Lecture 1

# EECS 483: COMPILER CONSTRUCTION

# Administrivia

- **Instructor:** Max S. New
  **Office hours:** Mondays 4:00-5:00pm,
  Thu 10-11am
  & by appointment.
  Beyster 4628 or Zoom, use oh.eecs.umich.edu for queuing

- **GSI:**
  - Eric Giovannini OH TBD
- **IA:**
  - Tingting Ding OH TBD

- **Web site:** http://www.maxsnew.com/teaching/eecs-483-wn24
- **Piazza:** https://piazza.com/class/lpwv6v8j1st66o
- **Gradescope**:

Links are on Canvas

# Announcements

- MLK Day next week
  - No class on Monday, January 15th
- HW1: Hellocaml
  - available on the course web site soon
  - due Tuesday, January 23rd at 11:59pm
- Note: New curriculum this semester adapted from course at UPenn

# Why EECS 483?


- You will learn:
  - Practical applications of theory
  - Lexing/Parsing/Interpreters
  - How high-level languages are implemented in machine language
  - (A subset of) Intel x86 architecture
  - More about common compilation tools like GCC and LLVM
  - A deeper understanding of code
  - A little about programming language semantics & types
  - Functional programming in OCaml
  - How to manipulate complex data structures
  - How to be a better programmer

- Expect this to be a *very challenging*, implementation-oriented course.
  - Programming projects can take *tens* of hours per week…

# The EECS 483 Compiler

- Course projects
  - HW1: Hellocaml! (OCaml programming)
  - HW2: X86lite interpreter
  - HW3: LLVMlite compiler
  - HW4: Lexing, Parsing, simple compilation
  - HW5: Higher-level Features
  - HW6: Analysis and Optimizations

- Goal: build a complete compiler from a high-level, type-safe language to x86 assembly.

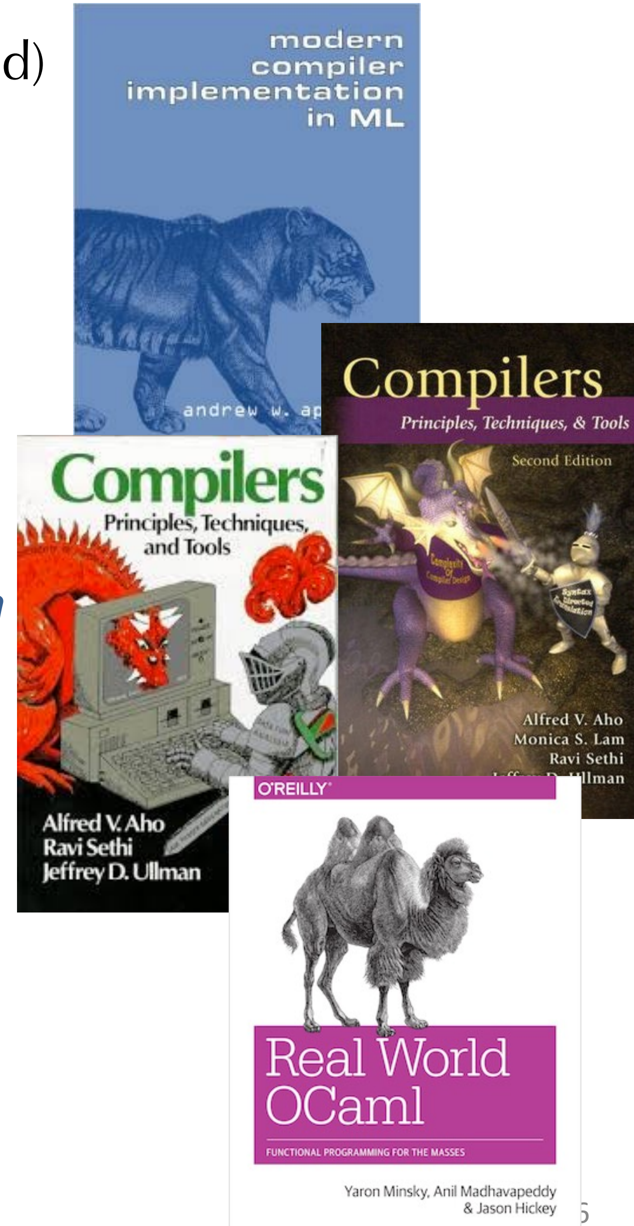
- Project will be graded on Linux (Gradescope's servers), assignments are most easily developed with it
  - Mac OS, including Apple Silicon should work too but might have minor issues
  - Windows users are recommended to use WSL (Windows Subsystem for Linux)

# Resources

- Course textbook: (recommended, not required)
  - *Modern compiler implementation in ML* (Appel)

- Additional compilers books:
  - *Compilers – Principles, Techniques & Tools* (Aho, Lam, Sethi, Ullman)
    - *a.k.a. "The Dragon Book"*
  - *Advanced Compiler Design & Implementation* (Muchnick)

- About Ocaml:
  - *Real World Ocaml* (Minsky, Madhavapeddy, Hickey)
    - realworldocaml.org
  - *Introduction to Objective Caml* (Hickey)



# Why OCaml?

- OCaml is a dialect of ML – "Meta Language"
  - It was designed to enable easy manipulation *abstract syntax trees*
  - Type-safe, mostly pure, functional language with support for polymorphic (generic) algebraic datatypes, modules, and mutable state
  - The OCaml compiler itself is well engineered
    - you can study its source!
  - It is the right tool for this job

- Never used OCaml?
  - Next couple lectures will (re)introduce it
  - First two projects will help you get up to speed programming
  - See "Introduction to Objective Caml" by Jason Hickey
    - book available on the course web pages, referred to in HW1

# HW1: Hellocaml

- Homework 1 available now on the course web site.
  - Individual project – no groups
  - Due: *Tuesday, 23 Jan. 2024 at 11:59pm*
  - Topic: OCaml programming, an introduction to interpreters
- We recommend using VSCode + OCaml Platform
  - See the course web pages about the tool chain to get started
- Quickstart guide:
  - open up the project in VSCode
  - start a "sandbox terminal" via OCaml Platform plugin
  - type `make test` at the command prompt

  - Please: Use Piazza to report any troubles with the toolchain!

# Homework Policies

- Homework (except HW1) should be done individually or in pairs
- Late projects:
  - up to 24 hours late:  10 point penalty
  - up to 48 hours late:  20 point penalty
  - after 48 hours: not accepted
- Submission policy:
  - Projects that don't compile will get no credit
  - Partial credit will be awarded according to the guidelines in the project description
- Academic integrity: don't cheat
  - This course will abide by the University's Honor Code
  - "low level" and "high level" discussions across groups are fine
  - "mid level" discussions / code sharing are not permitted
  - General principle: *When in doubt, ask!*

# Course Policies

Prerequisites: EECS 370

- Significant programming experience
- Some assembly experience
- If HW1 is a struggle, this class might not be a good fit for you (HW1 is significantly simpler than the rest…)

Grading:

- 70% Projects: *Compiler*
  - Groups of 2 students (besides HW1)
  - Implemented in Ocaml
  - Tentatively scheduled to be due on Tuesdays at 11:59pm

- 12% Midterm: in the evening… tentatively March 11$^{th}$
- 18% Final exam: cumulative, but more emphasis on latter half of the course

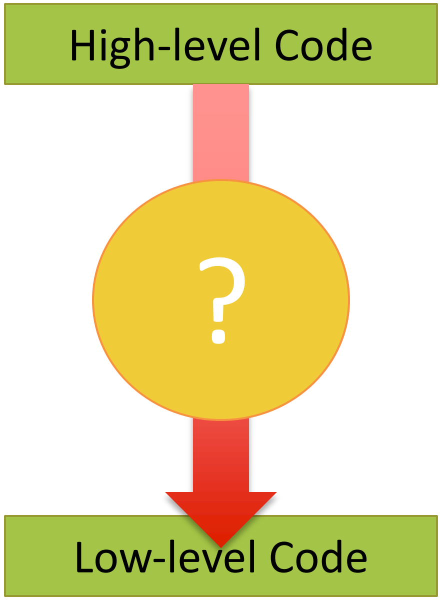
- Lecture attendance is crucial
  - Active participation (asking questions, etc.) is encouraged

What is a compiler?

# COMPILERS

# What is a Compiler?

- A compiler is a program that translates from one programming language to another.
- Typically: *high-level source code* to *low-level machine code* (object code)
  - Not always: Source-to-source translators, Java bytecode compiler, GWT Java ⇒ Javascript



# Historical Aside

- This is an old problem!
- Until the 1950's: computers were programmed in assembly.
- 1951—1952: Grace Hopper
  - developed the A-0 system for the UNIVAC I
  - She later contributed significantly to the design of COBOL
- 1957: FORTRAN compiler built at IBM
  - Team led by John Backus
- 1960's: development of the first bootstrapping compiler for LISP
- 1970's: language/compiler design blossomed
- Today: *thousands* of languages (most little used)
  - Some better designed than others...

1980s: ML / LCF
1984: Standard ML
1987: Caml
1991: Caml Light
1995: Caml Special Light
1996: Objective Caml
2005: F# (Microsoft)
2015: Reason ML
2020: OCaml Platform

# Source Code

- Optimized for human readability
  - *Expressive*: matches human ideas of grammar / syntax / meaning
  - *Redundant*: more information than needed to help catch errors
  - *Abstract*: exact computation possibly not fully determined by code
- Example C source:

```
#include <stdio.h>

int factorial(int n) {
  int acc = 1;
  while (n > 0) {
    acc = acc * n;
    n = n - 1;
  }
  return acc;
}

int main(int argc, char *argv[]) {
  printf("factorial(6) = %d\n", factorial(6));
}
```

# Low-level code

- Optimized for Hardware
  - Machine code hard for people to read
  - Redundancy, ambiguity reduced
  - Abstractions & information about intent is lost

- Assembly language
  - then machine language

- Figure at right shows (unoptimized) 32-bit code x86 for the factorial function

```
_factorial:
## BB#0:
        pushl   %ebp
        movl    %esp, %ebp
        subl    $8, %esp
        movl    8(%ebp), %eax
        movl    %eax, -4(%ebp)
        movl    $1, -8(%ebp)
LBB0_1:
        cmpl    $0, -4(%ebp)
        jle     LBB0_3
## BB#2:
        movl    -8(%ebp), %eax
        imull   -4(%ebp), %eax
        movl    %eax, -8(%ebp)
        movl    -4(%ebp), %eax
        subl    $1, %eax
        movl    %eax, -4(%ebp)
        jmp     LBB0_1
LBB0_3:
        movl    -8(%ebp), %eax
        addl    $8, %esp
        popl    %ebp
        retl
```

# How to translate?

- Source code – Machine code mismatch
- Some languages are farther from machine code than others:
  - Consider: C, C++, Java, Lisp, ML, Haskell, Ruby, Python, Javascript

- Goals of translation:
  - Source level expressiveness for the task
  - Best performance for the concrete computation
  - Reasonable translation efficiency ($< O(n^3)$)
  - Maintainable code
  - Correctness!

# Correct Compilation

- Programming languages describe computation precisely…
  - therefore, *translation* can be precisely described
  - a compiler can be correct with respect to the source and target language semantics.

- Correctness is important!
  - Broken compilers generate broken code.
  - Hard to debug source programs if the compiler is incorrect.
  - Failure has dire consequences for development cost, security, etc.

- This course: some techniques for building correct compilers
  - *Finding and Understanding Bugs in C Compilers*,
    Yang et al. PLDI 2011

  - There is much ongoing research about *proving* compilers correct.
    (Google for CompCert, Verified Software Toolchain, or Vellvm)

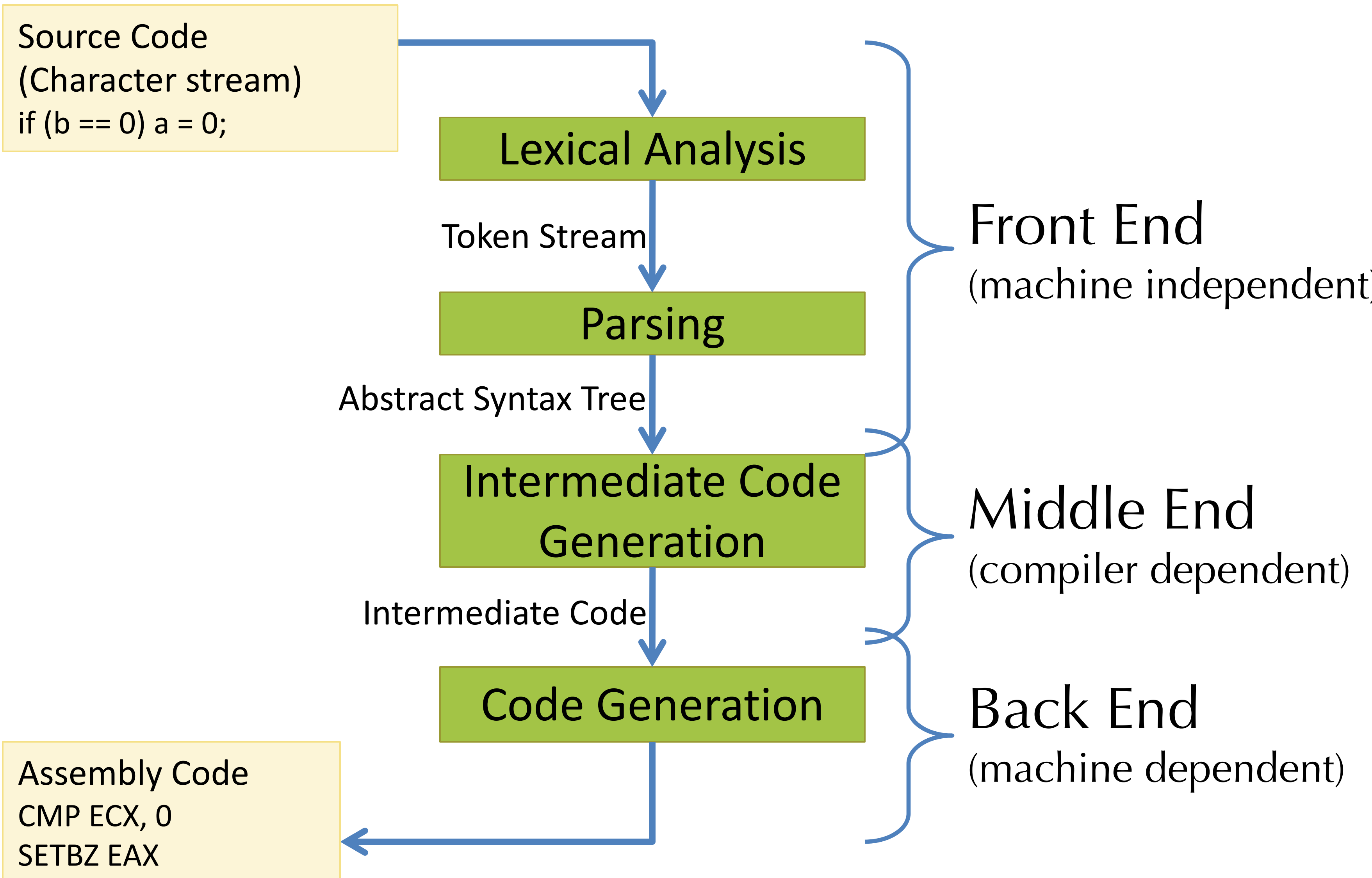
# Idea: Translate in Steps

- Compile via a series of program representations

- Intermediate representations are optimized for program manipulation of various kinds:
  - Semantic analysis: type checking, error checking, etc.
  - Optimization: dead-code elimination, common subexpression elimination, function inlining, register allocation, etc.
  - Code generation: instruction selection

- Representations are more machine specific, less language specific as translation proceeds

# (Simplified) Compiler Structure

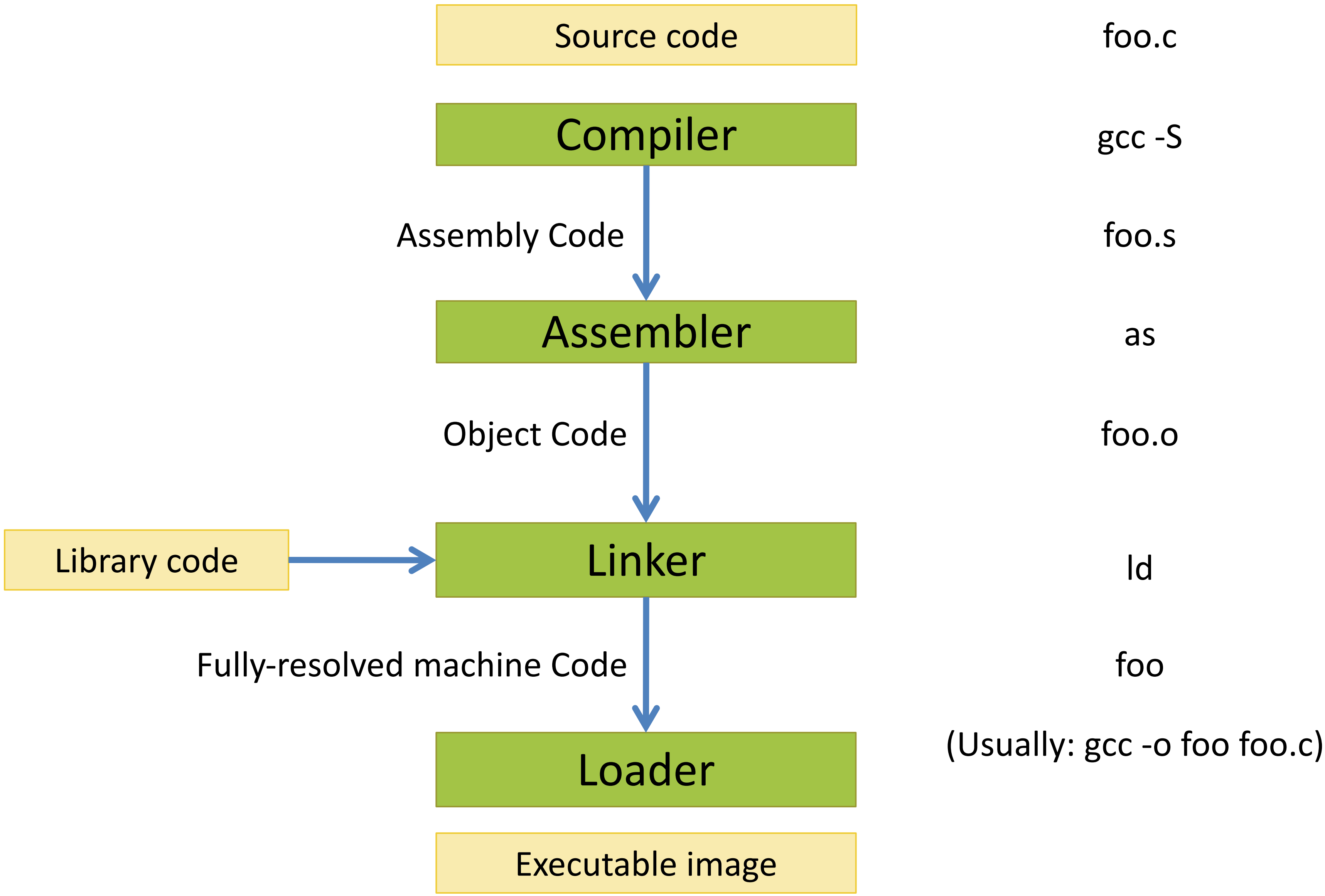


# Typical Compiler Stages

| Compiler Passes | | Representations of the program |
|---|---|---|
| Lexing | → | token stream |
| Parsing | → | abstract syntax |
| Disambiguation | → | abstract syntax |
| Semantic analysis | → | annotated abstract syntax |
| Translation | → | intermediate code |
| Control-flow analysis | → | control-flow graph |
| Data-flow analysis | → | interference graph |
| Register allocation | → | assembly |
| Code emission | | |

- Optimizations may be done at many of these stages
- Different source language features may require more/different stages
- Assembly code is not the end of the story

# Compilation & Execution



See lec01.zip

# COMPILER DEMO

# Short-term Plan

- Rest of today:
  - Intro to OCaml
  - “object language” vs. “meta language”
  - Build a simple interpreter

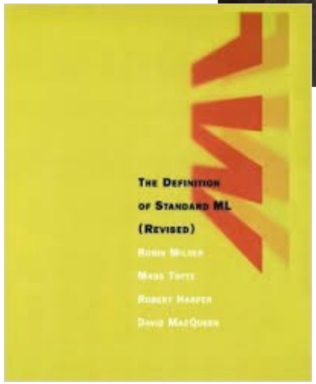
Introduction to OCaml programming

- A little background about ML
- Interactive tour of OCaml via UTop
- Writing simple interpreters

# OCAML

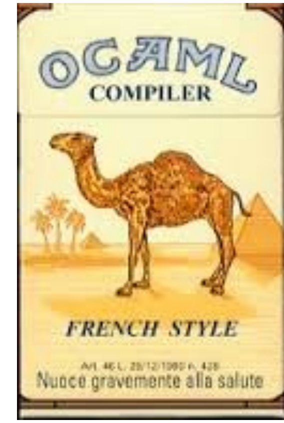
# ML's History

- **1971: Robin Milner** starts the LCF Project at Stanford
  - "logic of computable functions"
- **1973:** At Edinburgh, Milner implemented his theorem prover and dubbed it "Meta Language" – ML
- **1984:** ML escaped into the wild and became "Standard ML"
  - SML '97 newest version of the standard
  - There is a whole family of SML compilers:
    - SML/NJ – developed at AT&T Bell Labs
    - MLton – whole program, optimizing compiler
    - Poly/ML
    - Moscow ML
    - ML Kit compiler
    - MLj – SML to Java bytecode compiler
- ML 2000: failed revised standardization
- sML: successor ML – discussed intermittently
- **2014:** sml-family.org + definition on github



# OCaml's History

- The Formel project at the Institut National de Rechereche en Informatique et en Automatique (INRIA)
- **1987:** Guy Cousineau re-implemented a variant of ML
  - Implementation targeted the "Categorical Abstract Machine" (CAM)
  - As a pun, "CAM-ML" became "CAML"
- **1991:** Xavier Leroy and Damien Doligez wrote Caml-light
  - Compiled CAML to a virtual machine with simple bytecode (much faster!)
- **1996:** Xavier Leroy, Jérôme Vouillon, and Didier Rémy
  - Add an object system to create OCaml
  - Add native code compilation
- Many updates, extensions, since…
- **2005:** Microsoft's F# language is a descendent of OCaml
- **2013:** ocaml.org
- **2020:** OCaml Platform
- **2022:** Multicore OCaml




# OCaml Tools

- ocaml – the top-level interactive loop
- ocamlc – the bytecode compiler
- ocamlopt – the native code compiler
- ocamldep – the dependency analyzer
- ocamldoc – the documentation generator
- ocamllex – the lexer generator
- ocamlyacc – the parser generator

- menhir – a more modern parser generator
- dune – a compilation manager
- utop – a more fully-featured interactive top-level

- opam – package manager

# Distinguishing Characteristics

- Functional & (Mostly) "Pure"
  - Programs manipulate values rather than issue commands
  - Functions are first-class entities
  - Results of computation can be "named" using let
  - Has relatively few "side effects" (imperative updates to memory)

- Strongly & Statically typed
  - Compiler typechecks every expression of the program, issues errors if it can't prove that the program is type safe
  - Good support for type inference & generic (polymorphic) types
  - Rich user-defined "algebraic data types" with pervasive use of *pattern matching*
  - Very strong and flexible module system for constructing large projects

# Most Important Features for 483

- Types:
  - int, bool, int32, int64, char, string, built-in lists, tuples, records, functions
- Concepts:
  - Pattern matching
  - Recursive functions over algebraic (i.e. tree-structured) datatypes
- Libraries:
  - Int32, Int64, List, Printf, Format

How to represent programs as data structures.
How to write programs that process programs.

# INTERPRETERS

# Factorial: Everyone's Favorite Function

- Consider this implementation of factorial in a hypothetical programming language that we'll call "SIMPLE" (Simple IMperative Programming LanguagE):

```
X = 6;
ANS = 1;
whileNZ (x) {
        ANS = ANS * X;
        X = X + -1;
}
```

- We need to describe the constructs of this SIMPLE
  - ***Syntax***: which sequences of characters count as a legal "program"?
  - ***Semantics***: what is the meaning (behavior) of a legal "program"?

# ”Object” vs. “Meta” language

***Object language***:
the language (syntax / semantics) being described or manipulated

***Metalanguage***:
the language (syntax / semantics) used to *describe* some object language

| | |
|---|---|
| Today’s example: | |
| SIMPLE | interpreter written in OCaml |
| Course project: | |
| OAT ⇒ LLVM ⇒ x86_64 | compiler written in OCaml |
| Clang compiler: | |
| C/C++ ⇒ LLVM ⇒ x86_64 | compiler written in C++ |
| Metacircular interpreter: | |
| lisp | interpreter written in lisp |

# Grammar for a Simple Language

```
<exp> ::=
    |  <X>
    |  <exp> + <exp>
    |  <exp> * <exp>
    |  <exp> < <exp>
    |  <integer constant>
    |  (<exp>)

<cmd> ::=
    |  skip
    |  <X> = <exp>
    |  ifNZ <exp> { <cmd> } else { <cmd> }
    |  whileNZ <exp> { <cmd> }
    |  <cmd>; <cmd>
```

BNF grammars are themselves domain-specific metalanguages for describing the syntax of other languages…

- Concrete syntax (grammar) for a simple imperative language
  - Written in "Backus-Naur form"
  - *<exp>* and *<cmd>* are ***nonterminals***
  - '::=' , '|' , and <…> symbols are part of the *metalanguage*
  - keywords, like 'skip' and 'ifNZ' and symbols, like '{' and '+' are part of the *object language*

- Need to represent the ***abstract syntax*** (i.e. hide the irrelevant of the concrete syntax)
- Implement the ***operational semantics*** (i.e. define the behavior, or meaning, of the program)

# OCaml Demo

simple.ml