```
data FBAE where
       Num :: Int -> FBAE
       Plus :: FBAE -> FBAE -> FBAE
       Minus :: FBAE -> FBAE -> FBAE
       Bind :: String -> FBAE -> FBAE -> FBAE
       Id :: String -> FBAE
       App :: FBAE -> FBAE -> FBAE
       Lambda :: String -> FBAE -> FBAE
       Closure :: String -> FBAE -> Env -> FBAE
t ::= Num
       | True
       | False
       | t+t
       | t-t
       | bind id = t in t
       l id
       | if t then t else t
       | t <= t
       | t && t
       | isZero t
v ::= Num
       | True
       | False
eval :: AE -> Maybe Int
eval e (Lambda i b) = return (ClosureV i b e)
eval e (App f a) = do { (ClosureV i b j) <- eval e f;
               v <- eval e a;
               eval (i,v):j b}
elab (BindX i v b) = (App (Lambda i (elab b)) (elab v))
typeof :: ABE -> Maybe TABE
typeof (Plus I r) = do { TNum <- typeof I;
                       TNum <- typeof r;
                       return TNum}
```

```
First order: not in arguments. Not in returns
Higher order: arguments, returns,
First class: + function pointers
Lambda expression - function value
Application - call lambda with actuals
```

**Domain** - types of formals **Range** - type of return

```
bind inc = (lambda x in x+1) in (inc)(1)
```

**Dynamic scope** - depends on when evaluated **Static scope** - depends on where defined (closures)

Enrichment - adding features to a new language

Extension - adding new features by modifying the evaluation function

Derived form - defines new features in terms of existing expressions

Elaboration - translating one abstract syntax into another

External language - language that is input to an elaborator (bind)

Internal language - language that is output by an elaborator (lambda)

```
Dynamic:
```

```
bind f = lambda x in if x=0 then 1 else x^*(f)(x-1) in (f)(2)
Static:
```

eval e (Fix f) = do  $\{(ClosureV i b e') <- eval e f;$ 

eval e' (subst i (Fix (Lambda i b)) b)}

bind f = fix(lambda g in (lambda x in if x=0 then 1 else x\*(g)(x-1)))