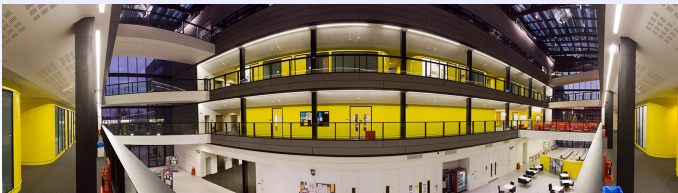


# How and How Not to Compute the Exponential of a Matrix

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# Outline

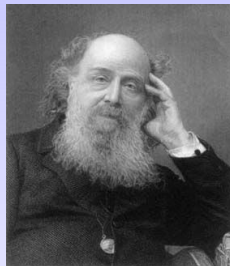
1 History & Properties

2 Applications

3 Methods

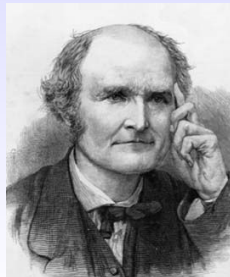
# Cayley and Sylvester

- Term “**matrix**” coined in 1850 by James Joseph Sylvester, FRS (1814–1897).



- **Matrix algebra** developed by Arthur Cayley, FRS (1821–1895).

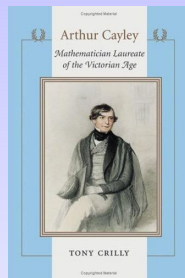
**Memoir on the Theory of Matrices (1858).**



# Cayley and Sylvester on Matrix Functions

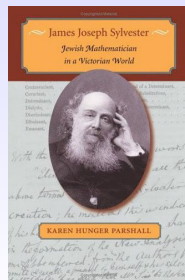
- Cayley considered matrix square roots in his 1858 memoir.

**Tony Crilly, *Arthur Cayley: Mathematician Laureate of the Victorian Age*, 2006.**



- Sylvester (1883) gave first definition of  $f(A)$  for general  $f$ .

**Karen Hunger Parshall, *James Joseph Sylvester. Jewish Mathematician in a Victorian World*, 2006.**



Laguerre (1867):

En particulier, si nous définissons  $e^X$ ,  $X$  étant un système d'ordre quelconque, comme étant la somme de la série

$$1 + X + \frac{X^2}{1.2} + \frac{X^3}{1.2.3} + \dots,$$

$e^X$  sera une fonction de la variable  $X$ ; mais il est à remarquer qu'en général on n'aura pas

$$e^X \cdot e^Y = e^{X+Y}.$$

Peano (1888):

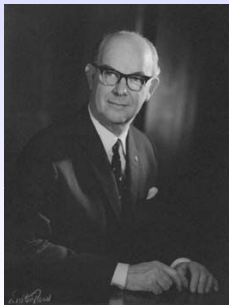
$$x = \left[ 1 + Rt + \frac{1}{2!} (Rt)^2 + \dots \right] a,$$

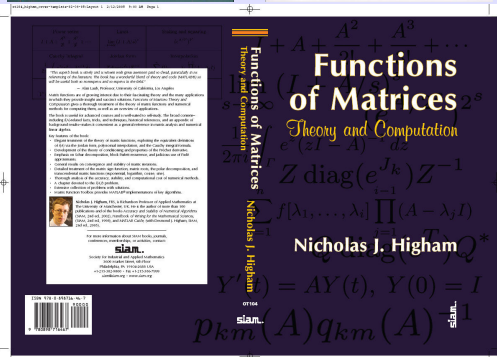
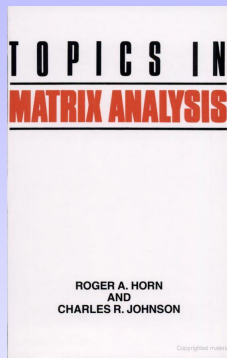
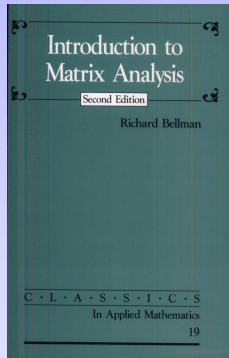
$$\text{ou, en posant } e^R = 1 + R + \frac{1}{2!} R^2 + \dots,$$

$$x = e^{Rt} a.$$

# Matrices in Applied Mathematics

- Frazer, Duncan & Collar, Aerodynamics Division of NPL: aircraft flutter, matrix structural analysis.
- **Elementary Matrices & Some Applications to Dynamics and Differential Equations, 1938.**  
Emphasizes importance of  $e^A$ .
- Arthur Roderick Collar, FRS (1908–1986): *“First book to treat matrices as a branch of applied mathematics”*.





# Formulae

$$\mathbf{A} \in \mathbb{C}^{n \times n}:$$

## Power series

$$\mathbf{I} + \mathbf{A} + \frac{\mathbf{A}^2}{2!} + \frac{\mathbf{A}^3}{3!} + \dots$$

## Limit

$$\lim_{s \rightarrow \infty} (\mathbf{I} + \mathbf{A}/s)^s$$

## Scaling and squaring

$$(e^{\mathbf{A}/2^s})^{2^s}$$

## Cauchy integral

$$\frac{1}{2\pi i} \int_{\Gamma} e^z (z\mathbf{I} - \mathbf{A})^{-1} dz$$

## Jordan form

$$\mathbf{Z} \text{diag}(e^{\mathbf{J}_k}) \mathbf{Z}^{-1}$$

## Interpolation

$$\sum_{i=1}^n f[\lambda_1, \dots, \lambda_i] \prod_{j=1}^{i-1} (\mathbf{A} - \lambda_j \mathbf{I})$$

## Differential system

$$\mathbf{Y}'(t) = \mathbf{A}\mathbf{Y}(t), \mathbf{Y}(0) = \mathbf{I}$$

## Schur form

$$\mathbf{Q} \text{diag}(e^{\mathbf{T}}) \mathbf{Q}^*$$

## Padé approximation

$$\mathbf{p}_{km}(\mathbf{A}) \mathbf{q}_{km}(\mathbf{A})^{-1}$$



# Properties (1)

## Theorem

*For  $A, B \in \mathbb{C}^{n \times n}$ ,  $e^{(A+B)t} = e^{At}e^{Bt}$  for all  $t$  if and only if  $AB = BA$ .*

## Theorem (Wermuth)

*Let  $A, B \in \mathbb{C}^{n \times n}$  have algebraic elements and let  $n \geq 2$ . Then  $e^A e^B = e^B e^A$  if and only if  $AB = BA$ .*

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## Theorem

*Let  $A \in \mathbb{C}^{n \times n}$  and  $B \in \mathbb{C}^{m \times m}$ . Then  $e^{A \oplus B} = e^A \otimes e^B$ , where  $A \oplus B = A \otimes I_m + I_n \otimes B$ .*

# Properties (2)

## Theorem (Suzuki)

For  $A \in \mathbb{C}^{n \times n}$ , let

$$T_{r,s} = \left[ \sum_{i=0}^r \frac{1}{i!} \left( \frac{A}{s} \right)^i \right]^s.$$

Then

$$\|e^A - T_{r,s}\| \leq \frac{\|A\|^{r+1}}{s^r(r+1)!} e^{\|A\|}$$

and  $\lim_{r \rightarrow \infty} T_{r,s}(A) = \lim_{s \rightarrow \infty} T_{r,s}(A) = e^A$ .

# Outline

1 History & Properties

2 Applications

3 Methods

# Application: Control Theory

Convert **continuous-time system**

$$\begin{aligned}\frac{dx}{dt} &= Fx(t) + Gu(t), \\ y &= Hx(t) + Ju(t),\end{aligned}$$

to **discrete-time state-space system**

$$\begin{aligned}x_{k+1} &= Ax_k + Bu_k, \\ y_k &= Hx_k + Ju_k.\end{aligned}$$

Have

$$A = e^{F\tau}, \quad B = \left( \int_0^\tau e^{Ft} dt \right) G,$$

where  $\tau$  is the sampling period.

MATLAB Control System Toolbox: **c2d** and **d2c**.

# Psi Functions: Definition

$$\psi_0(z) = e^z, \quad \psi_1(z) = \frac{e^z - 1}{z}, \quad \psi_2(z) = \frac{e^z - 1 - z}{z^2}, \dots$$

$$\psi_{k+1}(z) = \frac{\psi_k(z) - 1/k!}{z}.$$

$$\psi_k(z) = \sum_{j=0}^{\infty} \frac{z^j}{(j+k)!}.$$

# Psi Functions: Solving DEs

$$y \in \mathbb{C}^n, A \in \mathbb{C}^{n \times n}.$$

$$\frac{dy}{dt} = Ay, \quad y(0) = y_0 \quad \Rightarrow \quad y(t) = e^{At}y_0.$$

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$$\frac{dy}{dt} = Ay + b, \quad y(0) = 0 \quad \Rightarrow \quad y(t) = t\psi_1(tA)b.$$

$$\frac{dy}{dt} = Ay + ct, \quad y(0) = 0 \quad \Rightarrow \quad y(t) = t^2\psi_2(tA)c.$$

$$\vdots$$

# Exponential Integrators

Consider

$$y' = Ly + N(y).$$

$N(y(t)) \approx N(y(0))$  implies

$$y(t) \approx e^{tL}y_0 + t\psi_1(tL)N(y(0)).$$

**Exponential Euler method:**

$$y_{n+1} = e^{hL}y_n + h\psi_1(hL)N(y_n).$$

Lawson (1967); recent resurgence.

# The Average Eye

First order character of optical system characterized by transference matrix  $T = \begin{bmatrix} S & \delta \\ 0 & 1 \end{bmatrix} \in \mathbb{R}^{5 \times 5}$ , where  $S \in \mathbb{R}^{4 \times 4}$  is symplectic:  $S^T J S = J$ , where  $J = \begin{bmatrix} 0 & I_2 \\ -I_2 & 0 \end{bmatrix}$ .

Average  $m^{-1} \sum_{i=1}^m T_i$  is not a transference matrix.

Harris (2005) proposes the average  $\exp(m^{-1} \sum_{i=1}^m \log(T_i))$ .

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---

For Hermitian pos def  $A$  and  $B$ , Arsigny et al. (2007) define the log-Euclidean mean

$$E(A, B) = \exp(\frac{1}{2}(\log(A) + \log(B))).$$

# Beyond Matrices

- GluCat library: generic library of C++ templates for universal Clifford algebras: exp, log, square root, trig functions.

<http://glucat.sourceforge.net>.

- Group exponential of a diffeomorphism in computational anatomy to study variability among medical images (Bossa et al., 2008).

# Outline

1 History & Properties

2 Applications

3 **Methods**

# Sengupta (Adv. Appl. Prob., 1989)

*Note that  $e^T$  is a power series in  $T$ . This means that a **wide variety of methods** in linear algebra can also be used to evaluate  $e^T$ . . . . brute force evaluation of the power series, . . . matrix decomposition methods or polynomial methods based on the **Cayley-Hamilton theorem** . . .*

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*Since evaluation of functions of matrices may be **fraught with difficulties** (such as roundoff and truncation errors, ill conditioning, near confluence of eigenvalues, etc.), there is a distinct advantage in having a rich class of solution techniques available for finding  $e^T$ . **If one method fails to find an accurate answer, one can always fall back on a different method.***



# Cayley–Hamilton Theorem

## Theorem (Cayley, 1857)

*If  $A, B \in \mathbb{C}^{n \times n}$ ,  $AB = BA$ , and  $f(x, y) = \det(xA - yB)$  then  $f(B, A) = 0$ .*

- $p(t) = \det(tI - A)$  implies  $p(A) = 0$ .
- $A^n = \sum_{k=0}^{n-1} c_n A^k$ .
- $e^A = \sum_{k=0}^{n-1} d_n A^k$ .

# Walz's Method

Walz (1988) proposed computing

$$C_k = (I + 2^{-k}A)^{2^k}$$

with Richardson extrapolation to accelerate cgce of the  $C_k$ .

Numerically unstable in practice (Parks, 1994).

# Diagonalization (1)

$A = Z \text{diag}(\lambda_i) Z^{-1}$  implies  $f(A) = Z \text{diag}(f(\lambda_i)) Z^{-1}$ .

But

- $Z$  may be ill conditioned ( $\kappa(Z) = \|Z\| \|Z^{-1}\| \gg 1$ ).
- $A$  may not be diagonalizable.

# Diagonalization (2)

```
>> A = [3 -1; 1 1]; X = funm_ev(A,@exp)
```

```
X =
```

```
14.7781    -7.3891
```

```
7.3891      0
```

```
>> norm(X - expm(A))/norm(expm(A))
```

```
ans = 1.3431e-009
```

```
>> expm_cond(A)
```

```
ans = 3.4676
```

```
>> [Z,D]=eig(A)
```

```
Z =
```

```
0.7071
```

```
0.7071
```

```
0.7071
```

```
0.7071
```

```
D =
```

```
2.0000
```

```
0
```

```
0
```

```
2.0000
```

# Scaling and Squaring Method

- ▶  $B \leftarrow A/2^s$  so  $\|B\|_\infty \approx 1$
- ▶  $r_m(B) = [m/m]$  Padé approximant to  $e^B$
- ▶  $X = r_m(B)^{2^s} \approx e^A$

- Originates with **Lawson (1967)**.
- **Ward (1977)**: algorithm, with rounding error analysis and a posteriori error bound.
- **Moler & Van Loan (1978)**: give backward error analysis allowing choice of  $s$  and  $m$ .
- **H (2005)**: sharper analysis giving optimal  $s$  and  $m$ .  
MATLAB's **expm**, Mathematica, NAG Library Mark 22.

# Padé Approximants $r_m$ to $e^x$

$r_m(x) = p_m(x)/q_m(x)$  known explicitly:

$$p_m(x) = \sum_{j=0}^m \frac{(2m-j)! m!}{(2m)! (m-j)! j!} x^j$$

and  $q_m(x) = p_m(-x)$ . Error satisfies

$$e^x - r_m(x) = (-1)^m \frac{(m!)^2}{(2m)!(2m+1)!} x^{2m+1} + O(x^{2m+2}).$$

# Scaling and Squaring Method

$$h_{2m+1}(X) := \log(e^{-X} r_m(X)) = \sum_{k=2m+1}^{\infty} c_k X^k.$$

Then  $r_m(X) = e^{X+h_{2m+1}(X)}$ . Hence

$$r_m(2^{-s}A)^{2^s} = e^{A+2^s h_{2m+1}(2^{-s}A)} =: e^{A+\Delta A}.$$

Want  $\|\Delta A\|/\|A\| \leq u$ .

- Moler & Van Loan (1978): a priori bound for  $h_{2m+1}$ ;  $m = 6$ ,  $\|2^{-s}A\| \leq 1/2$  in MATLAB.
- H (2005): sharp normwise bound using symbolic arithmetic and high precision. Choose  $(s, m)$  to minimize computational cost.

# Scaling & Squaring Algorithm (H, 2005)

$m$	3	5	7	9	13
$\theta_m$	0.015	0.25	0.95	2.1	5.4

for  $m = [3 \ 5 \ 7 \ 9 \ 13]$

    if  $\|A\|_1 \leq \theta_m$ ,  $X = r_m(A)$ , quit, end  
end



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end

$A \leftarrow A/2^s$  with  $s \geq 0$  minimal s.t.  $\|A/2^s\|_1 \leq \theta_{13} = 5.4$

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$A \leftarrow A/2^s$  with  $s \geq 0$  minimal s.t.  $\|A/2^s\|_1 \leq \theta_{13} = 5.4$

$A_2 = A^2$ ,  $A_4 = A_2^2$ ,  $A_6 = A_2 A_4$

$U = A[A_6(b_{13}A_6 + b_{11}A_4 + b_9A_2) + b_7A_6 + b_5A_4 + b_3A_2 + b_1I]$

$V = A_6(b_{12}A_6 + b_{10}A_4 + b_8A_2) + b_6A_6 + b_4A_4 + b_2A_2 + b_0I$

Solve  $(-U + V)r_{13} = U + V$  for  $r_{13}$ .

$X = r_{13}^{2^s}$  by repeated squaring.

# Example

$$A = \begin{bmatrix} 1 & b \\ 0 & -1 \end{bmatrix}, \quad e^A = \begin{bmatrix} e & \frac{b}{2}(e - e^{-1}) \\ 0 & e^{-1} \end{bmatrix}.$$

$b$	$\text{expm}(A)$	$s$	$\text{expm}(A)^\dagger$	$s$	$\text{funm}(A)$
$10^3$	1.7e-15	8	1.9e-16	0	1.9e-16
$10^4$	1.8e-13	11	7.6e-20	0	3.8e-20
$10^5$	7.5e-13	15	1.2e-16	0	1.2e-16
$10^6$	1.3e-11	18	2.0e-16	0	2.0e-16
$10^7$	7.2e-11	21	1.6e-16	0	1.6e-16
$10^8$	3.0e-12	25	1.3e-16	0	1.3e-16

For  $b = 10^8$ ,  $r_m(x)^{2^{25}} \approx ((1 + \frac{1}{2}x)/(1 - \frac{1}{2}x))^{2^{25}}$  with  
 $x = \pm 2^{-25} \approx \pm 10^{-8}$ .

# Overscaling

Kenney & Laub (1998); Dieci & Papini (2000).

A large  $\|A\|$  causes a larger than necessary  $s$  to be chosen, with a harmful effect on accuracy.

$$\exp\left(\begin{bmatrix} A_{11} & A_{12} \\ 0 & A_{22} \end{bmatrix}\right) = \begin{bmatrix} e^{A_{11}} & \int_0^1 e^{A_{11}(1-s)} A_{12} e^{A_{22}s} ds \\ 0 & e^{A_{22}} \end{bmatrix}.$$

# Insight

## Al-Mohy & H (2009):

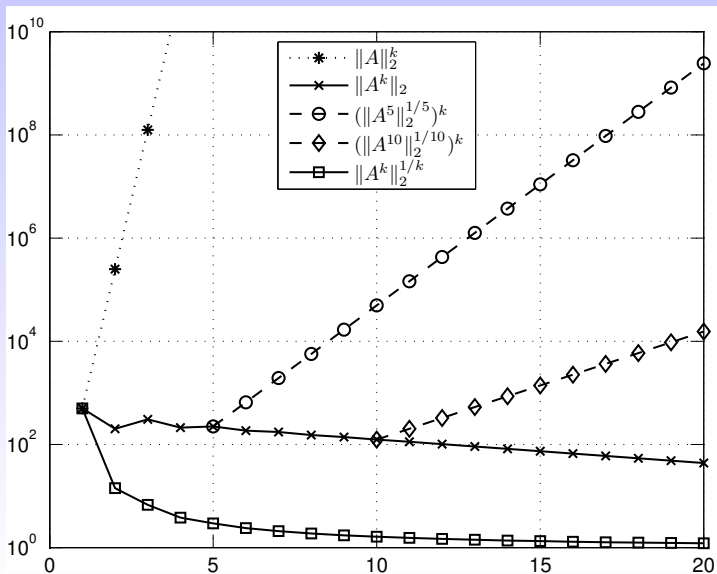
Existing method based on analysis in terms of  $\|A\|$ .

Why not instead use  $\|A^k\|^{1/k}$ ?

$$\rho(A) \leq \|A^k\|^{1/k} \leq \|A\|, \quad k = 1 : \infty,$$

$$\lim_{k \rightarrow \infty} \|A^k\|^{1/k} = \rho(A).$$

$$A = \begin{bmatrix} 0.9 & 500 \\ 0 & -0.5 \end{bmatrix}.$$



# Key Bounds (1)

## Theorem

For any  $A \in \mathbb{C}^{n \times n}$ ,

$$\left\| \sum_{k=\ell}^{\infty} c_k A^k \right\| \leq \sum_{k=\ell}^{\infty} |c_k| \left( \|A^t\|^{1/t} \right)^k$$

where  $\|A^t\|^{1/t} = \max\{\|A^k\|^{1/k} : k \geq \ell, c_k \neq 0\}$ .

**Proof.** Use  $\|A^k\| = (\|A^k\|^{1/k})^k \leq (\|A^t\|^{1/t})^k$ .  $\square$

# Key Bounds (2)

## Lemma

If  $k = pm_1 + qm_2$  with  $p, q \in \mathbb{N}$  and  $m_1, m_2 \in \mathbb{N} \cup \{0\}$ ,

$$\|A^k\|^{1/k} \leq \max(\|A^p\|^{1/p}, \|A^q\|^{1/q}).$$

**Proof.** Let  $\delta = \max(\|A^p\|^{1/p}, \|A^q\|^{1/q})$ . Then

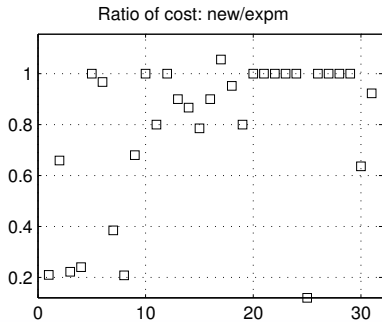
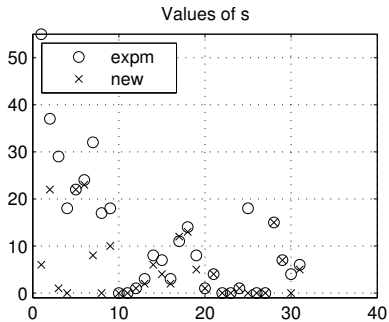
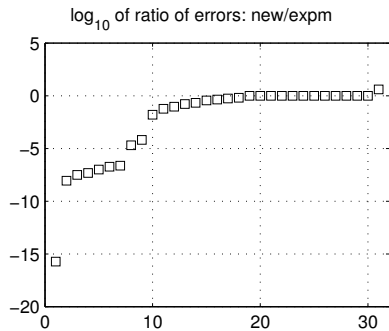
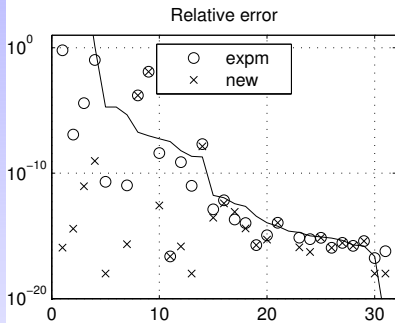
$$\begin{aligned} \|A^k\| &\leq \|A^{pm_1}\| \|A^{qm_2}\| \\ &\leq (\|A^p\|^{1/p})^{pm_1} (\|A^q\|^{1/q})^{qm_2} \\ &\leq \delta^{pm_1} \delta^{qm_2} = \delta^k. \quad \square \end{aligned}$$

■ Take  $\{p, q\} = \{r, r+1\}$  for  $k \geq r(r-1)$ .



# New Scaling and Squaring Algorithm

- Truncation bounds use  $\|A^k\|^{1/k}$  instead of  $\|A\|$ .
- Roundoff considerations: correction to chosen  $m$ .
- Use *estimates* of  $\|A^k\|$  where necessary (alg of H & Tisseur (2000)).
- Special treatment of triangular matrices to ensure accurate diagonal.
- New alg no slower than **expm**, potentially faster, potentially more accurate.



# Summary of New Alg

- Major benefits in speed and accuracy through using  $\|A^k\|^{1/k}$  in place of  $\|A\|^k$ .
- Overscaling problem “solved”.
- Stability of squaring phase remains an open question.

# Frechét Derivative

$$f(A + E) - f(A) - L(A, E) = o(\|E\|).$$

$$L(A, E) = \int_0^1 e^{A(1-s)} E e^{As} ds.$$

- Method based on

$$f\left(\begin{bmatrix} X & E \\ 0 & X \end{bmatrix}\right) = \begin{bmatrix} f(X) & L(X, E) \\ 0 & f(X) \end{bmatrix}.$$

- Kenney & Laub (1998): Kronecker–Sylvester alg, Padé of  $\tanh(x)/x$ :  $538n^3$  (complex) flops.
- **Al-Mohy & H (2009)**:  $e^A$  and  $L(A, E)$  in only  $48n^3$  flops.

# In Conclusion

- Many applications of  $f(A)$ , e.g. control theory, Markov chains, theoretical physics.
- Need better understanding of conditioning of  $f(A)$ .
- How to exploit structure?
- Need “factorization-free” methods for large, sparse  $A$ .
- Specialize to  $f(A)b$  problem.

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