Task-1

Algorithm of Array-based accumulation approach

void print_array_results (Index *index, int n_results, int n_documents) Check if n_results are not zero // 0(1) Initialise a float array of size n_documents // 0(1) Run a loop from 0 to n_documents-1 // 0(n) Set all the array index to zero // 0(1) // 0(n) Run a loop from 0 to num_of_terms-1 Traverse through each list and look for a document // 0(n) Add Score of the document to the array // 0(1) Initialise a min-heap with size n_results // 0(1) Run a loop from 0 to n_documents-1 // 0(n) If array index is less than n_results Insert score and id into min-heap // 0(logn) Else If score in array is greater than score in heap Remove min id with min score from heap // 0(logn) Insert array score into heap // O(logn) // 0(n) Print the heap Free heap // 0(1) Total time complexity, $T(n) = O(1) + O(n) + O(n^2) + O(1) + O(n) + O(n) + O(n\log n) + O(n\log n)$

+ 0(nlogn) + 0(1) + 0(1)

Hence Time complexity = $O(n^2)$ for worst case scenario.

Task-2

void print_merge_results(Index *index, int n_results)

Algorithm of Priority queue-based multi-way merge approach

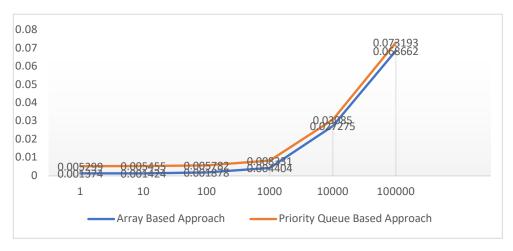
```
Check if n_results are not zero
                                                                            // 0(1)
Initialise min-heaplist of size num_terms
                                                                            // 0(1)
Initialise a node variable
                                                                            // 0(1)
Initialise a new list
                                                                            // 0(1)
Initialise n_size variable to 1
                                                                            // 0(1)
Initialise array to store the pointer pointing to first node of doclists.
                                                                           // 0(1)
Run a loop from 0 to num_terms-1
                                                                            // 0(n)
       Assign all the first nodes of doclist to array address
                                                                            // 0(1)
       Insert into heaplist document id, score, first node, doclist index // O(logn)
Run for loop for the heap size greater than zero
                                                                           // 0(n)
       Peek the node of min id from heap and assign it to node
                                                                           // 0(1)
       Peak the doclist index from which the id belongs to
                                                                           // 0(1)
       Remove score of min id
                                                                           // O(logn)
       Assign pointer of document id to Document type
                                                                           // 0(1)
       Update the n_size of total no of document id
                                                                           // 0(1)
       Add min id and score to new list l
                                                                           // 0(1)
       Update the array of pointer with next node of list
                                                                           // 0(1)
       If the next node of doclist is not NULL
               Assign its data to Document type
                                                                           // 0(1)
               Insert into heaplist document id, score, next node
               doclist index
                                                                           // 0(logn)
```

Initialise float array of size n_size	// 0(1)
Run loop from 0 to n_size-1 and set each index to 0	// 0(n)
Traverse through new list 1 and look at each document	// 0(n)
Add Score of the document to the float array	// 0(1)
Initialise a min-heap with size n_results	// 0(1)
Run a loop from 0 to n_size-1	// 0(n)
If array index is less than n_results	
Insert score and id into min-heap	// O(logn)
Else	
If score in array is greater than score in heap	
Remove min id with min score from heap	// O(logn)
Insert array score into heap	// O(logn)
Print the heap	// 0(n)
Free heap, heaplist, list	// 0(n) + 0(1) + 0(1)

Hence Time complexity = O(nlogn) for worst case scenario.

Based on the asymptotic time complexity that we had calculated, the array-based approach comes out to be $O(n^2)$ and the priority-based approach comes out to be $O(n\log n)$. This means that the array-based approach dominates the priority-based approach when n results $\rightarrow \infty$.

Analysis of computation time of algorithms with increase in n_results



Graph 1: Where X axis is the n_results and Y axis is the time computation. The code was run in macbook pro (16gb ram).

Task 1 is the array-based approach and Task 2 is the priority-based approach. The Graph 1 shows the time computation of each algorithm in different sets of n_results ranging from 1 to 100000. "Hello" and "world" were used as our query terms. Here we can see that for different values of n_results there is an increase time computation for both the algorithms. The Task 2 takes more time than Task 1 however both start to converge when n_results crosses 1000. The time computation for both approaches appears quadratic as seen from the graph.

Conclusion

Based on the asymptotic time complexity analysis, priority queue-based approach should perform better. However, experimental analysis of Task 1 shows less computation time than Task 2 (refer to graph above). Hence, array-based approach is better than priority queue-based approach though the computation time difference is very small. However, as per the graph the computation time will be less for priority-queue based approach when compared to the array-based approach as n_results -> ∞ . Which means when large computations are required, priority queue-based approach will be better.