

ARITHMETICA UNIVERSALIS

PROOF OF FUNDAMENTAL THEOREM OF ARITHMETICA UNIVERSALIS

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To Misfits...

Proof of Arithmetic Boolean Satisfiability Equation...

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1 ARITHMETIC BOOLEAN SATISFIABILITY EQUATION

Theorem.

$$\sum_{j=0}^{m-1} 2^{\sum_{i=0}^{n-1} x_{(n-1-i)j} 2^i} = (x_{0,0}^{\pm} \vee \dots \vee x_{n-1,0}^{\pm}) \wedge \dots \wedge (x_{0,m}^{\pm} \vee \dots \vee x_{n-1,m-1}^{\pm})$$

Proof. A CNF clause, is *False* only when all elements are *False*, then if *False* $\rightarrow 1$ and *True* $\rightarrow 0$ you can write any CNF clause how $(2^{|X|} - 1) - 1$ where $X = \{x_0 \dots x_{n-1}\}$ the set of variables and the lowest significant bit, correspond to the assignment $\{x_0 = \text{False} \rightarrow 1 \dots x_{n-1} = \text{False} \rightarrow 1\}$, repeating this process, covering all combinations:

$$\begin{array}{rcl} 0 & \underbrace{00 \dots 00}_{|X|} & \{x_0 = \text{True} \rightarrow 0 \dots x_{n-1} = \text{True} \rightarrow 0\} \\ 1 & \underbrace{00 \dots 01}_{|X|} & \{x_0 = \text{True} \rightarrow 0 \dots x_{n-1} = \text{False} \rightarrow 1\} \\ \vdots & \vdots & \vdots \\ 2^{|X|} - 2 & \underbrace{11 \dots 10}_{|X|} & \{x_0 = \text{False} \rightarrow 1 \dots x_{n-1} = \text{True} \rightarrow 0\} \\ 2^{|X|} - 1 & \underbrace{11 \dots 11}_{|X|} & \{x_0 = \text{False} \rightarrow 1 \dots x_{n-1} = \text{False} \rightarrow 1\} \end{array}$$

The *True* values of a CNF clause, is given by the positive assignments of each variable, then a particular clause have *True* values according to $\sum_{i=0}^{n-1} x_{(n-1-i)j} 2^i$ this is the binary number representation of their assignment, and then the entire CNF formula is given by the sum of this assignments over all clauses of the formula.

This proof the theorem. □

2 IMPLEMENTATION (Python)

```
def bits(n, p):
    s = []
    while n:
        s = [n % 2 == 0] + s
        n //= 2
    s = (p - len(s)) * [True] + s
    return s

def sat_equation(cnf, n, m):
    sat = 0
    for j in range(m):
        v = 0
        for i in range(n):
            v += int(cnf[j][n - 1 - i] > 0) * 2 ** i
        sat += 2 ** v
    return sat

if __name__ == '__main__':
    cnf = [(1, 2, 3, 4), (1, -2, 3, -4), (-1, -2, 3, 4)]

    n = 4
    m = len(cnf)

    print(bits(sat_equation(cnf, n, m), 2 ** n))
```

3 DISCLAIMER

All this work is the effort of 5 years of research, everything is completely original and of my authority, no reference exist, and no exist collaborators, only elementary concepts of the theory of complexity were used. Special thanks to my wife, my favorite persons, all my friends and followers, thank for all...

4 LINKS

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<http://twitter.com/maxtuno>
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<http://www.reverbnation.com/maxtuno>
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http://github.com/maxtuno/SAT_EQUATION

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