Merge Sort

Merge sort is a sorting technique based on divide and conquer technique. With worst-case time complexity being **O(n log n)**, it is one of the most used and approached algorithms. Merge sort first divides the array into equal halves and then combines them in a sorted manner.

What Is a Divide and Conquer Algorithm?

Divide-and-conquer recursively solves subproblems; each subproblem must be smaller than the original problem, and each must have a base case. A divide-and-conquer algorithm has three parts:

Divide: Divide the list or array recursively into two halves until it can no more be divided.

Conquer: Each subarray is sorted individually using the merge sort algorithm.

Merge: The sorted subarrays are merged back together in sorted order. The process continues until all elements from both subarrays have been merged.

Merge Sort Algorithm

```
MERGE-SORT(A, p, r)

1 if p < r

2 q = \lfloor (p+r)/2 \rfloor

3 MERGE-SORT(A, p, q)

4 MERGE-SORT(A, q+1, r)

5 MERGE(A, p, q, r)
```

```
MERGE(A, p, q, r)
 1 \quad n_1 = q - p + 1
 2 \quad n_2 = r - q
 3 let L[1...n_1 + 1] and R[1...n_2 + 1] be new arrays
 4 for i = 1 to n_1
         L[i] = A[p+i-1]
 5
    for j = 1 to n_2
 6
        R[j] = A[q+j]
 7
 8 L[n_1 + 1] = \infty
 9 \quad R[n_2+1] = \infty
10 i = 1
11 j = 1
12 for k = p to r
        if L[i] \leq R[j]
13
             A[k] = L[i]
14
             i = i + 1
15
         else A[k] = R[j]
16
             j = j + 1
17
```

Example: A=[2, 4, 5, 7, 1, 2, 3, 6]

Quick Sort

Quick Sort is a sorting algorithm based on the Divide and Conquer algorithm that picks an element as a pivot and partitions the given array around the picked pivot by placing the pivot in its correct position in the sorted array.

The following procedure implements quicksort.

```
QUICKSORT (A, p, r)

1 if p < r

2 then q \leftarrow PARTITION(A, p, r)

3 QUICKSORT (A, p, q - 1)

4 QUICKSORT (A, q + 1, r)
```

To sort an entire array A, the initial call is QUICKSORT(A, 1, length[A]).

Partitioning the array

The key to the algorithm is the PARTITION procedure, which rearranges the subarray $A[p \ | \ r]$ in place.

```
PARTITION(A, p, r)

1  x \leftarrow A[r]

2  i \leftarrow p - 1

3  for j \leftarrow p to r - 1

4  do if A[j] \leq x

5  then i \leftarrow i + 1

6  exchange A[i] \leftrightarrow A[j]

7  exchange A[i + 1] \leftrightarrow A[r]

8  return i + 1
```

Example: 2, 8, 7, 1, 3, 5, 6, 4

Heap Sort

```
Heapsort(A)
      Build-Max-Heap(A)
      for i = A.length downto 2
             exchange A[1] with A[i]
             A.heap-size = A.heap-size - 1
 5
             Max-Heapify(A, 1)
// Implementation of Heap Sort in C
 #include <stdio.h>
// Function to swap the position of two elements in an array
 void swap(int *a, int *b) {
 int tempvar = *a;
  *a = *b;
  *b = tempvar;
 }
 void heapify(int arr[], int n, int i) {
 // Finding the greatest among root, leftSide child, and rightSide child of
the tree
 int greatest = i;
 int leftSide = 2 * i + 1;
 int rightSide = 2 * i + 2;
```

```
if (leftSide < n && arr[leftSide] > arr[greatest])
  greatest = leftSide;
 if (rightSide < n && arr[rightSide] > arr[greatest])
  greatest = rightSide;
 // Swap and continue heapifying if the root is not the greatest
 if (greatest != i) {
  swap(&arr[i], &arr[greatest]);
  heapify(arr, n, greatest);
}
}
// Main function
void heapSort(int arr[], int n) {
 // Build max heap
 for (int i = n / 2 - 1; i \ge 0; i--)
  heapify(arr, n, i);
 for (int i = n - 1; i \ge 0; i--) {
  swap(&arr[0], &arr[i]);
  heapify(arr, i, 0);
 }
}
```

```
// Printing the array
void printArray(int arr[], int n) {
 for (int i = 0; i < n; ++i)
  printf("%d", arr[i]);
 printf("\n");
}
int main() {
 int arr[] = \{1, 12, 9, 5, 6, 10\};
 int n = sizeof(arr) / sizeof(arr[0]);
 heapSort(arr, n);
 printf("Sorted array is \n");
 printArray(arr, n);
}
```