

Probabilistic Programming Languages

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MPRI 2023-2024

Reminders

Probabilistic Programming Languages

Probabilistic programming

Programming and reasoning with uncertainty

- Sample from probability distributions
- Condition on observed data

Bayesian Inference: learn parameters from data

- Latent parameter θ
- Observed data x_1, \dots, x_n

$$\underbrace{p(\theta \mid x_1, \dots, x_n)}_{\text{posterior}} = \frac{p(\theta) p(x_1, \dots, x_n \mid \theta)}{p(x_1, \dots, x_n)} \quad (\text{Bayes' theorem})$$
$$\propto \underbrace{p(\theta)}_{\text{prior}} \underbrace{p(x_1, \dots, x_n \mid \theta)}_{\text{likelihood}} \quad (\text{Data are constants})$$



Thomas Bayes (1701-1761)

Example: Coin



Consider a series of coin tosses

- Observations: head or tail
- Each toss is independant
- What is the probability of getting head at the next toss?

Probabilistic model

- Prior: $z \sim \text{Uniform}(0, 1)$
- Observations: for $i \in [1, n]$, $x_i \sim \text{Bernoulli}(z)$
- Posterior: $p(z \mid x_1, \dots, x_n)$?

Example: Coin



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- Observations: for $i \in [1, n]$, $x_i \sim \text{Bernoulli}(z)$
- Posterior: $p(z \mid x_1, \dots, x_n)$?

$$\begin{aligned} p(z \mid x_1, \dots, x_n) &= \frac{p(x_1, \dots, x_n \mid z)p(z)}{p(x_1, \dots, x_n)} \\ &= \frac{p(x_1, \dots, x_n \mid z)p(z)}{\int_z p(x_1, \dots, x_n \mid z)} \end{aligned}$$

$$\begin{aligned} p(x_1, \dots, x_n \mid z) &= \prod_{i=1}^n p(x_i \mid z) \\ &= \prod_{i=1}^n z^{x_i} (1 - z)^{1-x_i} \\ &= z^{\sum_{i=1}^n x_i} (1 - z)^{\sum_{i=1}^n (1-x_i)} \\ &= z^{\text{\#heads}} (1 - z)^{\text{\#tails}} \end{aligned}$$

$$\begin{aligned} p(z \mid x_1, \dots, x_n) &= \frac{z^{\text{\#heads}} (1 - z)^{\text{\#tails}}}{\int_z z^{\text{\#heads}} (1 - z)^{\text{\#tails}}} \\ &= \frac{z^{\text{\#heads}} (1 - z)^{\text{\#tails}}}{B(\text{\#heads} + 1, \text{\#tails} + 1)} \\ &= \text{pdf}(\text{Beta}(\text{\#heads} + 1, \text{\#tails} + 1)) \end{aligned}$$

Example: Coin



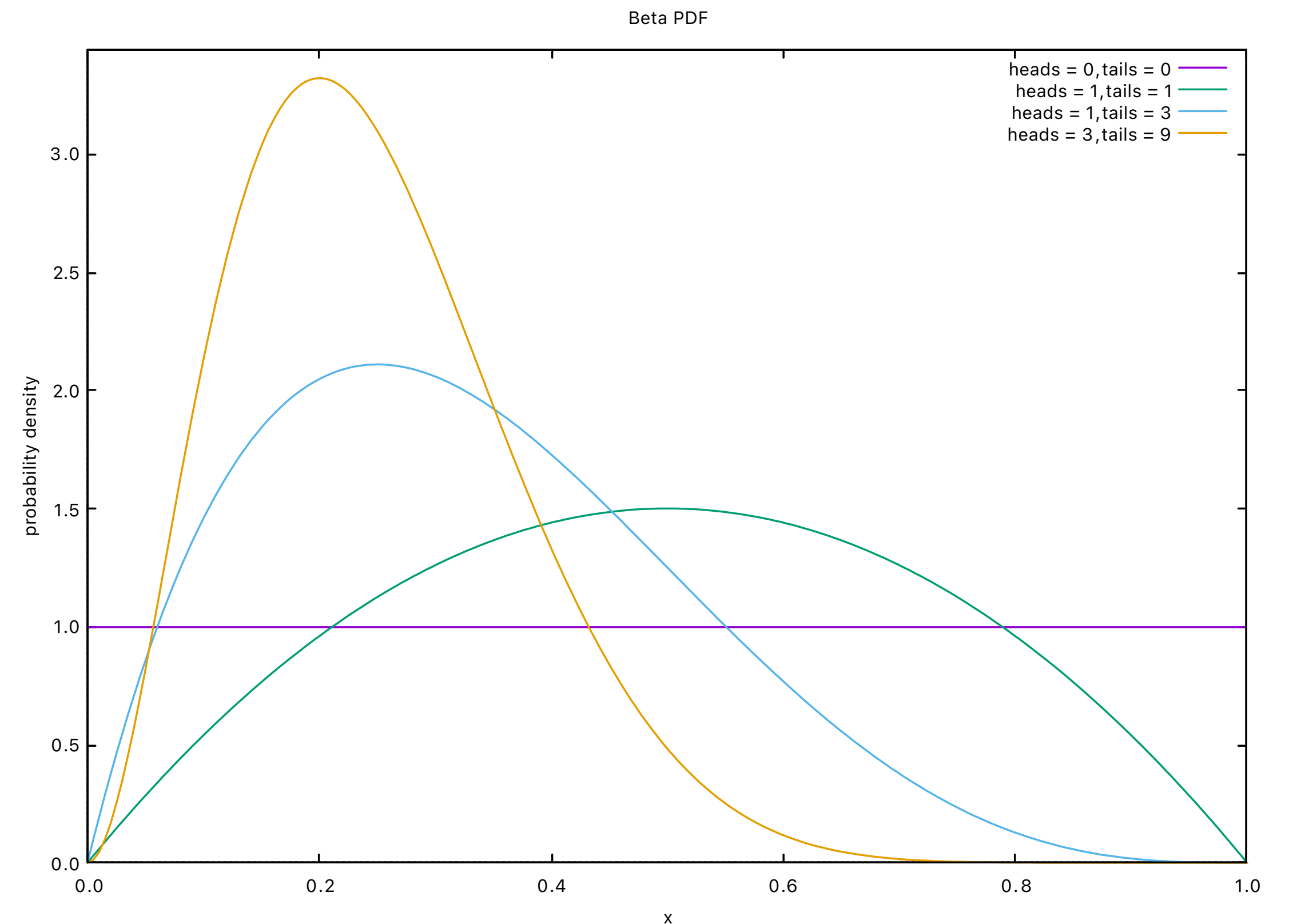
Consider a series of coin tosses

- Observations: head or tail
- Each toss is independant
- What is the probability of getting head at the next toss?

Probabilistic model

- Prior: $z \sim \text{Uniform}(0, 1)$
- Observations: for $i \in [1, n]$, $x_i \sim \text{Bernoulli}(z)$
- Posterior: $p(z \mid x_1, \dots, x_n)$?

$$z \sim \text{Beta}(\text{\#heads} + 1, \text{\#tails} + 1)$$



Example: Coin



Consider a series of coin tosses

- Observations: head or tail
- Each toss is independant
- What is the probability of getting head at the next toss?

Probabilistic model

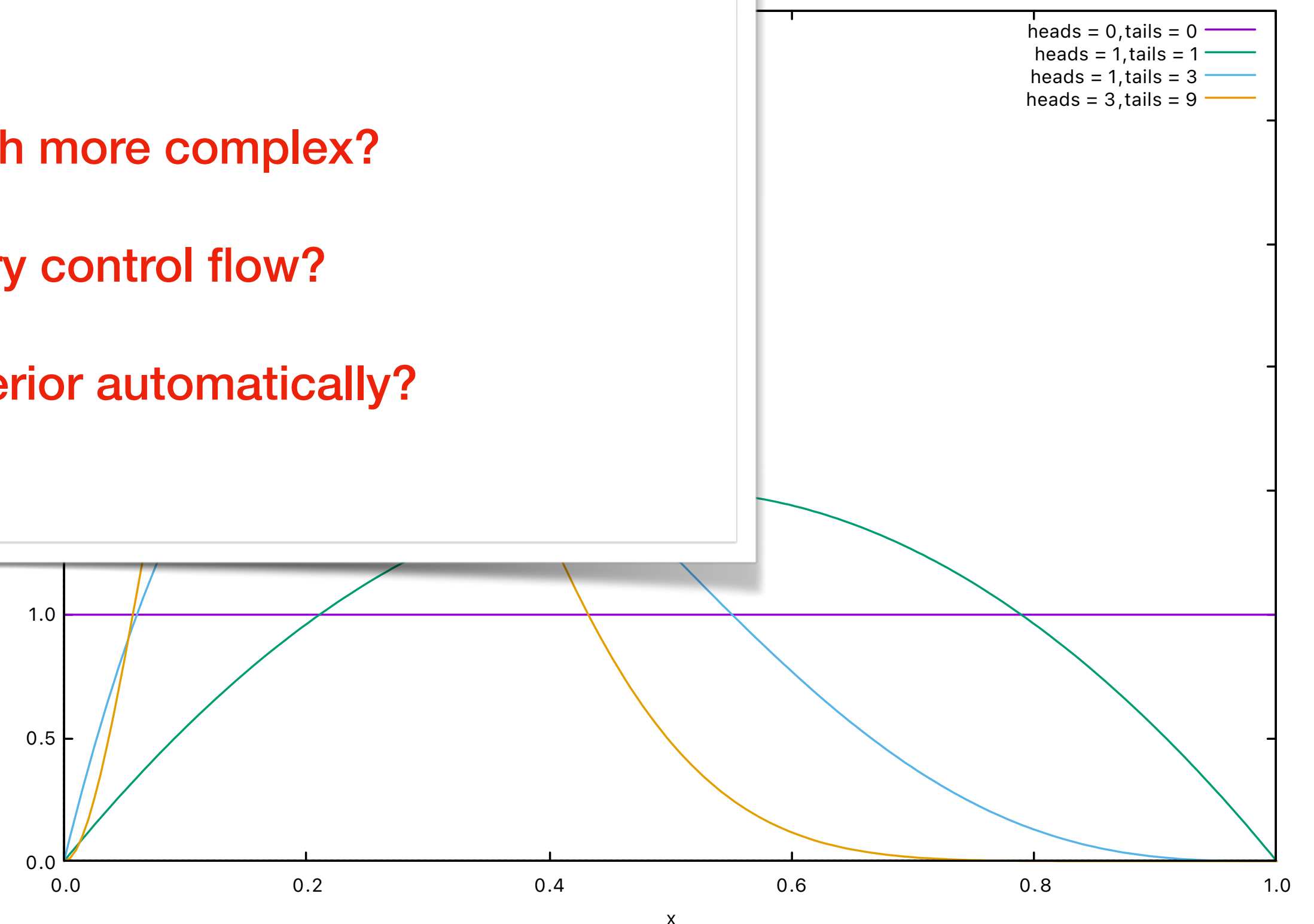
- Prior: $z \sim \text{Uniform}()$
- Observations: for i
- Posterior: $p(z \mid x_1, \dots, x_i)$

$$z \sim \text{Beta}(\text{\#heads} + 1, \text{\#tails} + 1)$$

What if the model is much more complex?

What if we use arbitrary control flow?

Can we compute the posterior automatically?



Probabilistic programming languages

General purpose programming languages extended with probabilistic constructs

- **sample**: draw a sample from a distribution
- **assume**, **factor**, **observe**: condition the model on inputs (e.g., observed data)
- **infer**: compute the posterior distribution of a model given the inputs

Multiple examples:

- Church, Anglican (lisp, clojure), 2008
- WebPPL (javascript), 2014
- Pyro/NumPyro (python), 2017/2019
- Gen (julia), 2018
- ProbZelus (Zelus), 2019
- ...

More and more, incorporating new ideas:

- New inference techniques, e.g., stochastic variational inference (SVI)
- Interaction with neural nets (deep probabilistic programming)

Outline

For a given inference algorithm, how to implement `sample`, `assume`, `factor`, `observe`, and `infer`?

I - Basic inference

- Rejection sampling
- Importance sampling

II - Typing

- Syntax: language and types
- Deterministic vs. probabilistic
- Guarding probabilistic constructs

II - Kernel Semantics

- Reminders: measure theory
- Types as measurable spaces
- Expressions as measures

TP : A short introduction to Stan

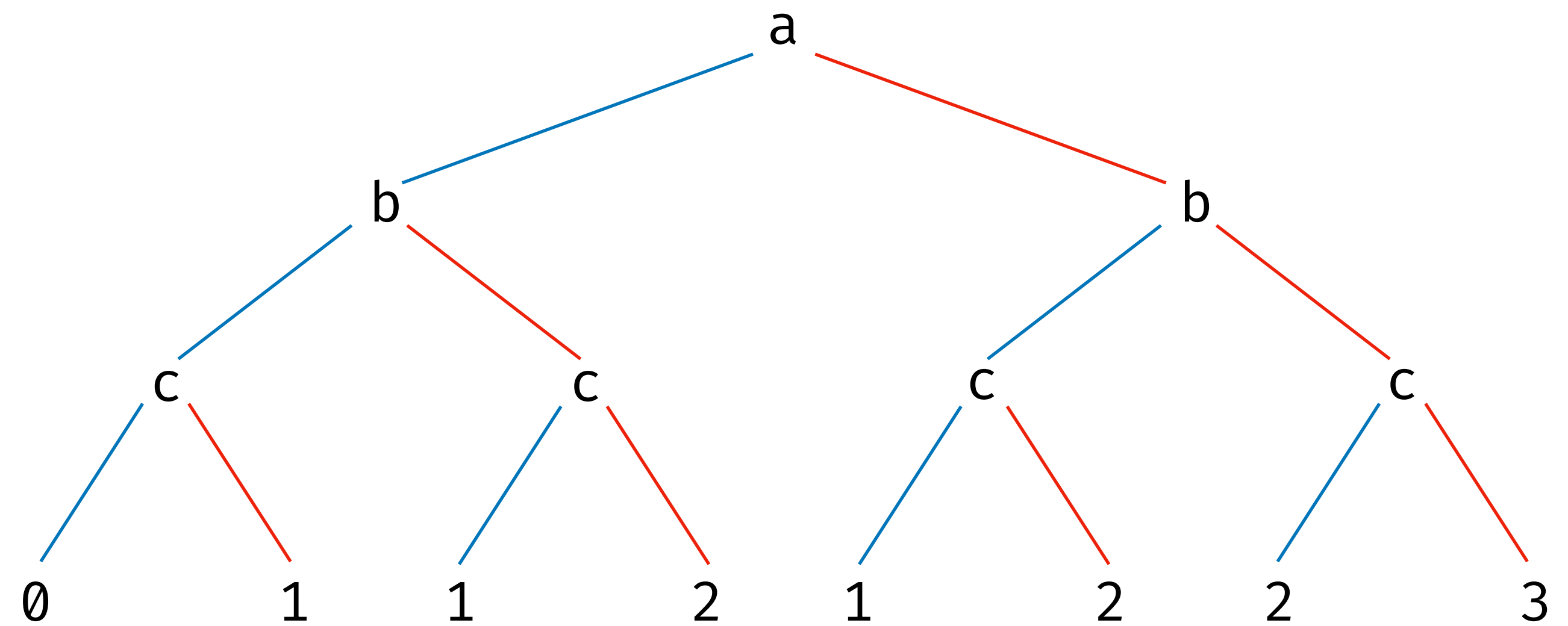
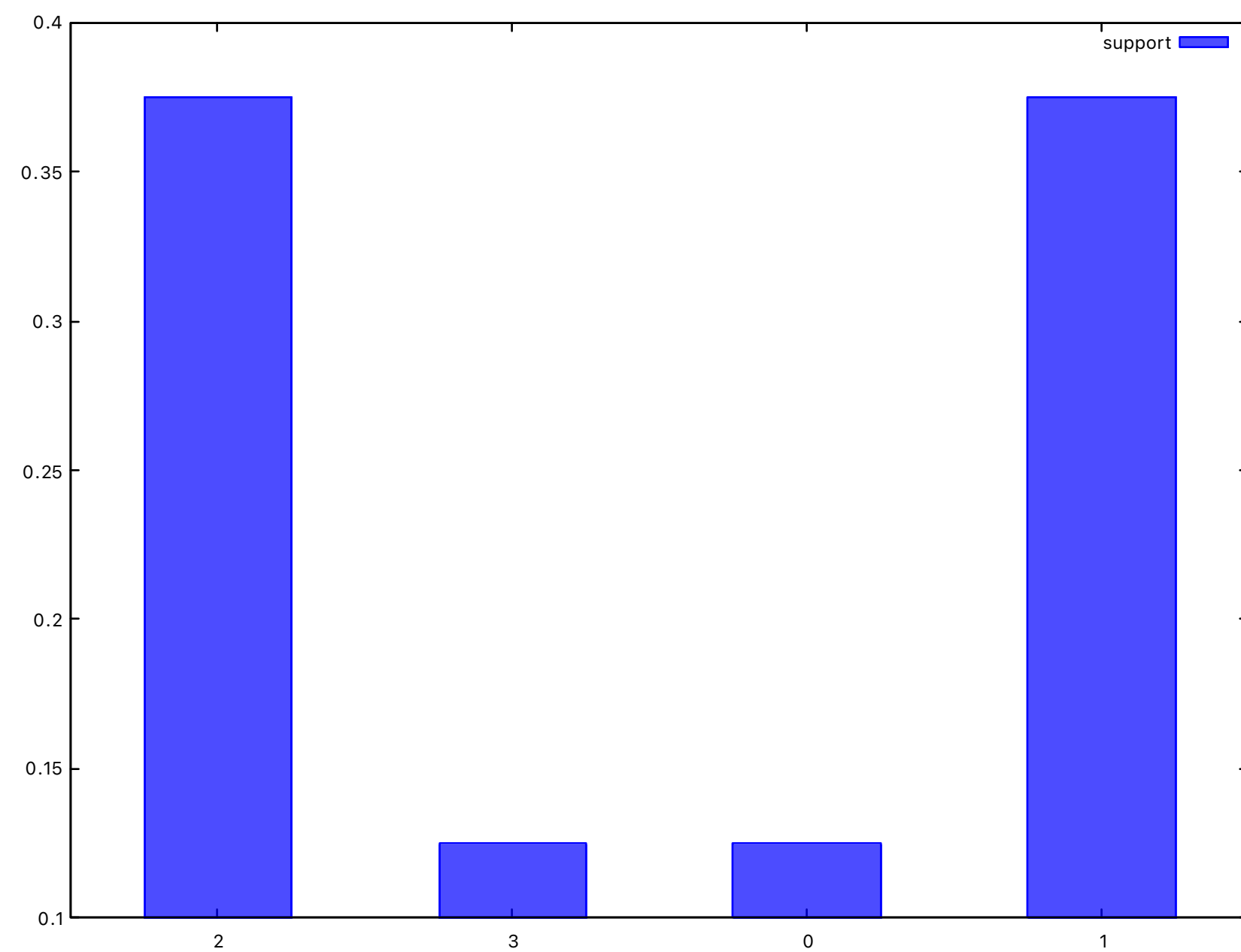
Rejection Sampling

Probabilistic Programming Languages

Example: Funny Bernoulli

funny_bernoulli.ml

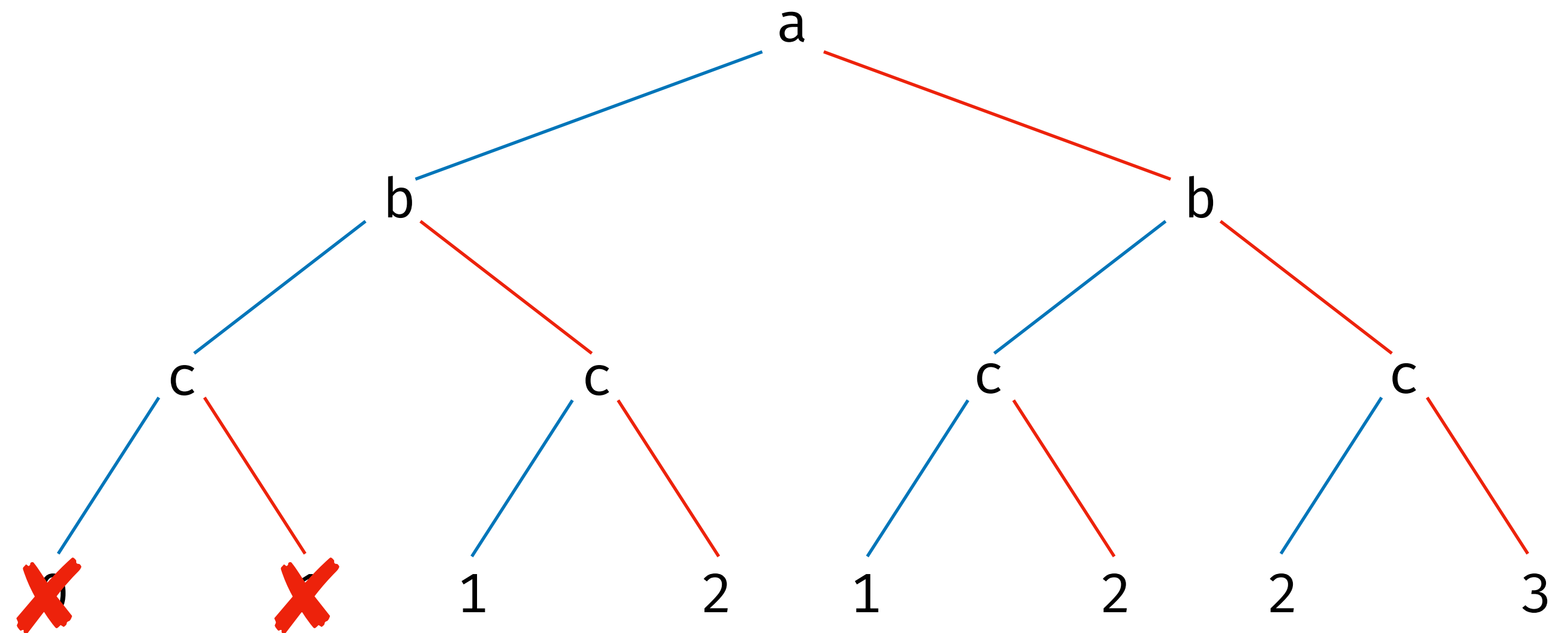
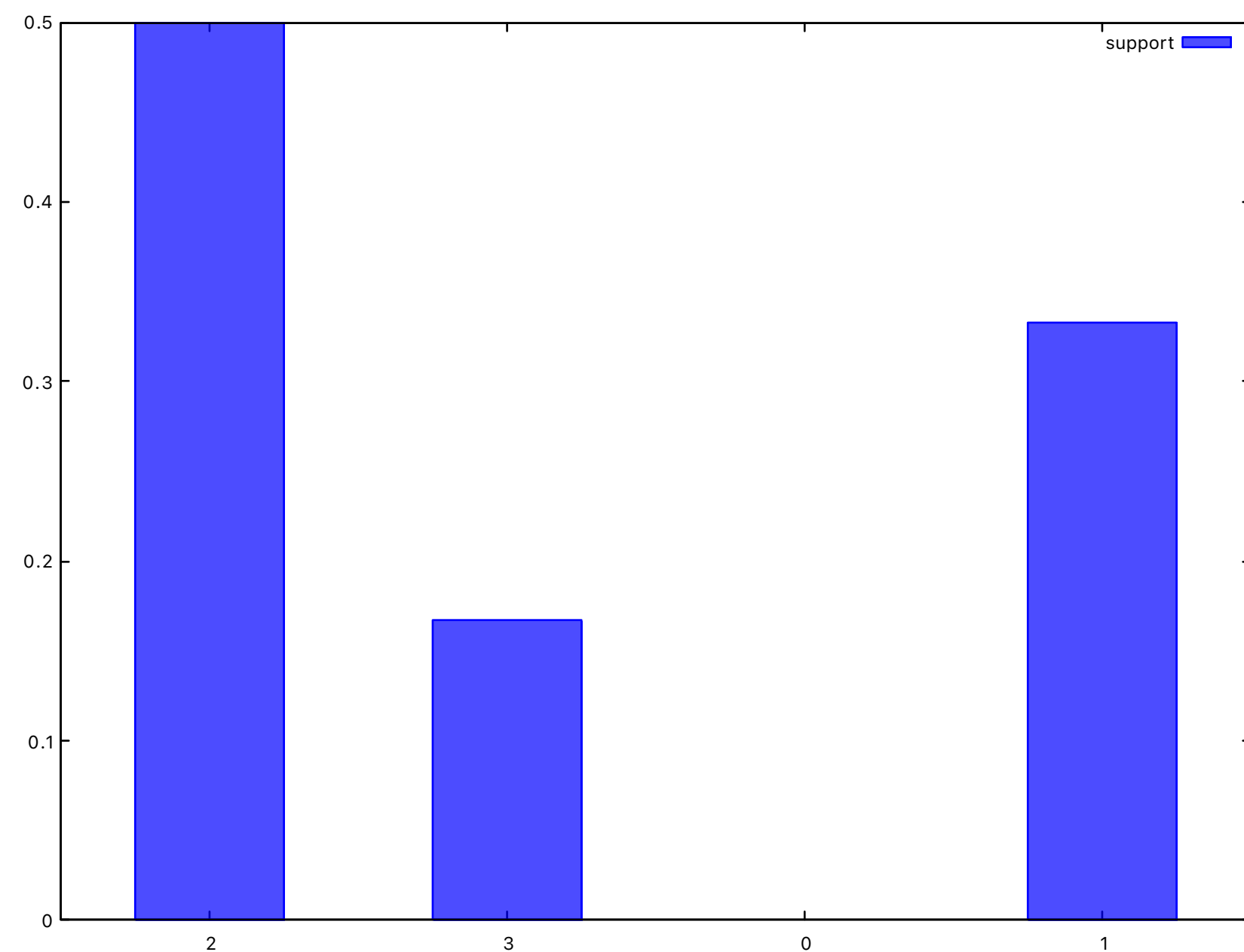
```
let funny_bernoulli () =  
  let a = sample (bernoulli ~p:0.5) in  
  let b = sample (bernoulli ~p:0.5) in  
  let c = sample (bernoulli ~p:0.5) in  
  a + b + c
```



Example: Funny Bernoulli

funny_bernoulli.ml

```
let funny_bernoulli () =  
  let a = sample (bernoulli ~p:0.5) in  
  let b = sample (bernoulli ~p:0.5) in  
  let c = sample (bernoulli ~p:0.5) in  
  let () = assume (a = 1 || b = 1) in  
  a + b + c
```



Rejection sampling

basic.ml

```
module Rejection_sampling : sig
  val sample : 'a Distribution.t → 'a
  val assume : bool → unit
  val infer : ?n:int → ('a → 'b) → 'a → 'b Distribution.t
end = struct ... end
```

Inference algorithm

- Run the model to get a sample
- **sample** : draw a value from a distribution
- **assume** : accept / reject a sample
- If a sample is rejected, re-run the model to get another sample

Rejection sampling

basic.ml

```
module Rejection_sampling = struct

  let sample d = assert false
  let assume p = assert false

  let infer ?(n = 1000) model obs = assert false
end
```


Rejection sampling

basic.ml

```
module Rejection_sampling = struct
  exception Reject

  let sample d = Distribution.draw d
  let assume p = if not p then raise Reject

  let infer ?(n = 1000) model obs =
    let rec exec i = try model obs with Reject → exec i in
    let values = Array.init n exec in
    Distribution.uniform_support ~values
end
```

The type `prob` trick

```
module Rejection_sampling : sig
  type prob
  val sample : prob → 'a Distribution.t → 'a
  val assume : prob → bool → unit
  val infer : ?n:int → (prob → 'a → 'b) → 'a → 'b Distribution.t
end = struct ... end
```

Forbid the use of probabilistic construct outside a model

- Define a simple abstract type `prob`
- Probabilistic constructs and models all require an argument of type `prob`
- Such a value can only be build by `infer`

Rejection sampling

basic.ml

```
module Rejection_sampling = struct
  type prob = Prob

  exception Reject

  let sample _prob d = Distribution.draw d
  let assume _prob p = if not p then raise Reject

  let infer ?(n = 1000) model obs =
    let rec exec i = try model Prob obs with Reject → exec i in
    let values = Array.init n exec in
    Distribution.uniform_support ~values
end
```

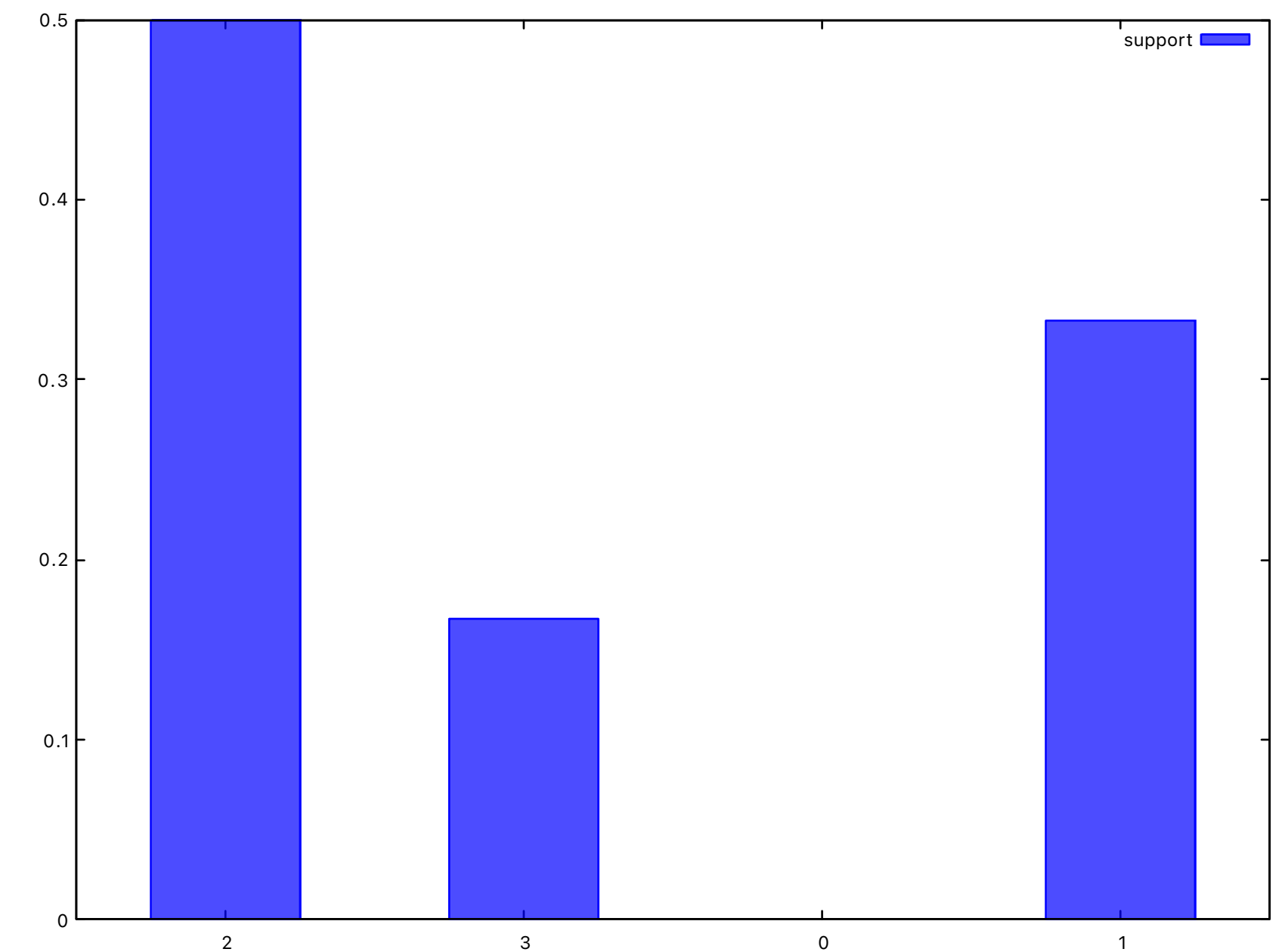
Example: Funny Bernoulli

funny_bernoulli.ml

```
open Byoppl
open Distribution
open Basic.Rejection_sampling
```

```
let funny_bernoulli prob () =
  let a = sample prob (bernoulli ~p:0.5) in
  let b = sample prob (bernoulli ~p:0.5) in
  let c = sample prob (bernoulli ~p:0.5) in
  let () = assume prob (a = 1 || b = 1) in
  a + b + c
```

```
let _ =
  let dist = infer funny_bernoulli () in
  let { values; probs; _ } = get_support ~shrink:true dist in
  Array.iteri (fun i x → Format.printf "%d %f@" x probs.(i)) values
```



› dune exec ./examples/funny_bernoulli.exe

Example: Coin

coin.ml

```
open Basic.Rejection_sampling
```

```
let coin prob x =  
  let z = sample prob (uniform ~a:0. ~b:1.) in  
  let () = List.iter (observe prob (bernoulli ~p:z)) x in  
  z
```

```
let _ =  
  let dist = infer coin [1; 1; 0; 0; 0; 0; 0; 0; 0; 0] in  
  let m, s = Distribution.stats dist in  
  Format.printf "Coin bias, mean:%f std:%f@." m s
```

```
> dune exec ./examples/coin.exe
```

```
Coin bias, mean:0.246161, std:0.119687
```

Example: Coin

coin.ml

```
open Basic.Rejection_sampling
```

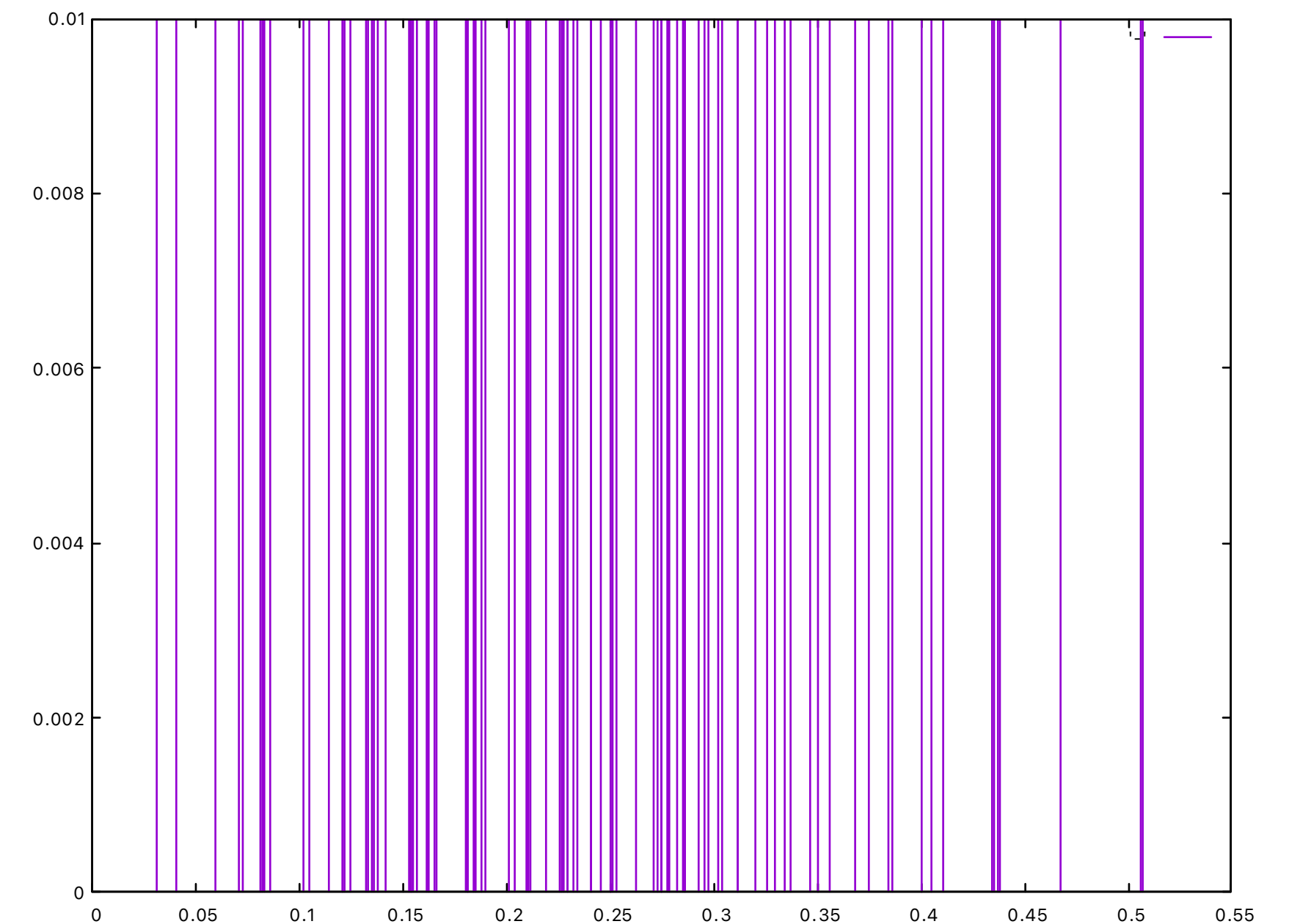
```
let coin prob x =  
  let z = sample prob (uniform ~a:0. ~b:1.) in  
  let () = List.iter (observe prob (bernoulli ~p:z)) x in  
  z
```

```
let _ =  
  let dist = infer coin [1; 1; 0; 0; 0; 0; 0; 0; 0; 0] in  
  let m, s = Distribution.stats dist in  
  Format.printf "Coin bias, mean:%f std:%f@." m s
```

```
> dune exec ./examples/coin.exe
```

```
Coin bias, mean:0.246161, std:0.119687
```

100 particles



Example: Coin

coin.ml

```
open Basic.Rejection_sampling
```

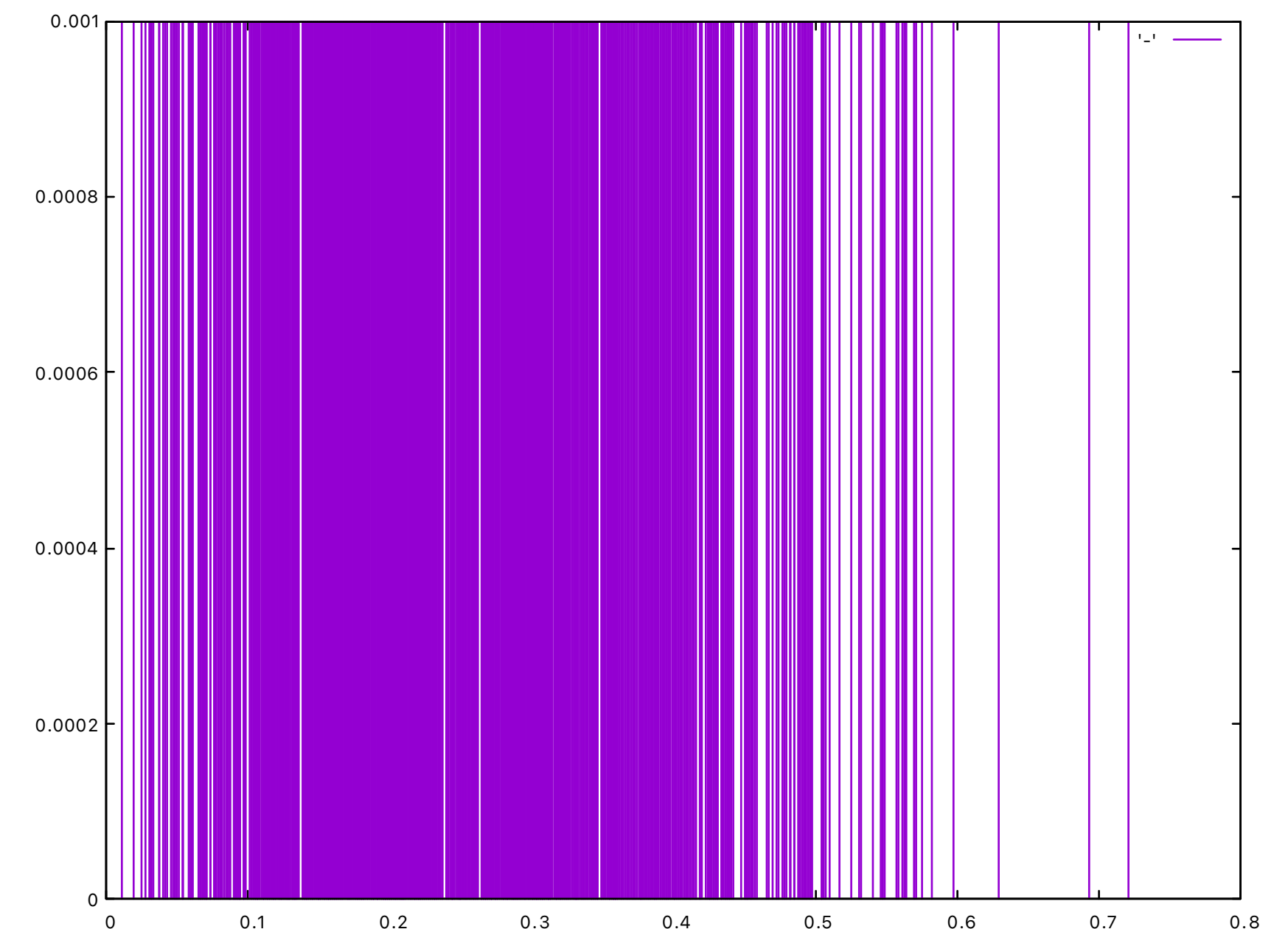
```
let coin prob x =  
  let z = sample prob (uniform ~a:0. ~b:1.) in  
  let () = List.iter (observe prob (bernoulli ~p:z)) x in  
  z
```

```
let _ =  
  let dist = infer coin [1; 1; 0; 0; 0; 0; 0; 0; 0; 0] in  
  let m, s = Distribution.stats dist in  
  Format.printf "Coin bias, mean:%f std:%f@." m s
```

```
> dune exec ./examples/coin.exe
```

```
Coin bias, mean:0.246161, std:0.119687
```

1000 particles



Example: Coin

coin.ml

```
open Basic.Rejection_sampling
```

```
let coin prob x =  
  let z = sample prob (uniform ~a:0. ~b:1.) in  
  let () = List.iter (observe prob (bernoulli ~p:z)) x in  
  z
```

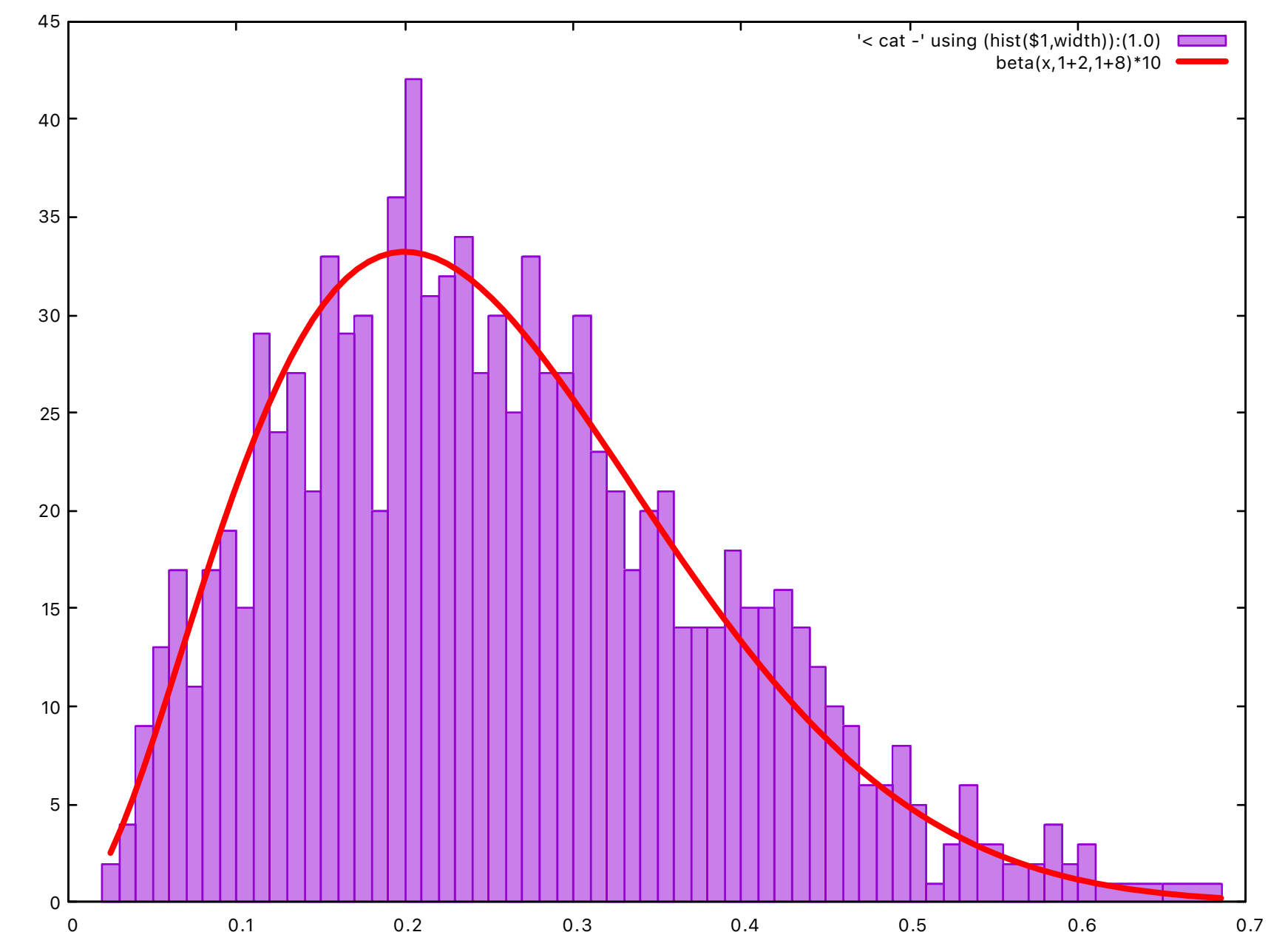
```
let _ =  
  let dist = infer coin [1; 1; 0; 0; 0; 0; 0; 0; 0; 0] in  
  let m, s = Distribution.stats dist in  
  Format.printf "Coin bias, mean:%f std:%f@." m s
```

```
> dune exec ./examples/coin.exe
```

```
Coin bias, mean:0.246161, std:0.119687
```

Slow!

1000 particles



Example: Laplace and gender bias

laplace.ml

```
open Basic.Rejection_sampling
```

```
let laplace prob () =
```

```
  let p = sample prob (uniform ~a:0. ~b:1.) in
```

```
  let g = sample prob (binomial ~p ~n:493_472) in
```

```
  let () = assume prob (g = 241_945) in
```

```
  p
```



```
let () = observe prob
```

```
(binomial ~p ~n:493_472) 241_945
```

```
let _ =
```

```
  let dist = infer ~n:1000 laplace () in
```

```
  let m, s = Distribution.stats dist in
```

```
  Format.printf "Gender bias, mean:%f std:%f@." m s
```

```
> dune exec ./examples/laplace.exe
```

Never terminate!

Importance Sampling

Probabilistic Programming Languages

Importance sampling

basic.ml

```
module Importance_sampling : sig
  type prob
  val sample : prob → 'a Distribution.t → 'a
  val factor : prob → float → unit
  val infer : ?n:int → (prob → 'a → 'b) → 'a → 'b Distribution.t
end = struct ... end
```

Inference algorithm

- Run a set of n independent executions
- **sample**: draw a sample from a distribution
- **factor**: associate a score to the current execution
- Gather output values and score to approximate the posterior distribution

Likelihood weighting

- **observe** $d\ x := \text{factor} (\text{logpdf } d\ x)$

Importance sampling

basic.ml

```
module Importance_sampling = struct
  type prob = ...

  let sample prob d = assert false
  let factor prob s = assert false
  let observe prob d x = factor prob (Distribution.logpdf d x)

  let infer ?(n = 1000) model obs = assert false
end
```


Importance sampling

basic.ml

```
module Importance_sampling = struct
  type prob = { id : int; scores : float array }

  let sample _prob d = Distribution.draw d
  let factor prob s = prob.scores.(prob.id) ← prob.scores.(prob.id) +. s
  let observe prob d x = factor prob (Distribution.logpdf d x)

  let infer ?(n = 1000) model obs =
    let scores = Array.make n 0. in
    let values = Array.make n (fun i → model { id = i; scores } obs) in
    Distribution.support ~values ~logits:scores
end
```

Example: Coin

coin.ml

```
open Basic.Importance_sampling
```

```
let coin prob x =  
  let z = sample prob (uniform ~a:0. ~b:1.) in  
  let () = List.iter (observe prob (bernoulli ~p:z)) x in  
  z
```

```
let _ =  
  let dist = infer coin [1; 1; 0; 0; 0; 0; 0; 0; 0; 0] in  
  let m, s = Distribution.stats dist in  
  Format.printf "Coin bias, mean:%f, std:%f@." m s
```

```
> dune exec ./examples/coin.exe
```

```
Coin bias, mean:0.247876, std:0.118921  
Beta(2+1, 8+1), mean:0.250000, std:0.120096
```

Example: Coin

coin.ml

```
open Basic.Importance_sampling
```

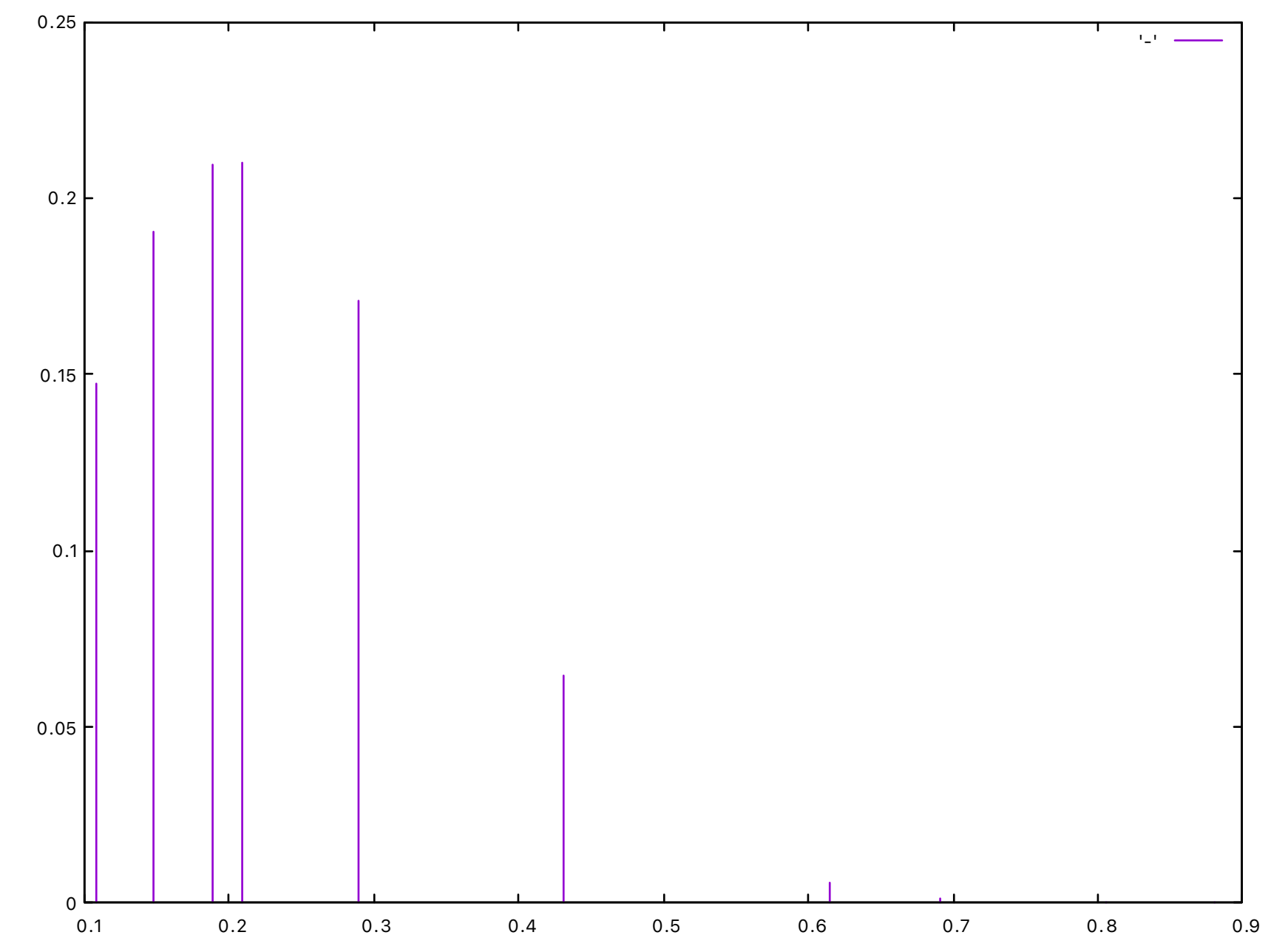
```
let coin prob x =  
  let z = sample prob (uniform ~a:0. ~b:1.) in  
  let () = List.iter (observe prob (bernoulli ~p:z)) x in  
  z
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let _ =  
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```

```
> dune exec ./examples/coin.exe
```

```
Coin bias, mean:0.247876, std:0.118921  
Beta(2+1, 8+1), mean:0.250000, std:0.120096
```

10 particles



Example: Coin

coin.ml

```
open Basic.Importance_sampling
```

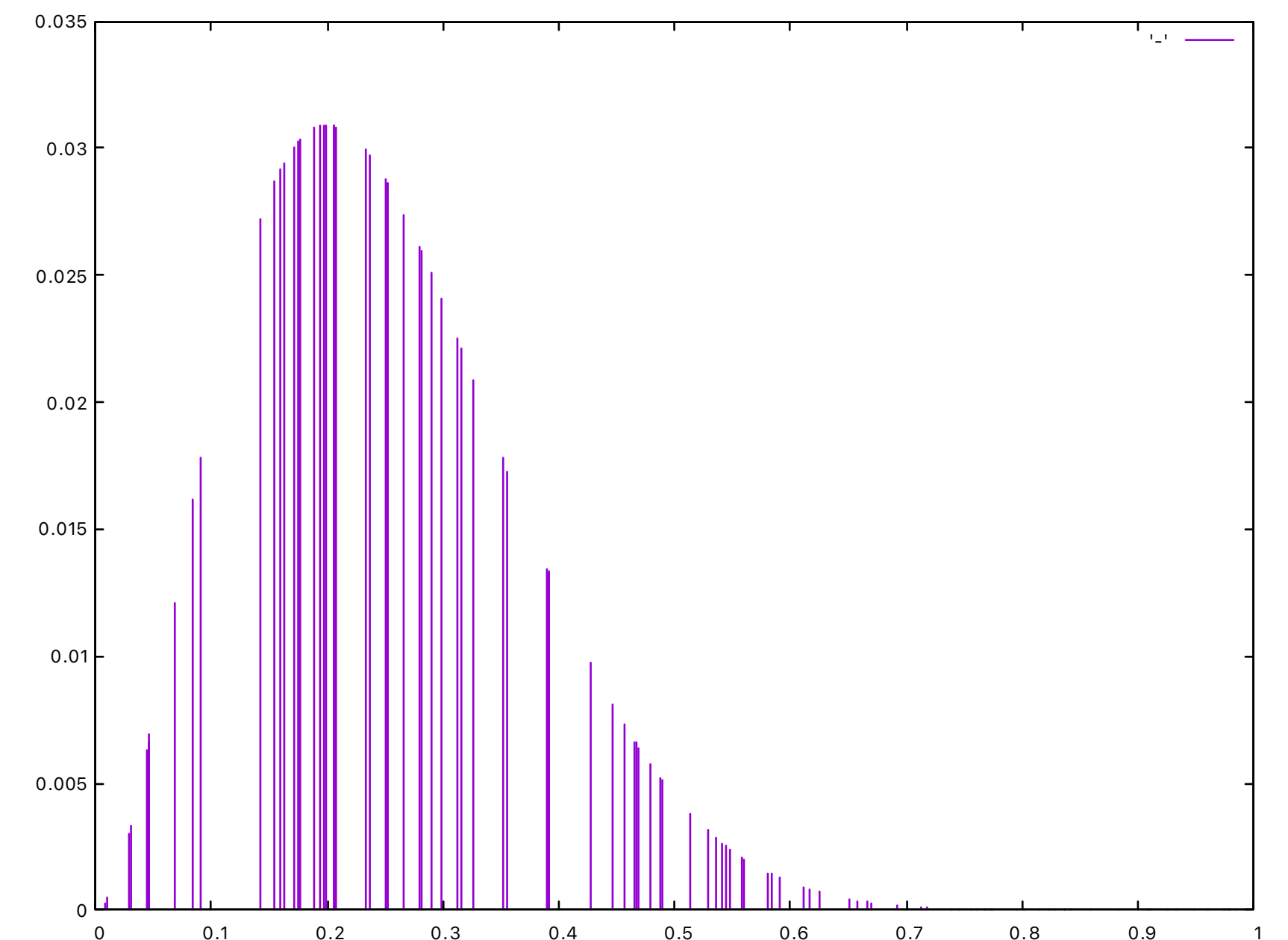
```
let coin prob x =  
  let z = sample prob (uniform ~a:0. ~b:1.) in  
  let () = List.iter (observe prob (bernoulli ~p:z)) x in  
  z
```

```
let _ =  
  let dist = infer coin [1; 1; 0; 0; 0; 0; 0; 0; 0; 0] in  
  let m, s = Distribution.stats dist in  
  Format.printf "Coin bias, mean:%f, std:%f@." m s
```

```
> dune exec ./examples/coin.exe
```

```
Coin bias, mean:0.247876, std:0.118921  
Beta(2+1, 8+1), mean:0.250000, std:0.120096
```

100 particles



Example: Coin

coin.ml

```
open Basic.Importance_sampling
```

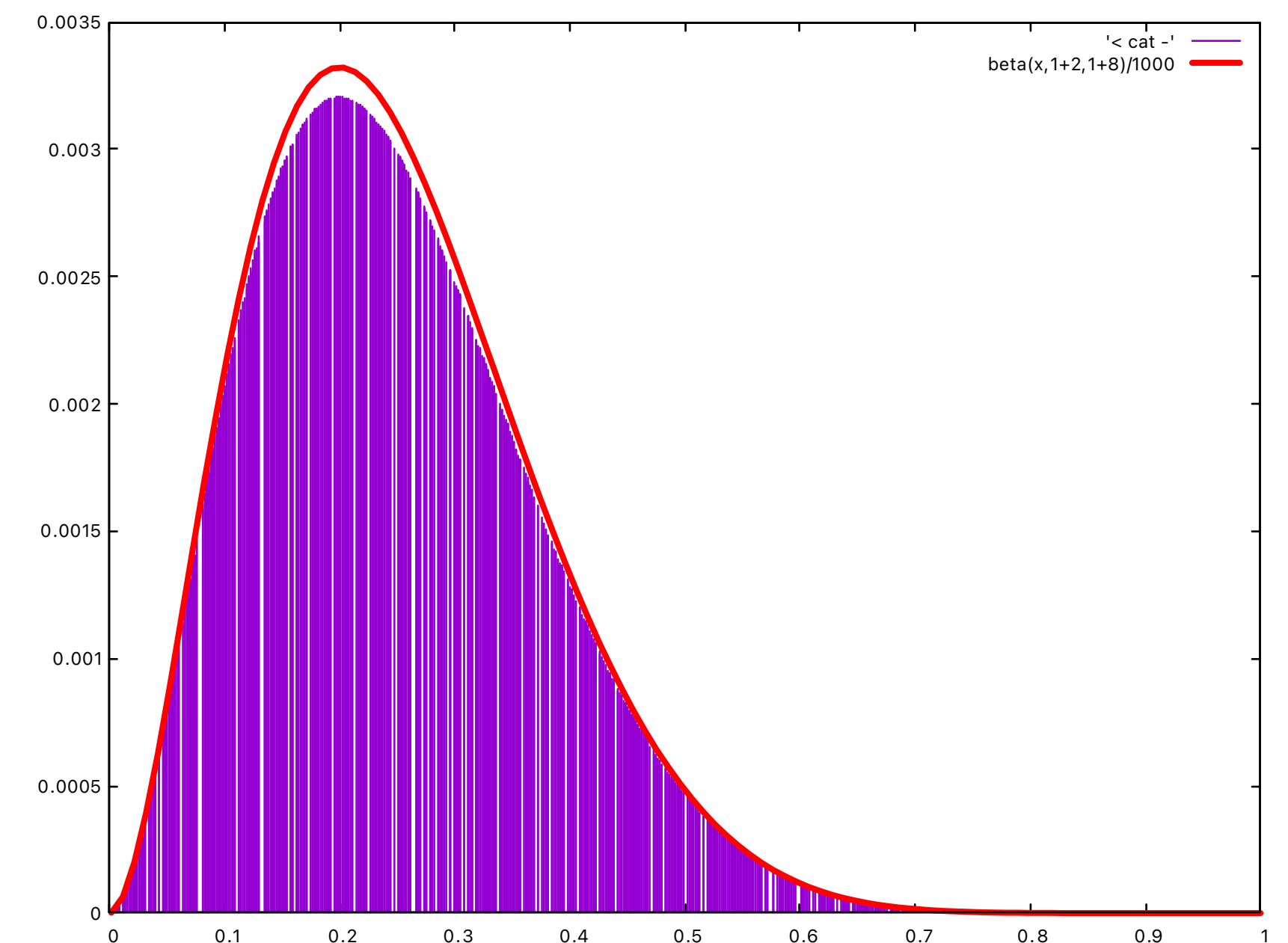
```
let coin prob x =  
  let z = sample prob (uniform ~a:0. ~b:1.) in  
  let () = List.iter (observe prob (bernoulli ~p:z)) x in  
  z
```

```
let _ =  
  let dist = infer coin [1; 1; 0; 0; 0; 0; 0; 0; 0; 0] in  
  let m, s = Distribution.stats dist in  
  Format.printf "Coin bias, mean:%f, std:%f@." m s
```

```
> dune exec ./examples/coin.exe
```

```
Coin bias, mean:0.247876, std:0.118921  
Beta(2+1, 8+1), mean:0.250000, std:0.120096
```

1000 particles



Conditioning

basic.ml

```
module Rejection_sampling = struct ...

  (* Reject if [p] is not true. *)
  let assume prob p =
    if not p then raise Reject

  (* Assume [x] was sampled from [d]. *)
  let observe prob d x =
    let v = sample d in
    assume prob (v = x)
```

Hard conditioning

Conditioning

basic.ml

```
module Rejection_sampling = struct ...
```

```
(* Reject if [p] is not true. *)
```

```
let assume prob p =
```

```
  if not p then raise Reject
```

```
(* Assume [x] was sampled from [d]. *)
```

```
let observe prob d x =
```

```
  let v = sample d in
```

```
  assume prob (v = x)
```

Hard conditioning

```
module Importance_sampling = struct ...
```

```
(* Update the (log)score. *)
```

```
let factor prob s =
```

```
  prob.(prob.id) ← prob.(prob.id) +. s
```

```
(* Assume [x] was sampled from [d]. *)
```

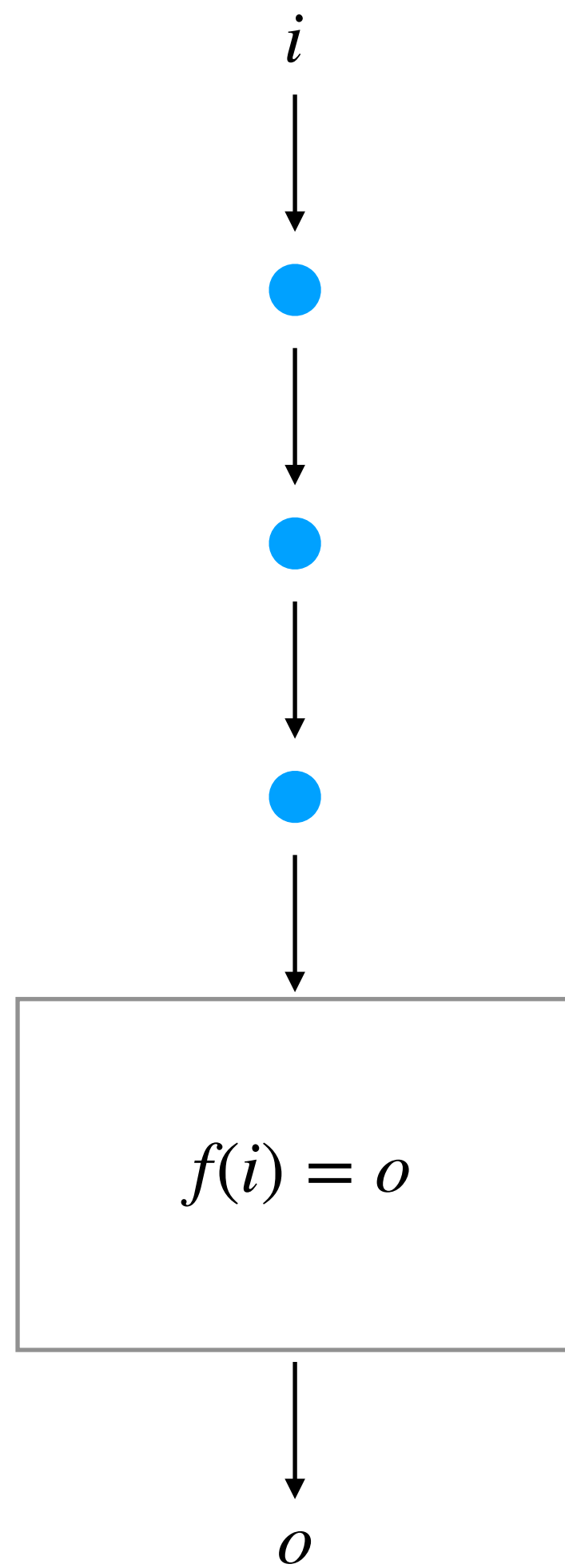
```
let observe prob d x =
```

```
  factor prob (logpdf d x)
```

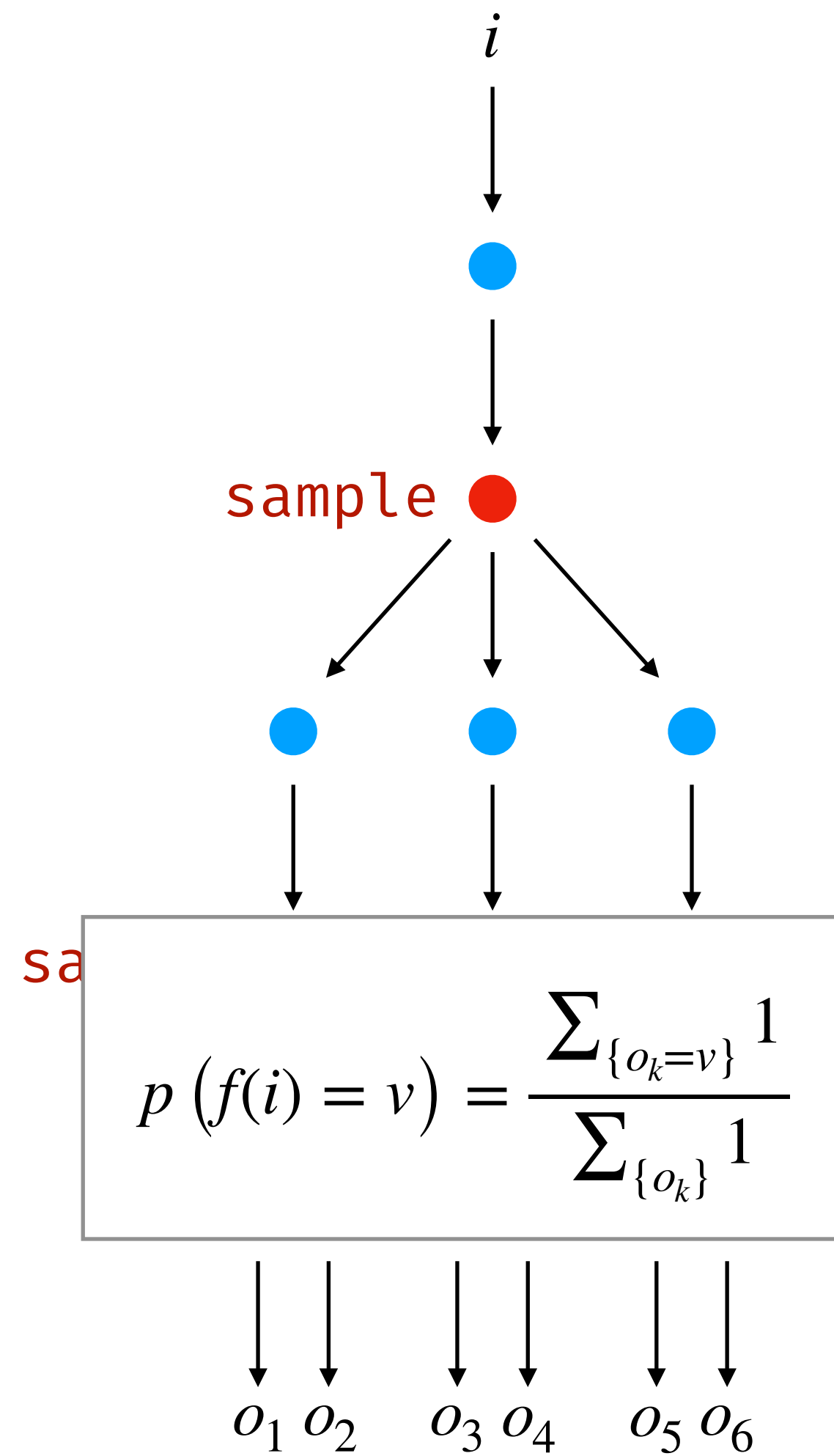
Soft conditioning

infer : $(\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \beta$ dist

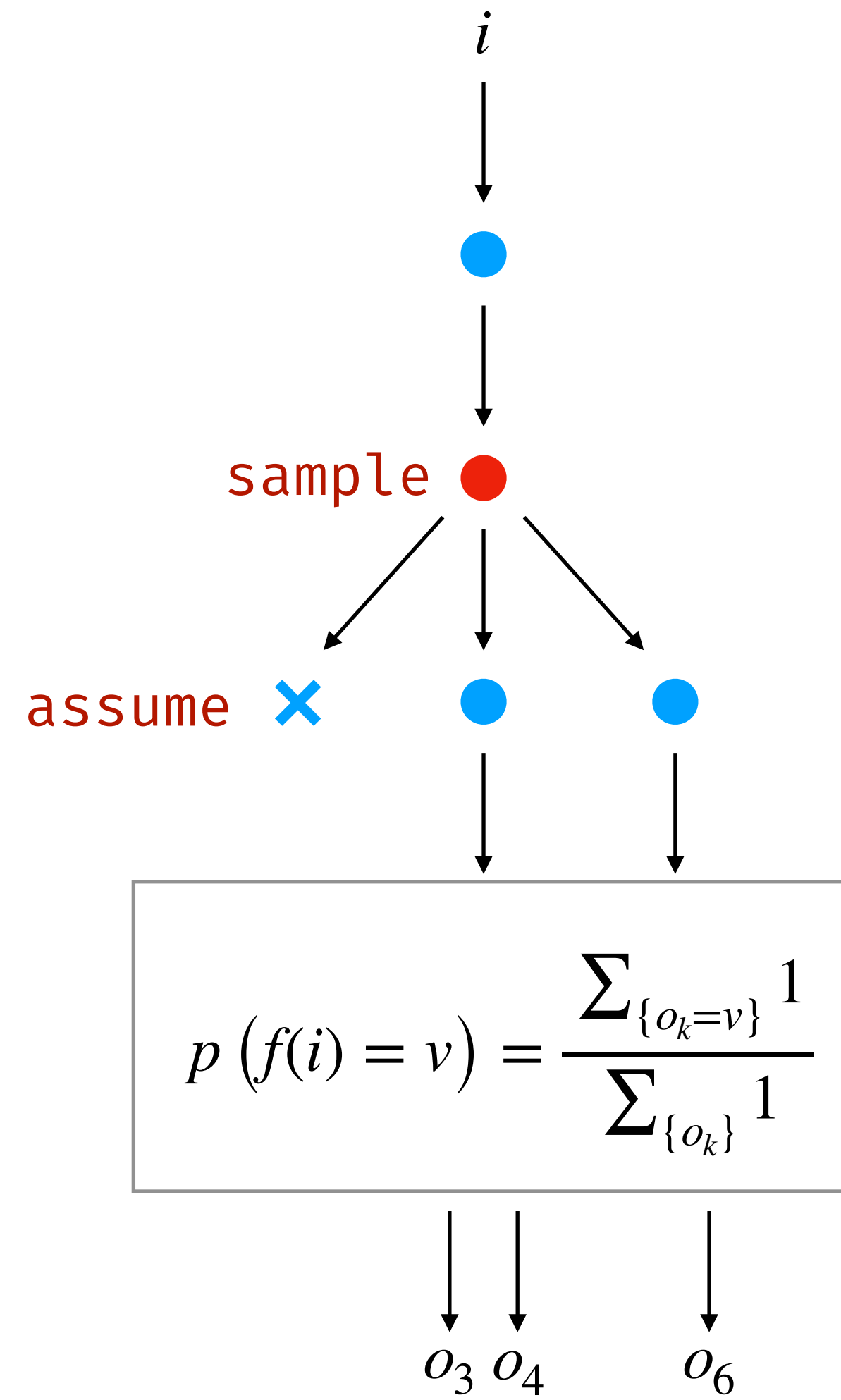
program



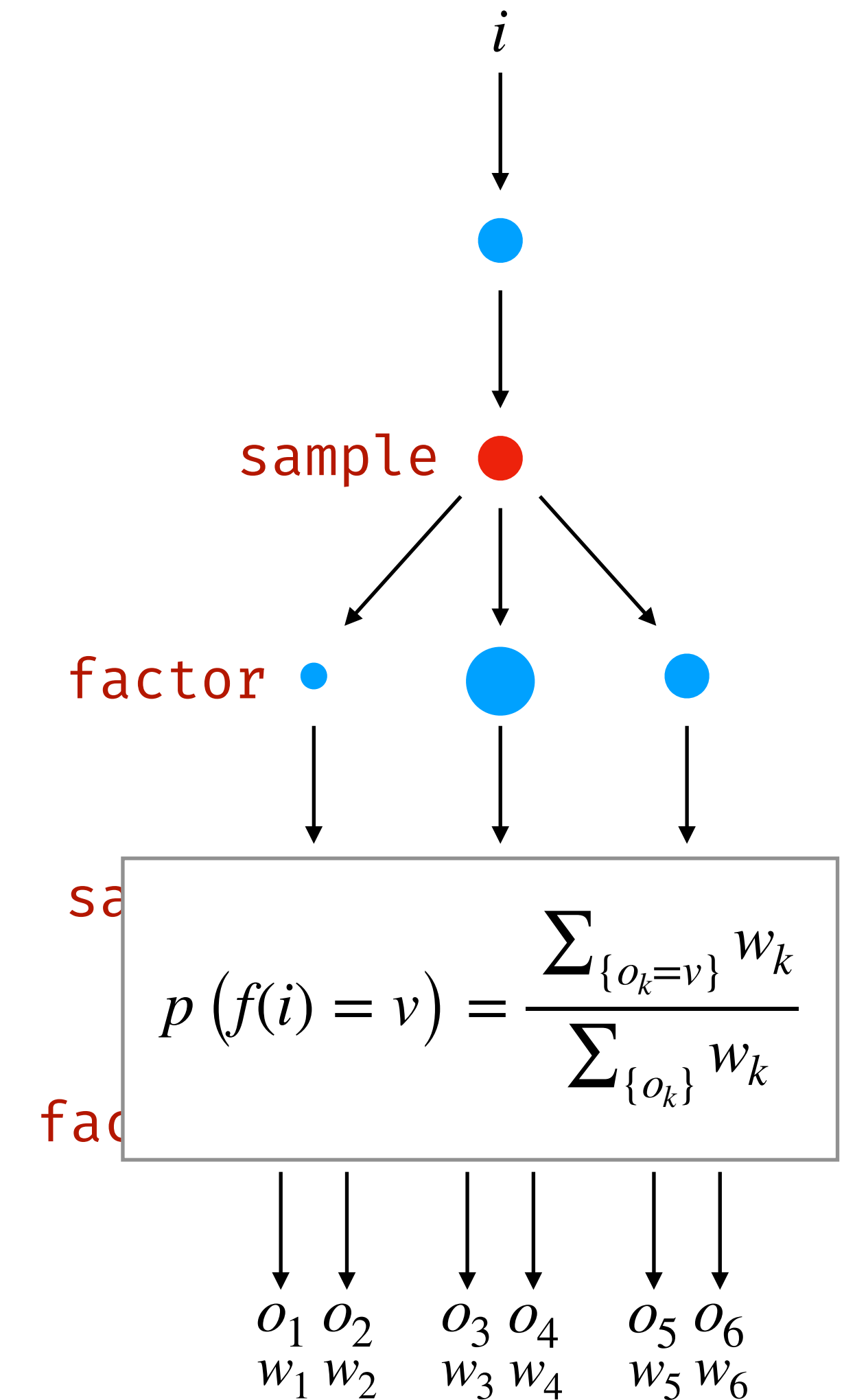
sample



assume



factor



Bayesian Reasoning

Bayesian Inference: learn parameters from data

- Latent parameter θ
- Observed data x_1, \dots, x_n

$$p(\theta \mid x_1, \dots, x_n) = \frac{p(\theta) p(x_1, \dots, x_n \mid \theta)}{p(x_1, \dots, x_n)} \quad (\text{Bayes' theorem})$$

posterior

$$\propto p(\theta) p(x_1, \dots, x_n \mid \theta) \quad (\text{Data are constants})$$

prior

likelihood

```
let model (x1, ..., xn) =  
  let theta = sample prior in  
  let () = observe (likelihood theta) (x1, ..., xn) in  
  theta
```

```
let posterior = infer model (x1, ..., xn)
```



Thomas Bayes (1701-1761)

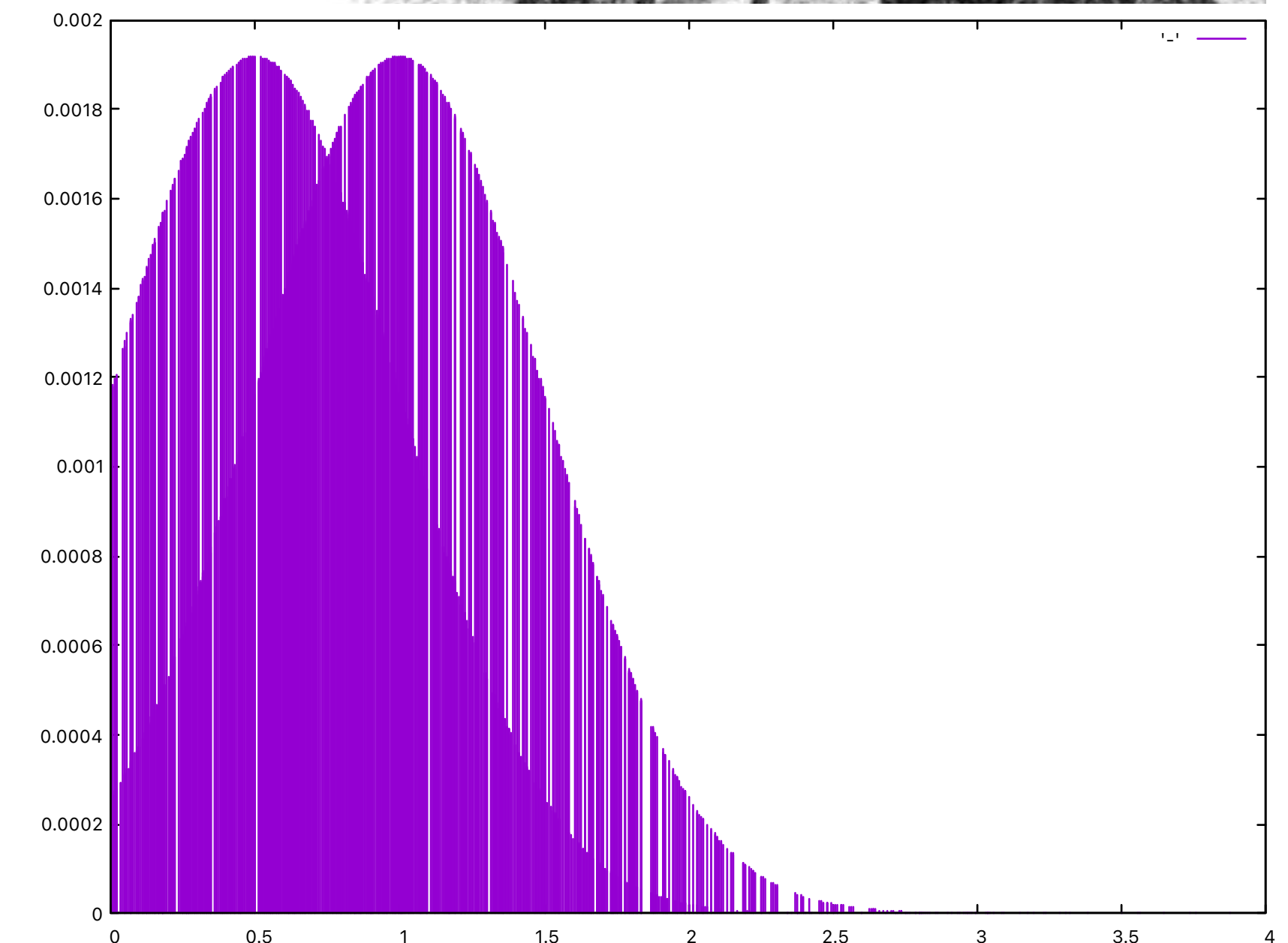
Probabilistic programming

Programming with probability distributions

- Draw a sample from a distribution
- Condition the model on assumption

More general than classic Bayesian Reasoning

```
let rec weird () =  
  let b = sample (bernoulli ~p:0.5) in  
  let mu = if (b = 1) then 0.5 else 1.0 in  
  let theta = sample (gaussian ~mu ~sigma:1.0) in  
  if theta > 0. then  
    observe (gaussian ~mu ~sigma:0.5) theta;  
    theta  
  else weird ()  
  
let weird_dist = infer weird ()
```



Typing

Probabilistic Programming Languages

Language and types

Simplified syntax

$x ::= \text{variables}$

$c ::= \text{constants}$

$d ::= \text{let } p = e \mid \text{let } f = \text{fun } p \rightarrow e \mid d \ d$

$p ::= x \mid (p, p)$

$e ::= c \mid x \mid (e, e) \mid \text{op } (e) \mid f(e)$

$\mid \text{if } e \text{ then } e \text{ else } e \mid \text{let } p = e \text{ in } e$

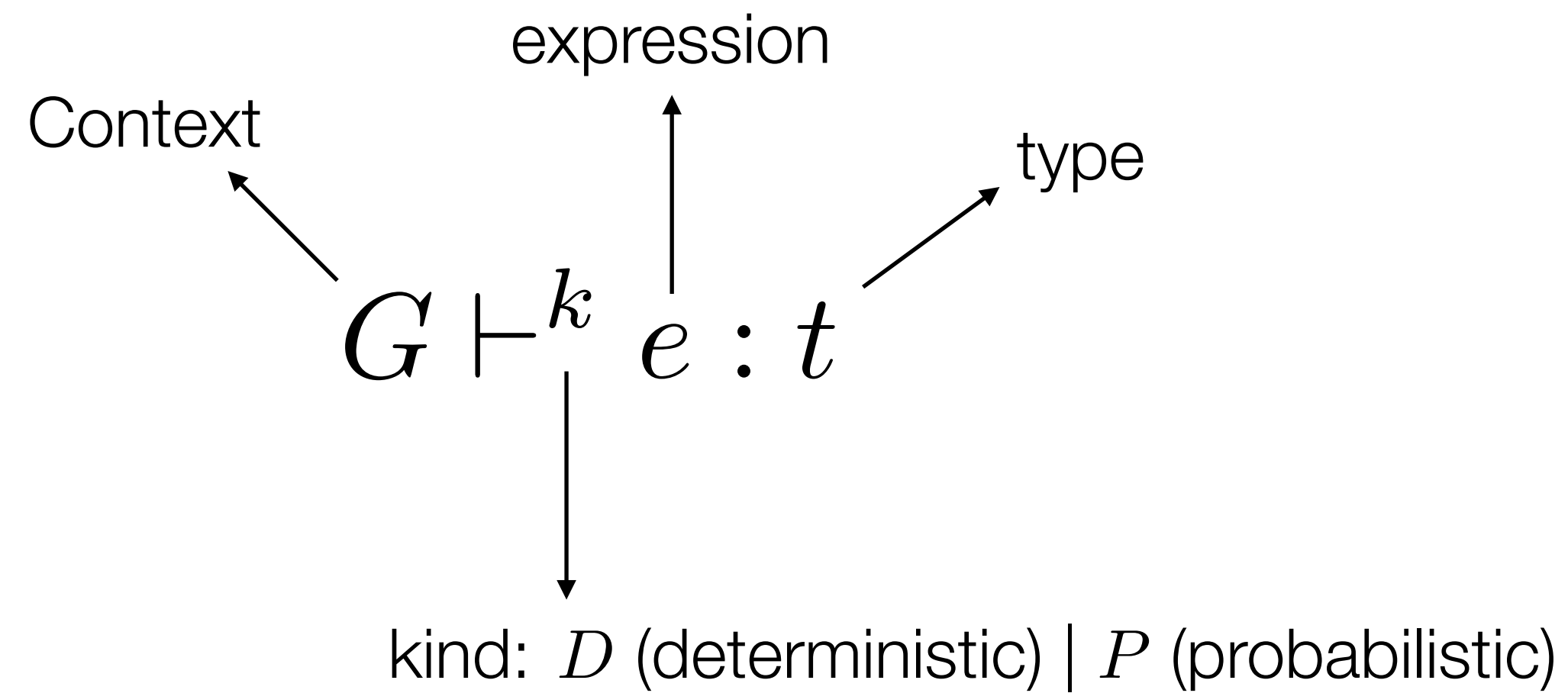
$\mid \text{sample } (e) \mid \text{factor } (e) \mid \text{observe } (e, e) \mid \text{infer } (e)$

Types

$t ::= \text{unit} \mid \text{bool} \mid \text{float} \mid t \text{ dist} \mid t \text{ dist}^* \mid t \times t \mid t \rightarrow t$

- $t \text{ dist}$: distribution over values of type t
- $t \text{ dist}^*$: distribution with densities ($\text{pdf}(d) : V \rightarrow [0, \infty)$ is defined)

Types and kinds



The type `prob` trick

```
module Rejection_sampling : sig
  type prob
  val sample : prob → 'a Distribution.t → 'a
  val assume : prob → bool → unit
  val infer : ?n:int → (prob → 'a → 'b) → 'a → 'b Distribution.t
end = struct ... end
```

Forbid the use of probabilistic construct outside a model

- Define a simple abstract type `prob`
- Probabilistic constructs and models all require an argument of type `prob`
- Such a value can only be build by `infer`

17

Kind `P` guards what can be expressed
in a probabilistic model

Typing declarations

$$\frac{G \vdash^D e : t}{G \vdash^D \text{let } p = e : G + [p \leftarrow t]}$$

$$\frac{k \in \{D, P\} \quad G + [p \leftarrow t_1] \vdash^k e : t_2}{G \vdash^D \text{let } f = \text{fun } p \rightarrow e : G + [f \leftarrow (t_1 \rightarrow^k t_2)]}$$

$$\frac{G \vdash^D d_1 : G_1 \quad G_1 \vdash^D d_2 : G_2}{G \vdash^D d_1 \ d_2 : G_2}$$

Declarations are deterministic
Functions can be D or P

Typing probabilistic constructs

$$\frac{G \vdash^P e : t}{G \vdash^D \text{infer}(e) : t \text{ dist}}$$

$$\frac{G \vdash^D e : t \text{ dist}}{G \vdash^P \text{sample}(e) : t}$$

$$\frac{G \vdash^D e : \text{float}}{G \vdash^P \text{factor}(e) : \text{unit}}$$

$$\frac{G \vdash^D e_1 : t \text{ dist}^* \quad G \vdash^D e_2 : t}{G \vdash^P \text{observe}(e_1, e_2) : \text{unit}}$$

$$\frac{G \vdash^D e : t}{G \vdash^P e : t}$$

$$\frac{G \vdash^D e : t \text{ dist}^*}{G \vdash^D e : t \text{ dist}}$$

Subtyping

Typing expressions

$$\frac{\text{typeOf}(c) = t}{G \vdash^D c : t}$$

$$\frac{G(x) = t}{G \vdash^D x : t}$$

$$\frac{G \vdash^D e_1 : t_1 \quad G \vdash^D e_2 : t_2}{G \vdash^D (e_1, e_2) : t_1 \times t_2}$$

$$\frac{\text{typeOf}(op) = t_1 \rightarrow^D t_2 \quad G \vdash^D e : t_1}{G \vdash^D op(e) : t_2}$$

$$\frac{G(f) = t_1 \rightarrow^k t_2 \quad G \vdash^D e : t_1}{G \vdash^k f(e) : t_2}$$

$$\frac{G \vdash^D e_1 : \text{bool} \quad G \vdash^k e_2 : t \quad G \vdash^k e_3 : t}{G \vdash^k \text{if } e_1 \text{ then } e_2 \text{ else } e_3 : t}$$

$$\frac{G \vdash^k e_1 : t_1 \quad G + [p \leftarrow t_1] \vdash^k e_2 : t_2}{G \vdash^k \text{let } p = e_1 \text{ in } e_2 : t_2}$$

Polymorphic kind

Example : Coin

coin.ml

```
let coin (x1, ... , xn) =  
  let z = sample (uniform (0., 1.)) in  
  observe (bernoulli (z), x1);  
  ... ;  
  observe (bernoulli (z), xn);  
  z  
  
let _ =  
  let d = infer (coin (1; 1; 0; 0; ... )) in  
  plot (d)
```

```
[coin : ???]  
[x1 :  $\alpha_1, \dots, x_n : \alpha_n$ ]  $\vdash^P$  z : float  
[x1 : int, ..., xn :  $\alpha_n$ , z : float]  $\vdash^P$  _ : unit  
  
[x1 : int, ..., xn : int, z : float]  $\vdash^P$  _ : unit  
[x1 : int, ..., xn : int, z : float]  $\vdash^P$  _ : float
```

Example : Coin

coin.ml

```
let coin (x1, ..., xn) =  
  let z = sample (uniform (0., 1.)) in  
  observe (bernoulli (z), x1);  
  ... ;  
  observe (bernoulli (z), xn);  
  z
```

```
let _ =  
  let d = infer (coin (1; 1; 0; 0; ... )) in  
  plot (d)
```

$[\text{coin} : (\text{int} \times \dots \times \text{int}) \rightarrow^P \text{float}]$

$[\text{x1} : \alpha_1, \dots, \text{xn} : \alpha_n] \vdash^P \text{z} : \text{float}$

$[\text{x1} : \text{int}, \dots, \text{xn} : \alpha_n, \text{z} : \text{float}] \vdash^P _ : \text{unit}$

$[\text{x1} : \text{int}, \dots, \text{xn} : \text{int}, \text{z} : \text{float}] \vdash^P _ : \text{unit}$

$[\text{x1} : \text{int}, \dots, \text{xn} : \text{int}, \text{z} : \text{float}] \vdash^P _ : \text{float}$

$[\text{coin} : (\text{int} \times \dots \times \text{int}) \rightarrow^P \text{float}] \vdash^D \text{d} : \text{float dist}$

$[\text{coin} : (\text{int} \times \dots \times \text{int}) \rightarrow^P \text{float}, \text{d} : \text{float dist}] \vdash^D _ : \text{unit}$

Example : Coin

coin.ml

```
let coin prob (x1, ..., xn) =  
  let z = sample prob (uniform (0., 1.)) in  
  observe prob (bernoulli (z), x1);  
  ... ;  
  observe prob (bernoulli (z), xn);  
  z
```

```
let _ =  
  let d = infer (coin (1; 1; 0; 0; ... )) in  
  plot (d)
```

```
[coin : (int × ... × int) →P float]  
[x1 : α1, ..., xn : αn] ⊢P z : float  
[x1 : int, ..., xn : αn, z : float] ⊢P _ : unit
```

```
[x1 : int, ..., xn : int, z : float] ⊢P _ : unit  
[x1 : int, ..., xn : int, z : float] ⊢P _ : float
```

```
[coin : (int × ... × int) →P float] ⊢D d : float dist  
[coin : (int × ... × int) →P float, d : float dist] ⊢D _ : unit
```

Reminders: Measure Theory

Probabilistic Programming Languages

Measurable spaces

A σ -algebra on a set X is a collection of subsets:

- containing \emptyset
- closed under complement
- closed under countable union

A measurable space is a pair (X, Σ_X)

- X : set
- Σ_X : measurable sets

Examples:

- $\emptyset \dots$
- Singleton: $\{\emptyset, \{()\}\}$
- Booleans: $\{\emptyset, \{\text{true}\}, \{\text{false}\}, \{\text{true}, \text{false}\}\}$
- Borel sets (intervals) on \mathbb{R}
- Product : $\Sigma_{A \times B} = \{(U, V) \mid U \in \Sigma_A, V \in \Sigma_B\}$

Measure

A measure maps a measurable set to a positive score: $\mu : \Sigma_X \rightarrow [0, \infty)$

- $\mu(U) \geq 0$ for all set $U \in \Sigma_X$
- $\mu(\emptyset) = 0$
- $\mu(\bigcup_{i \in I} U_i) = \sum_{i \in I} \mu(U_i)$ for all countable collection $\bigcup_{i \in I} U_i$ of pairwise disjoint sets in Σ_X

Probability distributions are normalized measure, i.e., $\mu(X) = 1$

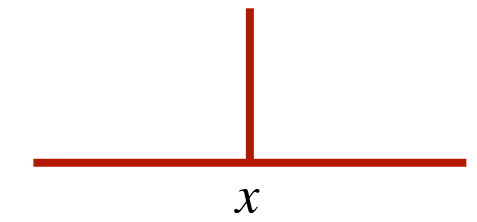
Examples

- Lebesgue measure
 $\lambda([a, b]) = b - a$
- Bernoulli distribution (discrete)
 $\mathcal{B}(0.3)(\{\text{true}\}) = 0.3$
 $\mathcal{B}(0.3)(\{\text{true}, \text{false}\}) = 1$

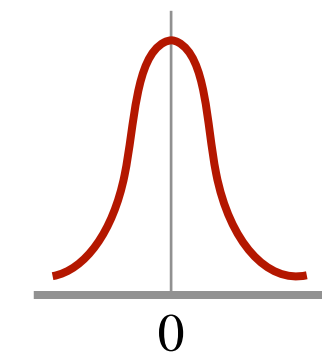


- Dirac delta measure

$$\delta_x(U) = \begin{cases} 1 & \text{if } x \in U \\ 0 & \text{otherwise} \end{cases}$$



- Normal distribution
 $\mathcal{N}(0, 1)([0, 1]) \approx 0.34$
 $\mathcal{N}(0, 1)((-\infty, 0]) = 0.5$



Measurable functions

$f : X \rightarrow Y$ is measurable if $f^{-1}(U) \in \Sigma_X$ for all $U \in \Sigma_Y$

Pushforward: transfer a measure from a measurable space to another one

- $f : X \rightarrow Y$ measurable
- $\mu : \Sigma_X \rightarrow [0, \infty)$
- $f_*(\mu) : \Sigma_Y \rightarrow [0, \infty)$

Examples

- $\mu : \Sigma_{A \times B} \rightarrow [0, \infty)$
- $\pi_{1*}(\mu) : \Sigma_A \rightarrow [0, \infty)$
- $\pi_{2*}(\mu) : \Sigma_B \rightarrow [0, \infty)$

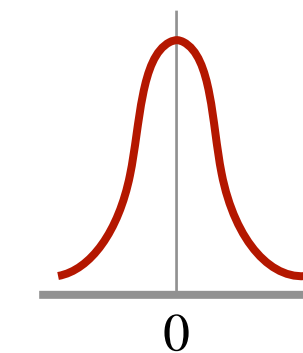
Integration

Given $f, g : X \rightarrow [0, \infty)$ measurable functions and a measure $\mu : \Sigma_X \rightarrow [0, \infty)$

- $\int f(x)\mu(dx) \in [0, \infty)$
- $\int \mu(dx)\delta_x(U) = \int_U \mu(dx) = \mu(U)$ for all $U \in \Sigma_X$
- $\int f(x)r\mu(dx) = r \int f(x)\mu(dx)$ for $r \in \mathbb{R}$
- $\int (f(x) + g(x))\mu(dx) = \int f(x)\mu(dx) + \int g(x)\mu(dx)$

We can also define new measures from integration

- $f(x) = \frac{1}{\sqrt{2\pi}}e^{-\frac{x^2}{2}}$
- $\int_U f(x)\lambda(dx) = \mathcal{N}(0, 1)(U)$ where λ is the Lebesgue measure



Kernel

A kernel $k : X \times \Sigma_Y \rightarrow [0, \infty)$ is a function such that:

- $k(x, _) : \Sigma_Y \rightarrow [0, \infty)$ for all $x \in X$ is a measure
- $k(_, U) : X \rightarrow [0, \infty)$ for all $U \in \Sigma_Y$ is measurable

A probability kernel is such that $k(x, Y) = 1$ for all $x \in X$

We now have everything to define
a kernel based denotational semantics!

Semantics

Probabilistic Programming Languages

Types as measurable spaces

A ground type t is interpreted as a measurable space $\llbracket t \rrbracket$

- $\llbracket \text{unit} \rrbracket$: discrete measurable space over the unique value $()$
- $\llbracket \text{bool} \rrbracket$: discrete measurable space with the two values $\text{true}, \text{false}$
- $\llbracket \text{float} \rrbracket$: reals with its Borel sets (intervals)
- $A \times B$ product space $\llbracket A \rrbracket \times \llbracket B \rrbracket$
with the rectangles $U \times V$ for $U \in \Sigma_A$ and $V \in \Sigma_B$
- $\llbracket t \text{ dist} \rrbracket$: set of probability measures on $\llbracket t \rrbracket$
with the sets $\{\mu \mid \mu(U) < r\}$ for $U \in \Sigma_{\llbracket t \rrbracket}$ and $r \in [0, 1]$ (Giry monad)
- A context $G = [x_1 : A_1, \dots, x_n : A_n]$ maps variables to types
 $\llbracket G \rrbracket = \prod_{i=1}^n \llbracket A_i \rrbracket$ is also a measurable space (product of all variables spaces)

What about function types?

Deterministic vs. probabilistic

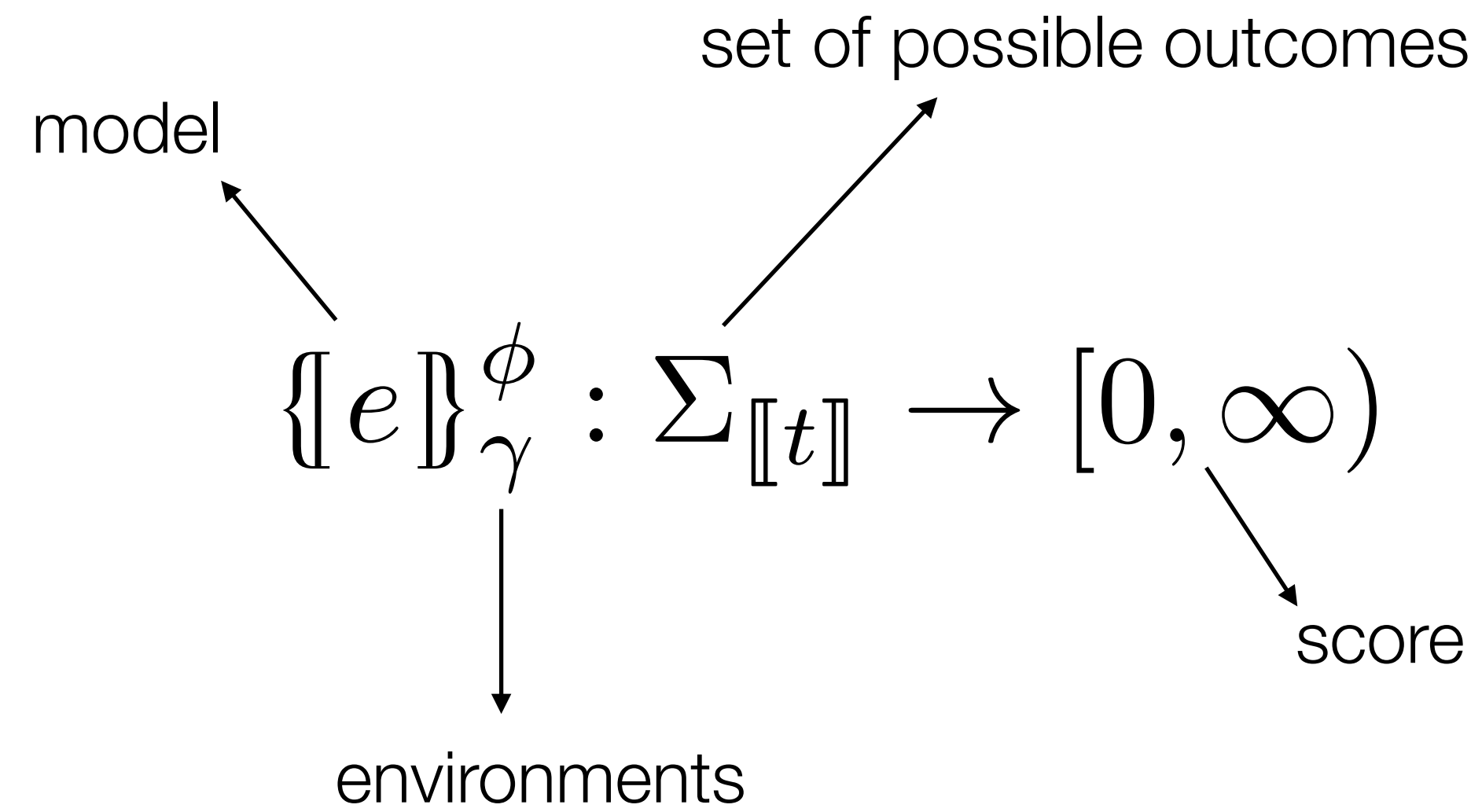
Deterministic semantics $G \vdash^D e : t$

- Classic denotational semantics
- Environments: ϕ (global declarations), γ (local variables)
- Given the declarations ϕ , $\llbracket e \rrbracket^\phi : \Gamma \rightarrow t$ is a measurable function
- $\llbracket e \rrbracket_\gamma^\phi$ is a value of type t

Probabilistic semantics $G \vdash^P e : t$

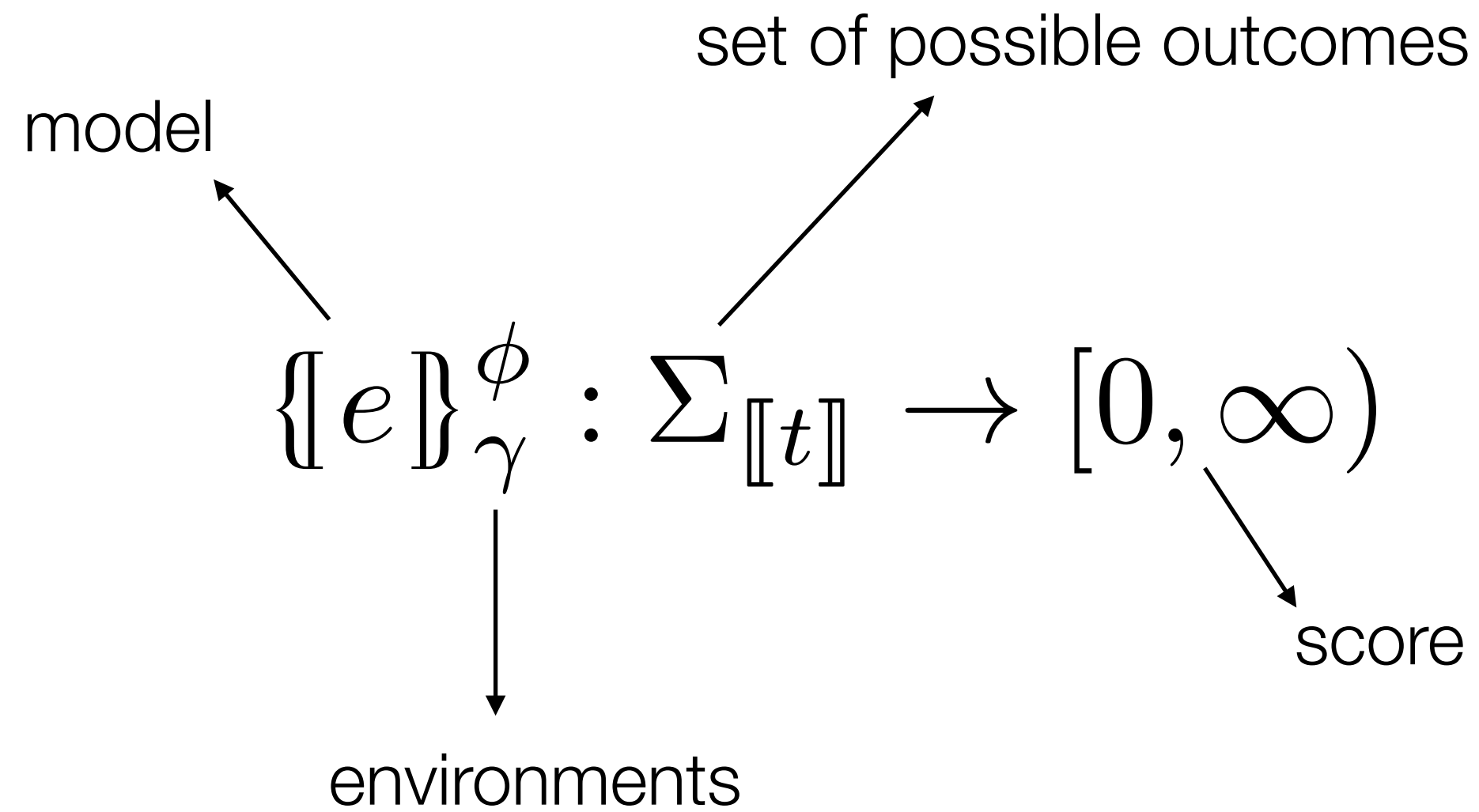
- Given the declarations ϕ , expressions are interpreted as kernels
- $\{e\}^\phi : \Gamma \times \Sigma_{\llbracket t \rrbracket} \rightarrow [0, \infty)$
- $\{e\}_\gamma^\phi$ is a measure on values of type t

(Un)normalized measures



Unnormalized measure

(Un)normalized measures



Unnormalized measure

$$\llbracket \text{infer}(e) \rrbracket_{\gamma}^{\phi} = \frac{\{\![e]\!\}_{\gamma}^{\phi}}{\{\![e]\!\}_{\gamma}^{\phi} (\llbracket \text{typeOf}(e) \rrbracket)}$$

Distribution

normalize over all possible values

Deterministic semantics

$$\begin{aligned}
 \llbracket \text{let } p = e \rrbracket^\phi &= \phi + \left[p \leftarrow \llbracket e \rrbracket_\emptyset^\phi \right] \\
 \llbracket \text{let } f = \text{fun } p \rightarrow e \rrbracket^\phi &= \phi + \left[f \leftarrow \lambda v. \llbracket e \rrbracket_{[p \leftarrow v]}^\phi \right] \\
 \llbracket d_1 \ d_2 \rrbracket^\phi &= \text{let } \phi_1 = \phi + \llbracket d_1 \rrbracket^\phi \text{ in } \llbracket d_2 \rrbracket^{\phi_1} \\
 \llbracket c \rrbracket_\gamma^\phi &= c \\
 \llbracket x \rrbracket_\gamma^\phi &= (\gamma + \phi)(x) \\
 \llbracket (e_1, e_2) \rrbracket_\gamma^\phi &= (\llbracket e_1 \rrbracket_\gamma^\phi, \llbracket e_2 \rrbracket_\gamma^\phi) \\
 \llbracket op(e) \rrbracket_\gamma^\phi &= op(\llbracket e \rrbracket_\gamma^\phi) \\
 \llbracket f(e) \rrbracket_\gamma^\phi &= \phi(f)(\llbracket e \rrbracket_\gamma^\phi) \\
 \llbracket \text{if } e_1 \text{ then } e_2 \text{ else } e_3 \rrbracket_\gamma^\phi &= \text{if } \llbracket e_1 \rrbracket_\gamma^\phi \text{ then } \llbracket e_2 \rrbracket_\gamma^\phi \text{ else } \llbracket e_3 \rrbracket_\gamma^\phi \\
 \llbracket \text{let } p = e_1 \text{ in } e_2 \rrbracket_\gamma^\phi &= \text{let } v = \llbracket e_1 \rrbracket_\gamma^\phi \text{ in } \llbracket e_2 \rrbracket_{\gamma + [p \leftarrow v]}^\phi
 \end{aligned}$$

Probabilistic semantics

$$\llbracket \text{let } f = \text{fun } p \rightarrow e \rrbracket^\phi = \phi + \left[f \leftarrow \lambda v. \llbracket e \rrbracket^\phi_{[p \leftarrow v]} \right] \text{ if } \text{kindOf}(e) = P$$

$$\llbracket e \rrbracket^\phi_\gamma = \lambda U. \delta_{\llbracket e \rrbracket^\phi_\gamma}(U) \text{ if } \text{kindOf}(e) = D$$

$$\llbracket f(e) \rrbracket^\phi_\gamma = \lambda U. \phi(f)(\llbracket e \rrbracket^\phi_\gamma)(U)$$

$$\llbracket \text{if } e_1 \text{ then } e_2 \text{ else } e_3 \rrbracket^\phi_\gamma = \lambda U. \text{if } \llbracket e_1 \rrbracket^\phi_\gamma \text{ then } \llbracket e_2 \rrbracket^\phi_\gamma(U) \text{ else } \llbracket e_3 \rrbracket^\phi_\gamma(U)$$

$$\llbracket \text{let } p = e_1 \text{ in } e_2 \rrbracket^\phi_\gamma = \lambda U. \int_{\llbracket \text{typeOf}(e_1) \rrbracket} \llbracket e_1 \rrbracket^\phi_\gamma(dv) \llbracket e_2 \rrbracket^\phi_{\gamma+[p \leftarrow v]}$$

$$\llbracket \text{sample}(e) \rrbracket^\phi_\gamma = \lambda U. \llbracket e \rrbracket^\phi_\gamma(U)$$

$$\llbracket \text{factor}(e) \rrbracket^\phi_\gamma = \lambda U. \llbracket e \rrbracket^\phi_\gamma \cdot \delta_{()}(U)$$

$$\llbracket \text{observe}(e_1, e_2) \rrbracket^\phi_\gamma = \lambda U. \text{pdf}(\llbracket e_1 \rrbracket^\phi_\gamma)(\llbracket e_2 \rrbracket^\phi_\gamma) \cdot \delta_{()}(U)$$

$$\llbracket \text{infer}(e) \rrbracket^\phi_\gamma = \begin{cases} \frac{\lambda U. \llbracket e \rrbracket^\phi_\gamma(U)}{\llbracket e \rrbracket^\phi_\gamma(\llbracket \text{typeOf}(e) \rrbracket)} & \text{if } 0 < \llbracket e \rrbracket^\phi_\gamma(\llbracket \text{typeOf}(e) \rrbracket) < \infty \\ \text{Error} & \text{otherwise} \end{cases}$$

Careful with 0, and ∞ ...

Example : Gaussian

my_gaussian.ml

```
let my_gaussian (mu, sigma) =  
  let x = sample (gaussian (mu, sigma)) in  
  x
```

$$\begin{aligned}\{\text{my_gaussian } (\mu, \sigma)\}_{\emptyset}(U) &= \int_{\mathbb{R}} \{\text{sample } (\text{gaussian } (\mu, \sigma))\}_{[\mu \leftarrow \mu, \sigma \leftarrow \sigma]}(dx) \{x\}_{[\mu \leftarrow \mu, \sigma \leftarrow \sigma, x \leftarrow x]} \\ &= \int_{\mathbb{R}} \text{Gaussian}(\mu, \sigma)(dx) \delta_x(U) \\ &= \int_U \frac{1}{\sigma\sqrt{2\pi}} e^{-\frac{(x-\mu)^2}{2\sigma^2}} dx \\ &= \text{Gaussian}(\mu, \sigma)(U)\end{aligned}$$

Example : Beta

my_gaussian.ml

```
let my_beta (a, b) =  
  let x = sample (uniform (0., 1.)) in  
  let () = observe (beta (a, b), x) in  
  x
```

$$\begin{aligned}\{\text{my_beta } (a, b)\}_{\emptyset}(U) &= \int_0^1 \{\text{sample } (\text{uniform } (0, 1))\}_{[a \leftarrow a, b \leftarrow b]}(dx) \\ &\quad \int_{()} \{\text{observe } (\text{beta } (a, b), x)\}_{[a \leftarrow a, b \leftarrow b, x \leftarrow x]}(du) \{\times\}_{[a \leftarrow a, b \leftarrow b, x \leftarrow x]}(U) \\ &= \int_0^1 \text{Uniform}(dx) \text{pdf}(\text{Beta}(a, b))(x) \delta_x(U) \\ &= \int_U \text{pdf}(\text{Beta}(a, b))(x) dx \\ &= \text{Beta}(a, b)(U)\end{aligned}$$

Example : Coin

coin.ml

```
let coin (x1, ..., xn) =
  let z = sample (uniform (0., 1.)) in
  observe (bernoulli (z), x1); ... ; observe (bernoulli (z), xn);
  z
```

$$\begin{aligned}
 \{\{\text{coin } (x_1, \dots, x_n)\}\}_{\emptyset}(U) &= \int_0^1 \{\{\text{sample } (\text{uniform } (0, 1))\}\}_{[x_1 \leftarrow x_1, \dots, x_n \leftarrow x_n]}(dz) \\
 &\quad \int_{(\cdot)} \{\{\text{observe } (\text{bernoulli } (z), x_1)\}\}_{[z \leftarrow z, x_1 \leftarrow x_1, \dots, x_n \leftarrow x_n]}(du_0) \\
 &\quad \int_{(\cdot)} \{\{\text{observe } (\text{bernoulli } (z), x_2)\}\}_{[z \leftarrow z, x_1 \leftarrow x_1, \dots, x_n \leftarrow x_n]}(du_1) \\
 &\quad \dots \\
 &\quad \int_{(\cdot)} \{\{\text{observe } (\text{bernoulli } (z), x_n)\}\}_{[z \leftarrow z, x_1 \leftarrow x_1, \dots, x_n \leftarrow x_n]}(du_n) \\
 &\quad \{\{z\}\}_{[z \leftarrow z, x_1 \leftarrow x_1, \dots, x_n \leftarrow x_n]}(U) \\
 &= \int_0^1 \text{Uniform}(0, 1)(dz) \prod_{i=1}^n \text{pdf}(\text{Bernoulli}(z))(x_i) \delta_z(U) \\
 &= \int_U z^{\#\text{heads}} (1 - z)^{\#\text{tails}} dz
 \end{aligned}$$

Unnormalized!

Example : Coin

coin.ml

```
let coin (x1, ..., xn) =  
  let z = sample (uniform (0., 1.)) in  
  observe (bernoulli (z), x1); ... ; observe (bernoulli (z), xn);  
  z
```

```
let d = infer (coin (data))
```

$$\{\{\text{coin } (x_1, \dots, x_n)\}\}_{\emptyset}(U) = \int_U z^{\#\text{heads}} (1 - z)^{\#\text{tails}} dz$$

$$\llbracket \text{infer } (\text{coin } (x_1, \dots, x_n)) \rrbracket_{[\text{coin}]} = \frac{\int_U z^{\#\text{heads}} (1 - z)^{\#\text{tails}} dz}{\int_0^1 z^{\#\text{heads}} (1 - z)^{\#\text{tails}} dz} = \frac{\int_U z^{\#\text{heads}} (1 - z)^{\#\text{tails}} dz}{B(\#\text{heads} + 1, \#\text{tails} + 1)} = \text{Beta}(\#\text{heads} + 1, \#\text{tails} + 1)(U)$$

Exercises

Prove the following properties

- `sample mu (* where mu is defined on [a, b] *)`

≡

```
let x = sample (uniform (a, b)) in
let () = observe (mu, x) in
x
```

- `observe (mu, x) (* where mu is a discrete distribution *)`

≡

```
let y = sample mu in
assume x = y
```

- `sample (bernoulli (0.5))`

≡

```
let x = sample (gaussian (0., 1.)) in
x > 0
```

Example: Laplace and gender bias

laplace.ml

```
open Basic.Rejection_sampling
```

```
let laplace prob () =
  let p = sample prob (uniform ~a:0. ~b:1.) in
  let g = sample prob (binomial ~p ~n:493_472) in
  let () = assume prob (g = 241_945) in
  p
```

→ let () = observe prob
(binomial ~p ~n:493_472) 241_945

```
let _ =
  let dist = infer ~n:1000 laplace () in
  let m, s = Distribution.stats dist in
  Format.printf "Gender bias, mean:%f std:%f@." m s
```

```
> dune exec ./examples/laplace.exe
```

Never terminate!

25

Improper priors

Uniform priors on bounded domains

- If $\mu : t \text{ dist}^*$ is defined on $[a, b]$ and has a density
- $\{\text{sample}(\mu)\} = \{\text{let } x = \text{sample}(\text{Uniform}(a, b)) \text{ in observe}(\mu, x); x\}$

Improper priors

```
let improper =  
  let x = sample (gaussian 0 1) in  
  factor (1. /. (pdf (gaussian 0 1) x));  
x
```

$$\begin{aligned}\{\text{improper}\}_{\emptyset}(U) &= \int_U \text{Gaussian}(0, 1)(dx) \frac{1}{f(x)} dx \\ &= \int_U f(x) \frac{1}{f(x)} dx \\ &= \lambda(U)\end{aligned}$$

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Semantics of probabilistic programs.

Dexter Kozen

Journal of Computer and System 1981

Commutative semantics for probabilistic programming

Sam Staton

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Semantics of Probabilistic Programs using s-Finite Kernels in Coq

Reynald Affeldt, Cyril Cohen, Ayumu Saito

CPP 2023

TP : A short introduction to Stan

Everything is on Github: <https://github.com/mpri-probprog/probprog-23-24>

- Go to td/td2
- Launch jupyter notebook (or jupyter lab)

Requirements

- Pandas
- CmdStanPy
- Jupyter
- Matplotlib



<https://mc-stan.org/>

