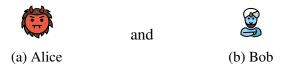
The Tiling Method

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Consider a two-party communication problem, in which the participants



participate to compute a function:

$$f: \underbrace{\mathbb{B}^n}_{\text{Alice's Bob's global}} \rightarrow \underbrace{\mathbb{B}}_{\text{output}}$$

$$\text{input input output}$$

The players can come up with a *protocol* $\Pi = (p_1, ..., p_t)$, namely, for some natural $t \in \mathbb{N}$, a sequence of t-many functions $p_i : \mathbb{B}^* \to \mathbb{B}^*$ such that the communication between the players looks like this:

Alice is given input x.

Hi Bob. I'm not divulging x, but, $p_1(x) = P1$.

Bob is given input y.

Thanks Alice. I'm not divulging y, but, $p_2(y,P1) = P2$.

Thanks Bob. Don't tell anyone, but: $p_3(x,P1,P2) = P3$.

Is that so? Well, $p_4(y,P1,P2,P3) = P4$.

Wicked. In that case, $p_5(x,P1,P2,P3,P4) = P5$.

...yada yada yada...

In that case, the last thing you need to know is that $p_1(x,P1,P2,P3,P4,...,P(t-1)) = Pt$.

Suppose that there is a protocol Π for f consisting of t messages, but, there does not exist any protocol Π' for f consisting of fewer than t messages. Then we say t is the *communication complexity* of f, and we write C(f) = t.

Given some such function f, it would be nice if we could automatically compute a reasonable lower bound on its communication complexity. One way to do this is with the *tiling method*. We will give the method immediately, and in tandem, we will illustrate the method using the function f(x,y) = x < y where x,y are integers in $\{0,1,2,3\}$, encoded in \mathbb{B} oolean. First, let M(f) be the *matrix of f*, namely, the $2^n \times 2^n$ matrix whose (x,y)th entry is the value f(x,y).

	0	1	2	3	4
0	0	1	1	1	1
1	0	0	1	1	1
2	0	0	0	1	1
3	0	0	0	0	1
4	0	0	0	0	0

Table 1: The matrix f(<) for inputs $x, y \in \{0, 1, 2, 3\}$. Values of x are given in the rows, while values of y are given in the columns. False (i.e. 0) values are marked red for clarity.

TODO...