Animation of Functional Specifications with PVSio

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¹Based on material by César A. Muñoz.

PVS as a Functional Programming Language



Most specifications in PVS are functional, e.g.,

```
sqrt_newton(a:nnreal,n:nat): RECURSIVE posreal =
   IF n=0 THEN a+1
   ELSE LET r=sqrt_newton(a,n-1) IN
     (1/2)*(r+a/r)
   ENDIF
   MEASURE n+1
```

■ What is the value of sqrt_newton(2,10)?

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Animation

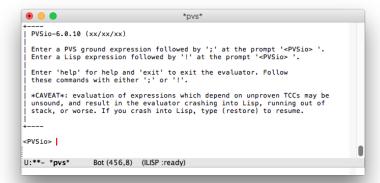


- Animation is the process of executing a specification to validate its intended semantics.
- Why: It is cheaper, faster, and more fun to test a specification than to prove it.
- How: PVSio.

PVSio is ...



- an read-eval-loop interface to the PVS Ground Evaluator;
- an efficient and sound mechanism to compute within the theorem prover;
- part of the PVS distribution;
- available as the standalone Unix command pvsio or through the Emacs command M-x pvsio.



<PVSio> sqrt_newton(2,10);

==>

<PVSio> sqrt_newton(2,10); ==>

48471731623224703430657/7555721707723979449648392625272648 2751814809171189000327511471361082396689939198281075027256...

PVSio Capabilities



- A predefined set of PVS functions for input/output operations, side-effects, unbounded-loops, exceptions, string manipulations, and floating point arithmetic
- 2 A high level interface for extending PVS programming language features.
- A tool for rapid prototyping.
- 4 An efficient strategy for evaluating ground expressions.

Contents



- 1 Input/Output Operations
- 2 Loops and Iteration
- 3 Exceptions
- 4 Local and Global Variables
- 5 PVS Parsing and Typechecking
- **6** Extending PVSio Programming Features
- Rapid Prototyping
- B PVSio and the PVS Theorem Prover



Output to screen

Basic output: print & println

println(s:string): void
println(r:real): void
println(b:bool): void

void

- just a rename for bool
- used to mark possible occurence of side effects

PVS source: [prelude] stdio & [prelude] stdlang



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PVS source: [prelude] stdio & [prelude] stdlang



```
<PVSio> print("sqrt_newton of 2: " + sqrt_newton(2,10));
```



```
<PVSio> print("sqrt_newton of 2: " + sqrt_newton(2,10));
sqrt_newton of 2: 1.4142135
```



```
<PVSio> print("sqrt_newton of 2: " + sqrt_newton(2,10));
sqrt_newton of 2: 1.4142135
<PVSio> print(sq(sqrt_newton(2,10)));
```



```
<PVSio> print("sqrt_newton of 2: " + sqrt_newton(2,10));
sqrt_newton of 2: 1.4142135
<PVSio> print(sq(sqrt_newton(2,10)));
2.0
```



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<PVSio> print("sqrt_newton of 2: " + sqrt_newton(2,10));
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<PVSio> sq(sqrt_newton(2,10)) = 2.0;
```



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==>
FALSE
```



```
<PVSio> print("sqrt_newton of 2: " + sqrt_newton(2,10));
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==>
FALSE
<PVSio> sq(sqrt_newton(2,10));
```



```
<PVSio> print("sqrt_newton of 2: " + sqrt_newton(2,10));
sqrt_newton of 2: 1.4142135
<PVSio> print(sq(sqrt_newton(2,10)));
2.0
<PVSio> sq(sqrt_newton(2,10)) = 2.0;
==>
FALSE
<PVSio> sq(sqrt_newton(2,10));
==>
11417786104914273667156938719275144887173049593916168
1541156469937779509690251478413369...
```

- Similar to Lisp's format function ≈ (format nil s t)
- s is the control string
 - a program in a syntax-based language
 - optimized for compactness rather than easy comprehension
- t are the values to print
 - given as a single value or as a tuple of values
- Additionally, the function
 printf(s:string,t:T):void
 prints the result of format(s,t)



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- Additionally, the function
 printf(s:string,t:T):void
 prints the result of format(s,t)



Output to screen – Example

```
<PVSio> format("The half of four is ~a.".4/2):
```

530617008867600505092775584086034866316307624567599571273
090520553619648095761832386318805390738103277561823284281
325003132706371396517146582357529867417615905908665879066
853985666554028115870511326582300341866167304359343960...



```
<PVSio> format("The half of four is ~a.",4/2);
==>
"The half of four is 2."
```

<PVSio> format("The half of four is ~a.",4/2);



```
==>
"The half of four is 2."
<PVSio> format("The half of ~a is ~a.",("four",4/2));
```



```
<PVSio> format("The half of four is ~a.",4/2);
==>
"The half of four is 2."
<PVSio> format("The half of ~a is ~a.",("four",4/2));
==>
"The half of four is 2."
```



```
<PVSio> format("The half of four is ~a.",4/2);
==>
"The half of four is 2."
<PVSio> format("The half of ~a is ~a.",("four",4/2));
==>
"The half of four is 2."
<PVSio> format("sqrt_newton of 2: ~a",sqrt_newton(2,10));
```



```
<PVSio> format("The half of four is ~a.".4/2):
==>
"The half of four is 2."
<PVSio> format("The half of ~a is ~a.".("four".4/2));
==>
"The half of four is 2."
<PVSio> format("sqrt_newton of 2: ~a",sqrt_newton(2,10));
==>
"sqrt_newton of 2: 10685404112580054249577309962027702517
530617008867600505092775584086034866316307624567599571273
090520553619648095761832386318805390738103277561823284281
325003132706371396517146582357529867417615905908665879066
853985666554028115870511326582300341866167304359343960...
```

Basic directives

- ~% new line, ~& fresh line, ~~ a tilde (no data consumption)
- ~a outputs next data in human-readable form
- ~d ~x ~o ~b allow to format integer values
- ~r ~:r print numbers as English words
- ~@r ~:@r print numbers as Roman numerals

Conditional Formatting

" ~[··· ~; ··· ~] uses next datum to index such list

```
format("\sim[zero\sim;um\sim;dois\sim]", 0) \rightarrow "zero" format("\sim[zero\sim;um\sim;dois\sim]", 1) \rightarrow "um" format("\sim[zero\sim;um\sim;dois\sim]", 2) \rightarrow "dois"
```



PVSio provides the outfix operator {| |} to use format directives on
 PVS lists and PVS boolean values

Lists

```
<PVSio> LET numbers = (:1,2,3:) IN
format("-{-a-^, -}",{|numbers|});
==>
"1, 2, 3"
```

■ Boolean values

```
<PVSio> LET b = true IN
     format("~:[falso~;verdade~]",{|b|})
==>
"verdade"
```



- PVSio provides the outfix operator {| |} to use format directives on
 PVS lists and PVS boolean values
- Lists

```
<PVSio> LET numbers = (:1,2,3:) IN
format("~{~a~^, ~}",{|numbers|});
==>
"1, 2, 3"
```

Boolean values

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<PVSio> LET b = true IN
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- PVSio provides the outfix operator {| |} to use format directives on
 PVS lists and PVS boolean values
- Lists

```
<PVSio> LET numbers = (:1,2,3:) IN
format("~{~a~^, ~}",{|numbers|});
==>
"1, 2, 3"
```

Boolean values

```
<PVSio> LET b = true IN
format("~:[falso~;verdade~]",{|b|});
==>
"verdade"
```



PVSio supports a special directive to print numbers in decimal form

- ~-n/pvs:d/
 - \blacksquare where n is the amount of fractional digits
- Example



- PVSio supports a special directive to print numbers in decimal form
- ~-n/pvs:d/
 - where *n* is the amount of fractional digits
- Example

Special cases

- PVSio supports a special directive to print numbers in decimal form
- ~-n/pvs:d/
 - where *n* is the amount of fractional digits
- Example



Input from stdin

| no prompt | prompt | Description |
|----------------|---------------------|--|
| read_real | query_real(msg) | Reads a real number |
| read_int | query_int(msg) | Reads an integer |
| read_word | query_word(msg) | Reads a word |
| read_bool(ans) | query_bool(msg,ans) | Checks if the entered word is equal to ans |
| read_line | query_line(msg) | Reads the whole line |
| read_token(s) | query_token(msg,s) | Returns the smaller prefix that ends in any of the chars in \boldsymbol{s} |

msg,s,ans: VAR string



















Streams

- As usual, PVSio Streams have kind & direction
 - Standard, File, String
 - Input, Ouput
- Ad-Hoc Datatypes & Constants

```
Stream : TYPE+ stdin : IStream

IStream: TYPE+ FROM Stream stdout : OStream

OStream: TYPE+ FROM Stream stderr : OStream
```

■ Functions, being f:VAR Stream

| fopen?(f):bool | Checks if the stream is open | |
|---------------------|-----------------------------------|--|
| strstream?(f):bool | | |
| filestream?(f):bool | Check the kind of stream | |
| sdtstream?(f):bool | | |
| finput?(f):bool | Check the direction of the stream | |
| fouput?(f):bool | Check the direction of the stream | |



■ Kind and direction are represented by the enumerated type Mode

```
■ Mode : TYPE = {input,output,create,append,overwrite,rename,str}
```

More Functions

Streams

```
      fopenin(m:Mode,s:string)
      : IStream
      Opens an input stream in mode m

      fopenin(s:string)
      : string
      Opens an input stream from file s

      fopenout(m:Mode,s:string)
      : OStream
      Opens an output stream in mode m

      eof?(f:IStream)
      : bool
      Checks if the stream has been completely consumed

      flength(f:Stream)
      : nat
      Returns the length of the stream
```



Read

Streams

Write

```
fprint(f:OStream,s:string)
                                           Writes the string s to the stream f
                                  : void
                                           Writes the real number r to the stream f
fprint(f:OStream,r:real)
                                  : void
fprint(f:OStream,b:bool)
                                           Writes the boolean value b to the stream f
                                  · void
fprintln(f:OStream,s:string)
                                  : void
                                           Writes the string s on a new line in f
fprintln(f:OStream,r:real)
                                  : void
                                           Writes the real number r on a new line in f
                                           Writes the boolean value b on a new line in f
fprintln(f:OStream,b:bool)
                                  · void
```



Input from file - Examples



Input from file - Examples

If the content of the file "dez.txt" are

```
line 1 10
```



Input from file - Examples

If the content of the file "dez.txt" are

```
line 1 10
```

The contents of the file "sqdez.txt" will be

line 1 3.1622777



Even more functions

Streams

| fcheck(f:IStream) | : bool | Checks if the stream is open and did not reach eof |
|------------------------------------|----------|--|
| <pre>fname(f:Stream)</pre> | : string | Returns the full name of the file stream f |
| fgetpos(f:Stream,n:nat) | : nat | Returns current position of the file stream f |
| <pre>fsetpos(f:Stream,n:nat)</pre> | : void | Set current position of file stream f |
| echo(f:OStream,s:string) | : void | Prints f to s and echoes to stdout |
| echo(f:OStream,r:real) | : void | Prints r to s and echoes to stdout |
| echo(f:OStream,b:bool) | : void | Prints b to s and echoes to stdout |
| echoln(f:OStream,s:string) | : void | Prints f to s in a new line and echoes to stdout |
| echoln(f:OStream,r:real) | : void | Prints r to s in a new line and echoes to stdout |
| echoln(f:OStream,b:bool) | : void | Prints b to s in a new line and echoes to stdout |

Loops and Iteration



- Bounded loops
 - for i = n to m do <statement>
 - Support for proofs of correctness
- Unbounded loops
 - Pragmatic approach
 - while(true) do <statement>



■ Functional version $(m \le n)$

$$f(n, f(...f(m+1, f(m, a))...))$$

Imperative version

```
local a : T := init;
local i : int;
for (i := m; i <= n; i++) {
   a := f(i,a);
}
return a;</pre>
```

```
for[T:TYPE](m,n:int,init:T,f:[subrange(m,n),T]->T) : T
```



```
\%% a = 1;
%% for (i=1; i <= n; i++) {
% a = a*x;
%% }
expit(x:real,n:nat): real =
  for[real](1,n,1,LAMBDA(i:subrange(1,n),a:real):a*x)
<PVSio> expit(2,10);
==>
1024
```



■ Functional version $(m \le n)$

$$f(m, f(...f(n-1, f(n, a))...))$$

Imperative version

```
local a : T := init;
local i : int;
for (i := n; i >= m; i--) {
   a := f(i,a);
}
return a;
```

```
for_down[T:TYPE](n,m:int,init:T,f:[subrange(m,n),T]->T) : T
```



```
for_down function - example
```

```
\%% a = 1;
%% for (i=n; i >= 1; i-) {
%% a = a*i;
%% }
factit(n:nat) : nat =
  for_down[nat](n,1,1,LAMBDA(i:subrange(1,n),a:nat):a*i)
<PVSio> factit(10);
==>
3628800
```



■ Functional version $(m \le n)$

$$(\cdots((f(m)\circ f(m+1))\circ f(m+2))\circ\cdots f(n))$$

Imperative version

```
local a : T = f(m);
local i : int;
for (i := m+1; i <= n; i++) {
   a := a o f(i)
}
return a;</pre>
```

```
iterate_left[T:TYPE](m,n:int,f:subrange(upfrom,upto)->T,o:[[T,T]->T]) : T
```

```
\% a = nth(1,0);
%% for (i=1;i<=length(l)-1;i++) {</pre>
%%
  a = \max(a, nth(1, i))
%% }
maxit(l:(cons?[real])) : real =
  iterate_left(0,length(1)-1,
                 LAMBDA(i:below(length(1))):nth(1,i),max)
<PVSio> maxit((:2,3,4,1,2:));
==>
4
```



Functional version $(m \le n)$

$$f(m) \circ (\cdots (f(n-2) \circ (f(n-1) \circ f(n))) \cdots)$$

Imperative version

```
local a : T = f(n);
local i : int;
for (i := n-1; i >= m; i-) {
   a := f(i) o a
}
return a;
```

```
iterate_right[T:TYPE](m,n:int,f:subrange(upfrom,upto)->T,o:[[T,T]->T]) : T
```

```
\% a = nth(1,0);
%% for (i=1;i<=length(l)-1;i++) {</pre>
\% a = min(nth(1,i),a)
%% }
minit(1:(cons?[real])) : real =
  iterate_right(0,
                length(1)-1,
                LAMBDA(i:below(length(1))):nth(1,i),
                min )
<PVSio> minit((:2,3,4,1,2:));
==>
```

Unbounded Loops



Previous definitions are not suitable for unbounded calculations

```
while(b:bool,s:void) : void
```

- Example: reads a file one line at a time
 - As in the cat unix command

```
cat : void =
  let f=fopenin("pvsio_examples.pvs") in
  while(not eof?(f),println(fread_line(f)))
& fclose(f)
```

Exceptions



- PVSio also provides support for exception handling
- Mechanism to respond to the occurrence of exceptional events
 - often changing the normal flow of program execution
- Usually used in input/output operations



```
int_aux : int =
  let i = query_int("Enter a number less than 10") in
  if i > 10 then throw("GreaterThan10")
  else i endif
```



```
int_aux : int =
  let i = query_int("Enter a number less than 10") in
  if i > 10 then throw("GreaterThan10")
  else i endif
readupto10 : int =
  catch[int]((:NotAnInteger, "GreaterThan10":),
       int_aux,0)
```



```
int_aux : int =
  let i = query_int("Enter a number less than 10") in
  if i > 10 then throw("GreaterThan10")
  else i endif
readupto10 : int =
  catch[int]((:NotAnInteger, "GreaterThan10":),
       int_aux,0)
<PVSio> readupto10;
```

```
Enter a number less than 10
15
==>
```



```
int_aux : int =
  let i = query_int("Enter a number less than 10") in
  if i > 10 then throw("GreaterThan10")
  else i endif
readupto10 : int =
  catch[int]((:NotAnInteger, "GreaterThan10":),
       int_aux,0)
<PVSio> readupto10;
Enter a number less than 10
```



```
int_aux : int =
  let i = query_int("Enter a number less than 10") in
  if i > 10 then throw("GreaterThan10")
  else i endif
readupto10 : int =
  catch[int]((:NotAnInteger, "GreaterThan10":),
       int_aux,0)
<PVSio> readupto10;
Enter a number less than 10
15
```



```
int_aux : int =
  let i = query_int("Enter a number less than 10") in
  if i > 10 then throw("GreaterThan10")
  else i endif
readupto10 : int =
  catch[int]((:NotAnInteger, "GreaterThan10":),
       int_aux,0)
<PVSio> readupto10;
Enter a number less than 10
15
==>
0
```

Exceptions



Throw

```
throw[T:TYPE] (tag:ExceptionTag): T
where ExceptionTag : TYPE = string

    int_aux : int =
        let i = query_int("Enter a number less than 10") in
        if i > 10 then throw("GreaterThan10")
        else i endif

Catch
    catch[T:TYPE] (tag:ExceptionTag,program,valueOnException:T): T
    catch[T:TYPE] (tags:list[ExceptionTag],program,valOnExcep:T): T

    readupto10 : int =
        catch[int]((:NotAnInteger, "GreaterThan10":),
        int_aux,0)
```

Local and Global Variables



Locally Scoped Imperative Variables

Imperative-like variables

- Mutable : TYPE+
- ref(value:T) : Mutable
 - defines a local variable storing the value value
- val(var:Mutable): T
 - returns the value stored in the variable var
 - if var stores no value, UndefinedMutableVariable is thrown
- undef(var:Mutable) : bool
 - indicates if the variable var stores any value or not

PVS source: [prelude] stdprog

Local and Global Variables



Local Variables - example

```
woow(x:int) : void =
  let lvar = ref[int](x) in
  println("The value of lvar is: "+val(lvar)) &
  set[int](lvar,x+1) &
  print("The new value of lvar is: "+val(lvar))
```

```
<PVSio> woow(23);
The value of lvar is: 23
The new value of lvar is: 24
```

Local and Global Variables



Local Variables - example

```
woow(x:int) : void =
  let lvar = ref[int](x) in
  println("The value of lvar is: "+val(lvar)) &
  set[int](lvar,x+1) &
  print("The new value of lvar is: "+val(lvar))
```

```
<PVSio> woow(23);
The value of lvar is: 23
The new value of lvar is: 24
```

Local and Global Variables



Local Variables - example

```
woow(x:int) : void =
  let lvar = ref[int](x) in
  println("The value of lvar is: "+val(lvar)) &
  set[int](lvar,x+1) &
  print("The new value of lvar is: "+val(lvar))
```

```
<PVSio> woow(23);
The value of lvar is: 23
The new value of lvar is: 24
```

PVS source: [NASAlib] examples@pvsio_examples

Local and Global Variables Globally Scoped Imperative Variables



Provided as PVS constants of type Global
Global[T:TYPE+,initial_value:T]: TYPE+ = Mutable[T]

Example

```
gvar : Global[int,0]

WOOW(x:int) : void =
  println("The original of gvar is: "+val(gvar)) &
  set(gvar,x) &
  print("The value of gvar is: "+val(gvar))
```

```
<PVSio> WOOW(23);
The original of gvar is: 0
The value of gvar is: 23
```

PVS source: [prelude] stdglobal

PVS Parsing and Typechecking



PVS parsing features are accesible through the function str2pvs str2pvs[T:TYPE+](s:string):T

Example

```
Point : TYPE = [# x : real, y: real #]
zero : Point = str2pvs("(# x := 0, y:= 0 #)")

<PVSio> zero;
==>
(# x := 0, y := 0 #)
```

PVS Parsing and Typechecking



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Example

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Point : TYPE = [# x : real, y: real #]
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<PVSio> zero;
==>
(# x := 0, y := 0 #)
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Example

```
Point : TYPE = [# x : real, y: real #]
zero : Point = str2pvs("(# x := 0, y:= 0 #)")
<PVSio> zero;
==>
(# x := 0, y := 0 #)
```

PVS Parsing and Typechecking Printing arbitrary PVS expressions



pvs2str returns a string representation of a PVS element pvs2str[T:TYPE+](t:T) : string

Example
 <PVSio> print((:1,2,3:));
 first argument to print has the
 Found: (list_adt[real].cons

```
<PVSio> pvs2str((:1,2,3:));
==>
```

PVS Parsing and Typechecking Printing arbitrary PVS expressions



pvs2str returns a string representation of a PVS element pvs2str[T:TYPE+](t:T) : string

```
Example
 <PVSio> print((:1,2,3:));
 first argument to print has the wrong type
      Found: (list adt[real].cons?)
   Expected: booleans.bool
 Try again.
```

PVS Parsing and Typechecking Printing arbitrary PVS expressions



pvs2str returns a string representation of a PVS element pvs2str[T:TYPE+](t:T) : string

```
Example
  <PVSio> print((:1,2,3:));
  first argument to print has the wrong type
      Found: (list_adt[real].cons?)
    Expected: booleans.bool
  Try again.

  <PVSio> pvs2str((:1,2,3:));
  ==>
  "(: 1, 2, 3 :)"
```



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 <PVSio> pvs2str((:1,2,3:));
 ==>
  "(: 1, 2, 3 :)"
```



- PVSio provides a "user-friendly" mechanism for extending the ground evaluator.
- Semantic attachements: Lisp functions attached to uninterpreted PVS functions.

- Every uninterpreted function symbol f_i in a PVS theory Th
- Can be semantically attached to Lisp code
 - using the macro defattach
 - \blacksquare the name must be $|Th.f_i|$
 - as many parameters as the PVS function
 - in a file named "pvs-attachments"
 - located in the context directory
- PVSio executes the attachment code when the symbol is evaluated

```
Th : \texttt{THEORY} BEGIN \dots f_i(p_0 \colon T_0 , \cdots , p_n \colon T_n) : T \\ \dots \\ \texttt{END} \ Th (defattach |Th.f_i| \ (p_0' \cdots p_n') (Documentation string) (Lisp code))
```

Extending PVSio Programming Features User-defined Attachments - Example



cubic_root : THEORY BEGIN cubic(x:real) : real END cubic_root



User-defined Attachments - Example

```
cubic_root : THEORY
BEGIN
 cubic(x:real) : real
END cubic_root
Create the file pvs-attachments in context directory:
;; File: pvs-attachments
(defattach cubic root.cubic (x)
   "Cubic root of x"
   (expt x (/ 1 3))) ;;; <==== THIS IS LISP
```



User-defined Attachments - Example

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BEGIN
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In PVSio
<PVSio> cubic(10);
```



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 cubic(x:real) : real
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   "Cubic root of x"
   (expt x (/ 1 3))) ;;; <==== THIS IS LISP
In PVSio
<PVSio> cubic(10);
==>
2.1544347
```



User-defined Attachments - Limitations

- Name of the attachment and number of parameters
 - Given by the PVS definition of the function
- Data types
 - Parameters and return value
 - Only basic types have an automatic translation to Lisp
 - \blacksquare string \leftrightarrow string (simple-array character)
 - lacksquare nat, int, bool ightarrow immediate fixnum
 - bool \leftarrow bool

PVS theories and attachments do not share namespaces

- PVS global variables can be accessed through ad-hoc macros
 - (pvsio_get_gvar_by_name \langle var name \rangle)
 (pvsio_set_gvar_bv_name \langle var_name \rangle \langle value \rangle)
- For more general cases, PVSio provides macro "using"
 - It allows to refer PVS declarations in attachments
 - Similar in structure to Lisp's let macro

```
(using ((\langle name_0 \rangle \ "\langle pvs \ decl \ name \rangle") \dots (\langle name_n \rangle \ "\langle pvs \ decl \ name \rangle")) (body)
```

lacksquare to use $name_i$ in body, Lisp's function funcall must be used



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Using PVS definitions in attachments

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```
(using ((\langle name_0 \rangle \ "\langle pvs \ decl \ name \rangle") ... (\langle name_n \rangle \ "\langle pvs \ decl \ name \rangle")) \langle body \rangle)
```

lacktriangle to use $name_i$ in body, Lisp's function funcall must be used



Using PVS definitions in attachments - Example

```
PVS theory att_test
  ct0: real = 13
  add_fun(x,y: nat): nat = x + y
  addtoct0(x: nat): nat
```



Using PVS definitions in attachments - Example

```
PVS theory att_test
  ct.0: real = 13
  add_fun(x,y: nat): nat = x + y
  addtoct0(x: nat): nat
In pvs-attachment file:
(defattach | att_test.addtoct0| (x)
  "Example of 'using' macro"
  (using
   ((ct "ct0")
    (add "add_fun"))
   (funcall add (funcall ct) x)))
```



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   ((ct "ct0")
    (add "add_fun"))
   (funcall add (funcall ct) x)))
In PVSio
<PVSio> addtoct0(3);
  ==>
  16
```



Trigonometric operations

- Trigonometric constants and operations are defined in NASAlib/trig
 - pi, sin, cos, tan, atan, asin, acos
 - Example

```
<PVSio> printf("--70/pvs:d/~%",pi_def.pi);
3.141592653589793115997963468544185161590576171875
```

(48 digits)

- By default, they are attached to Lisp's implementations
- NASAlib/fast_approx provides more accurate implementations
 - Example
 - Adding IMPORTING fast_approx@top in the PVS theory

```
<PVSio> printf("--70/pvs:d/-%",pi_def.pi);
3.141592652514881125260318649174312330393139633761036393026404016143926
(70 digits)
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(70 digits)
```



Semantic attachments may produce surprising results

- RANDOM is a constant defined in the prelude (stdmath theory)
- attached to a Lisp implementation of a random number generator

```
<PVSio> RANDOM = RANDOM;
==>
FALSE
```

but...

```
<PVSio> let r=RANDOM in r = r;
==>
TRUE
```

and the following lemma is trivially true

```
thefact : LEMMA
RANDOM = RANDO
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and the following lemma is trivially true

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thefact : LEMMA
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Rapid Prototyping



```
maxl_th : THEORY
BEGIN
 IMPORTING list[real]
 maxl(1:list) : RECURSIVE real =
   cases 1 of
     null:0,
     cons(a,r) : max(a,maxl(r))
   endcases
   MEASURE 1 by <<
END maxl_th
```



```
maxl_io : THEORY
BEGIN
  IMPORTING max1 th
  test : void =
    println("Testing the function maxl") &
    LET s = query_line("Enter a list of real numbers: "),
        1 = str2pvs[list[real]](s),
        m = max(1) IN
      println("The max of "+s+" is "+m)
END maxl io
```



<PVSio> test;

```
Testing the function maxl
Enter a list of real numbers:
(: -1, -2, 5, 3, 2 :)
The max of (: -1, -2, 5, 3, 2 :) is 5

<PVSio> test;
Testing the function maxl
Enter a list of real numbers:
(: -1, -2, -3, -4 :)
The max of (: -1, -2, -3, -4 :) is 0
```



```
<PVSio> test;
Testing the function maxl
Enter a list of real numbers:
```



```
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```

Rapid Prototyping Create PVSio Applications



\$ pvsio maxl_io:test

```
Testing the function maxl
Enter a list of real numbers:
(: 5, 4 ,3 ,2 :)
The max of (: 5, 4 ,3 ,2 :) is 5
```

Rapid Prototyping Create PVSio Applications



```
$ pvsio maxl_io:test
Testing the function maxl
Enter a list of real numbers:
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```

PVSio and the PVS Theorem Prover



- PVSio safely enables the ground evaluator in the theorem prover.
- Ground expressions are translated into Lisp and evaluated in the PVS Lisp engine.
- The theorem prover only trusts the Lisp code automatically generated from PVS functional specifications.
- Semantic attachments are always considered harmful for the theorem prover.



Evaluation of ground expressions via the ground evaluator:

```
{1} 2 < sqrt_newton(2, 10) * sqrt_newton(2, 10)
Rule? (eval-formula 1)
Q.E.D.</pre>
```



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Fast and Furious



Well, as Sound as the PVS Lisp Engine

```
{1} RANDOM /= RANDOM
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Function stdmath.RANDOM is defined as a semantic attachment. It cannot be evaluated in a formal proof.

No change on: (eval-formula 1)

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PVSio and the PVS Theorem Prover



Ground Evaluation With (grind)

```
{1} 2 < sqrt_newton(2, 10) * sqrt_newton(2, 10)

Rule? (grind)
sqrt_newton rewrites sqrt_newton(2, 10)
to (1/2) * (2 / ((1/2) * (2 / (3 * ((1/2) * (1/2))) + (1/2) * (2/(3 * (1/2) + (1/2) * (2/3))) + (1/2) * (1/2) * (2/3)))
+ 3 * ((1/2) * (1/2) * (1/2))))
+
```

PVSio and the PVS Theorem Prover



Ground Evaluation With (grind)

References



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