

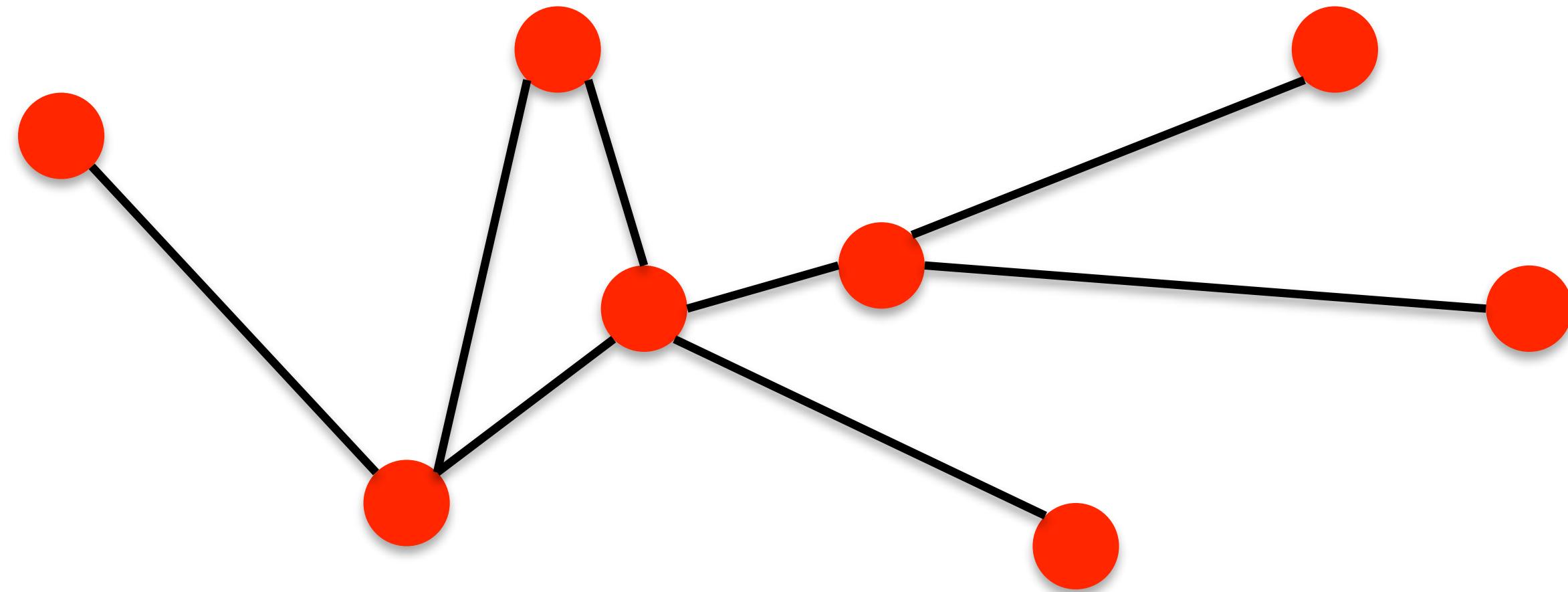
# Network theory

## Part I



Complexity in Social Systems  
AA 2023/2024  
Maxime Lucas  
Lorenzo Dall'Amico

# components



- **components:** nodes, vertices

N

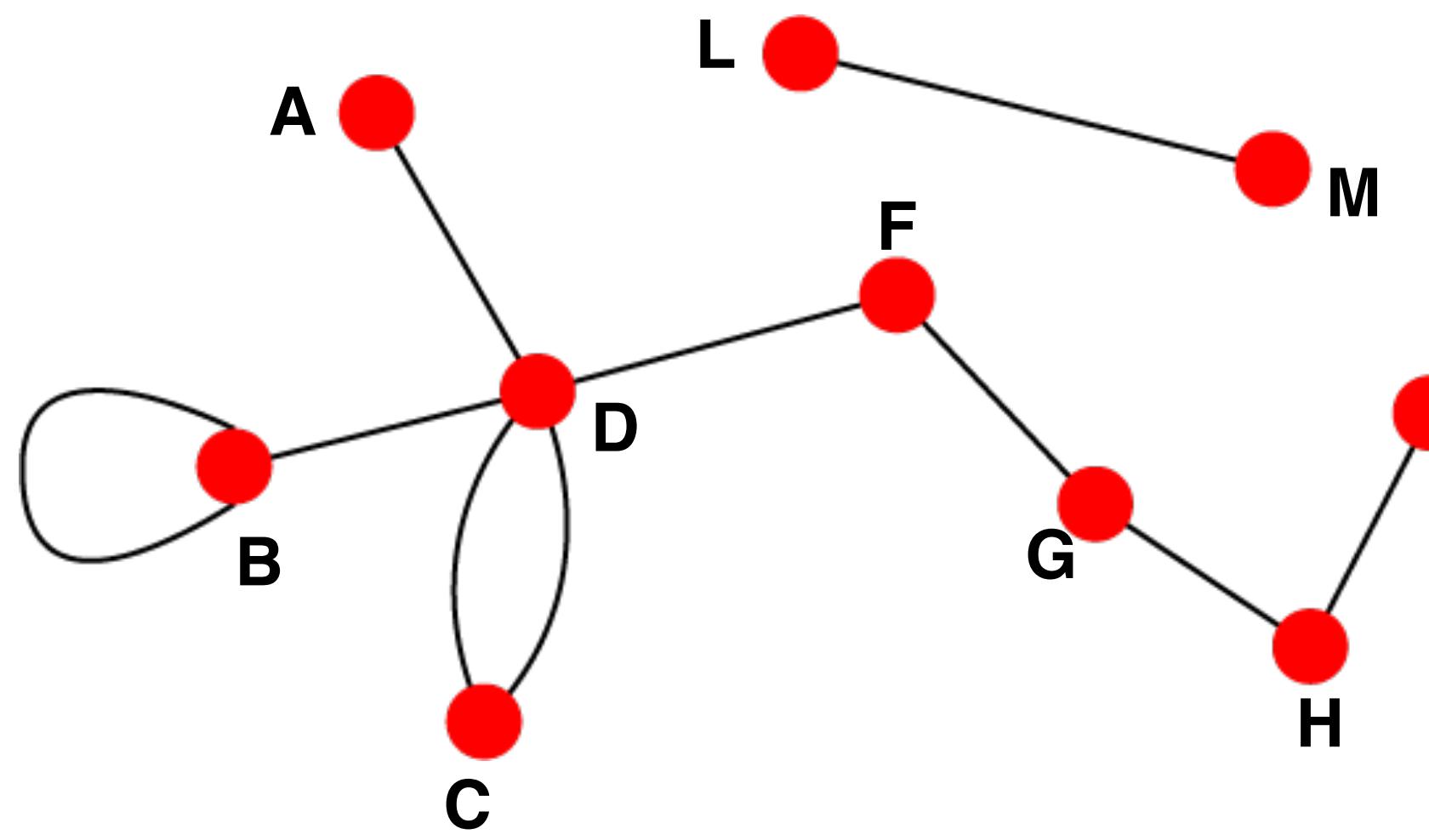
- **interactions:** links, edges

L

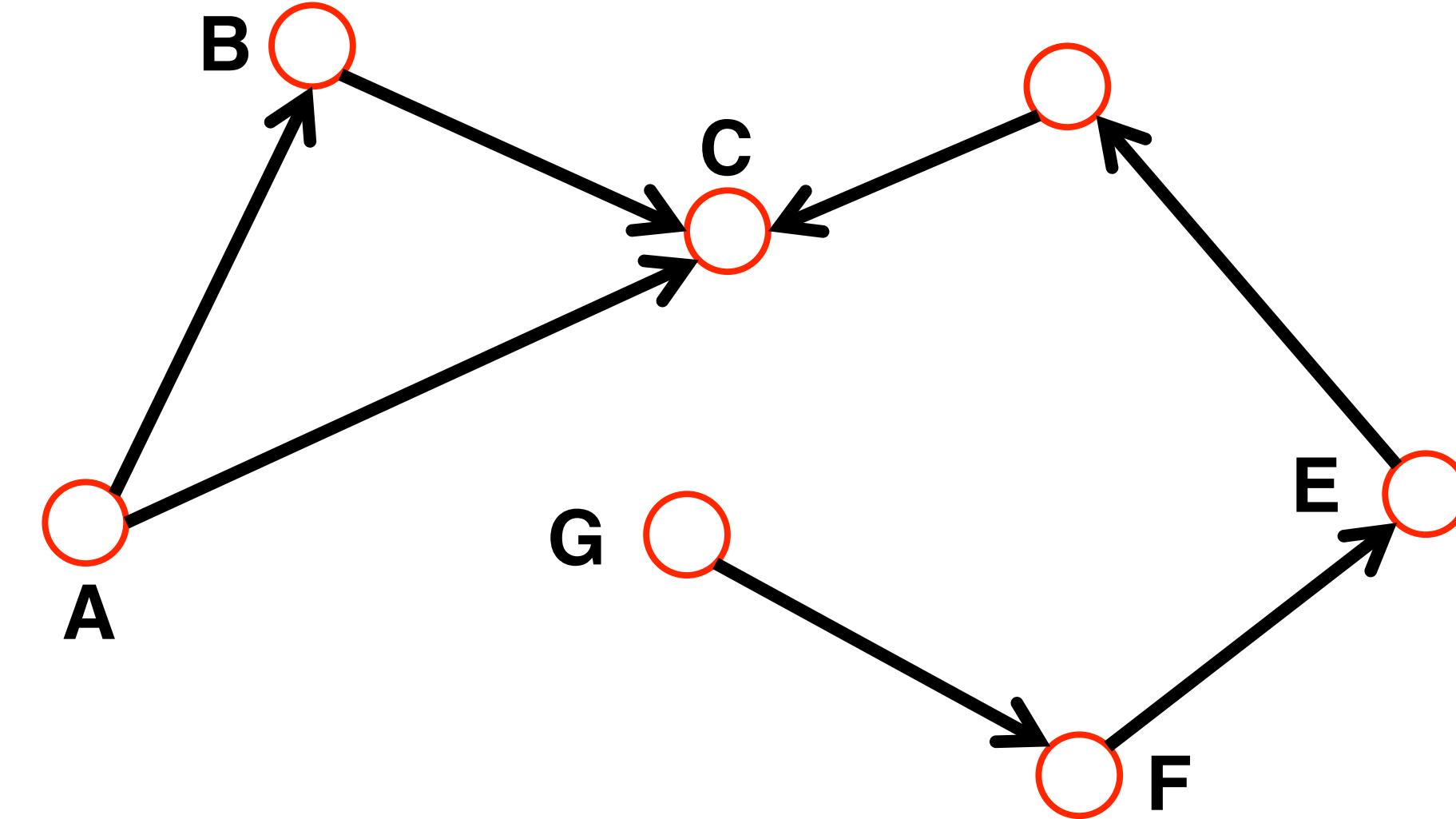
- **system:** network, graph

(N,L)

# undirected vs directed



*co-authorship  
actor networks  
co-occurrence*



*phone calls  
hyperlinks  
scientific citations*

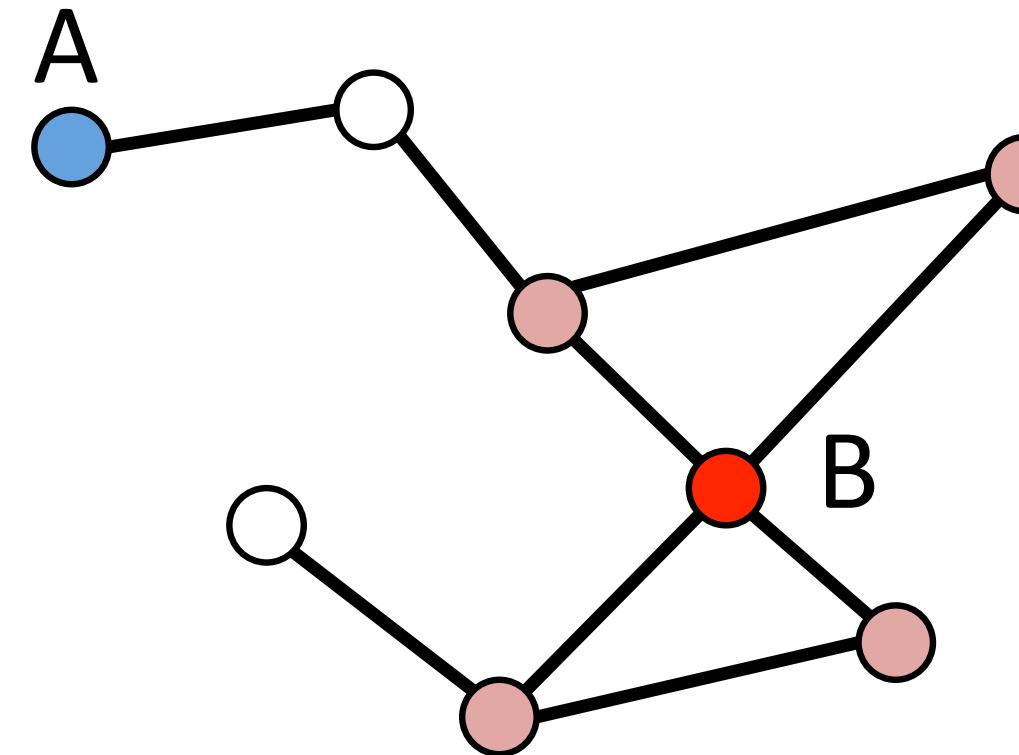
# reference networks

Network	Nodes	Links	Directed / Undirected	N	L	$\langle k \rangle$
Internet	Routers	Internet connections	Undirected	192,244	609,066	6.34
WWW	Webpages	Links	Directed	325,729	1,497,134	4.60
Power Grid	Power plants, transformers	Cables	Undirected	4,941	6,594	2.67
Mobile-Phone Calls	Subscribers	Calls	Directed	36,595	91,826	2.51
Email	Email addresses	Emails	Directed	57,194	103,731	1.81
Science Collaboration	Scientists	Co-authorships	Undirected	23,133	93,437	8.08
Actor Network	Actors	Co-acting	Undirected	702,388	29,397,908	83.71
Citation Network	Papers	Citations	Directed	449,673	4,689,479	10.43
E. Coli Metabolism	Metabolites	Chemical reactions	Directed	1,039	5,802	5.58
Protein Interactions	Proteins	Binding interactions	Undirected	2,018	2,930	2.90



degree and  
degree distribution

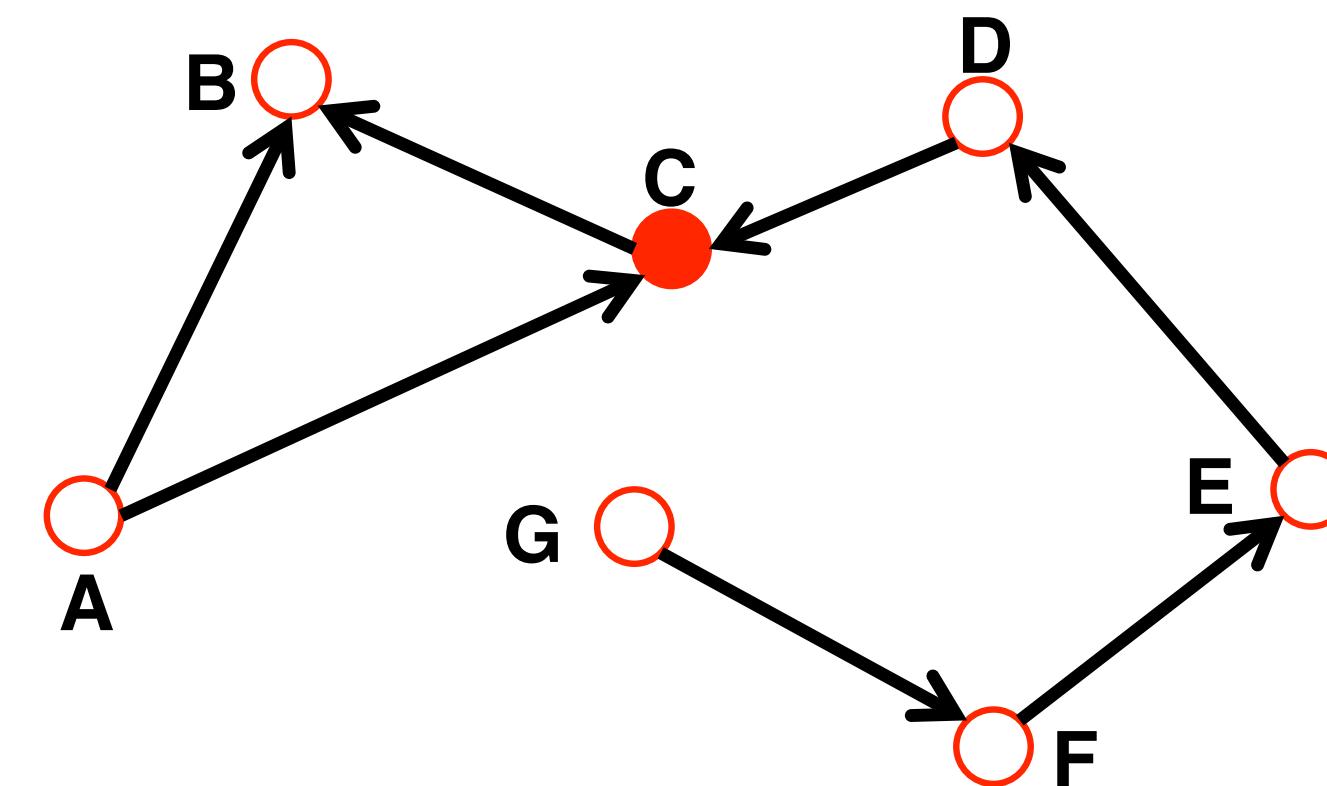
## Undirected



$$k_A = 1 \quad k_B = 4$$

**node degree:** the number of links connected to the node

## Directed



$$k_C^{in} = 2 \quad k_C^{out} = 1 \quad k_C = 3$$

**Source:** degree  $in = 0$

**Sink:** degree  $out = 0$

## BRIEF STATISTICS REVIEW

Four key quantities characterize a sample of  $N$  values  $x_1, \dots, x_N$ :

*Average (mean):*

$$\langle x \rangle = \frac{x_1 + x_2 + \dots + x_N}{N} = \frac{1}{N} \sum_{i=1}^N x_i$$

*The  $n^{\text{th}}$  moment:*

$$\langle x^n \rangle = \frac{x_1^n + x_2^n + \dots + x_N^n}{N} = \frac{1}{N} \sum_{i=1}^N x_i^n$$

*Standard deviation:*

$$\sigma_x = \sqrt{\frac{1}{N} \sum_{i=1}^N (x_i - \langle x \rangle)^2}$$

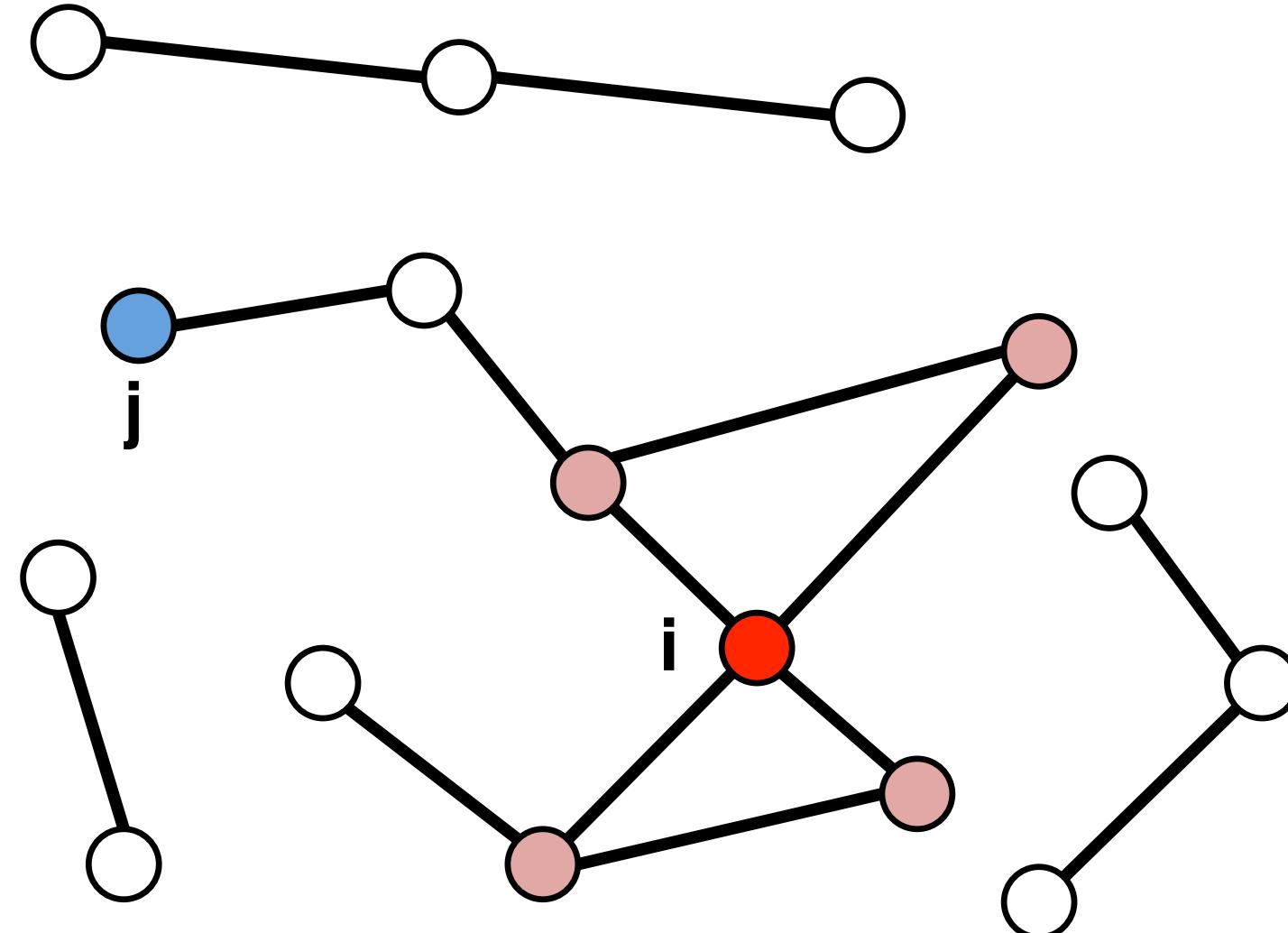
*Distribution of  $x$ :*

$$p_x = \frac{1}{N} \sum_i \delta_{x, x_i}$$

where  $p_x$  follows

$$\sum_i p_x = 1 \quad \left( \int p_x dx = 1 \right)$$

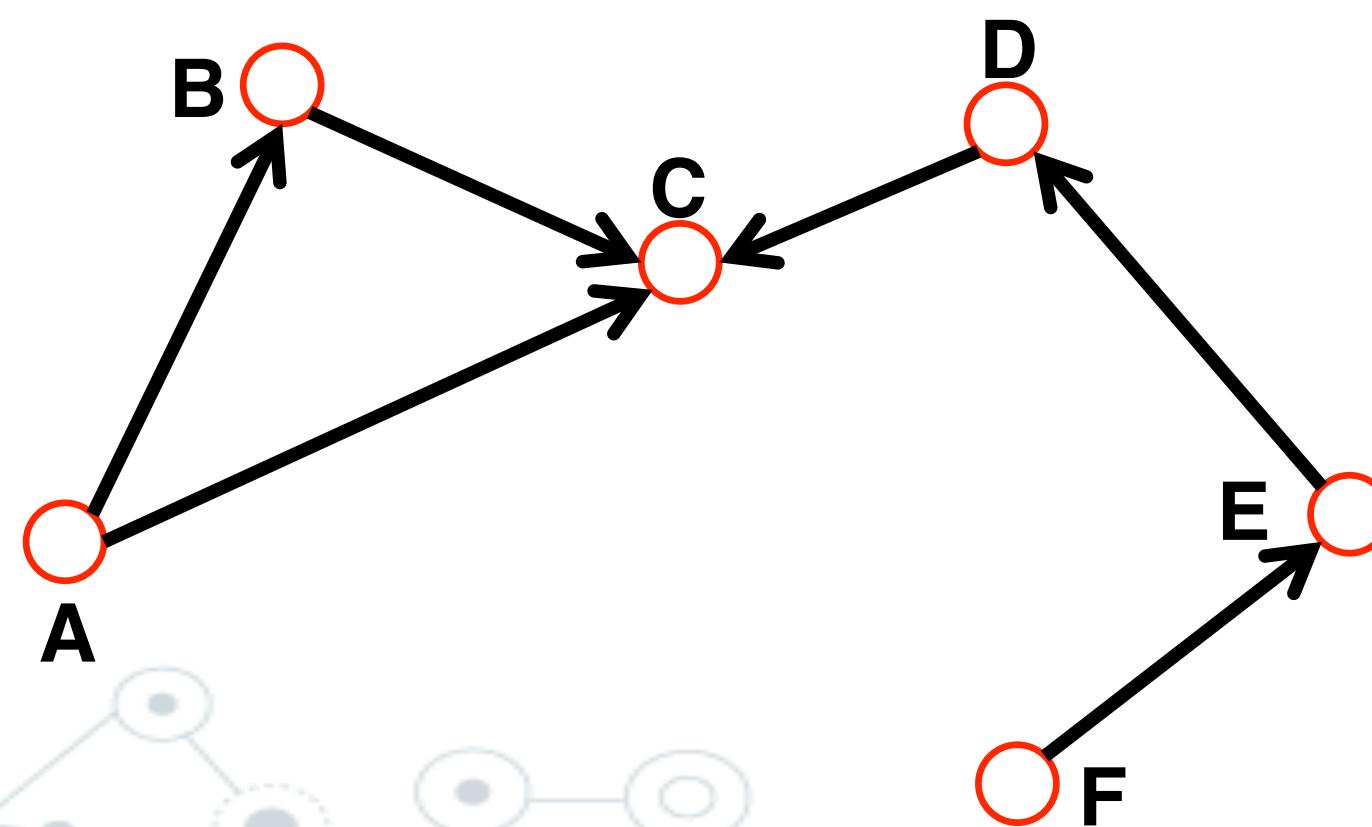
## Undirected



$$\langle k \rangle = \frac{1}{N} \sum_{i=1}^N k_i \quad \langle k \rangle = \frac{2L}{N}$$

N – the number of nodes in the graph

## Directed



$$\langle k^{in} \rangle = \frac{1}{N} \sum_{i=1}^N k_i^{in}, \quad \langle k^{out} \rangle = \frac{1}{N} \sum_{i=1}^N k_i^{out}, \quad \langle k^{in} \rangle = \langle k^{out} \rangle$$

$$\langle k \rangle = \frac{L}{N}$$

# degree distribution

$$P(k) = \frac{N_k}{N}$$

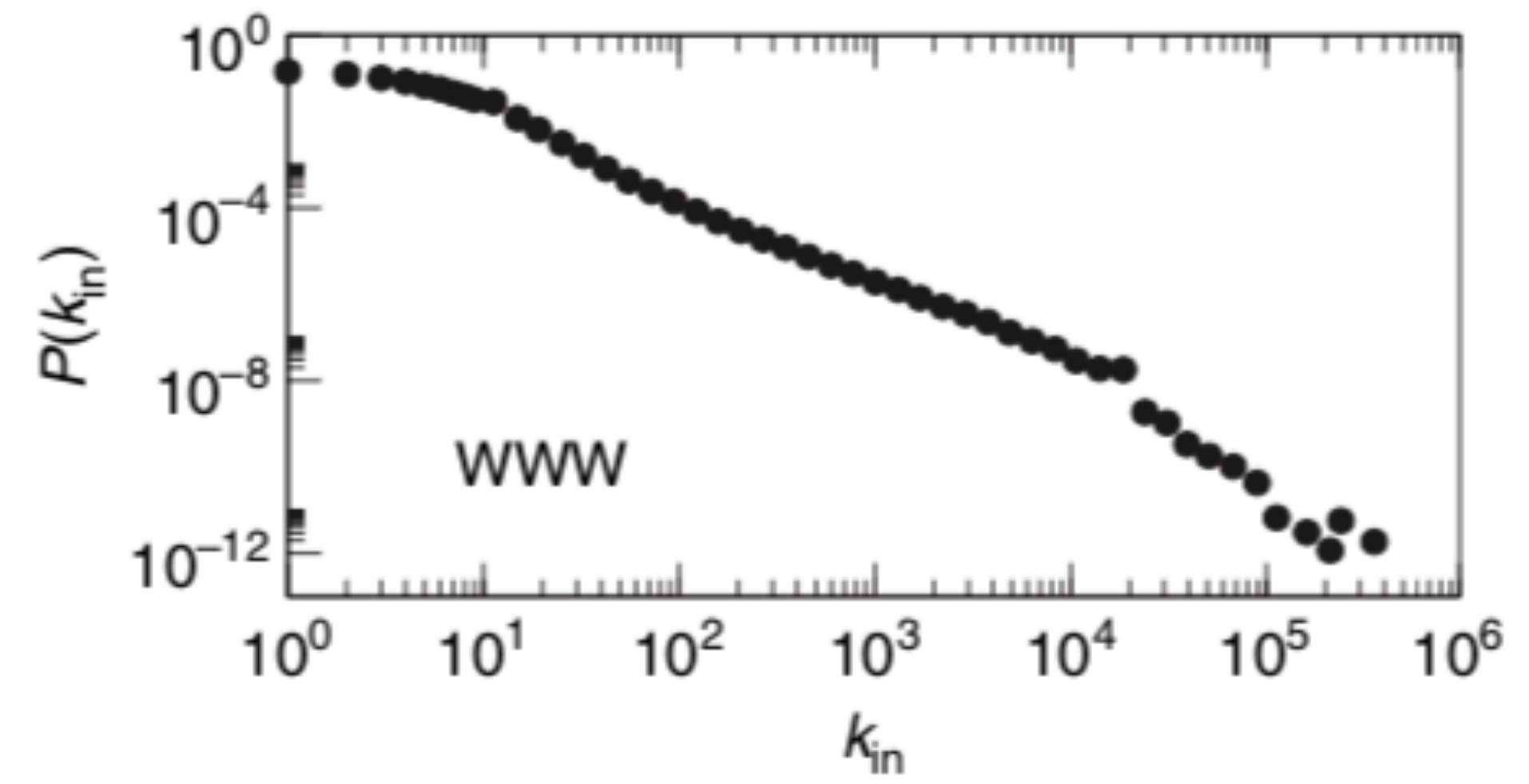
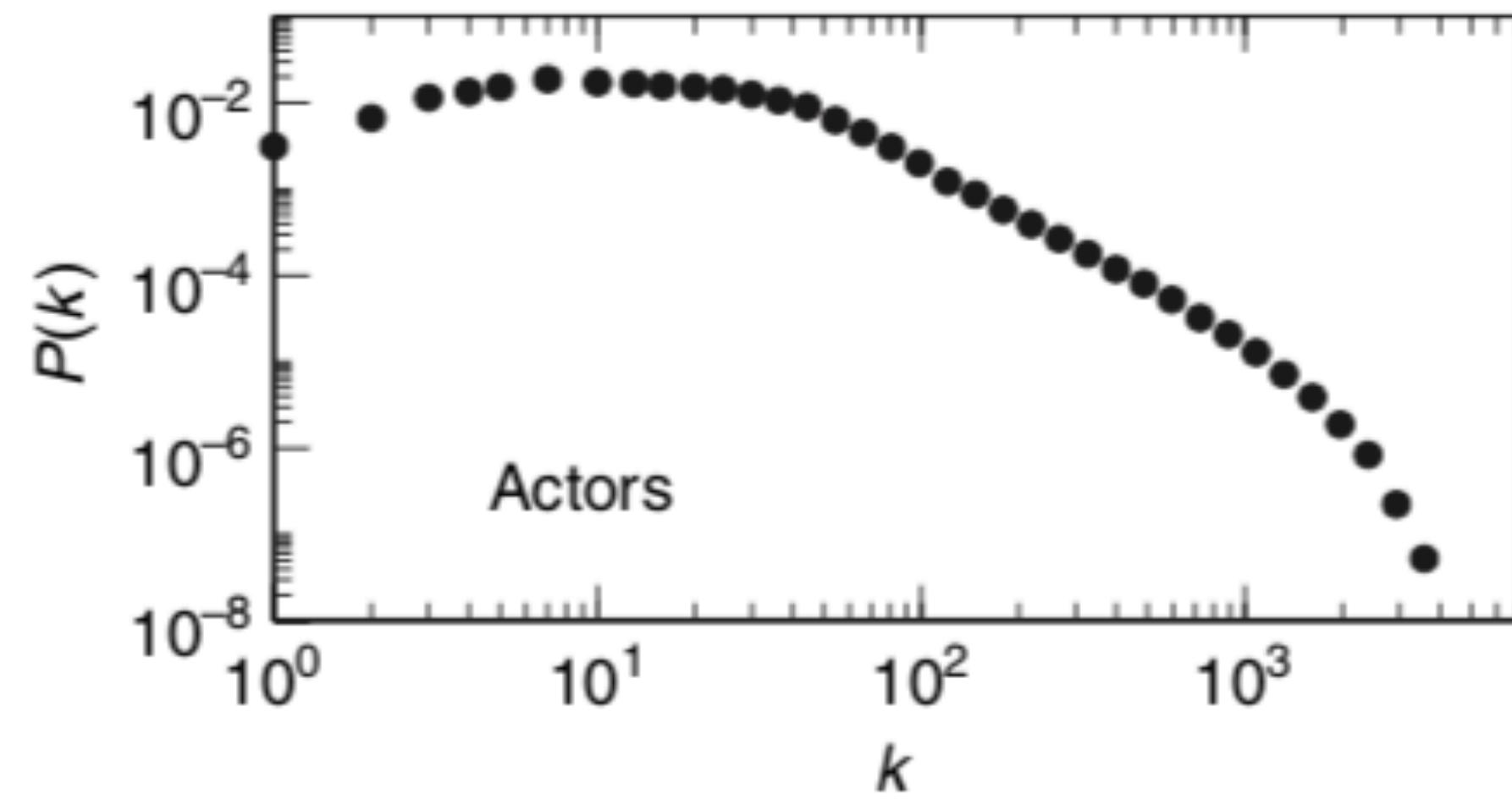
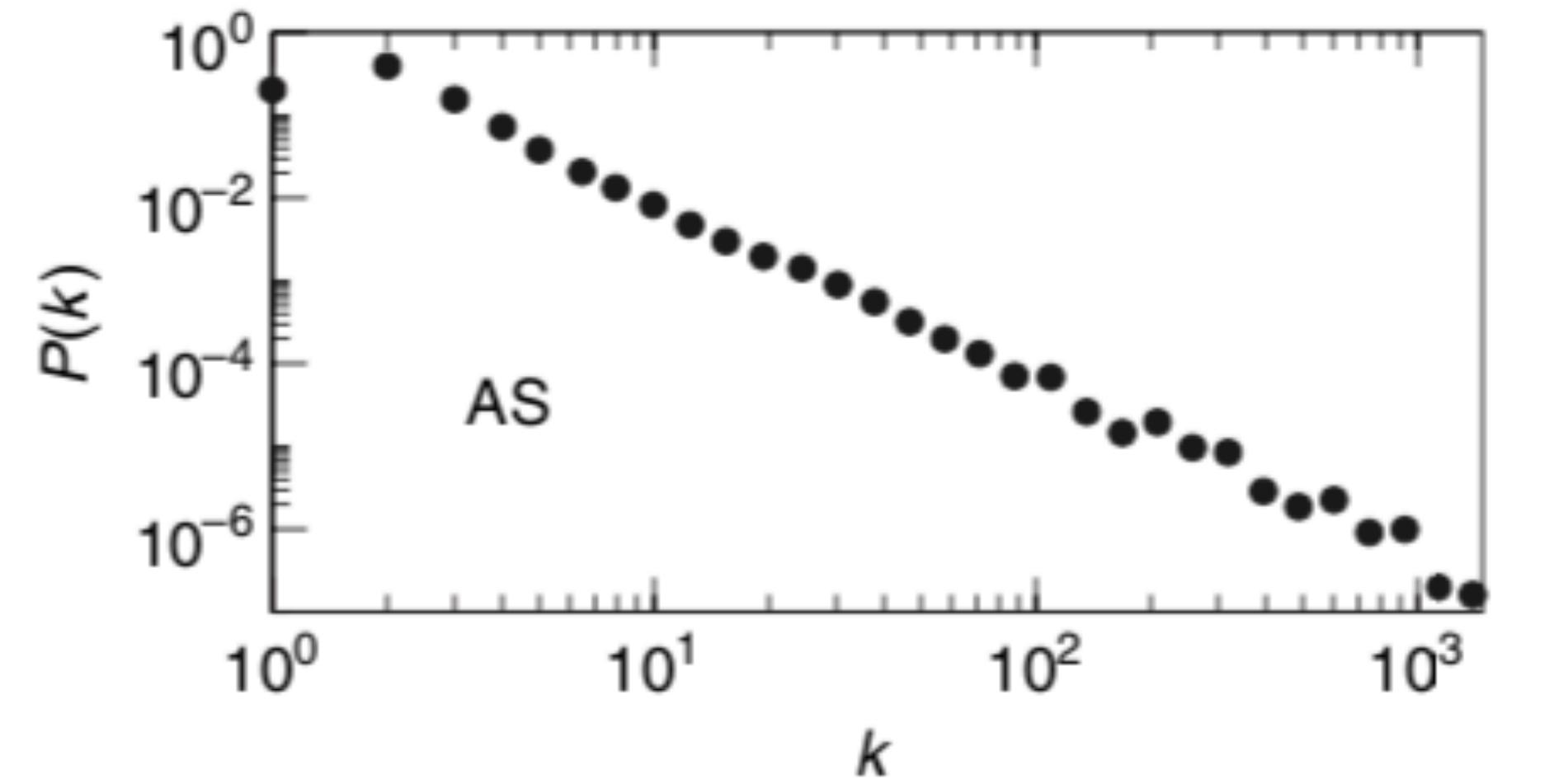
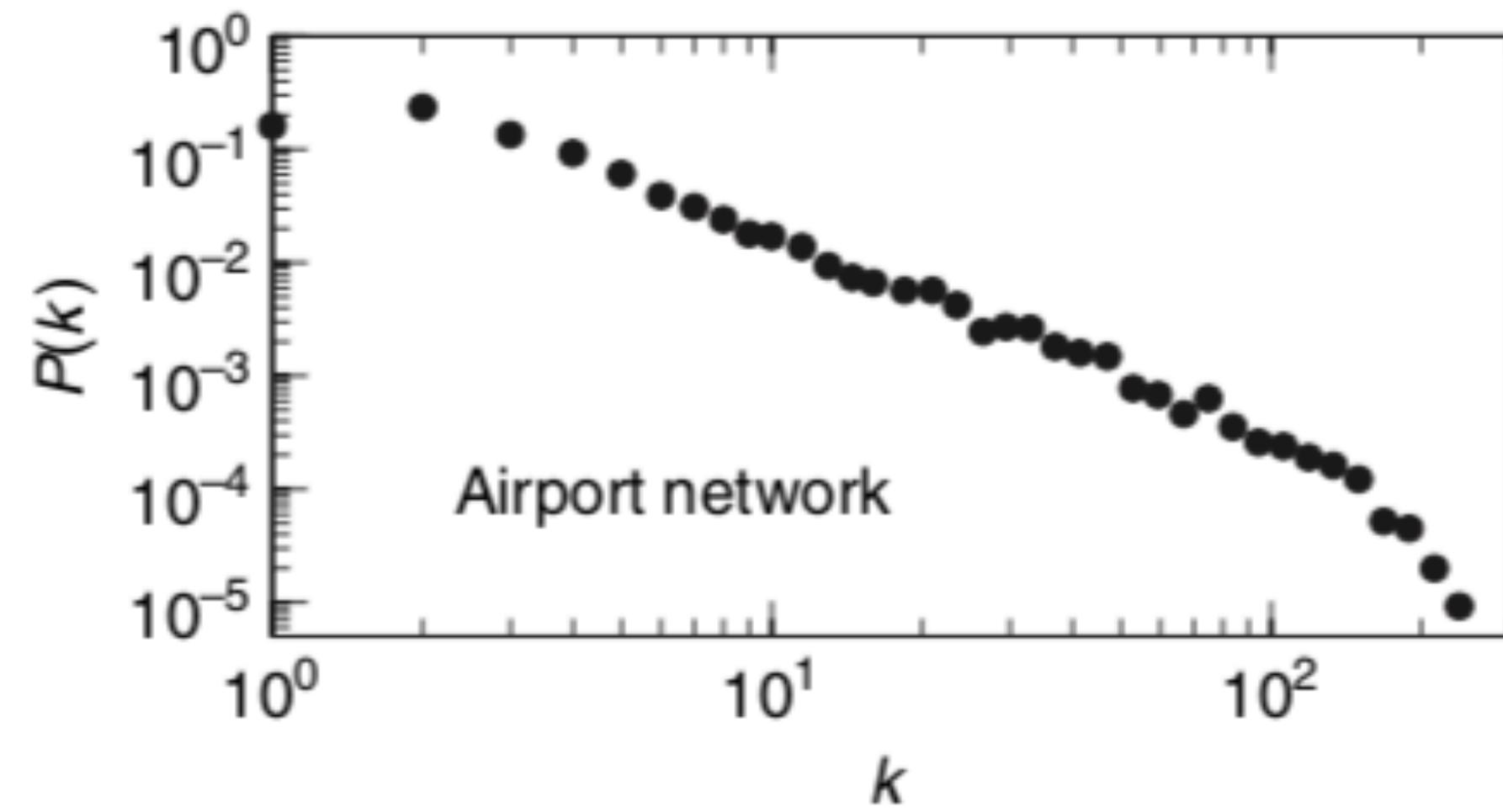
probability that a random chosen node has degree k

$$\langle k \rangle = \sum_k k P(k)$$

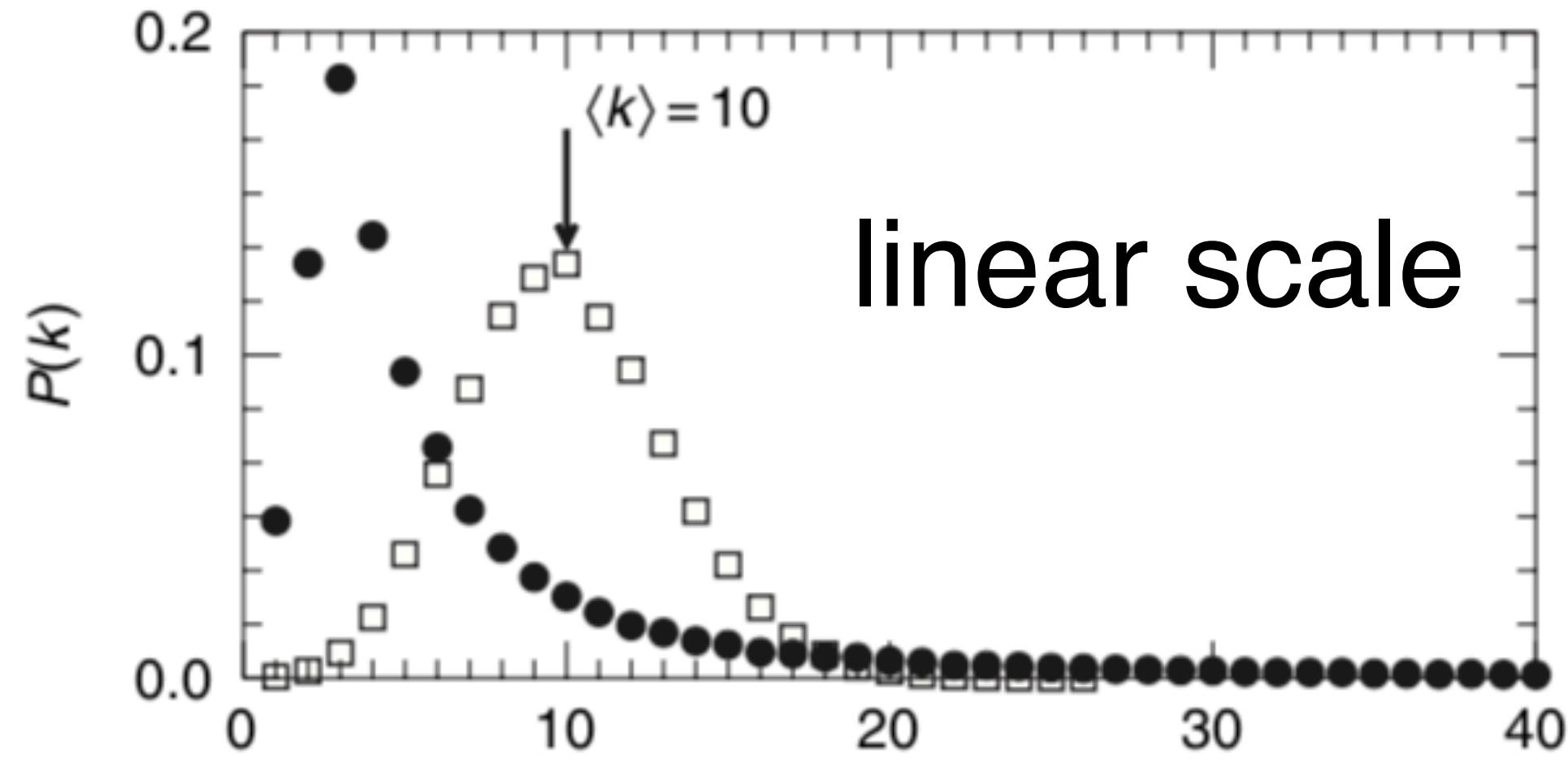
$$\sigma^2 = \langle k^2 \rangle - \langle k \rangle^2$$

$$\langle k^2 \rangle = \sum_k k^2 P(k)$$

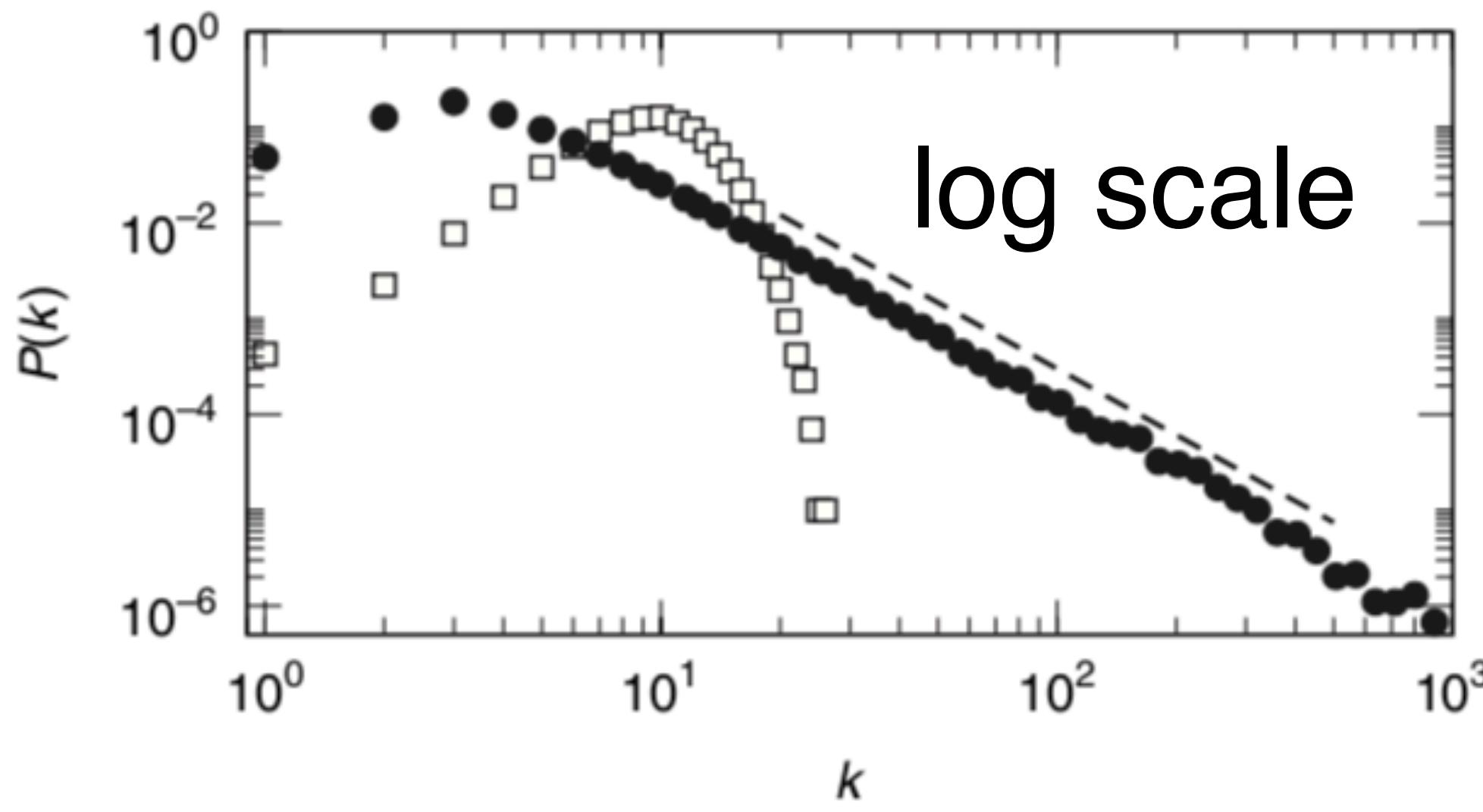
# real world networks



# real world networks



linear scale



log scale

Poisson  
(homogeneous)

**Broad** degree  
distributions

vs

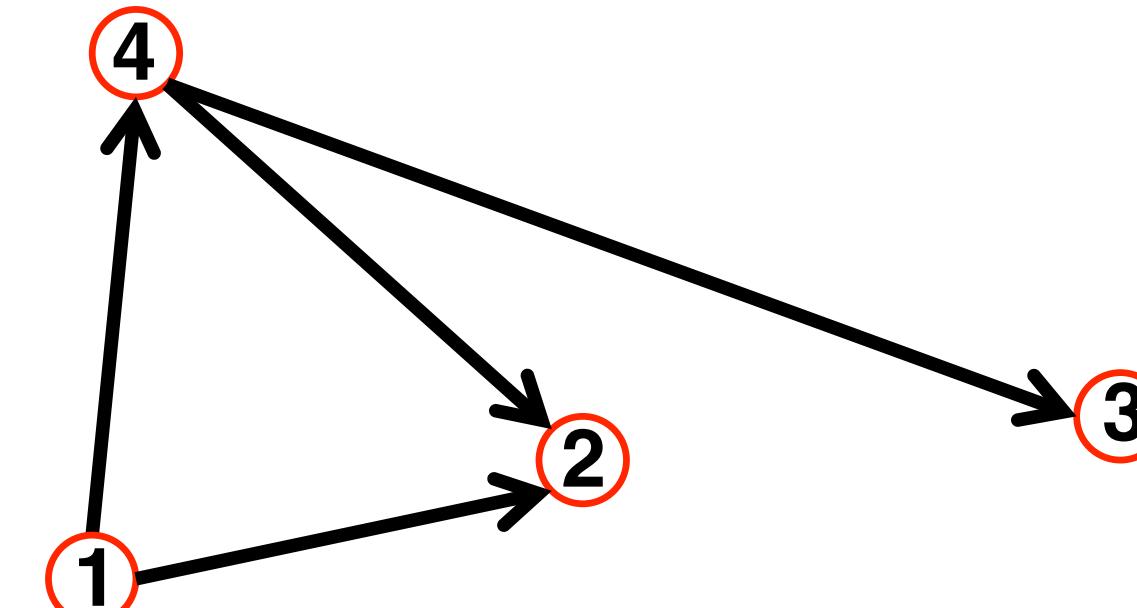
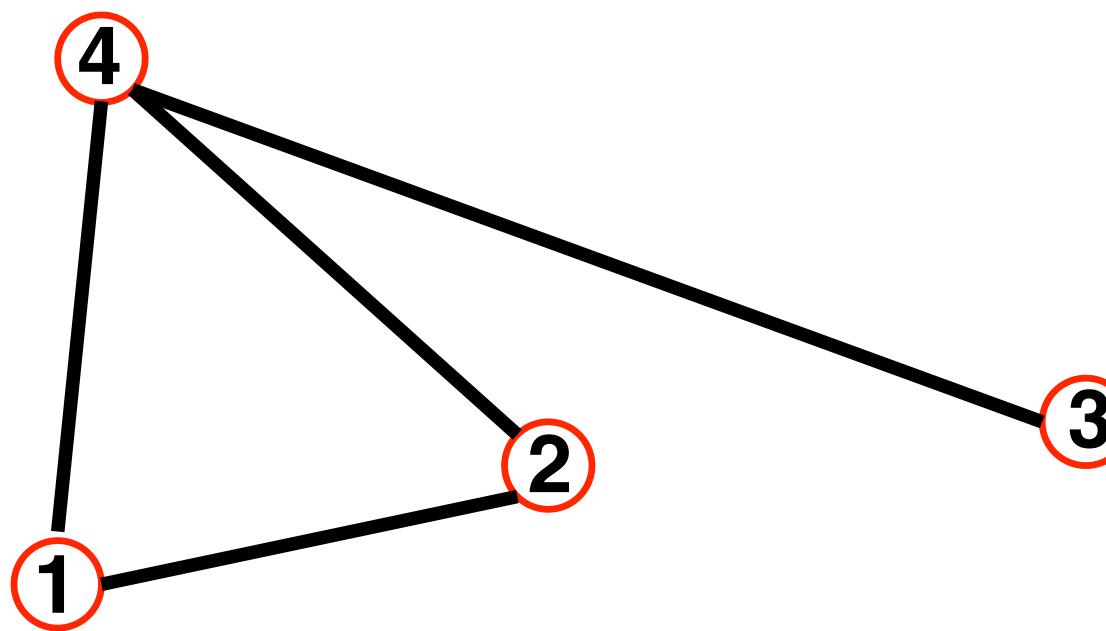
power-law  
(heterogeneous)

Power-law tails

$$P(k) \sim k^{-\gamma}, 2 < \gamma < 3$$

No characteristic scale

# adjacency matrix



$A_{ij}=1$  if there is a link between node  $i$  and  $j$

$A_{ij}=0$  if nodes  $i$  and  $j$  are not connected to each other.

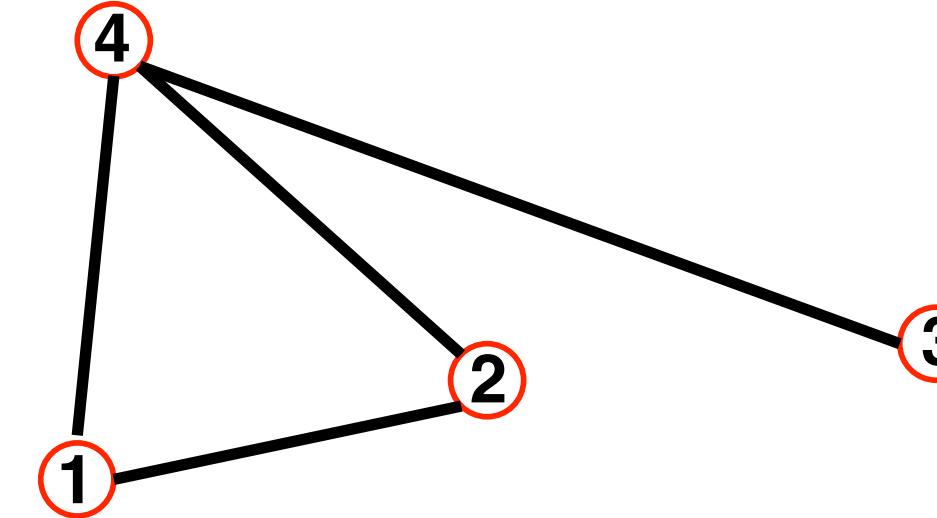
$$A_{ij} = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

$$A_{ij} = \begin{pmatrix} 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$

Note that for a directed graph (right) the matrix is not symmetric.

# adjacency matrix

**Undirected**



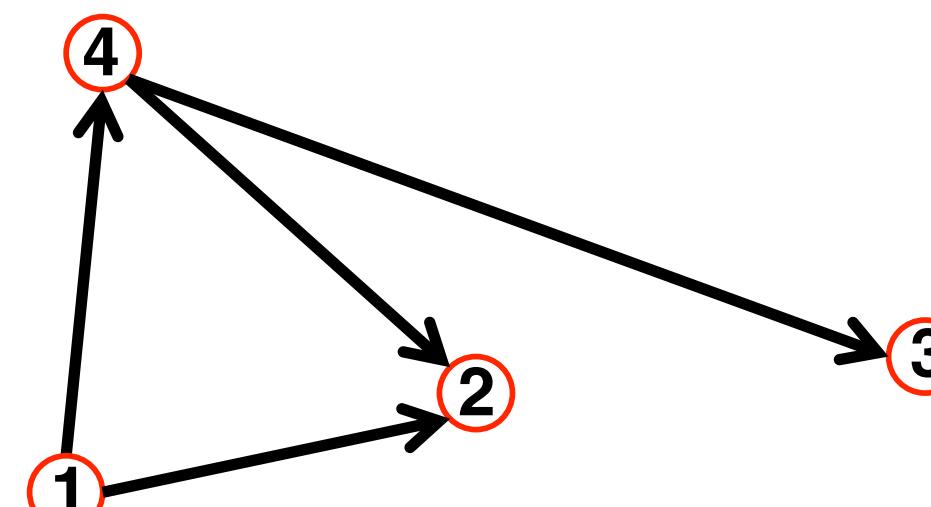
$$A_{ij} = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

$$\begin{aligned} A_{ij} &= A_{ji} \\ A_{ii} &= 0 \end{aligned}$$

$$k_j = \sum_{i=1}^N A_{ij}$$

$$L = \frac{1}{2} \sum_{i=1}^N k_i = \frac{1}{2} \sum_{i,j} A_{ij}$$

**Directed**



$$A_{ij} = \begin{pmatrix} 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$

$$\begin{aligned} A_{ij} &\neq A_{ji} \\ A_{ii} &= 0 \end{aligned}$$

$$k_i^{in} = \sum_{j=1}^N A_{ij}$$

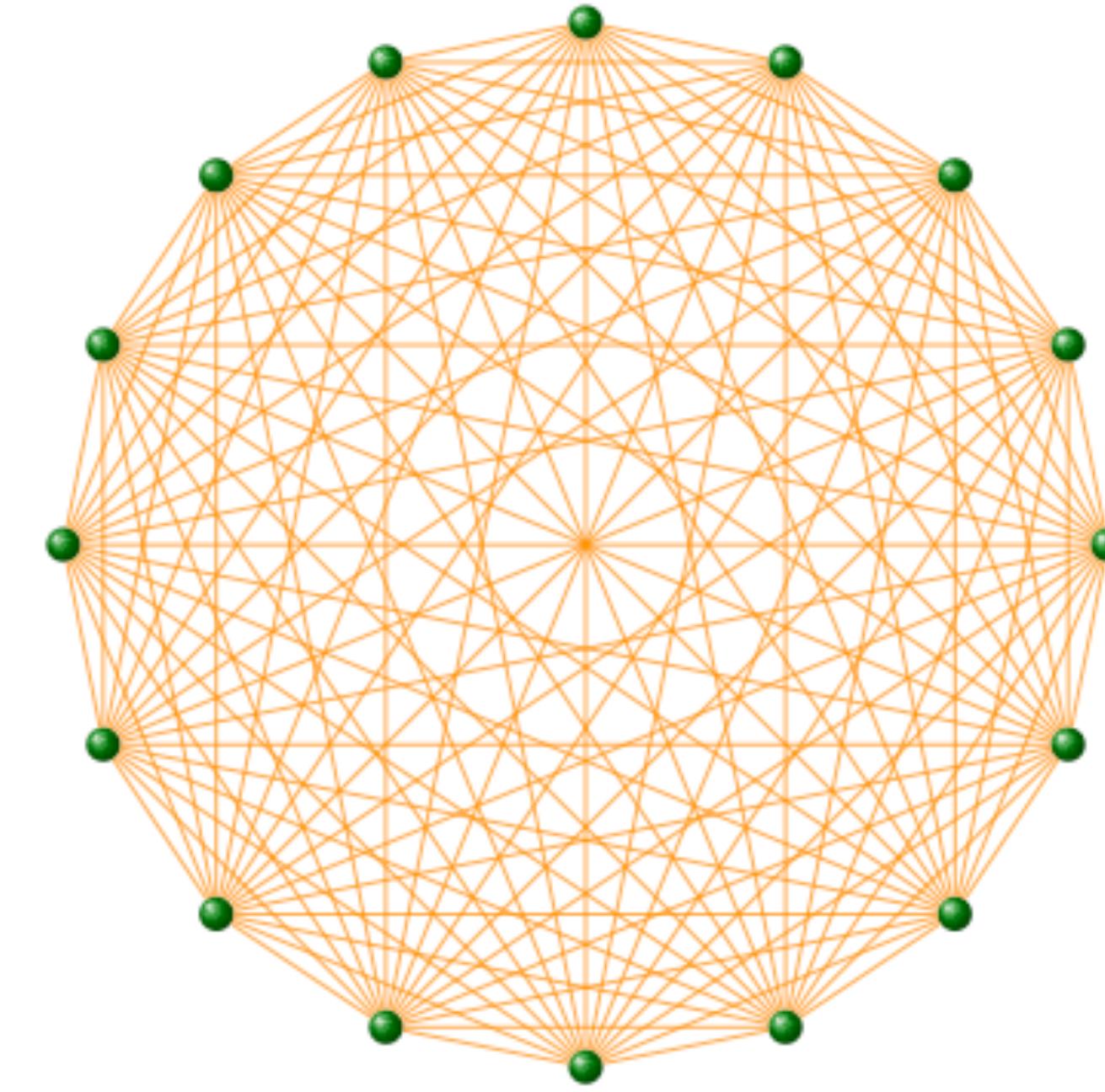
$$k_j^{out} = \sum_{i=1}^N A_{ij}$$

$$L = \sum_{i=1}^N k_i^{in} = \sum_{j=1}^N k_j^{out} = \sum_{i,j} A_{ij}$$



Real networks are sparse!

The maximum number of links a network of  $N$  nodes can have is:  $L_{\max} = \binom{N}{2} = \frac{N(N-1)}{2}$



A graph with degree  $L=L_{\max}$  is called a **complete graph**, and its average degree is  $\langle k \rangle = N-1$

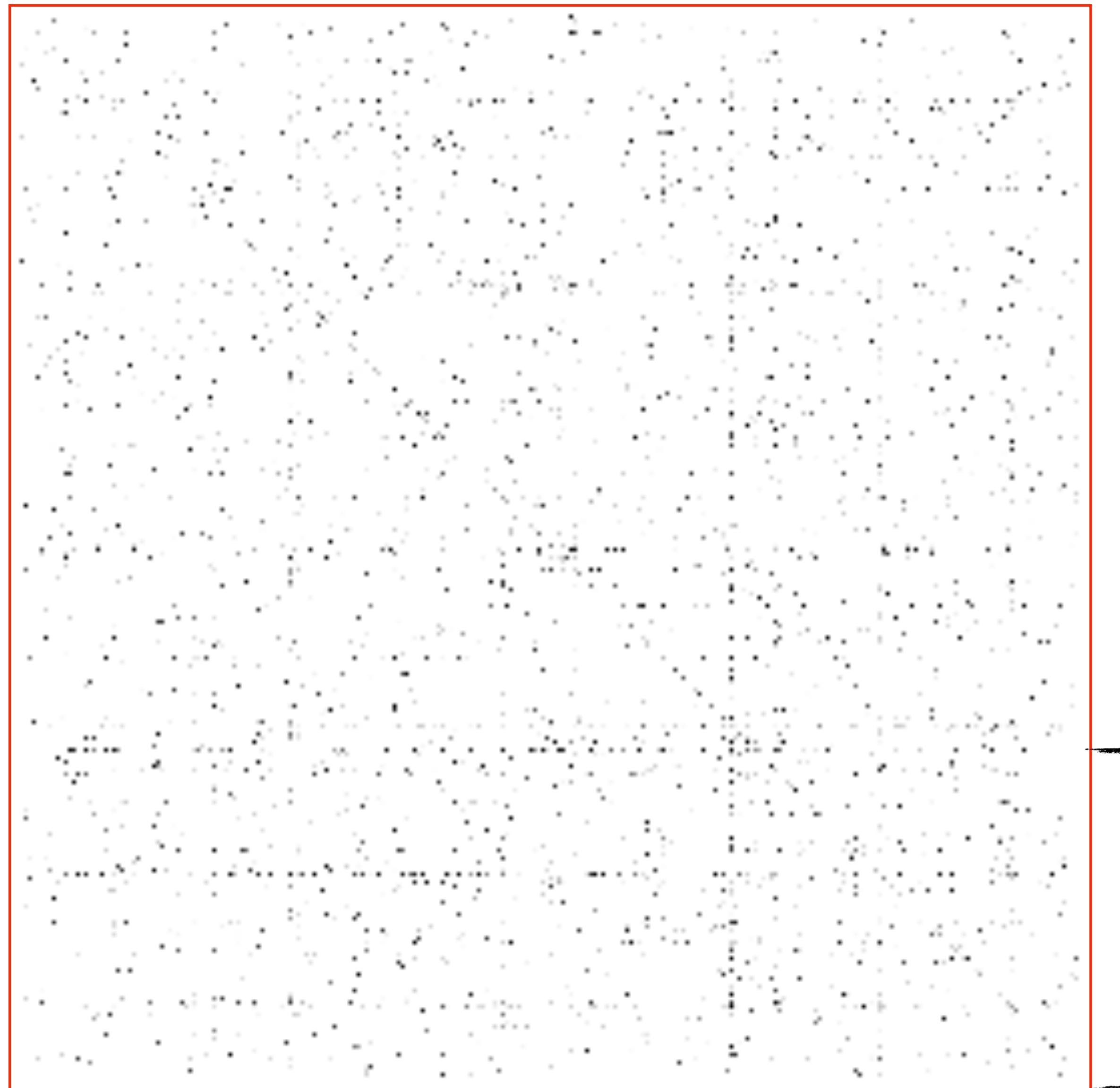
# Most networks observed in real systems are **sparse**

$$L \ll L_{max} \quad \langle k \rangle \ll N - 1$$

WWW (ND Sample):	$N=325,729;$	$L=1.4 \cdot 10^6$	$L_{max}=10^{12}$	$\langle k \rangle=4.51$
Protein ( <i>S. Cerevisiae</i> ):	$N=1,870;$	$L=4,470$	$L_{max}=10^7$	$\langle k \rangle=2.39$
Coauthorship (Math):	$N=70,975;$	$L=2 \cdot 10^5$	$L_{max}=3 \cdot 10^{10}$	$\langle k \rangle=3.9$
Movie Actors:	$N=212,250;$	$L=6 \cdot 10^6$	$L_{max}=1.8 \cdot 10^{13}$	$\langle k \rangle=28.78$

(Source: Albert, Barabasi, RMP2002)

The adjacency matrix of the yeast protein-protein interaction network, consisting of 2,018 nodes, each representing a yeast protein.



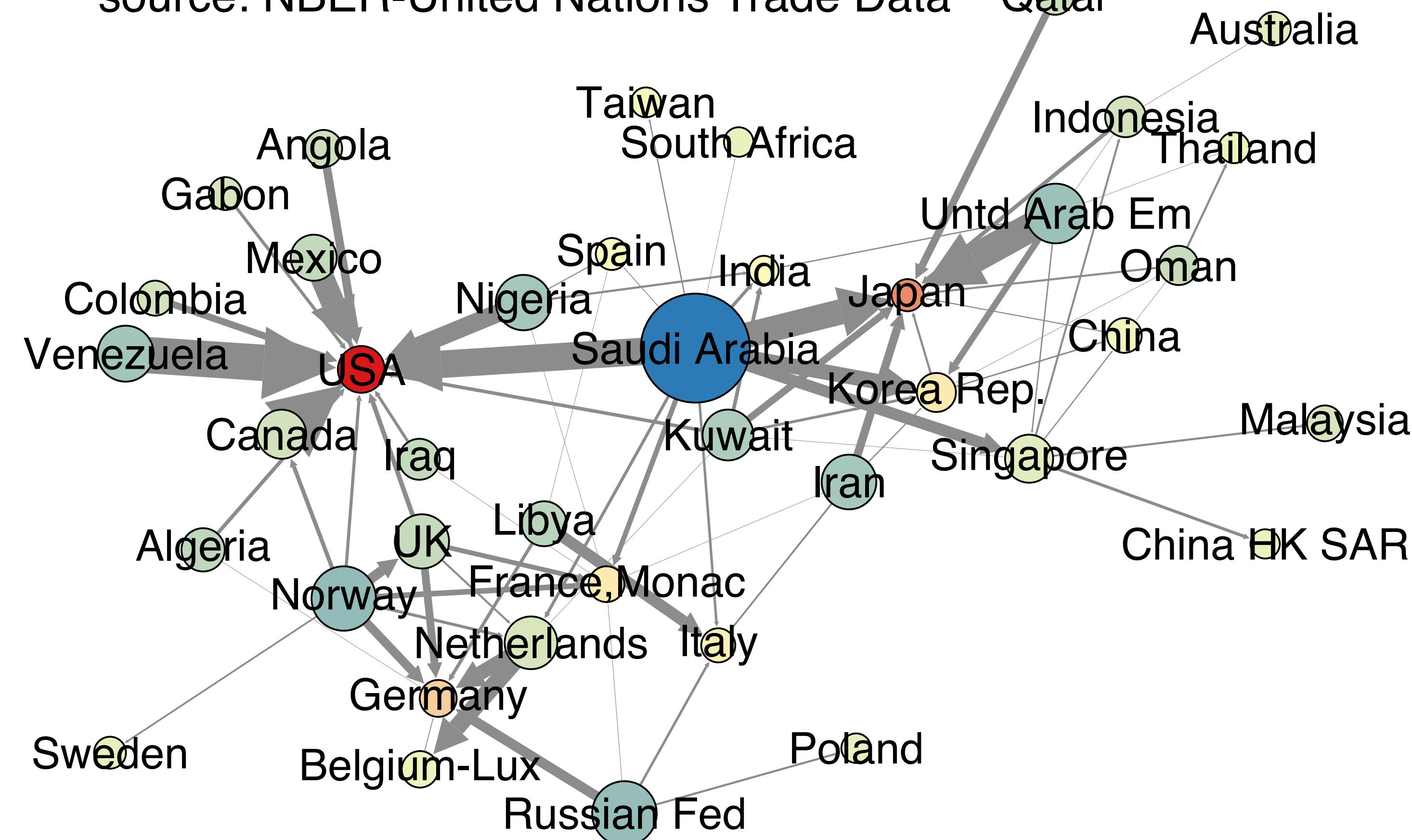
The adjacency matrix is not efficient to store the network

# weighted networks

$$A_{ij} = w_{ij}$$

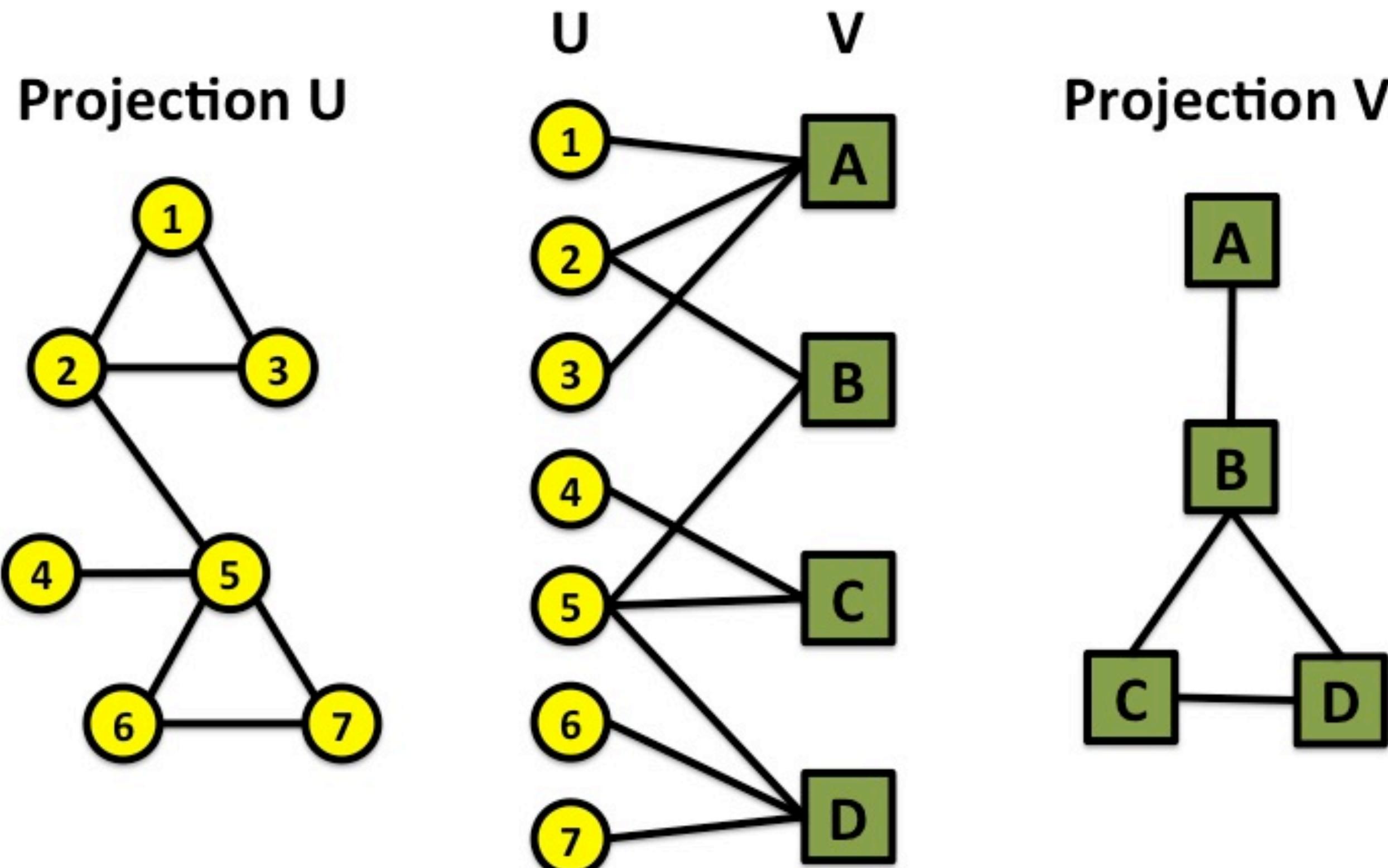
$$s_i = \sum_j w_{ij}$$

trade in petroleum and petroleum products, 1998,  
source: NBER-United Nations Trade Data    Qatar



# bipartite networks

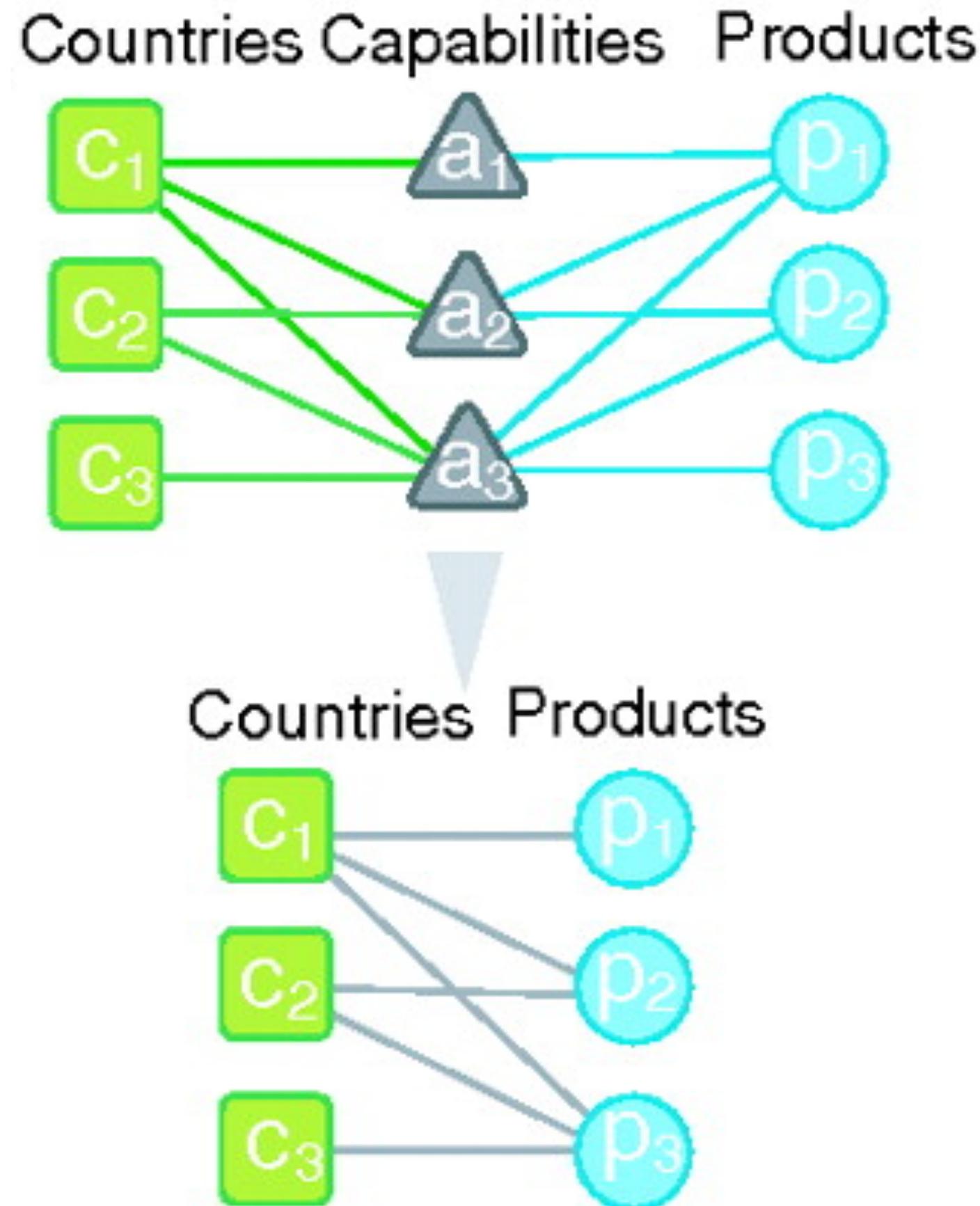
**bipartite graph** (or **bigraph**) is a [graph](#) whose nodes can be divided into two [disjoint sets](#)  $U$  and  $V$  such that every link connects a node in  $U$  to one in  $V$ ; that is,  $U$  and  $V$  are [independent sets](#).



## Examples

- actor network
- collaboration network
- host-pathogen networks

# bipartite networks

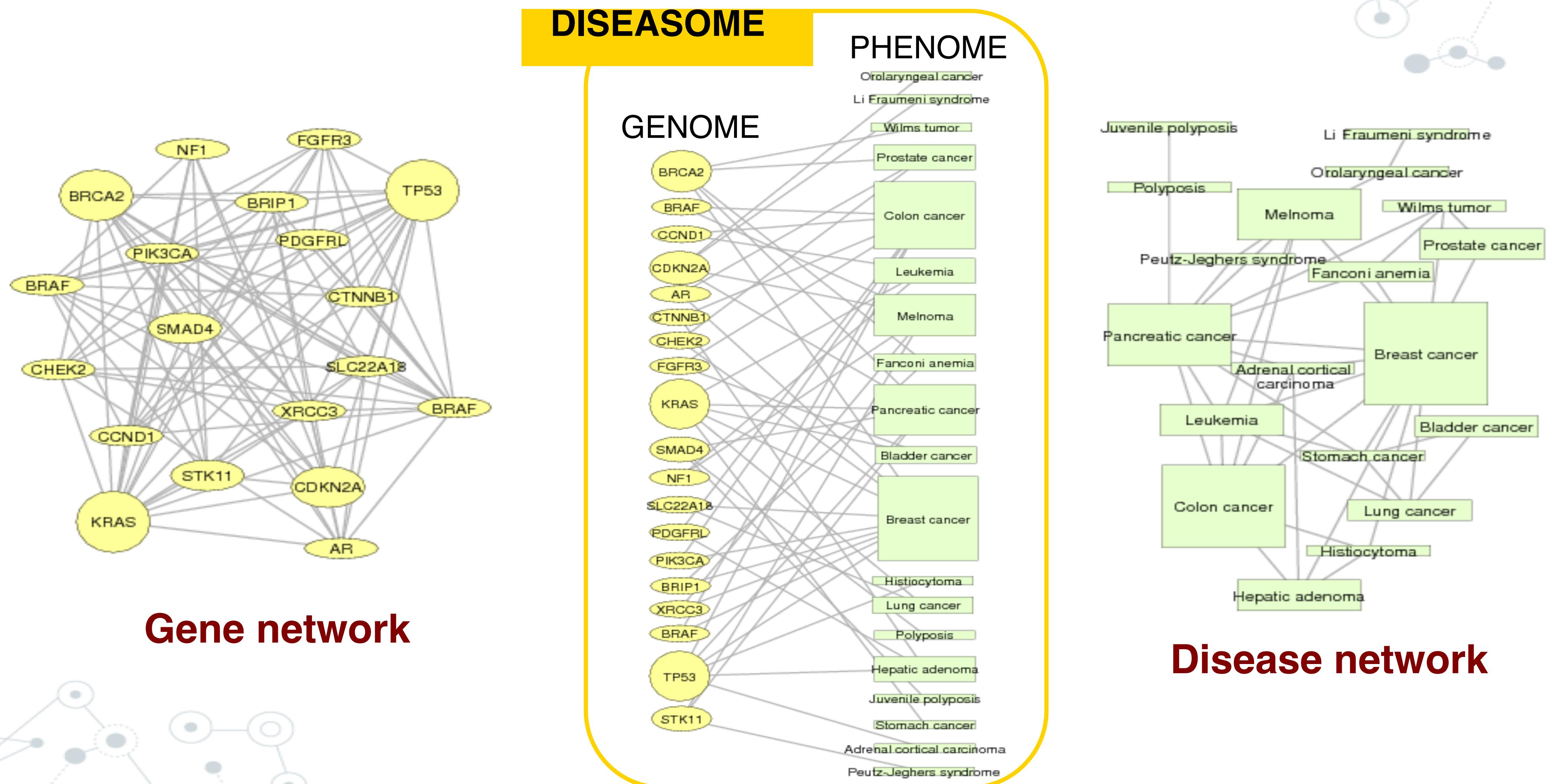


**The Atlas of  
Economic Complexity**

C. Hidalgo

<http://atlas.cid.harvard.edu/>

# bipartite networks



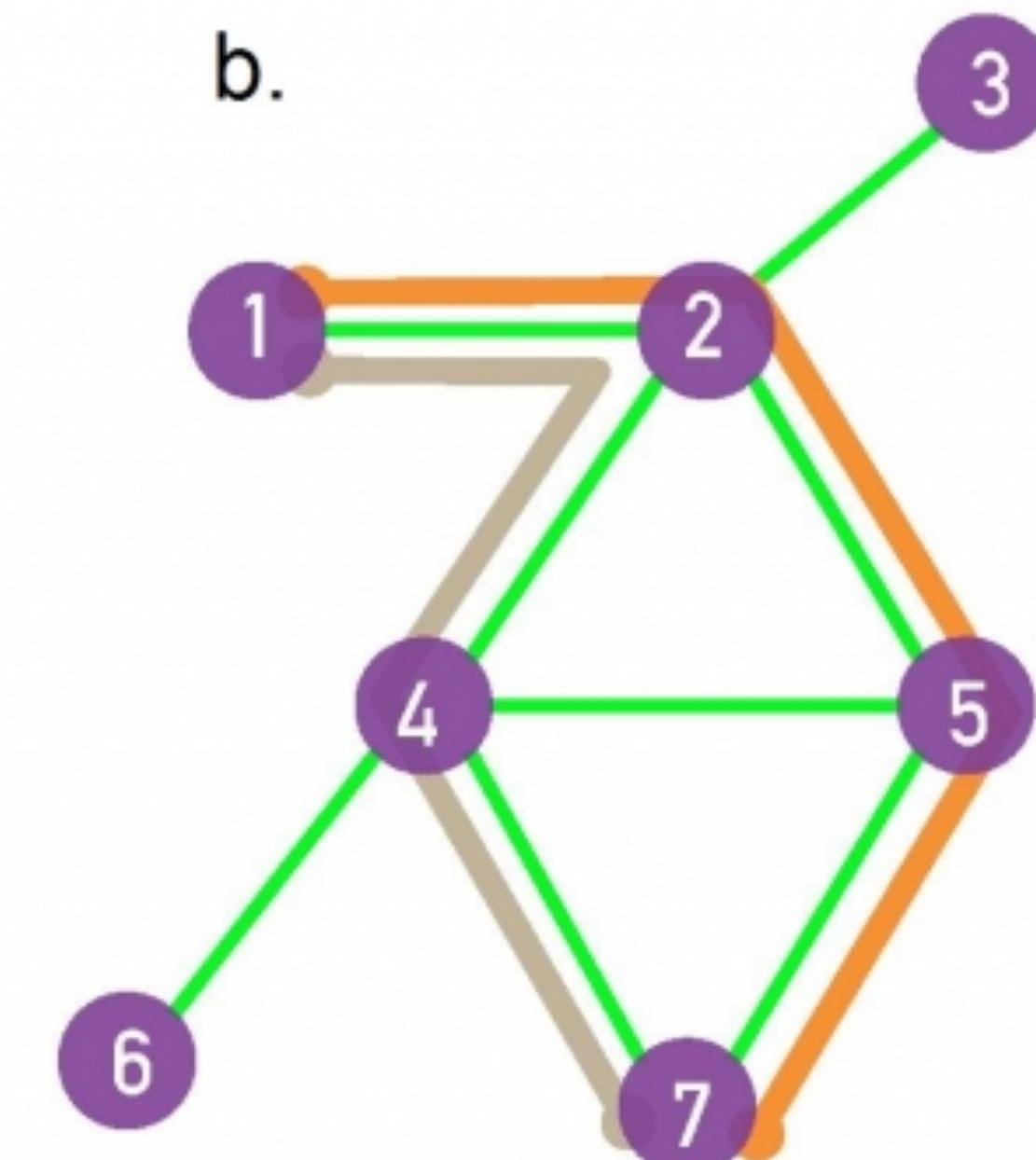
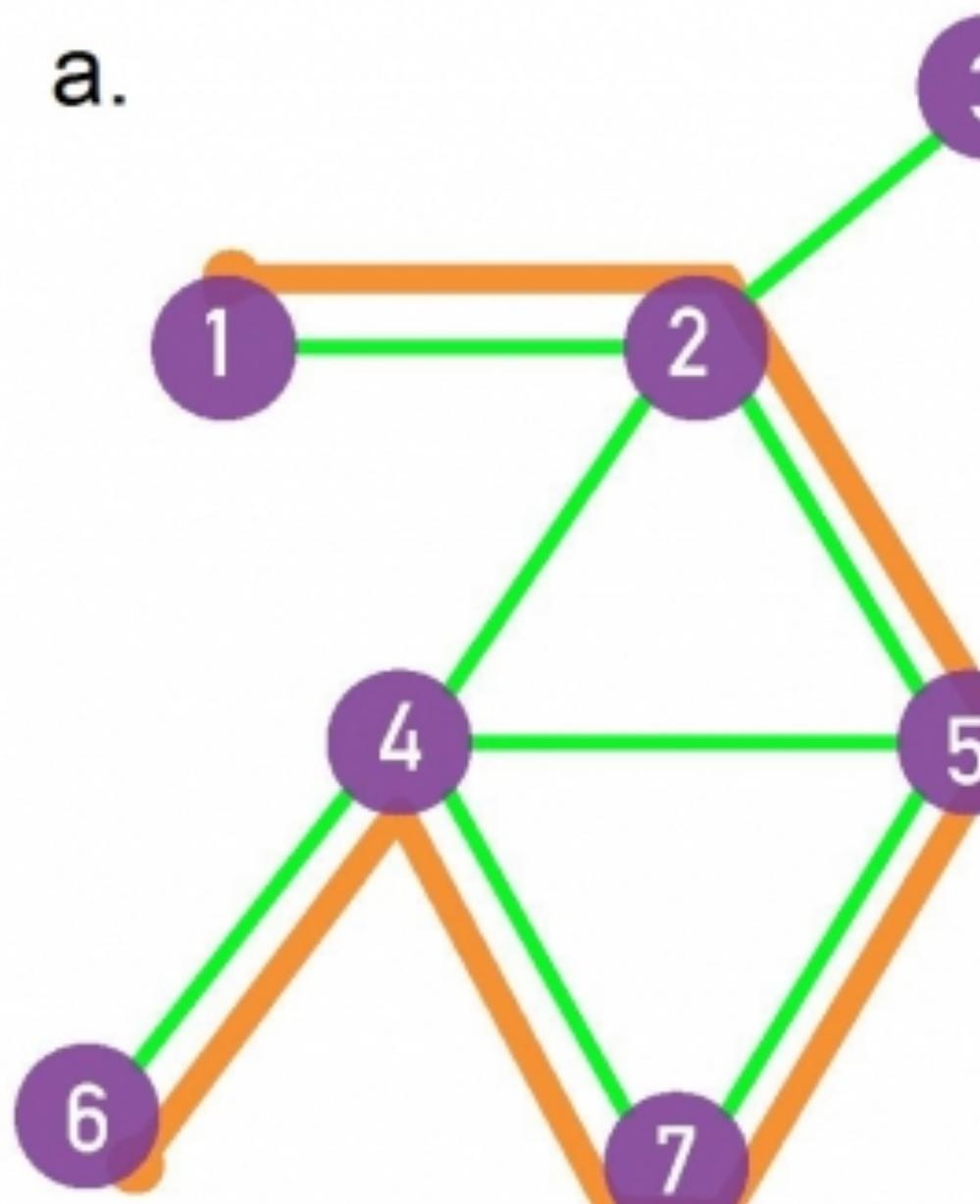
# paths

A *path* is a sequence of nodes in which each node is adjacent to the next one

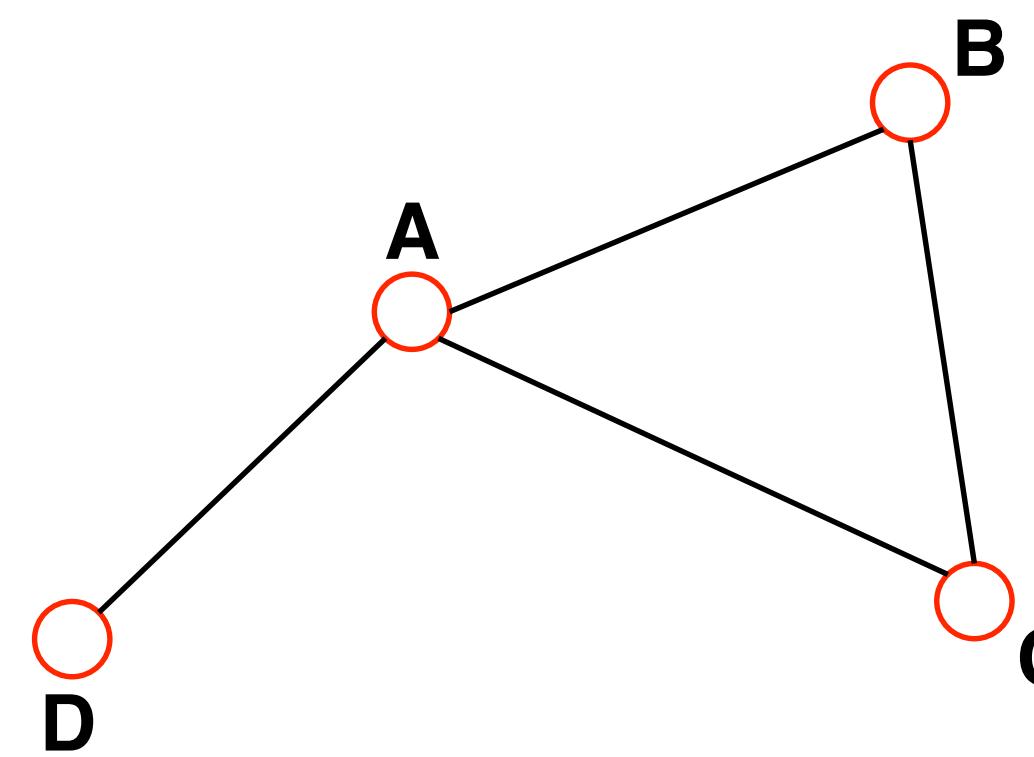
$P_{i_0, i_n}$  of length  $n$  between nodes  $i_0$  and  $i_n$  is an ordered collection of  $n+1$  nodes and  $n$  links

$$P_n = \{i_0, i_1, i_2, \dots, i_n\} \quad P_n = \{(i_0, i_1), (i_1, i_2), (i_2, i_3), \dots, (i_{n-1}, i_n)\}$$

- (a) path of length 5
- (b) two paths of equal length

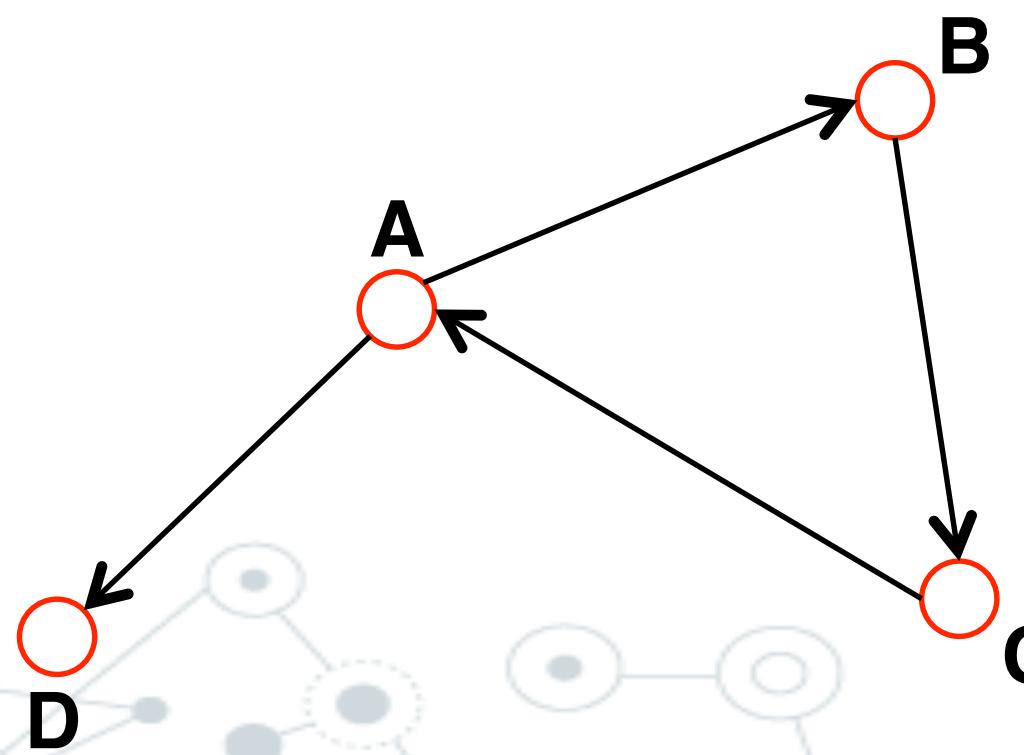


# distance



The *distance (shortest path, geodesic path)* between two nodes is defined as the number of edges along the shortest path connecting them. The *diameter of a graph* is the length of the longest geodesic path between any pair of vertices in the network for which a path actually exists.

\*If the two nodes are disconnected, the distance is infinity.



In *directed graphs* each path needs to follow the direction of the arrows.

Thus in a digraph the distance from node A to B (on an AB path) is generally different from the distance from node B to A (on a BCA path).

# paths

**$N_{ij}$ , number of paths between any two nodes  $i$  and  $j$ :**

**Length  $n=1$ :** If there is a link between  $i$  and  $j$ , then  $A_{ij}=1$  and  $A_{ij}=0$  otherwise.

**Length  $n=2$ :** If there is a path of length two between  $i$  and  $j$ , then  $A_{ik}A_{kj}=1$ , and  $A_{ik}A_{kj}=0$  otherwise.  
The number of paths of length 2:

$$N_{ij}^{(2)} = \sum_{k=1}^N A_{ik}A_{kj} = [A^2]_{ij}$$

**Length  $n$ :** In general, if there is a path of length  $n$  between  $i$  and  $j$ , then  $A_{ik}\dots A_{lj}=1$  and  $A_{ik}\dots A_{lj}=0$  otherwise.

The number of paths of length  $n$  between  $i$  and  $j$  is\*

$$N_{ij}^{(n)} = [A^n]_{ij}$$

# paths

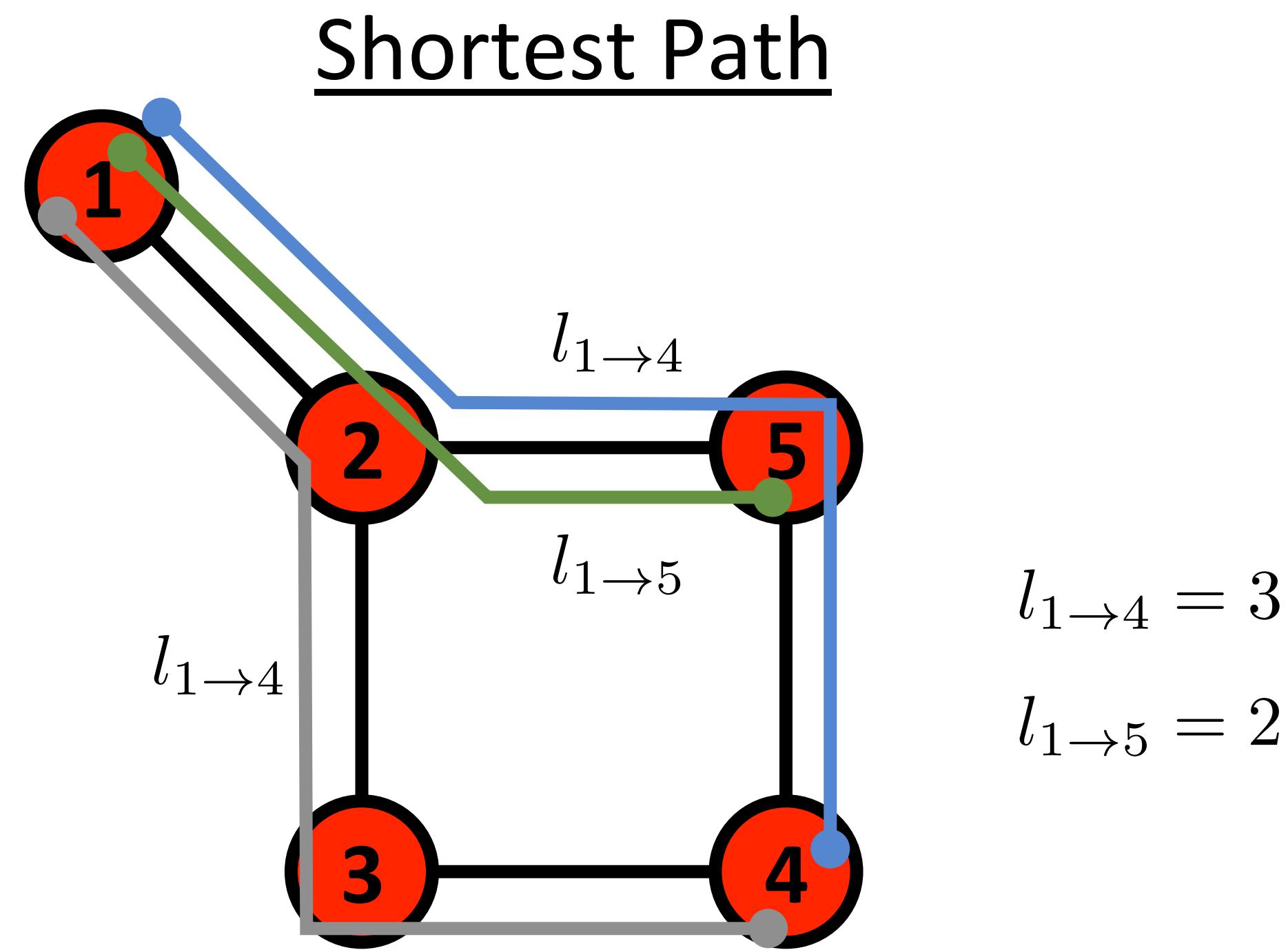
*Diameter:*  $d_{max}$  the maximum distance between any pair of nodes in the graph.

*Average path length/distance,*  $\langle d \rangle$ , for a **connected graph**: average distance between all pairs of nodes in the network, where  $d_{ij}$  is the distance from node  $i$  to node  $j$

$$\langle d \rangle = \frac{1}{2L_{max}} \sum_{i,j \neq i} d_{ij}$$

In an *undirected graph*  $d_{ij} = d_{ji}$ , so we only need to count them once:  $\langle d \rangle = \frac{1}{L_{max}} \sum_{i,j > i} d_{ij}$

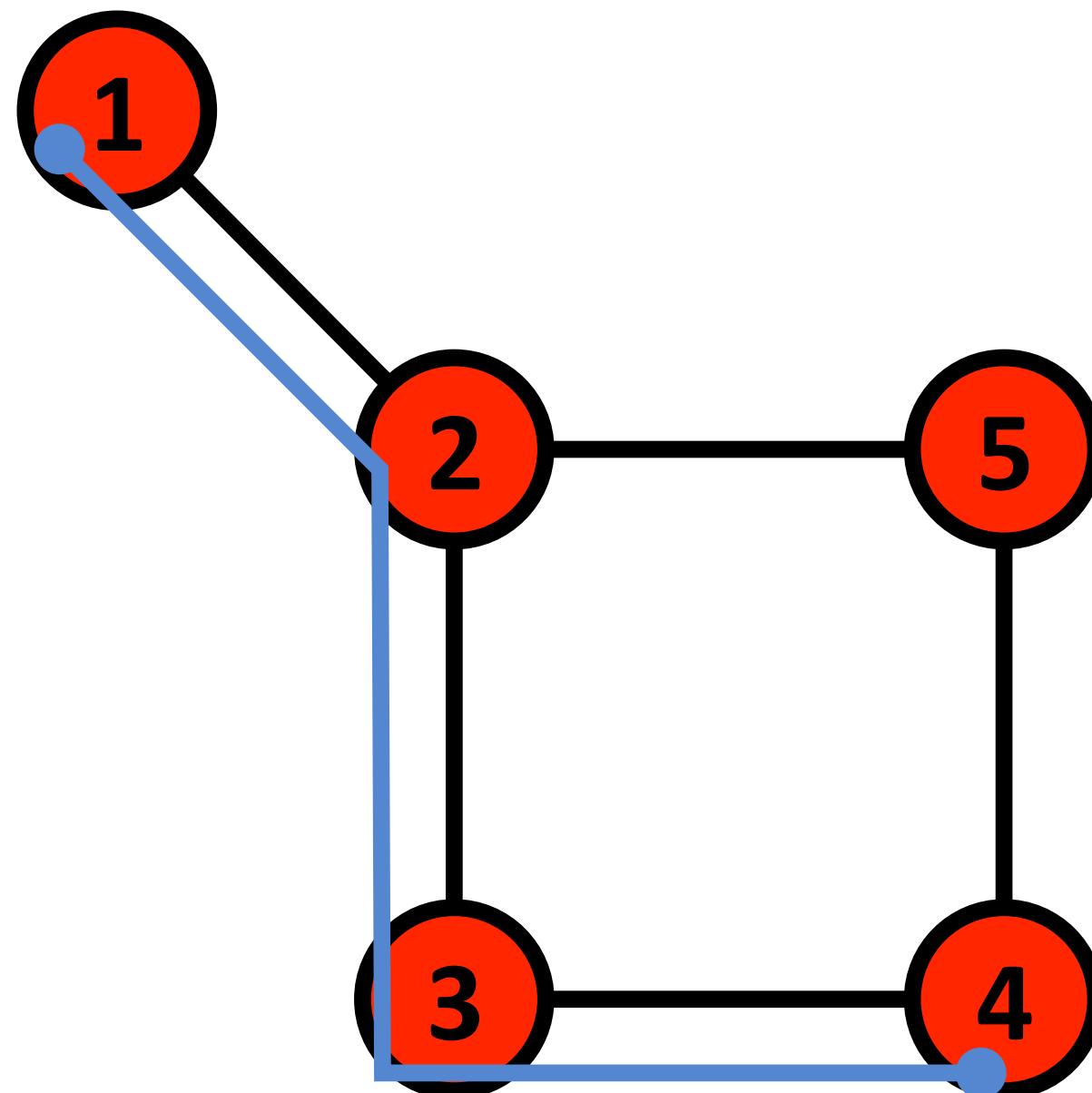
# paths: summary



The path with the shortest length between two nodes (distance).

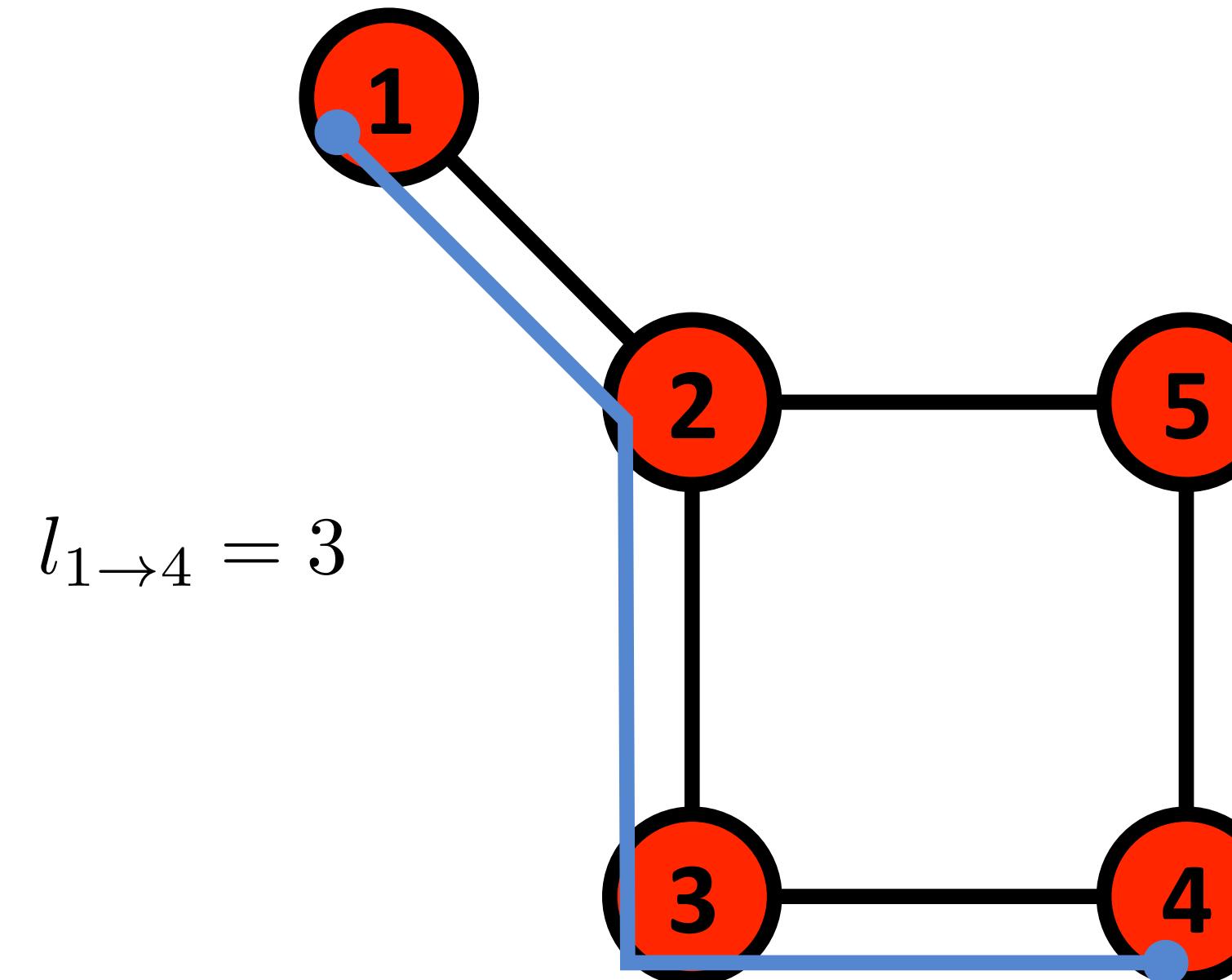
# paths: summary

## Diameter



The longest shortest path in  
a graph

## Average Path Length



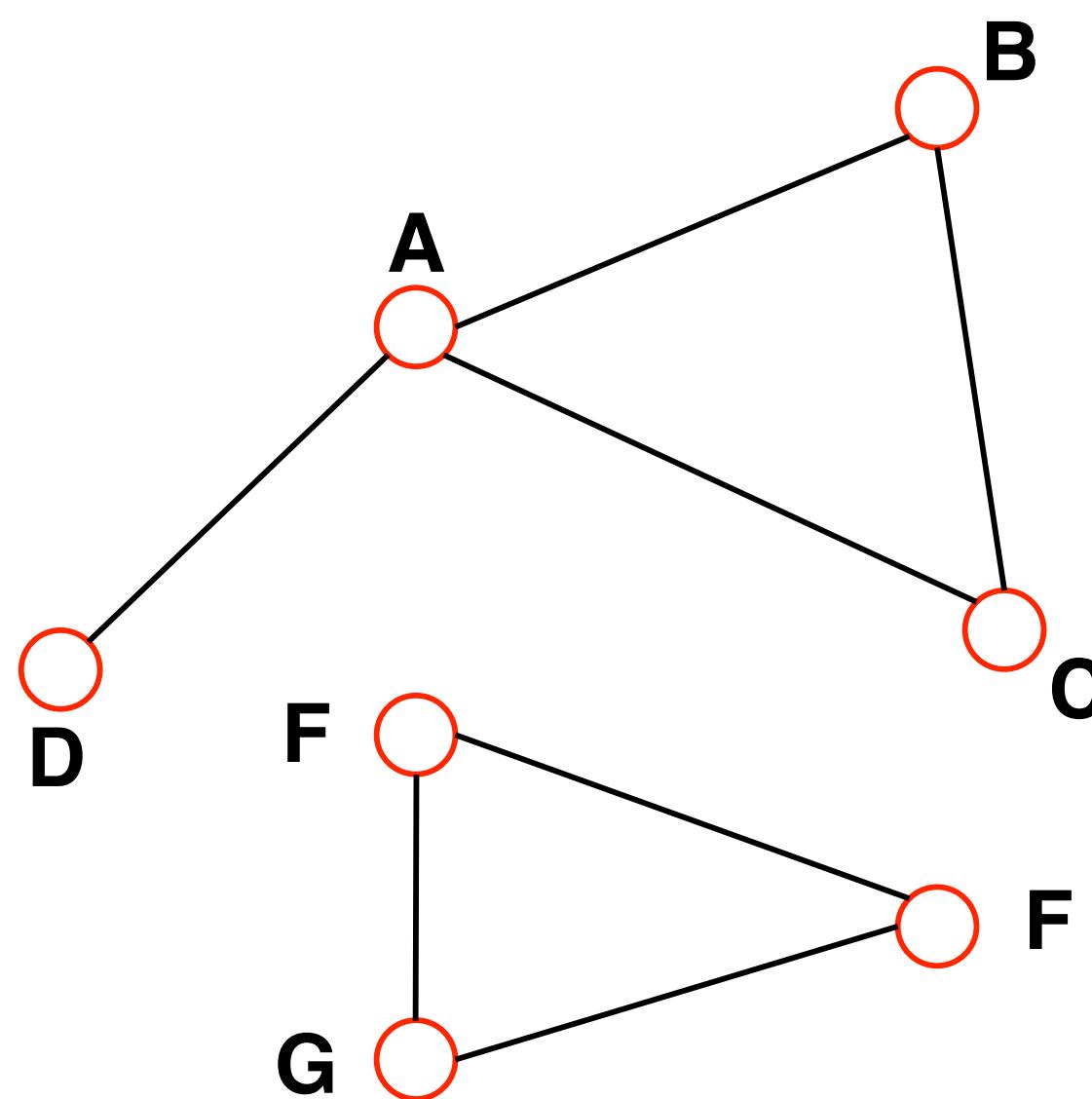
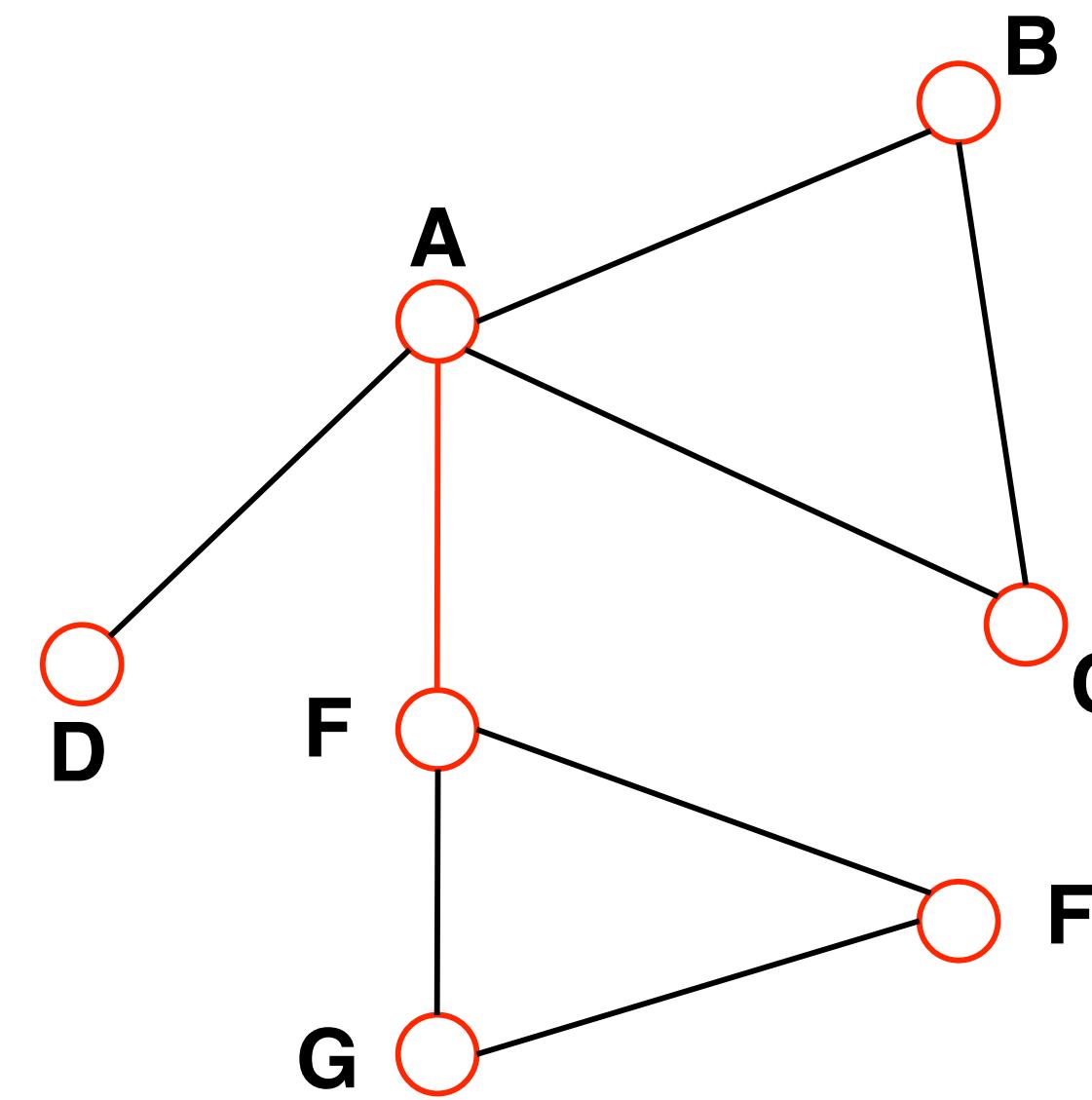
$$l_{1 \rightarrow 4} = 3$$

$$(l_{1 \rightarrow 2} + l_{1 \rightarrow 3} + l_{1 \rightarrow 4} + l_{1 \rightarrow 5} + l_{2 \rightarrow 3} + l_{2 \rightarrow 4} + l_{2 \rightarrow 5} + l_{3 \rightarrow 4} + l_{3 \rightarrow 5} + l_{4 \rightarrow 5}) / 10 = 1.6$$

The average of the shortest paths for  
all pairs of nodes.

# connectivity

Connected (undirected) graph: any two vertices can be joined by a path.  
A disconnected graph is made up by two or more connected components.



Largest Component:  
**Giant Component**

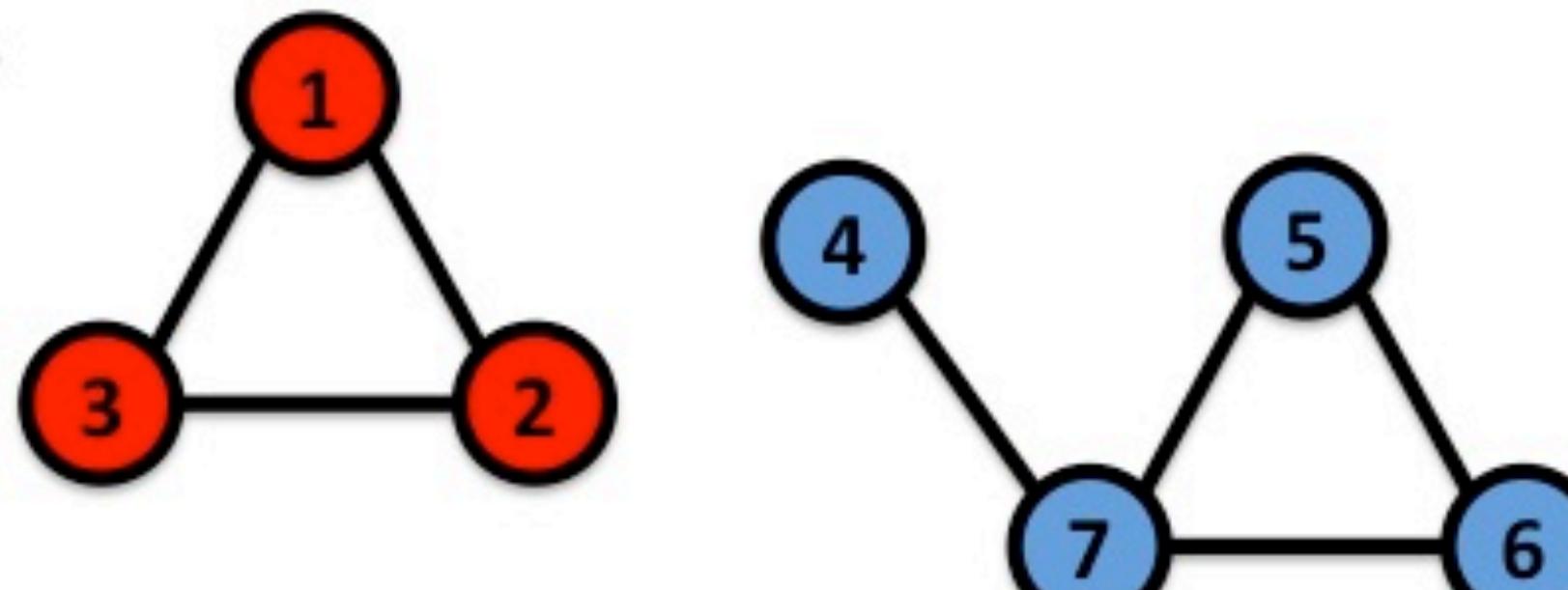
The rest: **Isolates**

Bridge: if we erase it, the graph becomes disconnected.

# connectivity

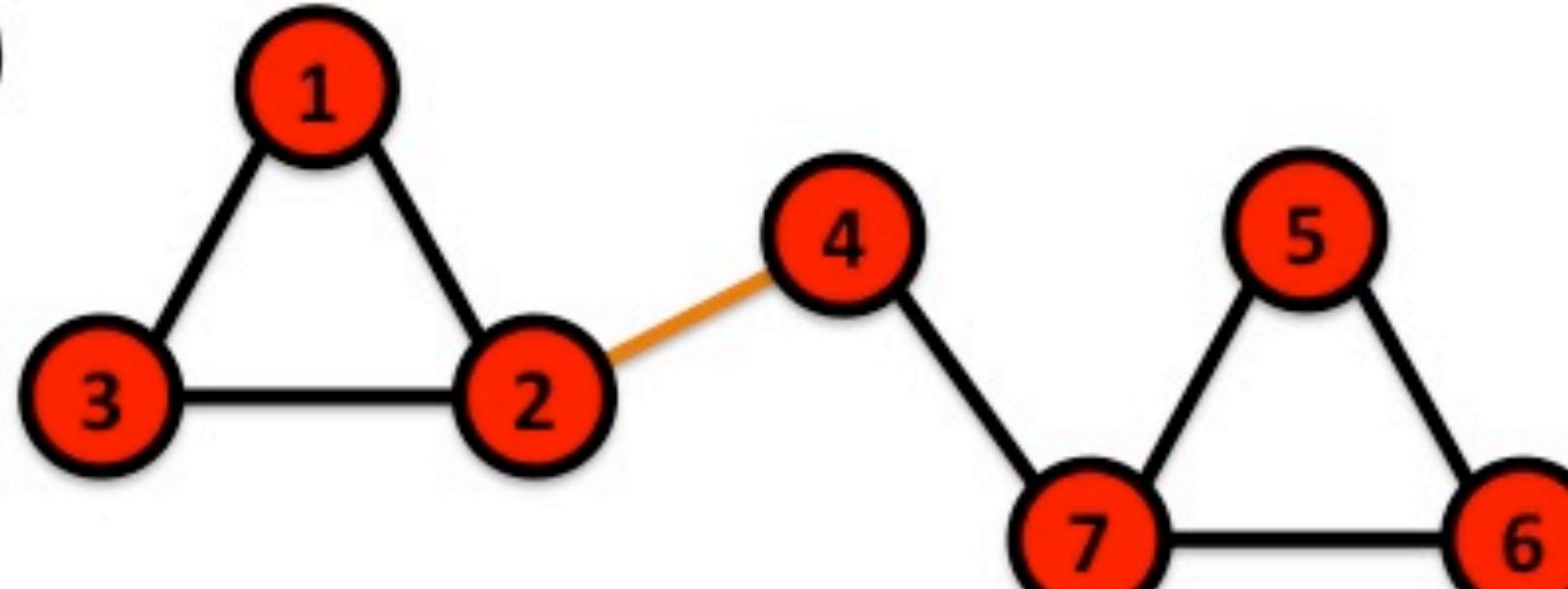
The adjacency matrix of a network with several components can be written in a block-diagonal form, so that nonzero elements are confined to squares, with all other elements being zero:

(a)



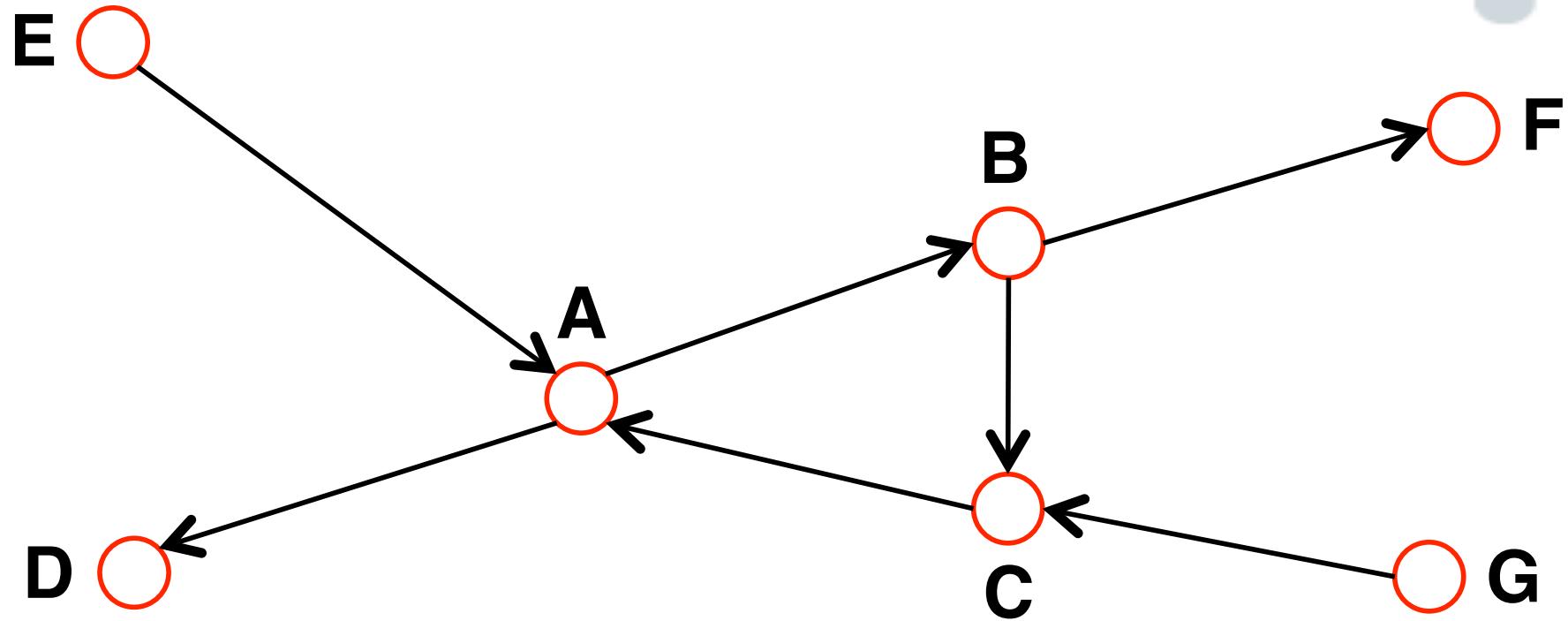
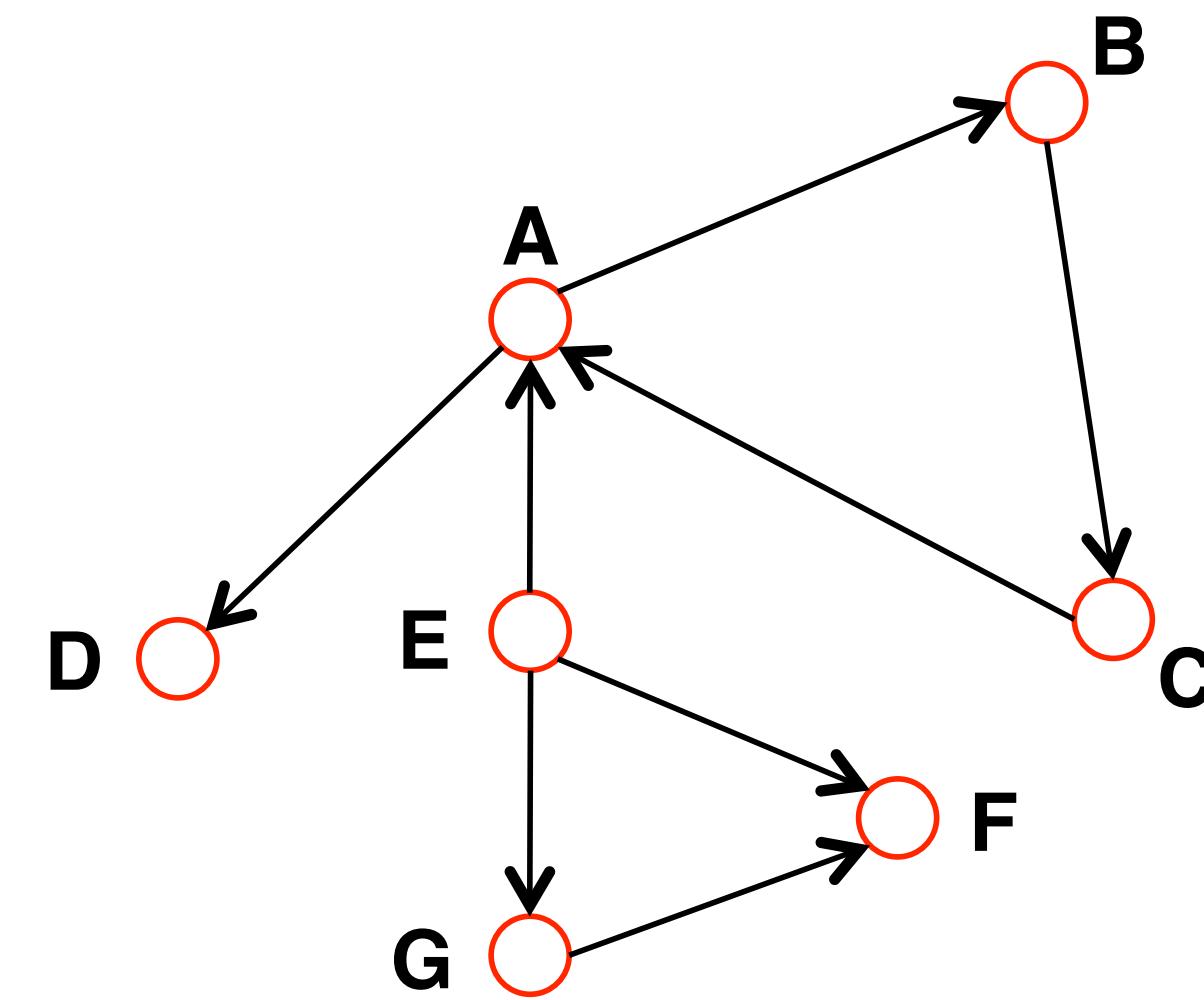
$$\left( \begin{array}{ccc|ccccc} 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 \end{array} \right)$$

(b)



$$\left( \begin{array}{ccc|ccccc} 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ \hline 0 & 1 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 \end{array} \right)$$

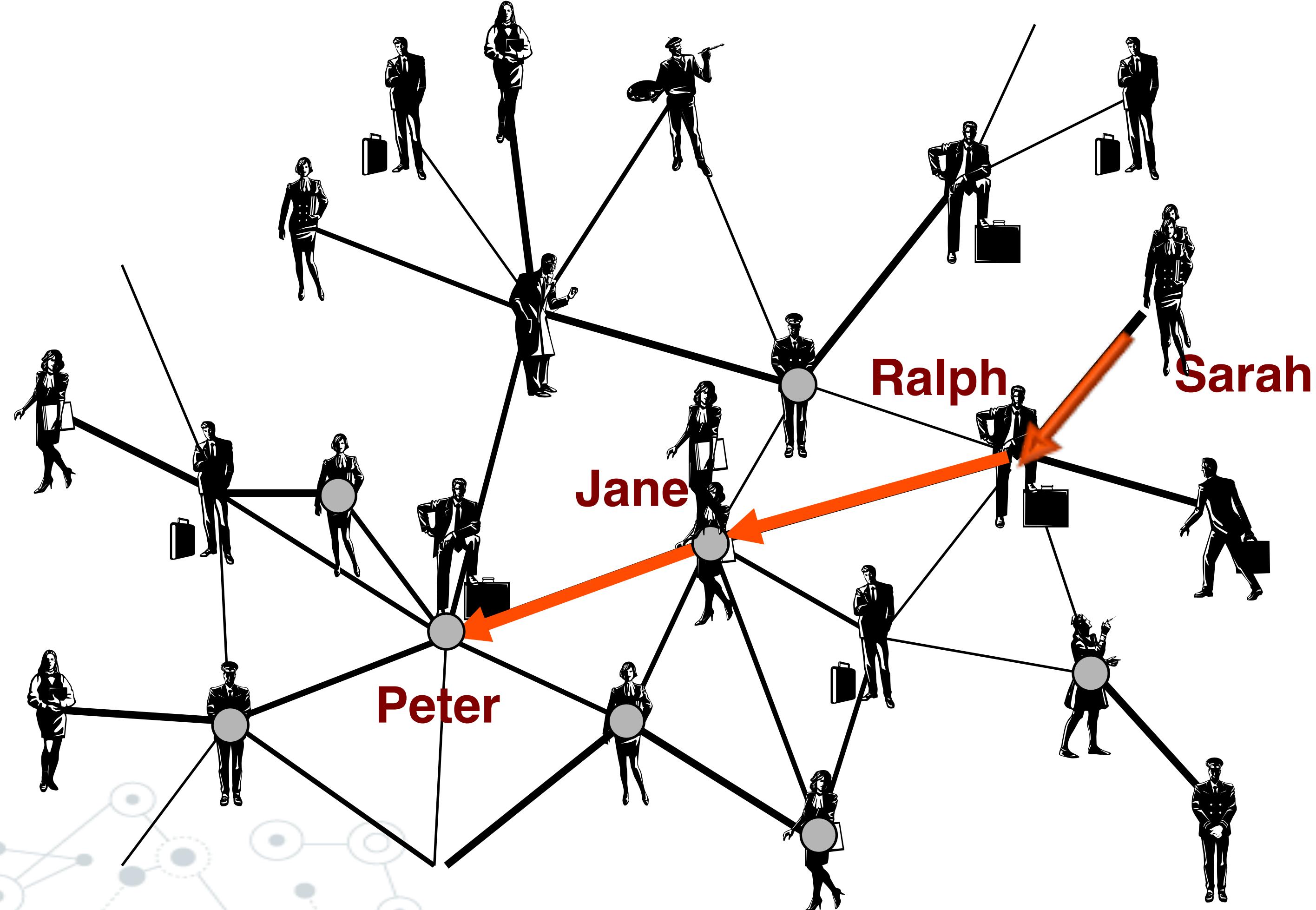
# connectivity in directed graphs



**Strongly connected directed graph:** has a path from each node to every other node and vice versa

**Weakly connected directed graph:** it is connected if we disregard the edge directions.

# real world networks



Small world effect

Frigyes Karinthy, 1929  
Stanley Milgram, 1967

# six degrees of separation

Stanley Milgram (1967)

Two targets in Boston and  
Sharon, MA.

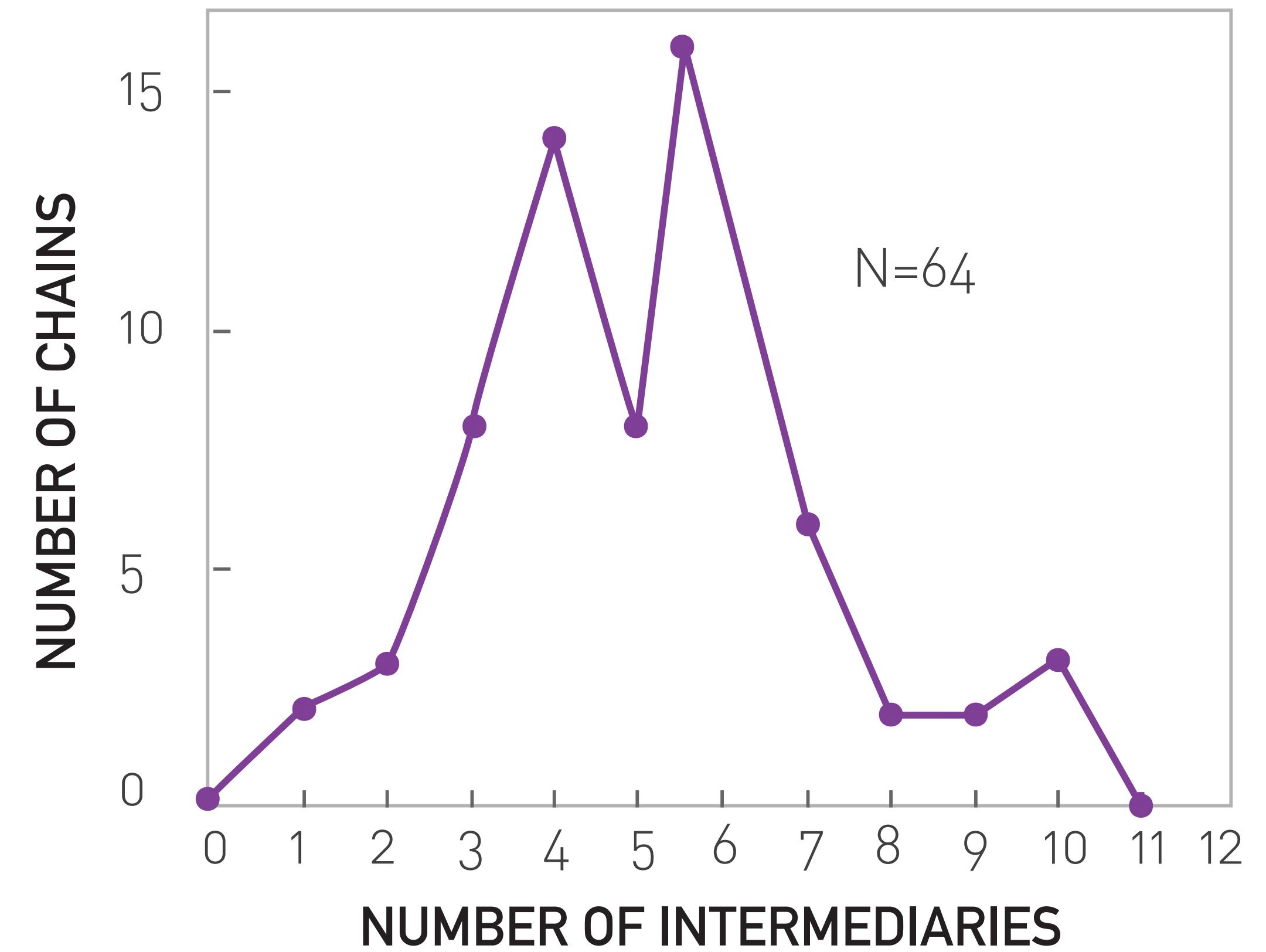
Randomly selected residents of  
Wichita and Omaha were asked  
to forward a letter to someone  
who is most likely to know the  
target person.

# six degrees of separation

Stanley Milgram (1967)

Two targets in Boston and Sharon, MA.

Randomly selected residents of Wichita and Omaha were asked to forward a letter to someone who is most likely to know the target person.



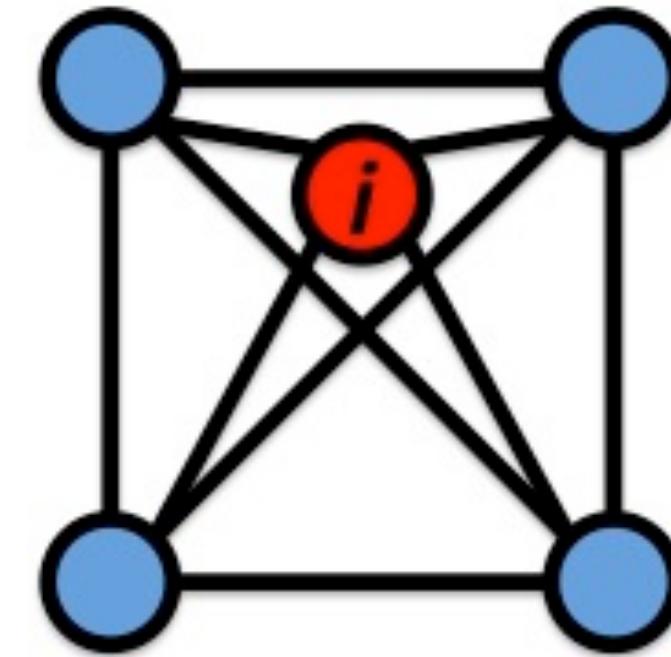
# clustering coefficient

The clustering coefficient of a node captures the degree to which the neighbors of a given node link to each other, i.e.  
**what fraction of your neighbors are connected?**

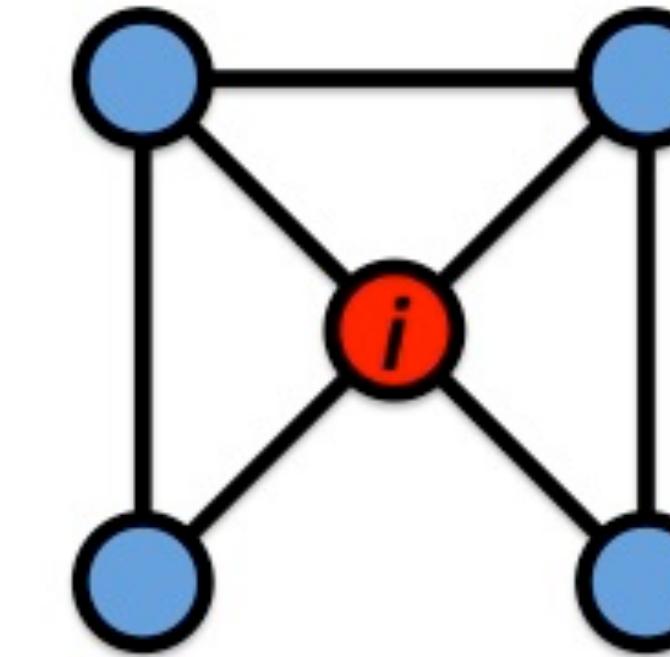
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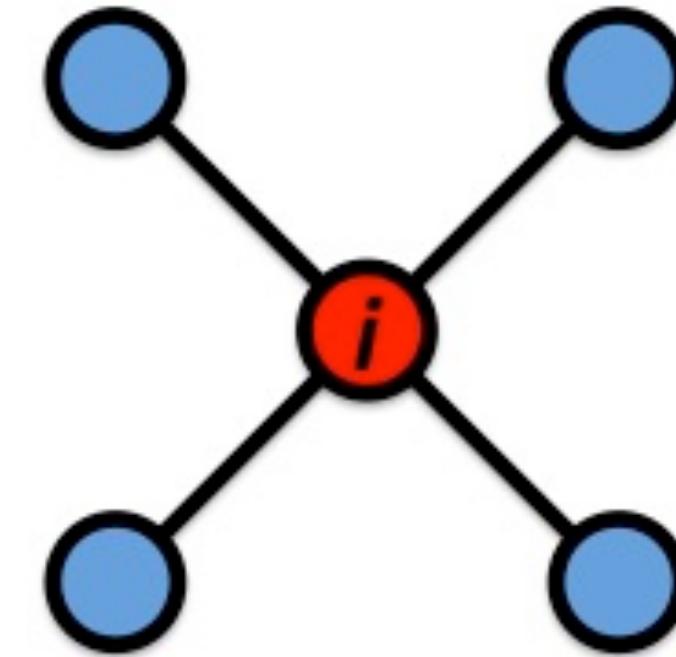
$$C_i = \frac{2L_i}{k_i(k_i - 1)}$$



$$C_i = 1$$



$$C_i = 1/2$$

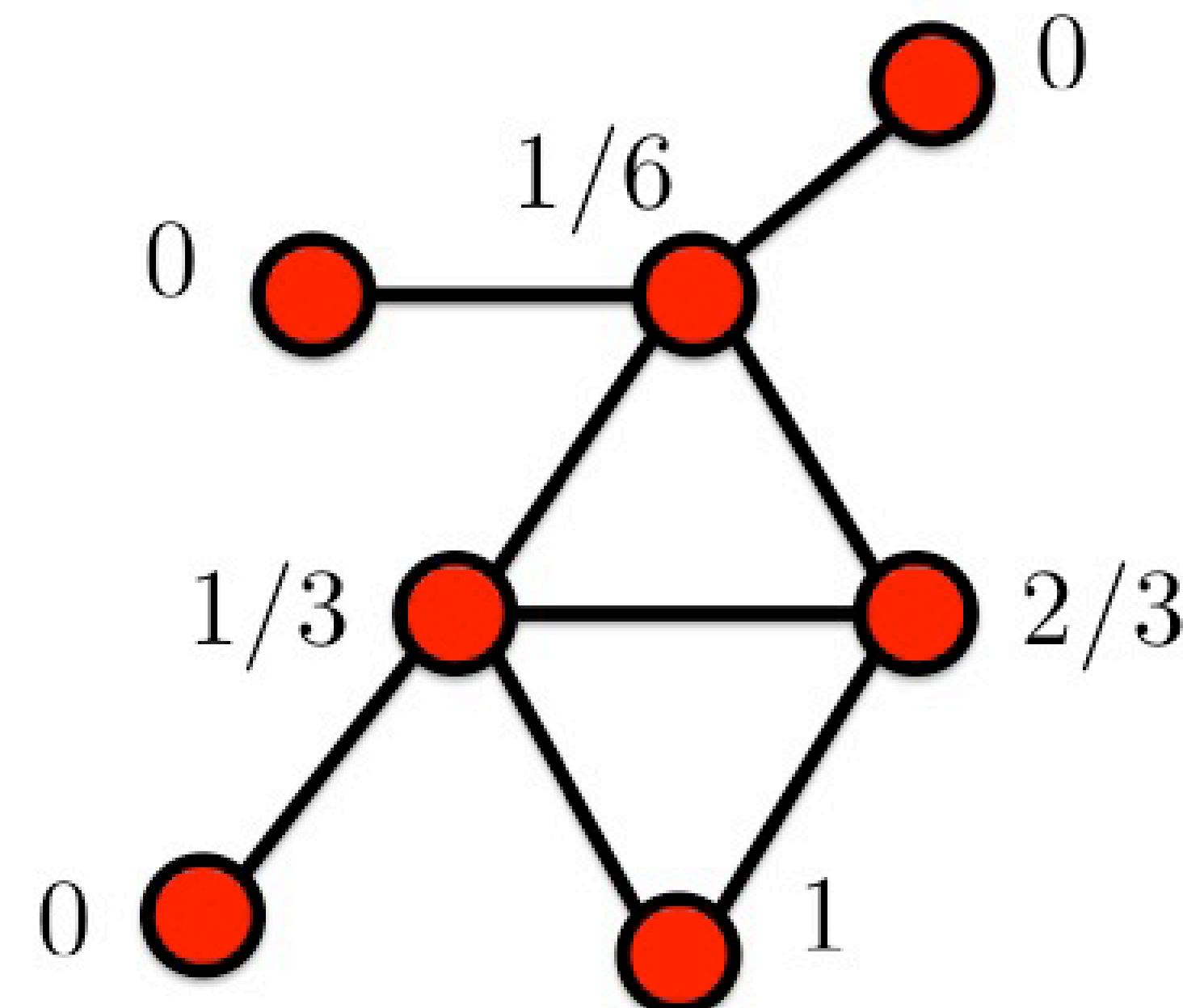


$$C_i = 0$$

# clustering coefficient

The degree of clustering of a whole network is captured by the **average clustering coefficient**, representing the average of  $C$  over all nodes  $i = 1, \dots, N$

$$\langle C_i \rangle = \frac{1}{N} \sum_i C_i$$



$$\langle C \rangle = \frac{13}{42} \approx 0.310$$

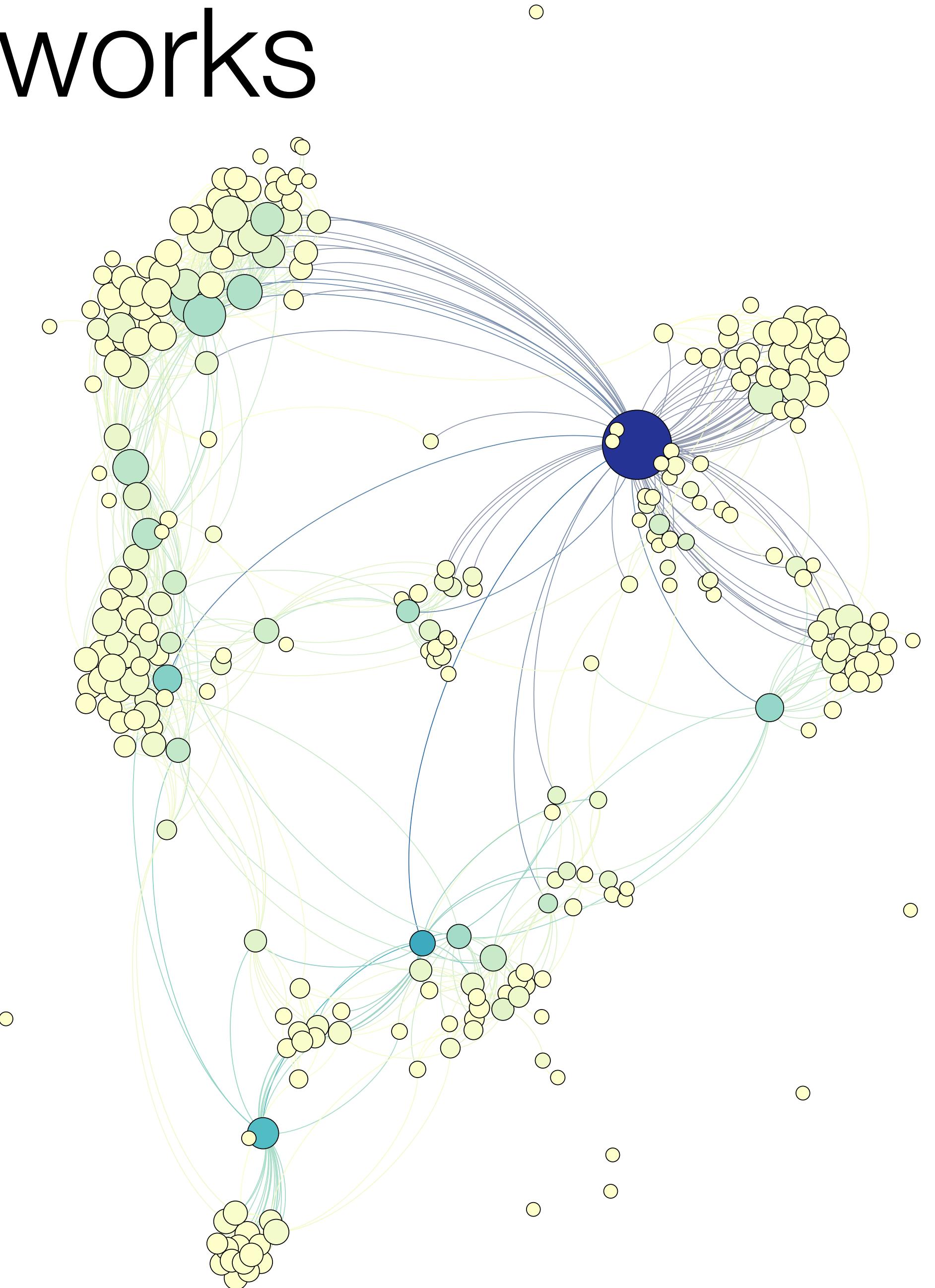
# clustering coefficient

the ***global clustering coefficient*** measures the total number of closed triangles in a network. Indeed,  $L_i$  in the previous equation is the number of triangles that node  $i$  participates in, as each link between two neighbors of node  $i$  closes a triangle.

$$C_{\Delta} = \frac{3 \times \text{Number Of Triangles}}{\text{Number Of Connected Triples}}$$

# real world networks

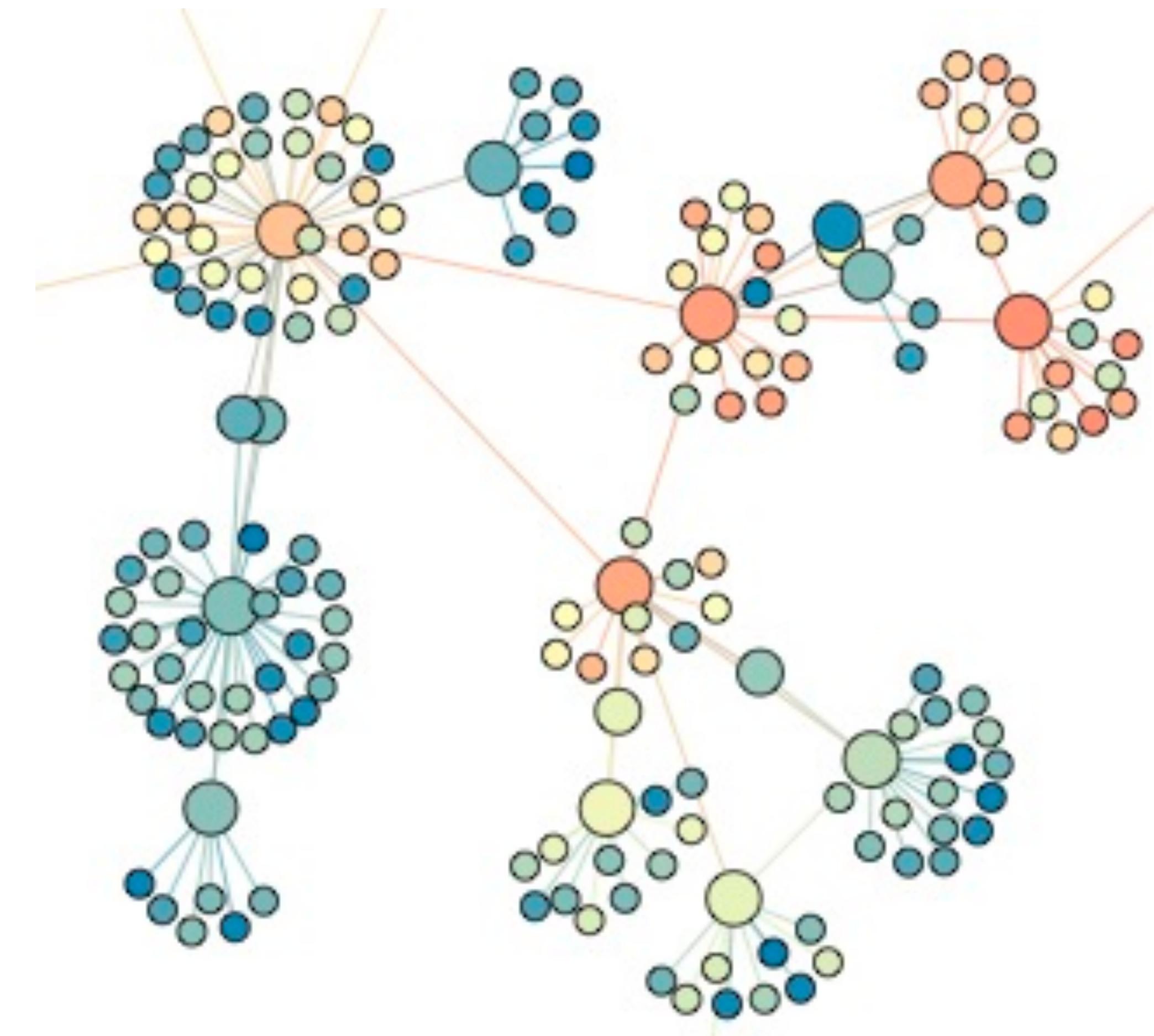
- Real world networks are highly clustered
- Average clustering coefficient can have values  $>0.5$
- **Triadic closure** in social networks is a common phenomenon



# measures of centrality

- Degree centrality
- Closeness centrality
- Betweenness centrality
- Eigenvector centrality
- Katz centrality
- Pagerank

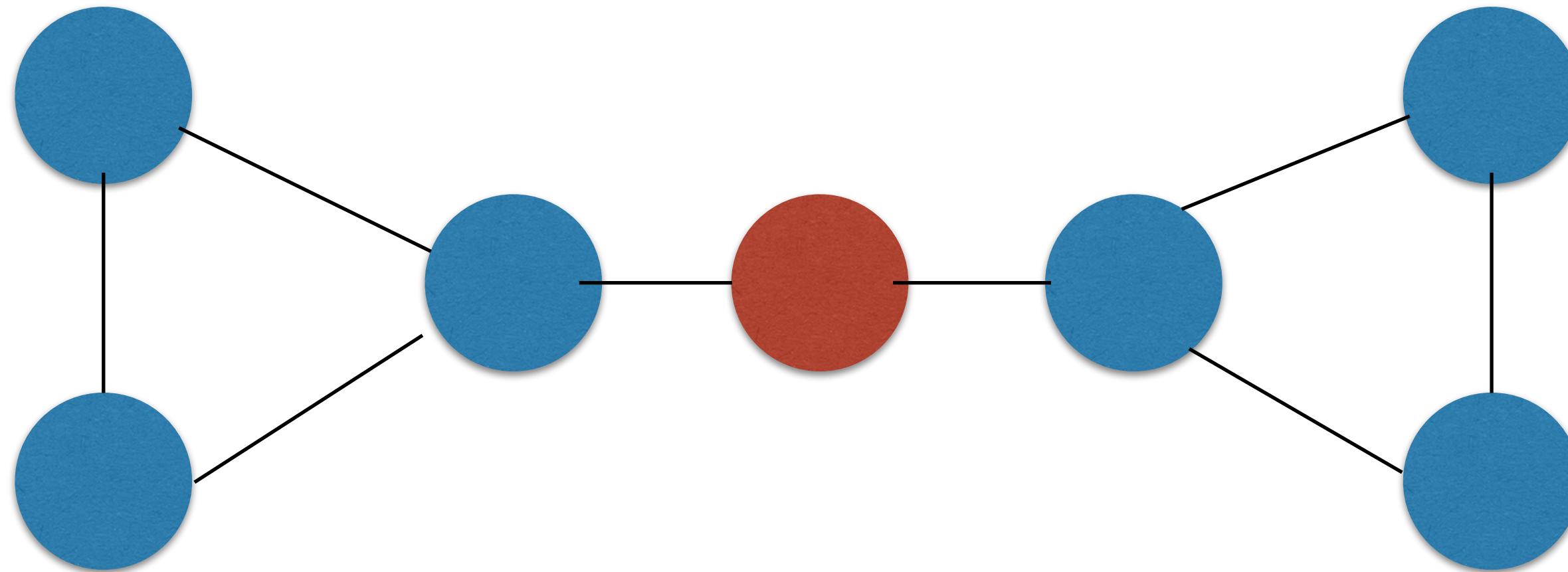
...



# betweenness centrality

**Betweenness captures a node's brokerage**

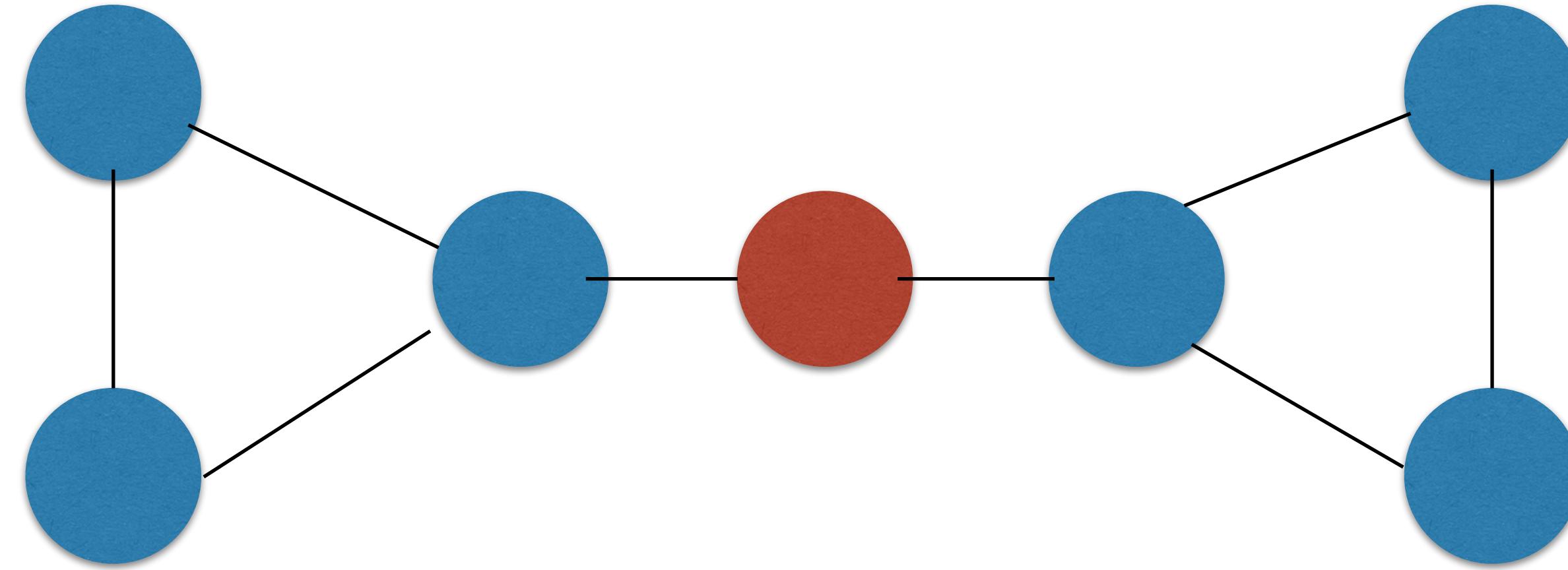
intuition: how many pairs of individuals would have to go through you in order to reach one another in the minimum number of hops?



# betweenness centrality

$$C_B(i) = \sum_{j < k} g_{jk}(i) / g_{jk}$$

$g_{jk}$  = #shortest paths connecting j and k  
 $g_{jk}(i)$  = #shortest paths connecting j and k through i

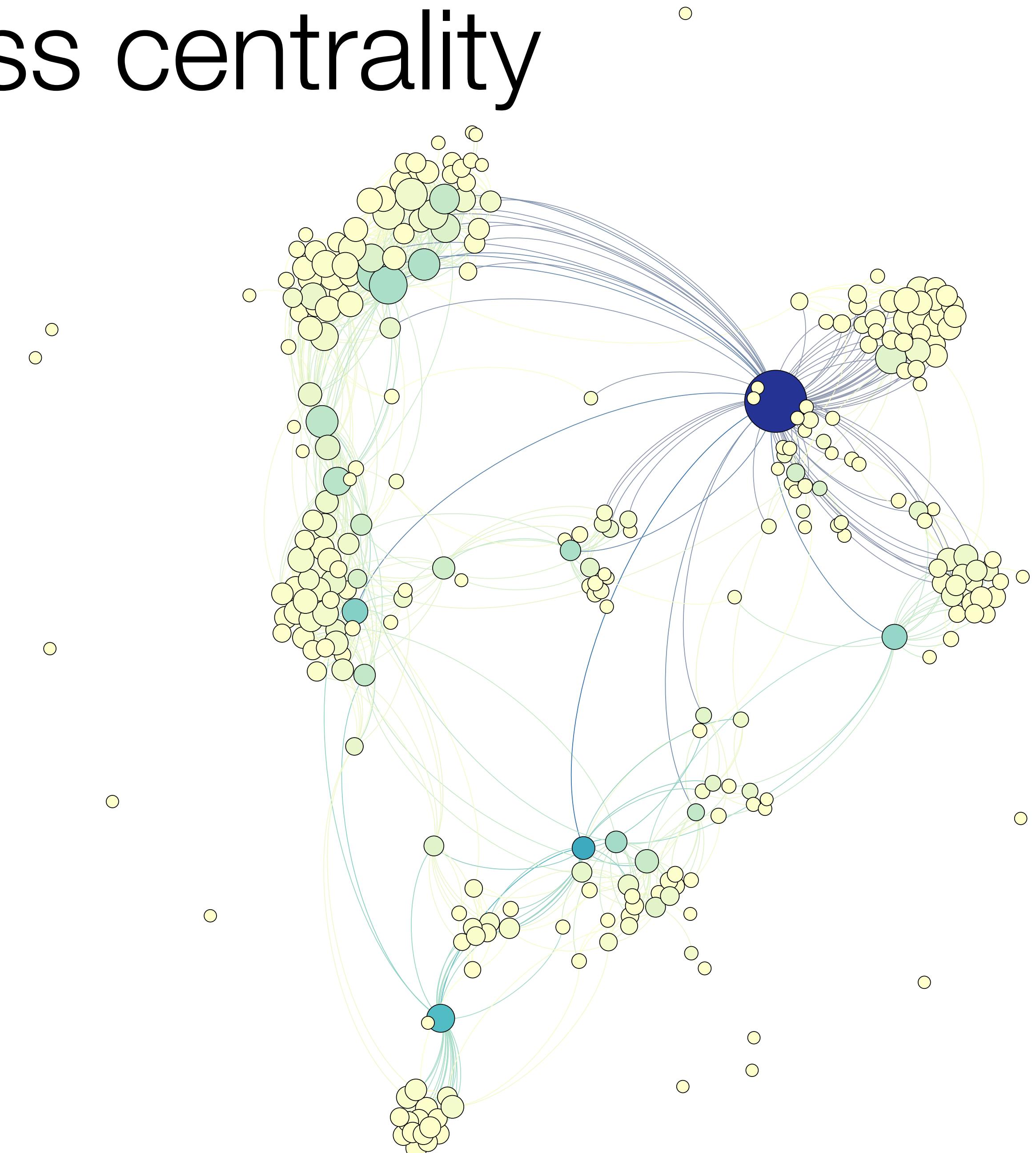


# betweenness centrality

My Facebook graph

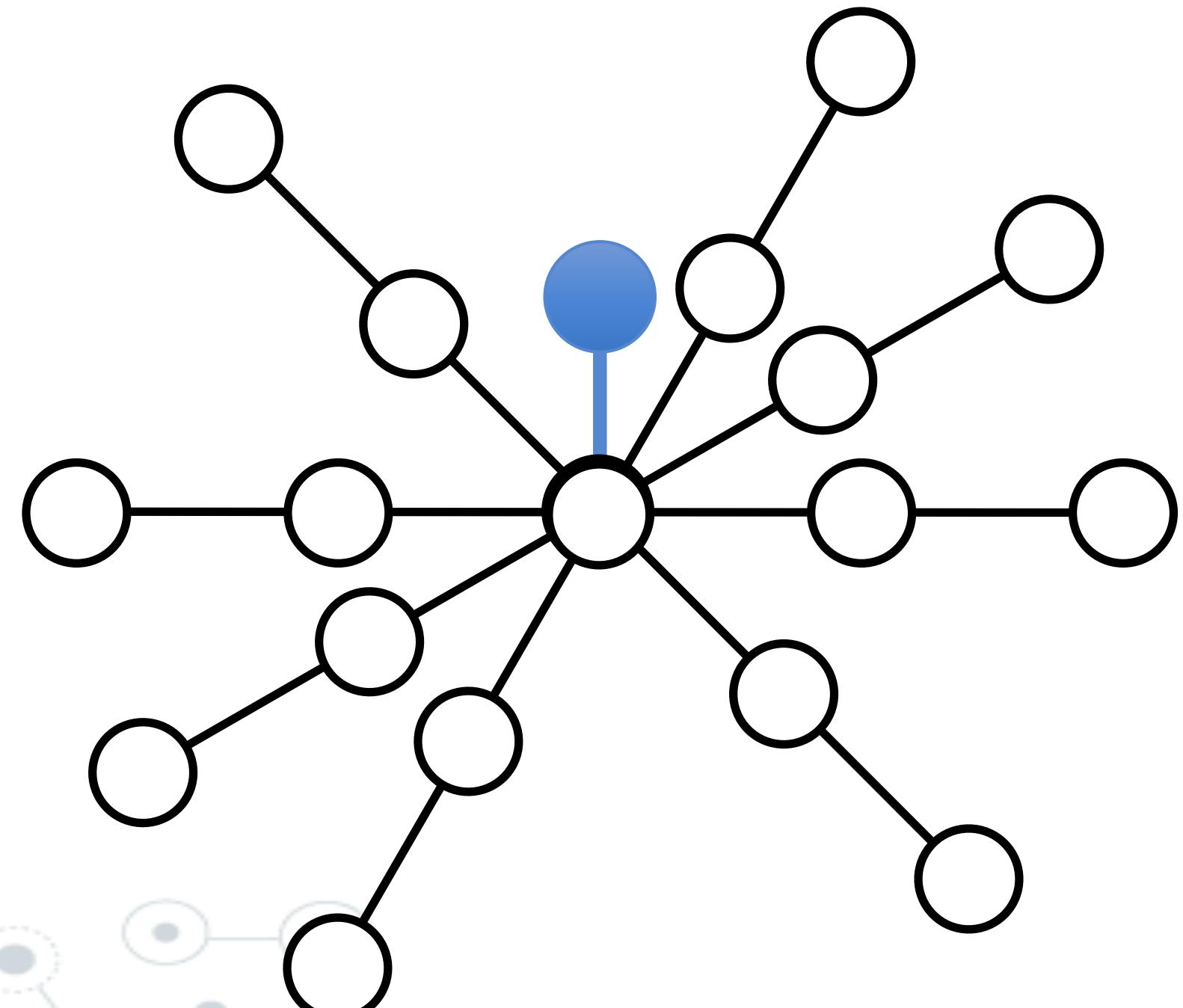
Node size is proportional  
to the degree

Node color is proportional  
to the betweenness



# closeness centrality

Closeness is based on the length of the average shortest path between a node and all other nodes in the network.  
It quantifies the reachability of a node.



$$C_c(i) = \left[ \sum_{j=1}^N d(i,j) \right]^{-1}$$

# Katz centrality

$$C_{KZ}(i) = \sum_{k=1}^{\infty} \sum_{j=1}^n \alpha^k (A^k)_{ji}$$

Katz centrality computes the relative influence of a node within a network by measuring the number of the immediate neighbors and also all other nodes in the network that connect to the node through these immediate neighbors. Connections made with distant neighbors are, however, penalized by an attenuation factor

# Katz centrality

$$C_{KZ}(i) = \alpha \sum_{j=1}^n A_{ij} C_{KZ}(j)$$

Katz centrality computes the relative influence of a node within a network by measuring the number of the immediate neighbors and also all other nodes in the network that connect to the node through these immediate neighbors. Connections made with distant neighbors are, however, penalized by an attenuation factor