step

FUNCTION SUMMARY (mean):

%Time	Exclusive msec	Inclusive total msec	#Call	#Subrs	Inclusive Name usec/call
100.0	6	2:41.619	1		161619151 .TAU application
100.0	0.911	2:41.612	1		161612979 main

97.6	0.537	2:37.808	1	109.272	157808979	run sim
91.2	0.0493	2:27.456	50	50		central2d_run
91.2	1	2:27.456	50	5382		central2d_xrun
71.2	184	1:55.142	2392	104785		central2d_step
37.1	415	1:00.023		1.79559E+06		central2d_predict
35.2	56,930		1.71666E+06	50023.1		limited_deriv1
35.2	56,909		1.71666E+06	49977.9		limited_derivk
33.9	508	54,821		1.73772E+06		central2d_correct
15.4	24,854	24,854	1196	0		MPI_Allreduce()
5.9	0.0624	9,552	51	101.37		gather_sol
5.8	0.0445	9,324	50.3704			send_full_u
5.8	9,324	9,324	50.3704			MPI_Send()
4.6	3	7,437	598			central2d_periodic
4.6	7,421	7,421	2392	0		MPI_Sendrecv()
1.6	2,592	2,592	1	0		MPI_Init()
0.7	1,210	1,210	1	0		MPI_Finalize()
0.3	548	548	1	0		MPI_Barrier()
0.1	0.0559	224	50.3704	100.741	4457	recv_full_u
0.1	220	220	51	769.988	4319	copy_u
0.1	167	167	0.62963	0	266331	solution_check
0.1	34	113	100001	100001	1	shallow2d_flux [THROTTLED]
0.0	79	79	100001	0	1	shallow2dv_flux [THROTTLED]
0.0	67	67	0.62963	0	107823	viz_frame
0.0	43	58	100001	100001	1	limdiff [THROTTLED]
0.0	22	22	100001	0	0	<pre>central2d_correct_sd [THROTTLED]</pre>
0.0	0.488	19	1196	1196	16	shallow2d_speed
0.0	19	19	1196	0	16	shallow2dv_speed
0.0	15	15	100001	0	0	xmin2s [THROTTLED]
0.0	9	14	1	37037	14556	lua_init_sim
0.0	11	11	14352	0	1	copy_subgrid
0.0	7	7	50.3704	0	142	MPI_Recv()
0.0	5	5	36572.4	0	0	central2d_offset
0.0	0.494	0.494	0.0123457	0	40006	viz_close
0.0	0.186	0.186	1234.58	0	0	central2d_offset [THROTTLED]
0.0	0.119	0.119	2	0	60	central2d_free
0.0	0.0135	0.0144	1	2	14	central2d_init
0.0	0.00979	0.00979	1	0	10	viz_open
0.0	0.00801	0.00801	1	0	8	copy_basic_info
0.0	0.000679	0.000679	1	0	1	MPI_Comm_size()
0.0	0.000247	0.000247	1	0	0	MPI_Comm_rank()