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MODULE *skeenAlgorithm*

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EXTENDS *TLC*, *Naturals*, *FiniteSets*, *Sequences*

CONSTANTS *NPROCESS*, *MESSAGES*

VARIABLES *pc*, *sent*, *pending*, *received*, *lc*, *messages*

ASSUME (*NPROCESS*  $\in$  *Nat*)  $\wedge$  (*MESSAGES*  $\neq$   $\{\}$ )

ASSUME (*NPROCESS*  $\geq$  2)

vars  $\triangleq$   $\langle pc, sent, pending, received, lc, messages \rangle$

*Processes*  $\triangleq$  1 .. *NPROCESS*

*Init*  $\triangleq$

$\wedge$  *messages* =  $[i \in \text{Processes} \mapsto \text{MESSAGES}]$   
 $\wedge$  *pending* =  $[i \in \text{Processes} \mapsto \{\}]$   
 $\wedge$  *received* =  $[i \in \text{Processes} \mapsto \{\}]$   
 $\wedge$  *sent* =  $[i \in \text{Processes} \mapsto [bc \mapsto \{\}, ts \mapsto \{\}, sn \mapsto \{\}]]$   
 $\wedge$  *pc*  $\in$  [*Processes*  $\rightarrow$  {"BCAST", ""}]  
 $\wedge$  *lc* =  $[i \in \text{Processes} \mapsto 0]$

*Broadcast*(*message*)  $\triangleq$

$[i \in \text{Processes} \mapsto [sent[i] \text{ EXCEPT } !.bc = sent[i].bc \cup \{message\}]]$

*UpponBCAST*(*self*)  $\triangleq$

$\wedge$  (*pc*[*self*] = "BCAST")  $\wedge$  (*messages*[*self*]  $\neq$   $\{\}$ )  
 $\wedge$  LET *currentMessage*  $\triangleq$  CHOOSE  $x \in \text{messages}[self] : \text{TRUE}$   
IN  $\wedge$  *sent'* = *Broadcast*( $\langle self, currentMessage \rangle$ )  
 $\wedge$  *messages'* = [*messages* EXCEPT  $![self] = \text{messages}[self] \setminus \{currentMessage\}$ ]  
 $\wedge$  UNCHANGED  $\langle lc, pending, pc, received \rangle$

$\langle \text{TYPE}, \text{SOURCE}, \text{DESTINATION}, \text{MESSAGE\_BODY}, \text{TIMESTAMP} \rangle$

*ReceivedMessage*(*self*)  $\triangleq$

$\wedge \exists msg \in sent[self].bc :$   
 $\wedge msg \notin pending[self]$   
 $\wedge pending' = [pending \text{ EXCEPT } ![self] = pending[self] \cup \{msg\}]$   
 $\wedge sent' = [sent \text{ EXCEPT } ![msg[1]].ts = sent[msg[1]].ts \cup \{\langle self, msg[2], lc[self] \rangle\}]$   
 $\wedge$  UNCHANGED  $\langle received, pc, messages, lc \rangle$

*MaxTSAllProcess*(*S*)  $\triangleq$  (CHOOSE  $t \in S : \forall s \in S : s[3] \leq t[3]$ )[3]

*SN*(*message*)  $\triangleq$

$[i \in \text{Processes} \mapsto [sent[i] \text{ EXCEPT } !.sn = sent[i].sn \cup \{message\}]]$

*ReceivedTSFromAll*(*self*)  $\triangleq$

$\wedge$  LET *msgs*  $\triangleq$   $\{m1 \in sent[self].ts : \forall m2 \in sent[self].ts : m1[2] = m2[2]\}$   
IN  $\wedge$  *Cardinality*(*msgs*) = *NPROCESS*  
 $\wedge$  LET *m*  $\triangleq$  CHOOSE  $x \in msgs : \text{TRUE}$

$$\text{IN} \quad \wedge \text{sent}' = \text{SN}(\langle \text{self}, m[2], \text{MaxTSAllProcess}(\text{msgs}) \rangle) \\ \wedge \text{UNCHANGED} \langle \text{received}, \text{pending}, \text{lc}, \text{messages}, \text{pc} \rangle$$

$$\text{ReceivedSN}(\text{self}) \triangleq \\ \wedge \exists \text{msg} \in \text{sent}[\text{self}].\text{sn} : \\ \wedge \langle \text{msg}[1], \text{msg}[2] \rangle \notin \text{received}[\text{self}] \\ \wedge \text{received}' = [\text{received} \text{ EXCEPT } ![\text{self}] = \text{received}[\text{self}] \cup \{ \langle \text{msg}[1], \text{msg}[2] \rangle \}] \\ \wedge \text{UNCHANGED} \langle \text{sent}, \text{pending}, \text{lc}, \text{messages}, \text{pc} \rangle$$

$$\text{FIX RECEIVED} \\ \text{Accept}(\text{self}) \triangleq \\ \wedge \text{Cardinality}(\text{Processes} \times \text{MESSAGES}) = \text{Cardinality}(\text{received}[\text{self}]) \\ \wedge \text{pc}' = [\text{pc} \text{ EXCEPT } ![\text{self}] = \text{"AC"}] \\ \wedge \text{UNCHANGED} \langle \text{sent}, \text{pending}, \text{lc}, \text{messages}, \text{received} \rangle$$

$$\text{Steps}(\text{self}) \triangleq \\ \vee \text{UponBCAST}(\text{self}) \\ \vee \text{ReceivedMessage}(\text{self}) \\ \vee \text{ReceivedTSFromAll}(\text{self}) \\ \vee \text{ReceivedSN}(\text{self}) \\ \vee \text{Accept}(\text{self}) \\ \vee \text{UNCHANGED vars}$$

$$\text{Next} \triangleq (\exists \text{self} \in \text{Processes} : \text{Steps}(\text{self}))$$

$$\text{Fairness} \triangleq \text{WF}_{\text{vars}}(\text{Next})$$

$$\text{Spec} \triangleq \text{Init} \wedge \Box[\text{Next}]_{\text{vars}} \wedge \text{Fairness}$$

$$\text{TypeOK} \triangleq \\ \wedge \text{pc} \in [\text{Processes} \rightarrow \{ \text{"BCAST"}, \text{"SN"}, \text{"AC"}, \text{""} \}]$$

Properties

$$\text{Agreement} \triangleq \Box((\exists \text{self} \in \text{Processes} : \text{pc}[\text{self}] = \text{"AC"}) \Rightarrow \Diamond\Box(\forall \text{self} \in \text{Processes} : \text{pc}[\text{self}] = \text{"AC"}))$$

$$\text{SelectedMessage} \triangleq \text{CHOOSE } m \in (\text{Processes} \times \text{MESSAGES}) : \text{TRUE}$$

$$\text{Validity} \triangleq \Box((\exists \text{self} \in \text{Processes} : \text{SelectedMessage} \in \text{sent}[\text{self}].\text{bc}) \Rightarrow \Diamond\Box(\forall p \in \text{Processes} : \text{SelectedMessage} \in \text{received}[p]))$$

$$\text{Integrity} \triangleq \Box((\exists \text{self} \in \text{Processes} : \text{SelectedMessage} \in \text{received}[\text{self}]) \Rightarrow \Diamond\Box(\forall \text{self} \in \text{Processes} : \text{Cardinality}(\text{received}[\text{self}]) \leq \text{Cardinality}(\text{Processes})))$$

$$\text{TotalOrder} \triangleq$$