

# On the inexact scaled gradient projection method

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The problem, definitions and preliminaries results

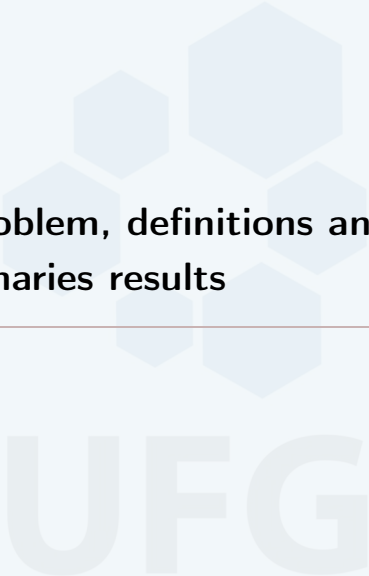
Inexact scaled gradient method

Partial asymptotic convergence

Full asymptotic convergence

Iteration-complexity bound

Numerical experiments



## The problem, definitions and preliminaries results

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## The main problem

We want to present an inexact version of the scaled gradient projection method for **constrained convex optimization problem** as follows

$$\min\{f(x) : x \in C\}, \tag{1}$$

where  $C$  is a closed and convex subset of  $\mathbb{R}^n$  and  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  is a continuously differentiable function.

# Scaled Gradient Projection Method<sup>1</sup>

**Step 0.** Choose  $\sigma, \tau \in (0, 1)$ ,  $0 < \alpha_{\min} \leq \alpha_{\max}$ . Let  $x^0 \in C$  and set  $k = 0$ ;

**Step 1.** Choose  $\alpha_k \in [\alpha_{\min}, \alpha_{\max}]$  and a positive definite matrix  $D_k$  and take  $w^k \in C$  as

$$w^k := \mathcal{P}_C^{D_k}(x^k - \alpha_k D_k^{-1} \nabla f(x^k))$$

If  $w^k = x^k$ , then **stop**; otherwise,

**Step 2.** Choose  $\tau_k$  and define the next iterate  $x^{k+1}$  as

$$x^{k+1} = x^k + \tau_k(w^k - x^k). \quad (2)$$

and go back to the Step 1.

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<sup>1</sup>S. Bonettini, R. Zanella, and L. Zanni. “A scaled gradient projection method for constrained image deblurring”. In: *Inverse Problems* 25.1 (2009), pp. 015002, 23.

## Scaled Gradient Projection Method

Let  $D$  be a  $n \times n$  positive definite matrix and  $\|\cdot\|_D : \mathbb{R}^n \rightarrow \mathbb{R}$  be the norm defined by

$$\|d\|_D := \sqrt{\langle Dd, d \rangle}, \quad \forall d \in \mathbb{R}^n.$$

For a fixed constant  $\mu \geq 1$ , denote by  $\mathcal{D}_\mu$  the set of symmetric positive definite matrices  $n \times n$  with all eigenvalues contained in the interval  $[\frac{1}{\mu}, \mu]$ .

- $\mathcal{D}_\mu$  is compact;
- If  $D \in \mathcal{D}_\mu$ , it follows that  $D^{-1}$  also belongs to  $\mathcal{D}_\mu$ ;
- $\forall D \in \mathcal{D}_\mu$ , we obtain

$$\frac{1}{\mu} \|d\|^2 \leq \|d\|_D^2 \leq \mu \|d\|^2, \quad \forall d \in \mathbb{R}^n.$$

## Definition

The **exact projection** of the point  $v \in \mathbb{R}^n$  onto  $C$  with respect to the norm  $\|\cdot\|_D$ , denoted by  $\mathcal{P}_C^D(v)$ , is defined by

$$\mathcal{P}_C^D(v) := \arg \min_{z \in C} \|z - v\|_D^2.$$

## Lemma

Let  $v, w \in \mathbb{R}^n$ . Then,  $w = \mathcal{P}_C^D(v)$  if and only if  $w \in C$  and

$$\langle D(v - w), y - w \rangle \leq 0,$$

for all  $y \in C$ .

### Definition

The **feasible inexact projection mapping**, with respect to the norm  $\|\cdot\|_D$ , onto  $C$  relative to a point  $u \in C$  and forcing parameter  $\zeta \in (0, 1]$ , denoted by  $\mathcal{P}_{C,\zeta}^D(u, \cdot) : \mathbb{R}^n \rightrightarrows C$ , is the set-valued mapping defined as follows

$$\mathcal{P}_{C,\zeta}^D(u, v) := \{w \in C : \|w - v\|_D^2 \leq \zeta \|\mathcal{P}_C^D(v) - v\|_D^2 + (1 - \zeta) \|u - v\|_D^2\}.$$

Each point  $w \in \mathcal{P}_{C,\zeta}^D(u, v)$  is called a **feasible inexact projection**, with respect to the norm  $\|\cdot\|_D$ , of  $v$  onto  $C$  relative to  $u$  and forcing parameter  $\zeta \in (0, 1]$ .

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<sup>2</sup>Ernesto G. Birgin, José Mario Martínez, and Marcos Raydan. “Inexact spectral projected gradient methods on convex sets”. In: *IMA J. Numer. Anal.* 23.4 (2003), pp. 539–559.



### Definition

The **feasible inexact projection mapping**, with respect to the norm  $\|\cdot\|_D$ , onto  $C$  relative to  $u \in C$  and forcing parameter  $\gamma \geq 0$ , denoted by  $\mathcal{R}_{C,\gamma}^D(u, \cdot) : \mathbb{R}^n \rightrightarrows C$ , is the set-valued mapping defined as follows

$$\mathcal{R}_{C,\gamma}^D(u, v) := \{w \in C : \langle D(v - w), y - w \rangle \leq \gamma \|w - u\|_D^2, \quad \forall y \in C\}.$$

Each point  $w \in \mathcal{R}_{C,\gamma}^D(u, v)$  is called a feasible inexact projection, with respect to the norm  $\|\cdot\|_D$ , of  $v$  onto  $C$  relative to  $u$  and forcing parameter  $\gamma \geq 0$ .

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<sup>3</sup>Fabiana R. de Oliveira, Orizon P. Ferreira, and Gilson N. Silva. “Newton’s method with feasible inexact projections for solving constrained generalized equations”. In: *Comput. Optim. Appl.* 72.1 (2019), pp. 159–177.

### Lemma

Let  $v \in \mathbb{R}^n$ ,  $u \in C$ ,  $\gamma \geq 0$  and  $\zeta \in (0, 1]$ . If  $0 \leq \gamma < 1/2$  and  $\zeta = 1 - 2\gamma$ , then

$$\mathcal{R}_{C,\gamma}^D(u, v) \subset \mathcal{P}_{C,\zeta}^D(u, v).$$

### Proposition

Let  $v \in \mathbb{R}^n$ ,  $u \in C$  and assume that  $C$  is a bounded set. Then, for each  $0 < \gamma < 1/2$ , there exist  $0 < \zeta < 1$  such that

$$\mathcal{P}_{C,\zeta}^D(u, v) \subseteq \mathcal{R}_{C,\gamma}^D(u, v)$$

### Lemma

Let  $x \in C$ ,  $\alpha > 0$  and  $z(\alpha) = x - \alpha D^{-1} \nabla f(x)$ . Take  $w(\alpha) \in \mathcal{P}_{C,\zeta}^D(x, z(\alpha))$  with  $\zeta \in (0, 1]$ . Then, there hold

- (i)  $\langle \nabla f(x), w(\alpha) - x \rangle \leq -\frac{1}{2\alpha} \|w(\alpha) - x\|_D^2 + \frac{\zeta}{2\alpha} [\|\mathcal{P}_C^D(z(\alpha)) - z(\alpha)\|_D^2 - \|x - z(\alpha)\|_D^2];$
- (ii) the point  $x$  is stationary for problem (1) if, and only if,  $x \in \mathcal{P}_{C,\zeta}^D(x, z(\alpha));$
- (iii) if  $x \in C$  is a nonstationary point for problem (1), then  $\langle \nabla f(x), w(\alpha) - x \rangle < 0$ .  
Equivalently, if there exists  $\bar{\alpha} > 0$  such that  $\langle \nabla f(x), w(\bar{\alpha}) - x \rangle \geq 0$ , then  $x$  is stationary for problem (1).



## Inexact scaled gradient method

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## InexProj-SGM employing nonmonotone line search

**Step 0.** Choose  $\sigma, \zeta_{\min} \in (0, 1)$ ,  $0 < \alpha_{\min} \leq \alpha_{\max}$  and  $\mu \geq 1$ . Let  $x^0 \in C$ ,  $\nu_0 \geq 0$  and set  $k \leftarrow 0$ .

**Step 1.** Choose positive real numbers  $\alpha_k$  and  $\zeta_k$  and a positive definite matrix  $D_k$  such that

$$\alpha_{\min} \leq \alpha_k \leq \alpha_{\max}, \quad 0 < \zeta_{\min} < \zeta_k \leq 1, \quad D_k \in \mathcal{D}_{\mu}.$$

Compute  $w^k \in C$  as any feasible inexact projection with respect to the norm  $\|\cdot\|_{D_k}$  of

$$z^k := x^k - \alpha_k D_k^{-1} \nabla f(x^k)$$

onto  $C$  relative to  $x^k$  with forcing parameter  $\zeta_k$ , i.e.,

$$w^k \in \mathcal{P}_{C, \zeta_k}^{D_k}(x^k, z^k).$$

If  $w^k = x^k$ , then **stop** declaring convergence.

## InexProj-SGM employing nonmonotone line search

**Step 2.** Set  $\tau_{\text{trial}} \leftarrow 1$ . If

$$f\left(x^k + \tau_{\text{trial}}(w^k - x^k)\right) \leq f(x^k) + \sigma \tau_{\text{trial}} \left\langle \nabla f(x^k), w^k - x^k \right\rangle + \nu_k, \quad (3)$$

then  $\tau_k \leftarrow \tau_{\text{trial}}$ , define the next iterate  $x^{k+1}$  as

$$x^{k+1} = x^k + \tau_k(w^k - x^k), \quad (4)$$

and go to **Step 3**. Otherwise, choose  $\tau_{\text{new}} \in [\underline{\omega}\tau_{\text{trial}}, \bar{\omega}\tau_{\text{trial}}]$ , set  $\tau_{\text{trial}} \leftarrow \tau_{\text{new}}$ , and repeat test (3).

**Step 3.** Take  $\delta_{k+1} \in [\delta_{\min}, 1]$  and choose  $\nu_{k+1} \in \mathbb{R}$  satisfying

$$0 \leq \nu_{k+1} \leq (1 - \delta_{k+1}) \left[ f(x^k) + \nu_k - f(x^{k+1}) \right].$$

Set  $k \leftarrow k + 1$  and go to **Step 1**.

### Remarks

There are several ways of choosing  $\nu_k$

(i) If  $\nu_k = 0$ , the line search (4) is the well-known Armijo line search.

(ii) If  $f_{\max} = \max\{f(x^{k-j}) \mid 0 \leq j \leq \min\{k, M\}\}$  and

$$\nu_k = f_{\max} - f(x^k) \tag{5}$$

the line search (4) is the same defined by Grippo, Lampariello and Lucidi<sup>4</sup>.

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<sup>4</sup>L. Grippo, F. Lampariello, and S. Lucidi. “A nonmonotone line search technique for Newton’s method”. In: *SIAM J. Numer. Anal.* 23.4 (1986), pp. 707–716.

(iii) Let  $0 \leq \eta_{\min} \leq \eta_{\max} < 1$ ,  $c_0 = f(x_0)$  and  $q_0 = 1$ . Choose  $\eta_k \in [\eta_{\min}, \eta_{\max}]$  and set

$$q_{k+1} = \eta_k q_k + 1, \quad c_{k+1} = (\eta_k q_k c_k + f(x^{k+1}))/q_{k+1}, \quad \forall k \in \mathbb{N}.$$

If  $\delta_{k+1} = 1/q_{k+1}$  and

$$\nu_k = c_k - f(x^k) \tag{6}$$

the line search (4) is the same defined by Zhang and Hager<sup>5</sup>.

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<sup>5</sup>H. Zhang and W. W. Hager. “A nonmonotone line search technique and its application to unconstrained optimization”. In: *SIAM J. Optim.* 14.4 (2004), pp. 1043–1056.





## Partial asymptotic convergence

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### Lemma

*There holds  $0 \leq \delta_{k+1} \left[ f(x^k) + \nu_k - f(x^{k+1}) \right] \leq \left( f(x^k) + \nu_k \right) - \left( f(x^{k+1}) + \nu_{k+1} \right)$ , for all  $k \in \mathbb{N}$ . As consequence the sequence  $\left( f(x^k) + \nu_k \right)_{k \in \mathbb{N}}$  is non-increasing.*

### Theorem

*Assume that  $\lim_{k \rightarrow +\infty} \nu_k = 0$ . Then, Algorithm InexProj-SGM stops in a finite number of iterations at a stationary point of problem (1), or generates an infinite sequence  $(x^k)_{k \in \mathbb{N}}$  for which every cluster point is stationary for problem (1).*

### Proposition

If  $\delta_{min} > 0$ , then  $\sum_{k=0}^{+\infty} \nu_k < +\infty$ . Consequently,  $\lim_{k \rightarrow +\infty} \nu_k = 0$ .<sup>6</sup>

### Remark

Armijo line search and nonmonotone line search strategy defined by (6) satisfies a condition  $\delta_{min} > 0$ . However, for the nonmonotone line search strategy proposed by (5), we can only guarantee that  $\delta_{min} \geq 0$ . Hence, we need deal with this case separately.

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<sup>6</sup>Geovani N. Grapiglia and Ekkehard W. Sachs. “On the worst-case evaluation complexity of non-monotone line search algorithms”. In: *Comput. Optim. Appl.* 68.3 (2017), pp. 555–577.

### Proposition

*Assume that the sequence  $(x^k)_{k \in \mathbb{N}}$  is generated by Algorithm InexProj-SGM with the nonmonotone line search (5), i.e.,  $\nu_k = f_{\max} - f(x^k)$  for all  $k \in \mathbb{N}$ . In addition, assume that the level set  $C_0 := \{x \in C : f(x) \leq f(x^0)\}$  is bounded and  $\nu_0 = 0$ . Then,*

$$\lim_{k \rightarrow +\infty} \nu_k = 0.$$



## Full asymptotic convergence

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## Full asymptotic convergence analysis

We will prove, under suitable assumptions, the full convergence of the sequence  $(x^k)_{k \in \mathbb{N}}$ . For this end, we assume that in **Step 1** of Algorithm InexProj-SGM:

**A1.** For all  $k \in \mathbb{N}$ , we take  $w^k \in \mathcal{R}_{C, \gamma_k}^{D_k}(x^k, z^k)$  with  $\gamma_k = (1 - \zeta_k)/2$ .

**A2.** For all  $k \in \mathbb{N}$ , we take  $0 \leq \nu_k$  such that  $\sum_{k=0}^{+\infty} \nu_k < +\infty$ .

Armijo line search and nonmonotone line search strategies defined by (6) satisfies the assumption **A2**.

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## Lemma

*For each  $x \in C$ , there holds*

$$\|x^{k+1} - x\|_{D_k}^2 \leq \|x^k - x\|_{D_k}^2 + 2\alpha_k \tau_k \left\langle \nabla f(x^k), x - x^k \right\rangle + \xi \left[ f(x^k) - f(x^{k+1}) + \nu_k \right], \quad \forall k \in \mathbb{N}.$$

*where  $\xi := \frac{2\alpha_{\max}}{\sigma} > 0$ .*

## Corollary

*Assume that  $f$  is a convex function. If  $U := \{x \in C : f(x) \leq \inf_{k \in \mathbb{N}} (f(x^k) + \nu_k)\}$  is not empty, then  $(x^k)_{k \in \mathbb{N}}$  converges to a stationary point of problem (1).*

### Theorem

*If  $f$  is a convex function and  $(x^k)_{k \in \mathbb{N}}$  has no cluster points, then  $\Omega^* = \emptyset$ ,  $\lim_{k \rightarrow \infty} \|x^k\| = +\infty$ , and  $\inf_{k \in \mathbb{N}} f(x^k) = \inf\{f(x) : x \in C\}$ .*



### Corollary

*If  $f$  is a convex function and  $(x^k)_{k \in \mathbb{N}}$  has at least one cluster point, then  $(x^k)_{k \in \mathbb{N}}$  converges to a stationary point of problem (1).*

### Theorem

*Assume that  $f$  is a convex function and  $\Omega^* \neq \emptyset$ . Then,  $(x^k)_{k \in \mathbb{N}}$  converge to an optimal solution of problem (1).*



## Iteration-complexity bound

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## Iteration-complexity bound

Besides assuming that in **Step 1** of Algorithm InexProj-SGM we take  $(x^k)_{k \in \mathbb{N}}$  satisfying **A1** and **A2**, we also need the following assumption.

**A3.** The gradient  $\nabla f$  of  $f$  is Lipschitz continuous with constant  $L > 0$ .

### Lemma

*The stepsize  $\tau_k$  in Algorithm InexProj-SGM satisfies  $\tau_k \geq \tau_{\min}$ ,*

where

$$\tau_{\min} := \min \left\{ 1, \frac{\tau(1 - \sigma)}{\alpha_{\max} \mu L} \right\}.$$

## Iteration-complexity bound

### Theorem

For every  $N \in \mathbb{N}$ , the following inequality holds

$$\min \left\{ \|w^k - x^k\| : k = 0, 1, \dots, N-1 \right\} \leq \sqrt{\frac{2\alpha_{\max}\mu(f(x^0) - f^* + \sum_{k=0}^{\infty} \nu_k)}{\sigma\tau_{\min}}} \frac{1}{\sqrt{N}}.$$

### Theorem

Let  $f$  be a convex function on  $C$ . Then, for every  $N \in \mathbb{N}$ , there holds

$$\min \left\{ f(x^k) - f^* : k = 0, 1, \dots, N-1 \right\} \leq \frac{\|x^0 - x^*\|_{D_0}^2 + \xi [f(x^0) - f^* + \sum_{k=0}^{\infty} \nu_k]}{2\alpha_{\min}\tau_{\min}} \frac{1}{N}.$$

### Lemma

Let  $N_k$  be the number of function evaluations after  $k \geq 1$  iterations of Algorithm InexProj-SGM. Then,

$$N_k \leq 1 + (k + 1) \left[ \frac{\log(\tau_{\min})}{\log(\tau)} + 1 \right].$$

### Theorem

*For a given  $\epsilon > 0$ , the number of function evaluations in Algorithm InexProj-SGM are at most*

$$1 + \left( \frac{2\alpha_{\max}\mu \left( f(x^0) - f^* + \sum_{k=0}^{\infty} \nu_k \right)}{\sigma\tau_{\min}} \frac{1}{\epsilon^2} + 1 \right) \left( \frac{\log(\tau_{\min})}{\log(\tau)} + 1 \right),$$

*to compute  $x^k$  and  $w^k$  such that  $\|w^k - x^k\| \leq \epsilon$ .*

### Theorem

Let  $f$  be a convex function on  $C$ . For a given  $\epsilon > 0$ , the number of function evaluations in Algorithm InexProj-SGM are at most

$$1 + \left( \frac{\|x^0 - x^*\|_{D_0}^2 + \xi (f(x^0) - f^* + \sum_{k=0}^{\infty} \nu_k)}{2\alpha_{\min} \tau_{\min}} \frac{1}{\epsilon} + 1 \right) \left( \frac{\log(\tau_{\min})}{\log(\tau)} + 1 \right),$$

to compute  $x^k$  such that  $f(x^k) - f^* \leq \epsilon$ .





## Numerical experiments

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## Numerical experiments

Given  $A$  and  $B$  two  $m \times n$  matrices, with  $m \geq n$ , and  $c \in \mathbb{R}$ , we consider the matrix function  $f : \mathbb{R}^{n \times n} \rightarrow \mathbb{R}$  given by:

$$f(X) := \frac{1}{2} \|AX - B\|_F^2 + \sum_{i=1}^{n-1} \left[ c (X_{i+1,i+1} - X_{i,i}^2)^2 + (1 - X_{i,i})^2 \right],$$

which combines a least squares term with a Rosenbrock-type function.  $X_{i,j}$  stands for the  $ij$ -element of the matrix  $X$  and  $\|\cdot\|_F$  denotes the Frobenius matrix norm, i.e.,  $\|A\|_F := \sqrt{\langle A, A \rangle}$  where the inner product is given by  $\langle A, B \rangle = \text{tr}(A^T B)$ .

**Problem I<sup>7</sup>:**

$$\begin{array}{ll}\min & f(X) \\ \text{s.t.} & X \in SDD^+, \\ & L \leq X \leq U,\end{array}$$

where  $SDD^+$  is the cone of symmetric and diagonally dominant real matrices with positive diagonal, i.e.,

$$SDD^+ := \{X \in \mathbb{R}^{n \times n} \mid X = X^T, X_{i,i} \geq \sum_{j \neq i} |X_{i,j}| \forall i\},$$

$L$  and  $U$  are given  $n \times n$  matrices, and  $L \leq X \leq U$  means that  $L_{i,j} \leq X_{i,j} \leq U_{i,j}$  for all  $i, j$ .

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<sup>7</sup>Ernesto G. Birgin, José Mario Martínez, and Marcos Raydan. “Inexact spectral projected gradient methods on convex sets”. In: *IMA J. Numer. Anal.* 23.4 (2003), pp. 539–559.

Problem II<sup>89</sup>:

$$\begin{array}{ll}\min & f(X) \\ \text{s.t.} & X \in \mathbb{S}_+^n, \\ & \text{tr}(X) = 1,\end{array}$$

where  $\mathbb{S}_+^n$  is the cone of symmetric and positive semidefinite real matrices. The feasible set of Problem II was known as *spectrahedron* and appears in several interesting applications.

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<sup>8</sup>Zeyuan Allen-Zhu et al. “Linear convergence of a Frank-Wolfe type algorithm over trace-norm balls”. In: *Advances in Neural Information Processing Systems*. 2017, pp. 6191–6200.

<sup>9</sup>D.S. Gonçalves, M.A. Gomes-Ruggiero, and C. Lavor. “A projected gradient method for optimization over density matrices”. In: *Optimization Methods and Software* 31.2 (2016), pp. 328–341.

## Numerical experiments

We are interested in the **spectral gradient version of the SPG method**, so we set  $D_k := I$  for all  $k$ ,  $\alpha_0 := \min(\alpha_{\max}, \max(\alpha_{\min}, 1/\|\nabla f(x^0)\|))$  and, for  $k > 0$ ,

$$\alpha_k := \begin{cases} \min(\alpha_{\max}, \max(\alpha_{\min}, \langle s^k, s^k \rangle / \langle s^k, y^k \rangle)), & \text{if } \langle s^k, y^k \rangle > 0 \\ \alpha_{\max}, & \text{otherwise,} \end{cases}$$

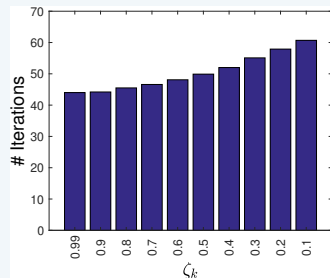
where  $s^k := X^k - X^{k-1}$ ,  $y^k := \nabla f(X^k) - \nabla f(X^{k-1})$ ,  $\alpha_{\min} = 10^{-10}$ , and  $\alpha_{\max} = 10^{10}$ .

Concerning the stopping criterion, all runs were stopped at an iterate  $X^k$  declaring convergence if

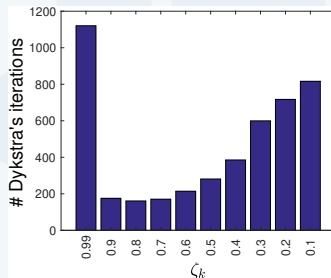
$$\max_{i,j} (|X_{i,j}^k - W_{i,j}^k|) \leq 10^{-6},$$

where  $W^k \in \mathcal{P}_{C, \zeta_k}^{D_k}(x^k, z^k)$ .

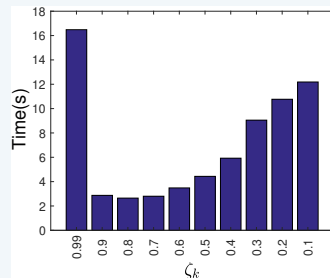
# Influence of the inexact projection



(a)



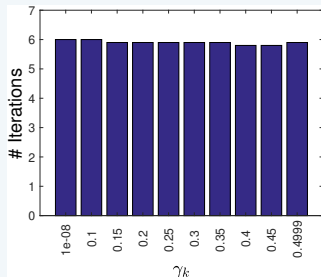
(b)



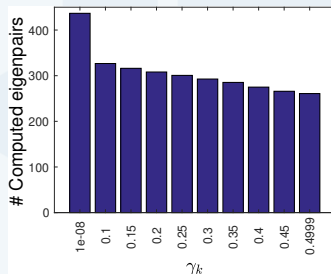
(c)

**Figure 1:** Results for 10 instances of Problem I using  $n = 100$ ,  $m = 200$ , and  $c = 10$ . Average number of: (a) iterations; (b) Dykstra's iterations; (c) CPU time in seconds needed to reach the solution for different choices of  $\zeta_k$ .

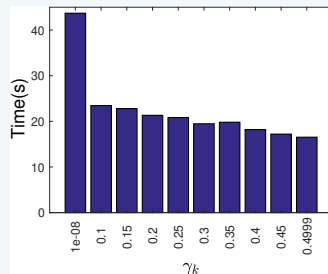
## Influence of the inexact projection



(a)



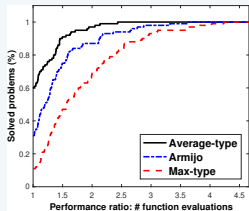
(b)



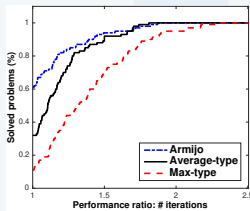
(c)

**Figure 2:** Results for 10 instances of Problem II using  $n = 800$ ,  $m = 1000$ , and  $c = 100$ . Average number of: (a) iterations; (b) computed eigenpairs; (c) CPU time in seconds needed to reach the solution for different choices of  $\gamma_k$ .

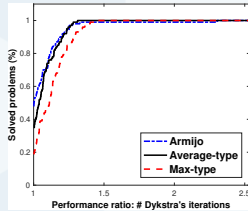
# Influence of the line search scheme



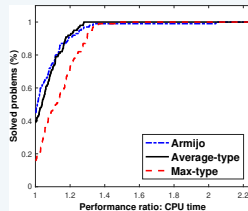
(a)



(b)



(c)

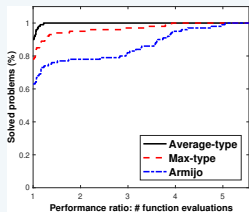


(d)

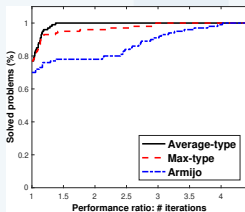
**Figure 3:** Performance profiles for Problem I considering the SPG method with the Armijo, the Average-type, and the Max-type line searches strategies using as performance measurement: (a) number of function evaluations; (b) number of (outer) iterations; (c) number of Dykstra's iterations; (d) CPU time.



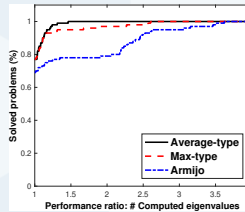
# Influence of the line search scheme



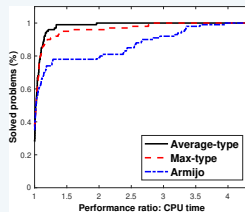
(a)



(b)



(c)



(d)

**Figure 4:** Performance profiles for Problem II considering the SPG method with the Armijo, the Average-type, and the Max-type line searches strategies using as performance measurement: (a) number of function evaluations; (b) number of (outer) iterations; (c) number of computed eigenpairs; (d) CPU time.



Thank you!

UFG