

2's Complement

Signed \rightarrow 1010 \rightarrow 1011 (only convert if negative)

0101 $+1$

0100 \rightarrow 1011 \rightarrow 1100 $+1$

Reversing:

Copy 0 ant 1st / bit

Compliment everything else

1's

As we will see with the other 2 representations, there is an easier way when you are dealing with binary numbers.

Table 5.1 Interpretation of four-bit signed integers.

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	-0	-7	-8
1001	-1	-6	-7
1010	-2	-5	-6
1011	-3	-4	-5
1100	-4	-3	-4
1101	-5	-2	-3
1110	-6	-1	-2
1111	-7	-0	-1

1's

$$\begin{array}{r} (+5) \\ + (+2) \\ \hline (+7) \end{array} \quad \begin{array}{r} 0101 \\ + 0010 \\ \hline 0111 \end{array} \quad \begin{array}{r} (-5) \\ + (+2) \\ \hline (-3) \end{array} \quad \begin{array}{r} 1010 \\ + 0010 \\ \hline 1100 \end{array}$$

$$\begin{array}{r} (+5) \\ + (-2) \\ \hline (+3) \end{array} \quad \begin{array}{r} 0101 \\ + 1101 \\ \hline 10010 \\ \text{Carry } 1 \rightarrow \\ \hline 0011 \end{array} \quad \begin{array}{r} (-5) \\ + (-2) \\ \hline (-7) \end{array} \quad \begin{array}{r} 1010 \\ + 1101 \\ \hline 10111 \\ \text{Carry } 1 \rightarrow \\ \hline 1000 \end{array}$$

2's

$$\begin{array}{r} (+5) \\ + (+2) \\ \hline (+7) \end{array} \quad \begin{array}{r} 0101 \\ + 0010 \\ \hline 0111 \end{array} \quad \begin{array}{r} (-5) \\ + (+2) \\ \hline (-3) \end{array} \quad \begin{array}{r} 1011 \\ + 0010 \\ \hline 1101 \end{array}$$

$$\begin{array}{r} (+5) \\ + (-2) \\ \hline (+3) \end{array} \quad \begin{array}{r} 0101 \\ + 1110 \\ \hline 10011 \\ \text{Carry } 1 \rightarrow \text{ignore} \end{array} \quad \begin{array}{r} (-5) \\ + (-2) \\ \hline (-7) \end{array} \quad \begin{array}{r} 1011 \\ + 1110 \\ \hline 11001 \\ \text{Carry } 1 \rightarrow \text{ignore} \end{array}$$

2's subtraction

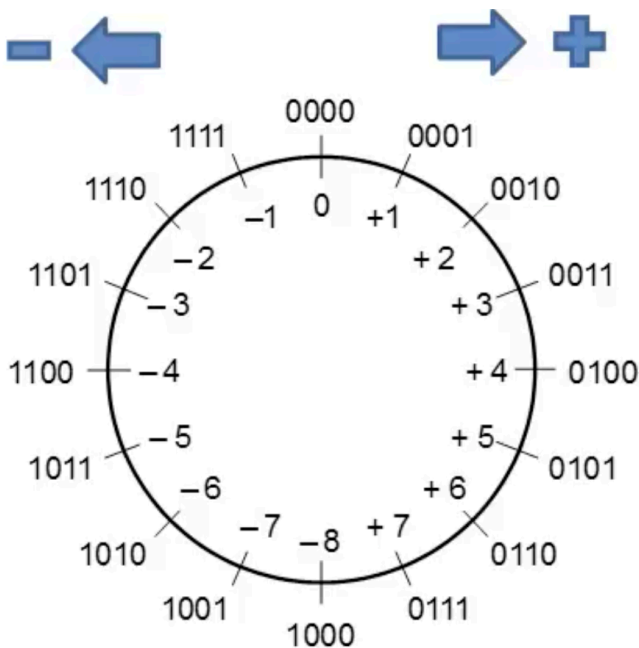
$$\begin{array}{r} (+5) \\ - (+2) \\ \hline (+3) \end{array} \quad \begin{array}{r} 0101 \\ - 0010 \\ \hline \end{array} \Rightarrow \begin{array}{r} 0101 \\ + 1110 \\ \hline 10011 \\ \text{Carry } 1 \rightarrow \text{ignore} \end{array}$$

$$\begin{array}{r} (-5) \\ - (+2) \\ \hline (-7) \end{array} \quad \begin{array}{r} 1011 \\ - 0010 \\ \hline \end{array} \Rightarrow \begin{array}{r} 1011 \\ + 1110 \\ \hline 11001 \\ \text{Carry } 1 \rightarrow \text{ignore} \end{array}$$

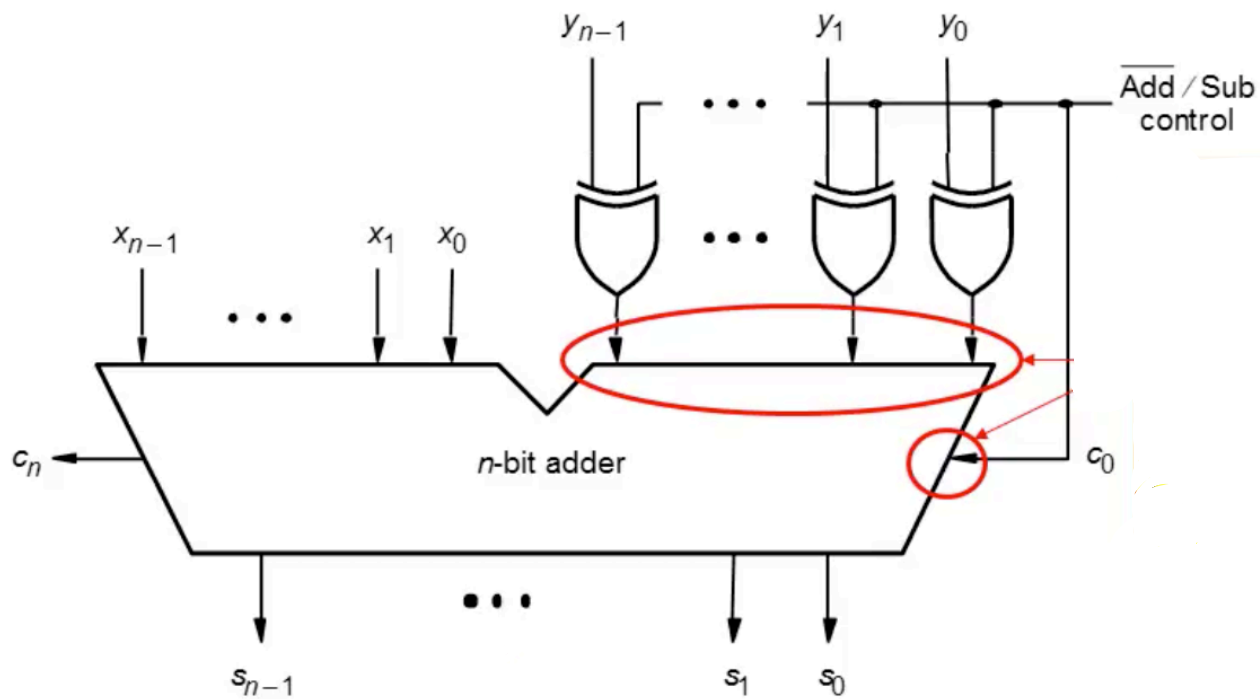
$$\begin{array}{r} (+5) \\ - (-2) \\ \hline (+7) \end{array} \quad \begin{array}{r} 0101 \\ - 1110 \\ \hline \end{array} \Rightarrow \begin{array}{r} 0101 \\ + 0010 \\ \hline 0111 \end{array}$$

$$\begin{array}{r} (-5) \\ - (-2) \\ \hline (-3) \end{array} \quad \begin{array}{r} 1011 \\ - 1110 \\ \hline \end{array} \Rightarrow \begin{array}{r} 1011 \\ + 0010 \\ \hline 1101 \end{array}$$

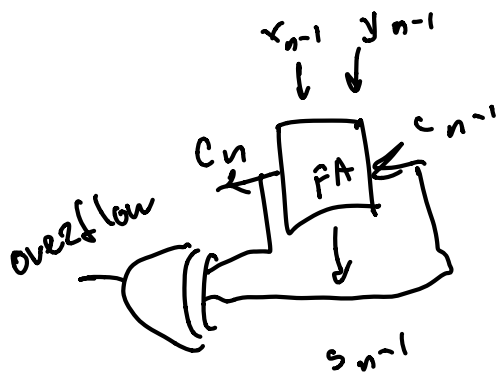
Take 2's



!Add / Sub - (bit) 0 - perform addition, 1 - subtraction



Result within: -2^{n-1} and $2^{n-1} - 1$
 Overflow when $c_{n-1} \oplus c_n$



Critical-path delay - for longest path of gates
 $2n$ - delay for n-bit adder
 $2n+1$ - for subtractor
 $2n+2$ - + overflow detector