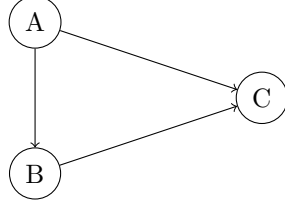


1 Trivial example 1

Consider the graph



where A is located at $(0, 1)$, B is located at $(0, -1)$ and C is located at $(h, 0)$. Denote the distance from A or B to C by l . The supplies are $s_A = (1, 0)$, $s_B = (0, 1)$ and $s_C = (-1, -1)$. Consider the trivial cost function $c(x, y) = |x + y|$ where x, y are the amounts of (signed) flow. Then

$$\partial c(0, 0) = \text{conv} \{(1, 1), (-1, -1), (1, -1), (-1, 1)\}$$

or with inequalities $\partial c(0, 0) = \{z : |e_i^T z| \leq 1 \forall i\}$. Obviously, an optimal flow is

$$f = \begin{pmatrix} 1 & 0 \\ 0 & 1 \\ 0 & 0 \end{pmatrix}.$$

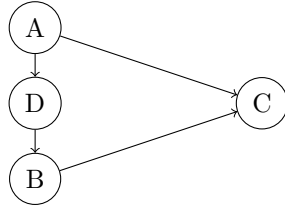
An optimal dual solution is $\phi_A = (l, l)$, $\phi_B = (l, l)$ and $\phi_C = (0, 0)$. It can be easily checked that this satisfies all constraints. The derivative of ϕ is

$$D\phi = \begin{pmatrix} -l/h & 0 \\ -l/h & 0 \end{pmatrix}.$$

We now check whether $D\phi$ meets the constraints, i.e. if $\|e_i^T D\phi\| \leq 1$. We see that $\|e_i^T D\phi\| = l/h > 1$ for all i . This is maybe surprising, as the flow f is globally optimal (with respect to all possible graph topologies).

Observation 1 *Even if the global optimum is found, the dual constraints might not be satisfied.*

Since $e_i^T D\phi = (-l/h, 0)$, this suggests that we should add an edge parallel to $(1, 0)$ to the graph. We get:



The optimal flow stays the same

$$f = \begin{pmatrix} 1 & 0 \\ 0 & 1 \\ 0 & 0 \\ 0 & 0 \end{pmatrix}$$

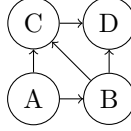
and we get optimal duals $\phi_A = (l, l)$, $\phi_B = (l, l)$, $\phi_C = (0, 0)$, $\phi_D = (h, h)$. One can again directly check that all constraints are satisfied. To see if $D\phi$ is feasible, by symmetry it suffices to consider the triangle DCA . Here,

$$D\phi = \begin{pmatrix} -1 & l-h \\ -1 & l-h \end{pmatrix}.$$

Then $\|e_i^T D\phi\|^2 = \|(-1, l-h)\|^2 = 1 + (l-h)^2 > 1$, so the dual solution is still not feasible.

2 Shortest path example

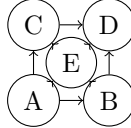
Consider



With supply 1 at A and -1 at D . An optimal solution is

$$f = \begin{pmatrix} 1 \\ 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}$$

(where edges are ordered lexicographically). An optimal dual solution is $\phi_A = 2$, $\phi_B = 1$, $\phi_C = 1$, $\phi_D = 0$. In the triangle BDC , $D\phi = (-1 - 1)$ which suggests an edge in the $(1, 1)$ -direction. In the triangle ABC , $D\phi = (-1 - 1)$ as well. Hence, we refine the traingulation and get



The optimal flow now is

$$f_{AE} = f_{ED} = 1$$

with duals $\phi_A = \sqrt{2}$, $\phi_B = 1/\sqrt{2}$, $\phi_C = 1/\sqrt{2}$, $\phi_D = 0$, $\phi_E = 1/\sqrt{2}$. This is obviously the “correct” optimal solution. Consider now $D\phi$ on all 4 triangles. In ABE , $D\phi = (-1/\sqrt{2}, -1/\sqrt{2})$ which has norm 1. By symmetry, this also holds for all other triangles. Hence in this example we indeed terminate with the optimal solution after one refinement.