

Problem Set 2: Simple Type Theory

Released: January 23, 2023
Due: February 03, 2022, 11:59pm

Submit your solutions to this homework on Canvas in a group of 2 or 3. Your solutions must be submitted in pdf produced using LaTeX.

Problem 1 Exponential Isomorphisms

We say that $x : A \vdash M : B$ and $y : B \vdash N : A$ form an *isomorphism* if $x : A \vdash N[M/y] = x$ and $y : B \vdash M[N/x] = y : B$. In this case we say A and B are isomorphic, written $A \cong B$.

Construct the following isomorphisms (with proof):

1. $A \Rightarrow B \Rightarrow C \cong (A \times B) \Rightarrow C$
2. $A \Rightarrow (B \times C) \cong (A \Rightarrow B) \times (A \Rightarrow C)$
3. $(A \Rightarrow 1) \cong 1$
4. $(A + B) \Rightarrow C \cong (A \Rightarrow C) \times (B \Rightarrow C)$
5. $0 \Rightarrow C \cong 1$

This gives an idea of why $A \Rightarrow B$ is in category theory sometimes called the *exponential* and written B^A .

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Problem 2 Admissible Rules in STT

When all variables are known to be distinct, substitution $M[N/x]$ can simply be defined as the replacement of x with N everywhere in the term M . This definition is the STT version of the admissibility of the substitution principle of IPL:

$$\frac{\Gamma \vdash M : A \quad \Gamma, x : A \vdash N : B}{\Gamma \vdash N[M/x]} \text{SUBST}(\ast)$$

The admissible principle of contraction can also be viewed as a textual substitution in the term:

$$\frac{\Gamma, x : A, y : A, \Delta \vdash M : C}{\Gamma, x : A, \Delta \vdash M[x/y] : C} \text{CONTRACTION}(\ast)$$

On the other hand, the use of variables to stand for assumptions means that exchange and weakening have no effect on the proof term:

$$\frac{\Gamma, y : B, x : A, \Delta \vdash M : C}{\Gamma, x : A, y : B, \Delta \vdash M : C} \text{EXCHANGE}(\ast) \qquad \frac{\Gamma, x : A, \Delta \vdash M : C}{\Gamma, \Delta \vdash M : C} \text{WEAKENING}(\ast)$$

In this exercise you will show all of these principles are admissible and additionally prove some *equations* about substitution.

The simplest inductive proofs involve an auxiliary notion. We define a substitution from Δ to Γ to be a function γ that for each variable $x : A \in \Gamma$ produces a term $\Delta \vdash \gamma(x) : A$. We write $\gamma : \Delta \rightarrow \Gamma$ to mean a substitution from Δ to Γ . We can then define an admissible action of substitution:

$$\frac{\gamma : \Delta \rightarrow \Gamma \quad \Gamma \vdash M : A}{\Delta \vdash M[\gamma] : A} \text{GENSUBST}$$

Defined by induction on M :

$$\begin{aligned} x[\gamma] &= \gamma(x) \\ f(M_1, \dots,)[\gamma] &= f(M_1[\gamma], \dots) \\ (M, N)[\gamma] &= (M[\gamma], N[\gamma]) \\ (\pi_j M)[\gamma] &= \pi_j M[\gamma] \\ ()[\gamma] &= \gamma \\ (i_j M)[\gamma] &= i_j M[\gamma] \\ (\text{case}_+ M \{i_1 x_1 \rightarrow N_1 \mid i_2 x_2 \rightarrow N_2\})[\gamma] &= (\text{case}_+ M[\gamma] \{i_1 x_1 \rightarrow N_1[\gamma, x_1/x_1] \mid i_2 x_2 \rightarrow N_2[\gamma, x_2/x_2]\}) \\ (\text{case}_0 M \{\})[\gamma] &= \text{case}_0 M[\gamma] \{\} \\ (\lambda x. M)[\gamma] &= \lambda x. M[\gamma, x/x] \\ (MN)[\gamma] &= M[\gamma] N[\gamma] \end{aligned}$$

Where the notation $\gamma[M/x]$ is the extension of the the function to map x to M :

$$\begin{aligned} \gamma[M/x](y) &= M & (\text{if } x = y) \\ \gamma[M/x](y) &= \gamma(y) & (\text{if } x \neq y) \end{aligned}$$

Define the *identity* substitution $\text{id}_\Gamma : \Gamma \rightarrow \Gamma$ to map each variable in Γ to itself: $\text{id}(x) = x$.

Given $\gamma : \Delta \rightarrow \Gamma$ and $\delta : \Xi \rightarrow \Delta$, define the *composition* $\gamma \circ \delta : \Xi \rightarrow \Gamma$ as $(\gamma \circ \delta)(x) = \gamma(x)[\delta]$.

Below assume $\gamma : \Delta \rightarrow \Gamma$, $\delta : \Xi \rightarrow \Delta$, $\xi : \Xi' \rightarrow \Xi$ and $\Gamma \vdash M : A$.

1. Show M that $\Delta \vdash M[\gamma] : A$, i.e., that the GenSubst typing rule is admissible.
2. Show that one-place substitution, weakening, exchange and contraction are all instances of GenSubst.
3. Show that $M[\text{id}_\Gamma] = M$. Note that this and the following equalities are exact syntactic equalities, you will not need to use any $\beta\eta$ rules to prove it.
4. Show that $\gamma \circ \text{id}_\Delta = \gamma$ and $\text{id}_\Gamma \circ \gamma = \gamma$.
5. Show that $M[\gamma \circ \delta] = M[\gamma][\delta]$
6. Show as a corollary that if $x_2 : A_2 \vdash N_1 : A_1$ and $x_3 : A_3 \vdash N_2 : A_2$ and $x_4 : A_4 \vdash N_3 : A_4$ then $(N_1[N_2/x_2])[N_3/x_3] = (N_1[N_2/x_2])[N_3/x_3]$.
7. Show that $(\gamma \circ \delta) \circ \xi = \gamma \circ (\delta \circ \xi)$

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