# ACM ICPC TEAM REFERENCE 2010

# Team Exceptions Universidade Federal do Espírito Santo

#### 1. Configuration files and scripts

#### 1.1. .vimrc.

```
set softtabstop=4
set shiftwidth=4
set tabstop=4
set expandtab
set ruler
set cindent
```

# 1.2. C++ template.

```
#include <cstdio>
#include <string>
#include <vector>
#include <vector>
#include <algorithm>
#include <algorithm>
#include <algorithm>
#include <cmath>
#include <queue>
#include <stack>
#include <set>
#include <deque>
#include <stock>
#include <set>
<algorithm>
#include <stack>
#include <stack>
#include <set>
<algorithm>
#include <stack>
#include <set>
<algorithm>
#include <stack>
<algorithm>
#include <sstream>
#include <sstream>
#include <sstream></a>
#include <sstream>
```

using namespace std;

```
set backspace=indent,eol,start
set showcmd
set nocp
syntax on
hi Normal guibg=black guifg=white
colors zellner
```

```
#define REP(i,n) for(int i=0;i<(int)n;++i)
#define EACH(i,c) for(_typeof((c).begin()) i=(c).begin(); i!=(c).end(); ++i)
#define ALL(c) (c).begin(), (c).end()
#define SIZE(x) (int((x).size()))
const int INF = 0x3F3F3F3F;
const double PI = 2*acos(0);
const double EPS = 1e-10;

inline int cmp(double x, double y=0, double tol=EPS){
    return (x<=y+tol) ? (x+tol<y) ? -1 : 0 : 1;
}
int main(){
    return 0;
}</pre>
```

# 2.1. Replace all occurrences.

```
string replaceAll(string s, string f, string t) {
   string r;
   for (int p = 0; (p = s.find(f)) != s.npos; ) {
      r += s.substr(0, p) + t;
}
```

# 2.2. Replace the first occurrence.

```
string replace(string s, string f, string t) {
   string r;
   int p = s.find(f);
   if (p != s.npos) {
      r += s.substr(0, p) + t;
}
```

# 2.3. Split on all occurrences.

```
vector<string> splitAll(string s, string t) {
  vector<string> v;
  for (int p = 0; (p = s.find(t)) != s.npos; ) {
    v.push_back(s.substr(0, p));
    s = s.substr(p + t.size());
```

# 2.4. Split on first occurrence.

```
vector<string> split(string s, string t) {
  vector<string> v;
  int p = s.find(t);
  if (p != s.npos) {
    v.push_back(s.substr(0, p));
}
```

# 3.1. Palindrome.

```
bool is_palindrome(string const &s){
    return equal(ALL(s), s.rbegin());
```

# 2. Useful functions

```
s = s.substr(p + f.size());
}
return r + s;
}

s = s.substr(p + f.size());
}
return r + s;
}

v.push_back(s);
return v;
}

s = s.substr(p + t.size());
}
v.push_back(s);
return v;
}
```

# 3. String

# 3.2. Longest palindrome.

```
int longest_palindrome(char *text, int n) {
   int rad[2*n], i, j, k;
   for(i=0,j=0;i<2*n;i+=k,j=max(j-k,0)) {
      while(i-j>=0&&i+j+1<2*n&&text[(i-j)/2]==text[(i+j+1)/2]) ++j;
      rad[i] = j;
      for(k=1;i-k>=0&&rad[i]-k>=0&&rad[i-k]!=rad[i]-k;++k)
```

# 3.3. Knuth-Morris-Pratt.

#### 3.4. Suffix Tree.

```
// Larsson-Sadakane's Suffix array Construction: O(n (log n)^2)
struct SAComp {
    const int h, *g;
    SAComp(int h, int* g) : h(h), g(g) {}
    bool operator() (int a, int b) {
        return a == b ? false : g[a] != g[b] ? g[a] < g[b] : g[a+h] < g[b+h];
    }
};
int *buildSA(char* t, int n) {
    int g[n+1], b[n+1], *v = new int[n+1];
    REP(i,n+1) v[i] = i, g[i] = t[i];
    b[0] = 0; b[n] = 0;

    sort(v, v+n+1, SAComp(0, g));
    for(int h = 1; b[n] != n ; h *= 2) {
        SAComp comp(h, g);
        sort(v, v+n+1, comp);
    }
}</pre>
```

```
rad[i+k] = min(rad[i-k], rad[i]-k);
   return *max_element(rad, rad+2*n);
   // ret. centre of the longest palindrome
   int k = -1; //k + 1 eh o tamanho do match atual
   for(int i = 0; i < text.size(); ++i) {</pre>
      while(k \ge 0 \& \& pattern[k + 1] != text[i])
         k = pi[k];
      if(pattern[k + 1] == text[i]) ++k;
      if(k + 1 == pattern.size()) {
         res[nres++] = i - k;
         k = pi[k];
      REP(i, n) b[i+1] = b[i] + comp(v[i], v[i+1]);
      REP(i, n+1) g[v[i]] = b[i];
   return v;
// Naive matching O(m log n)
int find(char *t, int n, char *p, int m, int *sa) {
  int a = 0, b = n;
   while (a < b) {
      int c = (a + b) / 2;
      if (strncmp(t+sa[c], p, m) < 0) a = c+1; else b = c;</pre>
   return strncmp(t+sa[a], p, m) == 0 ? sa[a] : -1;
// Kasai-Lee-Arimura-Arikawa-Park's simple LCP computation: O(n)
```

int \*buildLCP(char \*t, int n, int \*a) {

```
int h = 0, b[n+1], *lcp = new int[n+1];
   REP(i, n+1) b[a[i]] = i;
   REP(i, n+1) {
     if (b[i]) {
         for (int j = a[b[i]-1]; j+h<n && i+h<n && t[j+h] == t[i+h]; ++h);</pre>
         lcp[b[i]] = h;
     } else lcp[b[i]] = -1;
     if (h > 0) --h;
   return lcp;
// call RMO = buildRMO(lcp, n+1)
int *buildRMQ(int *a, int n) {
  int logn = 1;
   for (int k = 1; k < n; k *= 2) ++logn;
   int *r = new int[n * logn];
   int *b = r; copy(a, a+n, b);
   for (int k = 1; k < n; k *= 2) {
     copy (b, b+n, b+n); b += n;
     REP(i, n-k) b[i] = min(b[i], b[i+k]);
   return r:
// inner LCP computation with RMQ: O(1)
int minimum(int x, int y, int *rmq, int n)
  int z = y - x, k = 0, e = 1, s; // y - x >= e = 2^k
  s = ((z \& 0xffff0000) != 0) << 4; z >>= s; e <<= s; k |= s;
   s = ((z \& 0x0000ff00) != 0) << 3; z >>= s; e <<= s; k |= s;
   s = ((z \& 0x000000f0) != 0) << 2; z >>= s; e <<= s; k |= s;
   s = ((z \& 0x0000000c) != 0) << 1; z >>= s; e <<= s; k |= s;
   s = ((z \& 0x00000002) != 0) << 0; z >>= s; e <<= s; k |= s;
   return min( rmq[x+n*k], rmq[y+n*k-e+1] );
```

# 3.5. Regex JAVA.

```
import java.io.BufferedReader;
import java.io.InputStreamReader;
import java.util.regex.Matcher;
import java.util.regex.Pattern;

/*
    [abc] a, b, or c (simple class)
    [^abc] Any character except a, b, or c (negation)
    [a-zA-Z] a through z, or A through Z, inclusive (range)
    [a-d[m-p]] a through d, or m through p: [a-dm-p] (union)
```

```
// outer LCP computation: O(m - o)
int computeLCP(char *t, int n, char *p, int m, int o, int k) {
   int i = 0;
   for (; i < m && k+i < n && p[i] == t[k+i]; ++i);</pre>
   return i;
// Mamber-Myers's O(m + log n) string matching with LCP/RMQ
#define COMP (h, k) (h == m || (k+h < n \&\& p[h] < t[k+h]))
int find(char *t, int n, char *p, int m, int *sa, int *rmq)
   int 1 = 0, 1h = 0, r = n, rh = computeLCP(t, n+1, p, m, 0, sa[n]);
   if (!COMP(rh, sa[r])) return -1;
   for (int k = (1+r)/2; 1+1 < r; k = (1+r)/2) {
      int A = minimum(1+1, k, rmq, n+1), B = minimum(k+1, r, rmq, n+1);
      if (A >= B) {
         if (1h < A) 1 = k;
         else if (1h > A) r = k, rh = A;
         else {
            int i = computeLCP(t, n+1, p, m, A, sa[k]);
            if (COMP(i, sa[k])) r = k, rh = i; else l = k, lh = i;
      } else {
         if (rh < B) r = k;
         else if (rh > B) 1 = k, 1h = B;
         else {
            int i = computeLCP(t, n+1, p, m, B, sa[k]);
            if (COMP(i, sa[k])) r = k, rh = i; else l = k, lh = i;
   return rh == m ? sa[r] : -1;
      [a-z&&[def]] d, e, or f (intersection)
      [a-z&&[^bc]] a through z, except for b and c: [ad-z] (subtraction)
      [a-z\&\&[^nm-p]] a through z, and not m through p: [a-lq-z] (subtraction)
      . Any character (may or may not match line terminators)
      \d A digit: [0-9]
      \D A non-digit: [^0-9]
      \s A whitespace character: [ \t \n \x 0B \f \r]
      \S A non-whitespace character: [^\s]
      \w A word character: [a-zA-Z_0-9]
```

\W A non-word character: [^\w]

```
^ The beginning of a line
$ The end of a line
\b A word boundary
\B A non-word boundary
\A The beginning of the input
\G The end of the previous match
\Z The end of the input but for the final terminator, if any
\z The end of the input

*/

public class Regex {

   public static void main(String[] args) throws Exception {
        BufferedReader stdin;
   stdin = new BufferedReader(new InputStreamReader(System.in));
```

```
System.out.println("Digite_a_regex:_");
String regex = stdin.readLine();
System.out.println("Digite_a_palavra_a_ser_verificada_na_regex:_");
String word = stdin.readLine();
Pattern pattern = Pattern.compile(regex);
Matcher matcher = pattern.matcher(word);
if(matcher.find()) {
    System.out.println("Achou!");
}else{
    System.out.println("Nao_achou!");
}
```

#### 4. Data structures

#### 4.1. Union find.

```
struct UnionFind {
  vector<int> data;
  UnionFind(int size) : data(size, -1) { }
  bool unionSet(int x, int y) {
    x = root(x); y = root(y);
    if (x != y) {
        if (data[y] < data[x]) swap(x, y);
        data[x] += data[y]; data[y] = x;
    }
    return x != y;
}</pre>
```

#### 4.2. Interval Tree.

```
typedef int position;
typedef int contents;
struct interval {
   position low, high;
   interval(position low, position high) :
        low(low), high(high) {
};
struct interval_tree {
```

```
bool findSet(int x, int y) {
    return root(x) == root(y);
}
int root(int x) {
    return data[x] < 0 ? x : data[x] = root(data[x]);
}
int size(int x) {
    return -data[root(x)];
}
};</pre>
```

```
vector<position> pos;
struct node {
   vector<contents> values;
   position B, E, M;
   node *left, *right;
} *root;

template <class IN>
```

```
interval_tree(IN begin, IN end) : pos(begin, end) {
      root = build_tree(0, pos.size()-1);
~interval_tree() { release(root); }
node *build_tree(int i, int j) {
   int m = (i+j)/2;
   node *t = new node;
   t\rightarrow B = pos[i]; t\rightarrow E = pos[j]; t\rightarrow M = pos[m];
   t\rightarrow left = (i+1 < j ? build tree(i, m) : NULL);
   t->right = (i+1 < j ? build_tree(m, j) : NULL);
   return t;
// insert
void insert(const interval& I, contents c) { insert(root, I, c); }
void insert(node *v, const interval& I, contents c) {
   if (I.low <= v->B && v->E <= I.high) {</pre>
      v->values.push_back( c );
   } else {
      if (I.low < v->M) insert(v->left , I, c);
      if (I.high > v->M) insert(v->right, I, c);
```

# 4.3. Range Minimum Query.

```
// Given an array A[0, N-1] find the position of
// the element with the minimum value between two given indices.
// RMQ(2,7)=3
// a[0] | a[1] | a[2] | a[3] | a[4] | a[5] | a[6] | a[7] | a[8] | a[9]
// 2 | 4 | 3 | 1 | 6 | 7 | 8 | 9 | 1 | 7
int *buildRMQ(int *a, int n) {
  int logn = 1;
  for (int k = 1; k < n; k *= 2) ++logn;
  int *r = new int[n * logn];
  int *b = r; copy(a, a+n, b);
  for (int k = 1; k < n; k *= 2) {
    copy(b, b+n, b+n); b += n;
    REP(i, n-k) b[i] = min(b[i], b[i+k]);</pre>
```

# 4.4. Range Maximum Query.

```
int r[50010];
```

```
// query
   template <class OUT>
      void query(position p, OUT out) { query(root, p, out); }
   template <class OUT>
     void query(node *t, position p, OUT out) {
        if (!t || p < t->B || p >= t->E) return;
         copy(t->values.begin(), t->values.end(), out);
        if (p < t->M) query(t->left, p, out);
         else query(t->right, p, out);
   // release
   void release(node *t) {
     if (t->left) release(t->left);
     if (t->right) release(t->right);
      delete t;
};
 return r;
int minimum(int x, int y, int *rmq, int n) {
 int z = y - x, k = 0, e = 1, s; // y - x >= e = 2^k
 s = ((z \& 0xffff0000) != 0) << 4; z >>= s; e <<= s; k |= s;
 s = ((z \& 0x0000ff00) != 0) << 3; z >>= s; e <<= s; k |= s;
 s = ((z \& 0x000000f0) != 0) << 2; z >>= s; e <<= s; k |= s;
 s = ((z \& 0x0000000c) != 0) << 1; z >>= s; e <<= s; k |= s;
 s = ((z \& 0x000000002) != 0) << 0; z >>= s; e <<= s; k |= s;
 return min( rmq[x+n*k], rmq[y+n*k-e+1] );
```

int mm[50010][18]; // or n and log(n) +1

```
void construct() {
                                                                                                 int getmax(int a, int b) {
   int i, j, b;
                                                                                                    if(a>b) return -1;
   for (i=0; i<n; i++) mm[i][0]=r[i];</pre>
                                                                                                    for (int i=17; i>=0; i--)
   for(i=1;i<18;i++)
                                                                                                       if ((a\% (1 << i)) == 0 \& \& (a + (1 << i) - 1) <= b)
      for (j=0; (j+(1<<i)-1)<n; j+=(1<<i))
                                                                                                           return max(mm[a][i], getmax(a+(1<<i), b));</pre>
         mm[j][i]=max(mm[j][i-1], mm[j+(1<<(i-1))][i-1]);
4.5. Trie.
struct Trie {
                                                                                                    for (int i = 0; t[i]; ++i) {
   int value;
                                                                                                       char c = t[i];
   Trie *next[0x100];
                                                                                                       if (!r->next[c]) r->next[c] = new Trie;
   Trie() { fill(next, next+0x100, (Trie*)0); }
                                                                                                        r = r->next[c];
};
                                                                                                    return r;
Trie *find(char *t, Trie *r) {
```

#### 5. Dynamic Programming

#### 5.1. Edit distance.

typedef unsigned int ui;

```
template < class T > ui edit_distance (const T& s1, const T& s2) {
   const size_t len1 = s1.size(), len2 = s2.size();
   vector < vector < ui > > d(len1 + 1, vector < ui > (len2 + 1));

   d[0][0] = 0;
   for (ui i = 1; i <= len1; ++i) d[i][0] = i;

5.2. LCS.

// longest common sequence
vector < int > lcs (const vector < int > & a, const vector < int > & b) {
    const int n = a.size(), m = b.size();
    vector < vector < int > > X(n+1, vector < int > (m+1));
   vector < vector < int > > Y(n+1, vector < int > (m+1));
   for (int i = 0; i < n; ++i) {</pre>
```

X[i+1][j+1] = X[i][j+1];

```
Y[i+1][j+1] = +1;
} else {
    X[i+1][j+1] = X[i+1][j];
    Y[i+1][j+1] = -1;
}
}
vector<int> c;
```

#### 5.3. LIS.

```
// Longest increasing subsequence
vector<int> lis(vector<int>& seq) {
  int smallest_end[seq.size()+1], prev[seq.size()];
  smallest_end[1] = seq[0];

int sz = 1;
  for(int i = 1; i < seq.size(); ++i) {
    int lo = 0, hi = sz;
    while(lo < hi) {
        int mid = (lo + hi + 1)/2;
        if(seq[smallest_end[mid]] <= seq[i])
            lo = mid;
        else
            hi = mid - 1;
    }
}</pre>
```

#### 5.4. Matrix Chain.

```
// matrix chain
int matrix_chain(vector<int>& p, vector< vector<int>>& s) {
    const int n = p.size()-1;
    vector< vector<int>> X(n, vector<int>(n, inf));
    s.resize(n, vector<int>(n));
    for (int i = 0; i < n; ++i) X[i][i] = 0;
    for (int w = 1; w < n; ++w)
        for (int i = 0, j; j = i+w, j < n; ++i)
            for (int k = i, f; k < j; ++k) {</pre>
```

# 5.5. Counting change.

```
#define MAX 100
```

```
for (int i = n, j = m; i > 0 && j > 0; ) {
      if (Y[i][j] > 0) --i;
      else if (Y[i][j] < 0) --j;
      else { c.push_back(a[i-1]); --i; --j; }
   reverse(c.begin(), c.end());
   return c;
      prev[i] = smallest_end[lo];
      if(10 == sz)
         smallest\_end[++sz] = i;
      else if(seq[i] < seq[smallest_end[lo+1]])</pre>
         smallest_end[lo+1] = i;
  vector<int> ret;
  for(int cur = smallest_end[sz]; sz > 0; cur = prev[cur], --sz)
      ret.push_back(seq[cur]);
   reverse(ret.begin(), ret.end());
   return ret;
            int f = p[i]*p[k+1]*p[j+1];
            if (X[i][k] + X[k+1][j] + f < X[i][j]) {</pre>
               X[i][j] = X[i][k] + X[k+1][j] + f;
               s[i][j] = k;
   return X[0][n-1];
int s, n, coins[MAX], sol[MAX];
```

```
int solve() {
    REP(i,s+1) sol[i]=INF;
    sol[0] = 0;
    for(int i=1; i<=s; i++) {
        REP(j,n) {
            if(coins[j]<=i && sol[i-coins[j]]+1 < sol[i]) {
                 sol[i] = sol[i-coins[j]]+1;
            }
        }
    }
}</pre>
```

# 5.6. Knapsack.

```
#define MAXN 2000
#define MAXK 1000

int matrix[MAXN][MAXK+1];
int cost[MAXN];
int profit[MAXN];

int knapsack(int n,int k) {
   int c;
   REP(i,n) memset(matrix[i],0,sizeof(matrix[i]));
   for(int i=cost[0]; i<=k; i++)
        matrix[0][i] = profit[0];</pre>
```

#### 6.1. **GCD**.

```
// Greatest common divisor
// if gcd(a,b) then a and b are coprime
inline int gcd(int a, int b) {
   if(a<0) return gcd(-a,b);</pre>
```

# 6.2. **LCM.**

```
// Least common multiple
// requisits: gcd
inline int lcm(int a, int b) {
   if(a<0) return lcm(-a,b);</pre>
```

```
REP(i,s+1){
      printf("%d_",sol[i]);
  puts("");
   return sol[s];
void read() {
   scanf("%d_%d",&s,&n);
  REP(i,n) scanf("%d",&coins[i]);
  printf("%d",solve());
  for (int i=1;i<n;i++)</pre>
      REP (j, k+1) {
         if(j-cost[i] >= 0)
            c = matrix[i-1][j-cost[i]]+profit[i];
         else
            c = 0;
         if(c < matrix[i-1][j])
            c = matrix[i-1][j];
         matrix[i][j] = c;
   return matrix[n-1][k];
```

# 6. Матн

```
if(b<0) return gcd(a,-b);
return (b==0) ? a : gcd(b,a%b);
}

if(b<0) return lcm(a,-b);
return a*(b/gcd(a,b));
}</pre>
```

#### 6.3. **EXTGCD.**

```
// Extended Euclidean algorithm
// Find x and y that solve:
// a*x + b*y = gcd(a,b)
int extgcd(int a, int b, int &x, int &y) {
6.4. Binomial.
// Binomial - Pascal Triangle
#define DIM 80
unsigned long long C[DIM+1][DIM+1];
void run(){
   unsigned long long n,k;
6.5. Factorize.
// Factorize an integer N
// factorize (60) returns => (2,2), (3,1), (5,1)
template<class T> inline vector<pair<T,int> > factorize(T n) {
   vector<pair<T,int> > R;
   for(T i=2;n>1;) {
      if(n%i==0){
         int C=0;
         for (; n%i==0; C++, n/=i);
6.6. Fractional Library.
// Fractional library
struct frac {
   long long num, den;
   frac() : num(0), den(1) { };
   frac(long long num, long long den) { set_val(num, den); }
   frac(long long num) : num(num), den(1) { };
   void set_val(long long _num, long long _den) {
      num = _num/__gcd(_num, _den);
      den = _den/__gcd(_num, _den);
      if (den < 0) { num \star = -1; den \star = -1; }
```

```
int q = a; x = 1; y = 0;
  if (b!=0) g=extgcd(b,a\%b,y,x), y=(a/b)*x;
   return g;
   for (n=0; n<=DIM; ++n) {</pre>
      C[n][0] = C[n][n] = 1;
      for(k=1; k<n; ++k)
        C[n][k] = C[n-1][k-1] + C[n-1][k];
         R.push_back(make_pair(i,C));
      i++;
      if(i>n/i) i=n;
  if(n>1) R.push_back(make_pair(n,1));
  return R;
   void operator*=(frac f) { set_val(num * f.num, den * f.den); }
   void operator+=(frac f) { set_val(num * f.den + f.num * den, den * f.den); }
   void operator-=(frac f) { set_val(num * f.den - f.num * den, den * f.den); }
   void operator/=(frac f) { set_val(num * f.den, den * f.num); }
bool operator<(frac a, frac b) {</pre>
  if((a.den < 0) ^ (b.den < 0)) return a.num * b.den > b.num * a.den;
  return a.num * b.den < b.num * a.den;</pre>
```

```
std::ostream& operator<<(std::ostream& o, const frac f) {
   o << f.num << "/" << f.den;
   return o;
}
bool operator==(frac a, frac b) { return a.num * b.den == b.num * a.den; }
bool operator!=(frac a, frac b) { return ! (a == b); }</pre>
```

#### 6.7. Pollard rho.

```
// Pollard
// input: integer to be factored
// output: a non-trivial factor of n
long long pollard_r, pollard_n;
inline long long f(long long val) { return (val*val + pollard_r) % pollard_n; }
inline long long myabs(long long a) { return a >= 0 ? a : -a; }
long long pollard(long long n) {
    srand(unsigned(time(0)));
    pollard_n = n;
    long long d = 1;
```

#### 6.8. Miller Rabin.

```
// Prime numbers verify
int fastpow(int base, int d, int n) {
   int ret = 1;
   for(long long pow = base; d > 0; d >>= 1, pow = (pow * pow) % n)
        if(d & 1)
        ret = (ret * pow) % n;
   return ret;
}

bool miller_rabin(int n, int base) {
   if(n <= 1) return false;
   if(n % 2 == 0) return n == 2;

   int s = 0, d = n - 1;
   while(d % 2 == 0) d /= 2, ++s;

   int base_d = fastpow(base, d, n);</pre>
```

```
bool operator<=(frac a, frac b) { return (a == b) || (a < b); }</pre>
bool operator>=(frac a, frac b) { return !(a < b); }</pre>
bool operator>(frac a, frac b) { return ! (a <= b); }</pre>
frac operator/(frac a, frac b) { frac ret = a; ret /= b; return ret; }
frac operator*(frac a, frac b) { frac ret = a; ret *= b; return ret; }
frac operator+(frac a, frac b) { frac ret = a; ret += b; return ret; }
frac operator-(frac a, frac b) { frac ret = a; ret -= b; return ret; }
frac operator-(frac f) { return 0 - f; }
   do {
      d = 1;
      pollard_r = rand() % n;
      long long x = 2, y = 2;
      while(d == 1)
         x = f(x), y = f(f(y)), d = \underline{gcd(myabs(x-y), n)};
   } while(d == n);
   return d;
  if(base d == 1) return true;
   int base_2r = base_d;
   for(int i = 0; i < s; ++i) {</pre>
     if(base_2r == 1) return false;
      if(base_2r == n - 1) return true;
     base 2r = (long long)base 2r * base 2r % n;
   return false;
bool isprime(int n) {
  if(n == 2 || n == 7 || n == 61) return true;
   return miller_rabin(n, 2) && miller_rabin(n, 7) && miller_rabin(n, 61);
```

#### 6.9. **Sieve.**

```
// Sieve primes
const unsigned MAX = 1000000020/60, MAX_S = sqrt(MAX/60);
unsigned w[16] = {1, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41, 43, 47, 49, 53, 59};
unsigned short notprime[MAX];
vector<int> primes;

void sieve() {
   unsigned mod[16][16], di[16][16], num;
   for(int i = 0; i < 16; i++)
        for(int j = 0; j < 16; j++) {
        di[i][j] = (w[i]*w[j])/60;
        mod[i][j] = lower_bound(w, w + 16, (w[i]*w[j])%60) - w;
    }

   primes.push_back(2); primes.push_back(3); primes.push_back(5);</pre>
```

#### 6.10. Carmicheal Lambda function.

```
// Carmichael Lambda
// requisits: lcm
// Given n, returns a^m == 1 (mod n)
// Lamda (p^k) = p^{k-1} (p - 1) (p >= 3 or k <= 2)
// Lamda (2^k) = 2^{k-2} (k > 2)
// Lamda (p1^k1 ... pn^kn) = lcm(Lamda (p1^k1), ..., Lamda (pn^kn))
int carmichaelLambda (int n) {
  int ans = 1;
  if (n%8==0) n/=2;
  for (int d=2; d<=n; ++d) {</pre>
```

#### 6.11. Euler function.

```
// Euler Phi
// eulerPhi(n) returns the number
// of co-primes smaller than n
int eulerPhi(int n) {
   if(n==0) return 0;
   int ans=n;
   for(int x=2; x*x<=n; ++x) {
      if(n%x==0) {</pre>
```

```
memset(notprime, 0, sizeof notprime);
for(unsigned i = 0; i < MAX; i++)
  for(int j = (i==0); j < 16; j++) {
    if(notprime[i] & (1<<j)) continue;
    primes.push_back(num = 60*i + w[j]);

    if(i > MAX_S) continue;
    for(unsigned k = i, done = false; !done; k++)
        for(int 1 = 0; 1 < 16 && !done; l++) {
        unsigned mult = k*num + i*w[1] + di[j][1];
        if(mult >= MAX) done = true;
        else notprime[mult] |= 1<<mod[j][1];
    }
}</pre>
```

```
if (n%d==0) {
    int y=d-1;
    n/=d;
    while (n%d==0) n/=d, y*=d;
    ans = lcm(ans,y);
    }
} return ans;
}
```

```
ans-=ans/x;
while(n%x==0) n/=x;
}
if(n>1) ans-=ans/n;
return ans;
}
```

 $// (m/n) = (m/p_1)^e_1 * ... * (m/p_k)^e_k$ 

# 6.12. Floyd cycle detection.

```
// Floyd cycle detection
pair<int, int> floyd(int x0) {
   int t = f(x0), h = f(f(x0)), start = 0, length = 1;
   while(t != h)
      t = f(t), h = f(f(h));
   h = t; t = x0;
   while(t != h)
6.13. invMod.
// invMod
// requisits: extgcd
// x a == 1 (mod n
int invMod(Int a, Int m) {
6.14. sqrtMod.
// Quadratic residue
// requisits: invMod, powMod and jacobi
// x^2 == a \pmod{p}
int sqrtMod(int n, int p) {
   int S, Q, W, i, m = invMod(n, p);
   for (Q=p-1, S=0; Q%2==0; Q/=2, ++S);
   do{ W=rand()%p; } while(W== 0||jacobi(W,p)!=-1);
6.15. powMod.
// powMod
// x^k (mod n)
int powMod(int x, int k, int m) {
   if(k==0) return 1;
6.16. Jacobi function.
// Jacobi Symbol (m/n),
// For any m integer and n positive odd integer:
```

```
t = f(t), h = f(h), ++start;
  h = f(t);
  while(t != h)
     h = f(h), ++length;
   return make_pair(start, length);
  Int x, y;
  if (extgcd(a, m, x, y) == 1) return (x + m) % m;
  else return 0; // unsolvable
   for (int R = powMod(n, (Q+1)/2, p), V=powMod(W, Q, p);;)
      int z=R*R*m%p;
      for (i=0; i<S&&z%p!=1; z*=z,++i);
      if(i==0) return R;
      R = (R * powMod(V, 1 << (S-i-1), p)) %p;
  if(k%2==0) return powMod(x*x%m,k/2,m);
   else return x*powMod(x,k-1,m)%m;
// where n = p_1^e_1 * ... * p_k^e_k
// (m/n) returns:
// * 0 if a == 0 \pmod{p}
```

```
// * 1 if a != 0 (mod p) and for some integer x, a == x^2 (mod p)
// * -1 if there is no such x
#define NEGPOW(e) ((e) % 2 ? -1 : 1)
int jacobi(int a, int m) {
6.17. Möbius function.
// Mobius
// - mobiusMu(n) = 1 if n is a square-free positive integer
// with an even number of prime factors.
// - mobiusMu(n) = -1 if n is square-free positive integer
// with an odd number of prime factors.
// - mobiusMu(n) = 0 if n is not square-free.
// square-free: integer is one divisible by no perfect square, except 1.
// example: 10 is square-free but 18 is not, as it is divisible by 9 = 3^2
int N = 20000000000;
int mobiusMu(int n) {
   static int lookup = 0, p[N], f[N];
   if(!lookup) {
      REP(i, N) p[i]=1, f[i]=1;
6.18. Romberg.
// Romberg
// Assume F' = f
// input: interval [a,b] and a function f
// ouput: F(b)-F(a)
long double romberg(int a, int b, double(*func)(double)) {
   long double approx[2][50];
   long double *cur=approx[1], *prev=approx[0];
   prev[0] = 1/2.0 * (b-a) * (func(a) + func(b));
   for (int it = 1; it < 25; ++it, swap(cur, prev)) {
      if(it > 1 && cmp(prev[it-1], prev[it-2]) == 0)
         return prev[it-1];
6.19. Polynomials library.
// Polynomial library
typedef complex<double> cdouble;
int cmp(cdouble x, cdouble y = 0) {
   return cmp(abs(x), abs(y));
```

```
if(a==0) return m == 1 ? 1 : 0;
   if(a%2) return NEGPOW((a-1)*(m-1)/4)*jacobi(m%a, a);
   else return NEGPOW((m*m-1)/8)*jacobi(a/2, m);
      for(int i=2;i<N;++i) {</pre>
         if(p[i]){
            f[i] = -1;
            for(int j=i+i; j<N; j+=i) {</pre>
               p[j]=0;
               f[j] *= (j% (i*i) == 0) ? 0 : -1;
      lookup=1;
   return f[n];
      cur[0] = 1/2.0 * prev[0];
      long double div = (b-a)/pow(2, it);
      for(long double sample = a + div; sample < b; sample += 2 * div)</pre>
         cur[0] += div * func(a + sample);
      for(int j = 1; j <= it; ++j)</pre>
         cur[j] = cur[j-1] + 1/(pow(4, it) - 1)*(cur[j-1] + prev[j-1]);
   return prev[24];
const int TAM = 200;
struct poly {
   cdouble poly[TAM]; int n;
```

```
poly(int n = 0): n(n) { memset(p, 0, sizeof(p)); }
cdouble& operator [](int i) { return p[i]; }
poly operator ~() {
  poly r(n-1);
  for (int i = 1; i <= n; i++)</pre>
      r[i-1] = p[i] * cdouble(i);
  return r;
pair<poly, cdouble> ruffini(cdouble z) {
  if (n == 0) return make_pair(poly(), 0);
  poly r(n-1);
  for (int i = n; i > 0; i--) r[i-1] = r[i] * z + p[i];
  return make_pair(r, r[0] * z + p[0]);
cdouble operator ()(cdouble z) { return ruffini(z).second; }
cdouble find_one_root(cdouble x) {
  poly p0 = *this, p1 = ~p0, p2 = ~p1;
  int m = 1000;
  while (m--) {
      cdouble y0 = p0(x);
      if (cmp(y0) == 0) break;
      cdouble G = p1(x) / y0;
```

# 6.20. Simplex.

```
typedef long double T;
typedef vector<T> VT;
typedef vector<int> VI;
vector<VT> A;
VT b,c,res;
VI kt, N;
int m;
inline void pivot(int k, int 1, int e) {
   int x=kt[1]; T p=A[1][e];
   REP(i,k) A[1][i]/=p; b[1]/=p; N[e]=0;
   \label{eq:REP(i,m)} \textbf{ if } (i!=1) \ b[i]-=A[i][e]*b[1], A[i][x]=A[i][e]*-A[1][x];
   REP(j,k) if (N[j]) {
      c[j] -= c[e] *A[1][j];
      REP(i,m) if (i!=1) A[i][j]-=A[i][e]*A[1][j];
   kt[1]=e; N[x]=1; c[x]=c[e]*-A[1][x];
VT doit(int k) {
   VT res: T best:
```

```
cdouble H = G * G - p2(x) - y0;
         cdouble R = sqrt(cdouble(n-1) * (H * cdouble(n) - G * G));
         cdouble D1 = G + R, D2 = G - R;
        cdouble a = cdouble(n) / (cmp(D1, D2) > 0 ? D1 : D2);
        x -= a;
         if (cmp(a) == 0) break;
      return x;
   vector<cdouble> roots() {
      poly q = *this;
      vector<cdouble> r;
      while (q.n > 1) {
        cdouble z(rand() / double(RAND_MAX), rand() / double(RAND_MAX));
        z = q.find_one_root(z); z = find_one_root(z);
        q = q.ruffini(z).first;
        r.push_back(z);
      return r;
};
   while (1) {
      int e=-1, l=-1; REP(i,k) if (N[i] \&\& c[i]>EPS) \{e=i; break;\}
      if (e==-1) break;
      REP(i,m) if (A[i][e]>EPS && (l==-1 || best>b[i]/A[i][e]))
        best=b[ l=i ]/A[i][e];
      if (l==-1) /*ilimitado*/ return VT();
      pivot(k,l,e);
   res.resize(k,0); REP(i,m) res[kt[i]]=b[i];
   return res;
```

m=AA.size(); k=n+m+1; kt.resize(m); b=bb; c=cc; c.resize(n+m);

c=VT(k,0); c[k-1]=-1; pivot(k,pos,k-1); res=doit(k);

 $A=AA; REP(i,m) \{ A[i].resize(k); A[i][n+i]=1; A[i][k-1]=-1; kt[i]=n+i; \}$ 

VT simplex(vector<VT> &AA, VT &bb, VT &cc) {

N=VI(k,1); REP(i,m) N[kt[i]]=0;
int pos=min\_element(ALL(b))-b.begin();

int n=AA[0].size(),k;

**if** (b[pos] <-EPS) {

```
if (res[k-1]>EPS) /*impossivel*/ return VT();
REP(i,m) if (kt[i]==k-1)
    REP(j,k-1) if (N[j] && (A[i][j]<-EPS || EPS<A[i][j])){
      pivot(k,i,j); break;
}</pre>
```

```
c=cc; c.resize(k,0); REP(i,m) REP(j,k) if (N[j]) c[j]-=c[kt[i]]*A[i][j];
}
res=doit(k-1); if (!res.empty()) res.resize(n);
return res;
}
```

#### 7. Graph

#### 7.1. Base element.

7.1.1. Global definitions.

```
typedef int Weight;

struct Edge {
  int src, dst;
  Weight weight;
  Edge(int src, int dst, Weight weight) :
        src(src), dst(dst), weight(weight) { }
};
```

# 7.2. Connected component.

7.2.1. Joint point, connected component decomposition double peak.

```
struct UndirectionalCompare {
  bool operator() (const Edge& e, const Edge& f) const {
      if (min(e.src,e.dst) != min(f.src,f.dst))
         return min(e.src,e.dst) < min(f.src,f.dst);</pre>
      return max(e.src,e.dst) < max(f.src,f.dst);</pre>
};
typedef set<Edge, UndirectionalCompare> Edgeset;
void visit (const Graph &g, int v, int u,
      vector<int>& art, vector<Edgeset>& bcomp,
      stack<Edge>& S, vector<int>& num, vector<int>& low, int& time) {
   low[v] = num[v] = ++time;
   EACH(e, g[v]) {
      int w = e->dst;
      if (num[w] < num[v]) S.push(*e); // for bcomps</pre>
      if (num[w] == 0) {
         visit(g, w, v, art, bcomp, S, num, low, time);
```

```
bool operator < (const Edge &e, const Edge &f) {</pre>
   return e.weight != f.weight ? e.weight > f.weight :
        e.src != f.src ? e.src < f.src : e.dst < f.dst;
typedef vector<Edge> Edges;
typedef vector<Edges> Graph;
typedef vector<Weight> Array;
typedef vector<Array> Matrix;
         low[v] = min(low[v], low[w]);
         if ((num[v] == 1 \&\& num[w] != 2) || // for arts
               (num[v] != 1 \&\& low[w] >= num[v])) art.push_back(v);
         if (low[w] >= num[v]) { // for bcomps}
            bcomp.push_back(Edgeset());
            while (1) {
               Edge f = S.top(); S.pop();
               bcomp.back().insert(f);
               if (f.src == v && f.dst == w) break;
      } else low[v] = min(low[v], num[w]);
void articulationPoint (const Graph& q,
```

vector<int>& art, vector<Edgeset>& bcomp) {

const int n = q.size();

```
vector<int> low(n), num(n);
stack<Edge> S;
REP(u, n) if (num[u] == 0) {
```

7.2.2. Bridge, a double-edge connected component decomposition.

```
void visit (const Graph & q, int v, int u,
     Edges& brdg, vector< vector<int> >& tecomp,
     stack<int>& roots, stack<int>& S, vector<bool>& inS,
     vector<int>& num, int& time) {
   num[v] = ++time;
   S.push(v); inS[v] = true;
   roots.push(v);
   EACH(e, g[v]) {
     int w = e->dst;
     if (num[w] == 0)
        visit(q, w, v, brdq, tecomp, roots, S, inS, num, time);
     else if (u != w && inS[w])
         while (num[roots.top()] > num[w]) roots.pop();
  if (v == roots.top()) {
     brdg.push_back(Edge(u, v));
     tecomp.push_back(vector<int>());
     while (1) {
         int w = S.top(); S.pop(); inS[w] = false;
```

7.2.3. Strongly connected component decomposition.

```
int time = 0;
     visit(q, u, -1, art, bcomp, S, num, low, time);
        tecomp.back().push back(w);
        if (v == w) break;
     roots.pop();
void bridge(const Graph& g, Edges& brdg, vector< vector<int> >& tecomp) {
  const int n = q.size();
  vector<int> num(n);
  vector<bool> inS(n);
  stack<int> roots, S;
   int time = 0;
  REP(u, n) if (num[u] == 0) {
     visit(g, u, n, brdg, tecomp, roots, S, inS, num, time);
     brdg.pop_back();
      scc.push_back(vector<int>());
      while (1) {
        int w = S.top(); S.pop(); inS[w] = false;
        scc.back().push back(w);
        if (v == w) break;
void stronglyConnectedComponents(const Graph& q,
     vector< vector<int> >& scc) {
  const int n = q.size();
  vector<int> num(n), low(n);
   stack<int> S;
```

```
vector<bool> inS(n);
int time = 0;
REP(u, n) if (num[u] == 0)
```

#### 7.3. Shortest path.

7.3.1. Single-source shortest path (Dijkstra).

```
void shortestPath(const Graph &g,int s,vector<Weight> &dist,vector<int> &prev) {
   int n = g.size();
   dist.assign(n, INF); dist[s] = 0;
   prev.assign(n, -1);
   priority_queue<Edge> Q;
   for (Q.push(Edge(-2, s, 0)); !Q.empty(); ) {
      Edge e = Q.front(); Q.pop();
      if (prev[e.dst] != -1) continue;
      prev[e.dst] = e.src;
      EACH(f,g[e.dst]) {
        if (dist[f->dst] > e.weight+f->weight) {
            dist[f->dst] = e.weight+f->weight;
            Q.push(Edge(f->src, f->dst, e.weight+f->weight));
      }
}
```

7.3.2. Single-source shortest paths (Bellman-Ford).

```
bool shortestPath(const Graph g,int s,vector<Weight> &dist,vector<int> &prev){
   int n = g.size();
   dist.assign(n, INF+INF); dist[s] = 0;
   prev.assign(n, -2);
   bool negative_cycle = false;
   REP(k, n) REP(i, n) EACH(e,g[i]) {
     if (dist[e->dst] > dist[e->src] + e->weight) {
        dist[e->dst] = dist[e->src] + e->weight;
        prev[e->dst] = e->src;
        if (k == n-1) {
            dist[e->dst] = -INF;
            negative_cycle = true;
        }
}
```

7.3.3. k-shortest path.

```
Weight k_shortestPath(const Graph &g, int s, int t, int k) {
  const int n = g.size();
```

```
vector<int> buildPath(const vector<int> &prev, int t) {
  vector<int> path;
   for (int u = t; u >= 0; u = prev[u])
     path.push_back(u);
   reverse(path.begin(), path.end());
  return path;
   return !negative_cycle;
vector<int> buildPath(const vector<int> &prev, int t) {
   vector<int> path;
   for (int u = t; u >= 0; u = prev[u])
     path.push_back(u);
   reverse(path.begin(), path.end());
   return path;
```

visit(q, u, scc, S, inS, low, num, time);

Graph h(n); // make reverse graph

```
REP(u, n) EACH(e, q[u])
  h[e->dst].push_back(Edge(e->dst,e->src,e->weight));
vector<Weight> d(n, INF); d[t] = 0; // make potential
vector<int> p(n, -1); // using backward dijkstra
priority_queue<Edge> Q; Q.push(Edge(t, t, 0));
while (!Q.empty()) {
  Edge e = Q.top(); Q.pop();
  if (p[e.dst] >= 0) continue;
  p[e.dst] = e.src;
  EACH(f, h[e.dst]) if (d[f->dst] > e.weight + f->weight) {
     d[f->dst] = e.weight + f->weight;
     Q.push(Edge(f->src, f->dst, e.weight + f->weight));
  }
int 1 = 0; // forward dijkstra-like method
priority_queue<Edge> R; R.push(Edge(-1,s,0));
while (!R.empty()) {
  Edge e = R.top(); R.pop();
  if (e.dst == t && ++1 == k) return e.weight + d[s];
```

#### 7.3.4. Shortest path between all points for (Johnson).

```
bool shortestPath(const Graph &g,
      Matrix &dist, vector<vector<int> > &prev) {
   int n = q.size();
   Array h(n+1);
   REP(k,n) REP(i,n) EACH(e,g[i]) {
     if (h[e->dst] > h[e->src] + e->weight) {
         h[e->dst] = h[e->src] + e->weight;
         if (k == n-1) return false; // negative cycle
   dist.assign(n, Array(n, INF));
   prev.assign(n, vector<int>(n, -2));
   REP(s, n) {
      priority_queue<Edge> Q;
      Q.push(Edge(s, s, 0));
      while (!Q.empty()) {
         Edge e = Q.top(); Q.pop();
         if (prev[s][e.dst] != -2) continue;
         prev[s][e.dst] = e.src;
```

```
EACH(f, g[e.dst])
        R.push(Edge(f->src, f->dst, e.weight+f->weight-d[f->src]+d[f->dst]));
  return -1; // not found
// simpler method
Weight k shortestPath(const Graph &g, int s, int t, int k) {
  const int n = g.size();
  vector<Weight> dist[n];
  priority_queue<Edge> Q; Q.push(Edge(-1, s, 0));
  while (!Q.empty()) {
     Edge e = Q.top(); Q.top();
     if (dist[e.dst].size() >= k) continue;
     dist[e.dst].push_back(e.weight);
     EACH(f, g[e.dst]) Q.push(Edge(f->src, f->dst, f->weight+e.weight));
        EACH(f,q[e.dst]) {
           if (dist[s][f->dst] > e.weight + f->weight) {
              dist[s][f->dst] = e.weight + f->weight;
              Q.push(Edge(f->src, f->dst, e.weight + f->weight));
     REP(u, n) dist[s][u] += h[u] - h[s];
vector<int> buildPath(const vector< vector<int> >& prev, int s, int t) {
  vector<int> path;
   for (int u = t; u >= 0; u = prev[s][u])
     path.push_back(u);
  reverse (ALL (path));
  return path;
```

7.3.5. Shortest path between all points for (Floyd Warshall).

```
void shortestPath(const Matrix &g,Matrix &dist,vector<vector<int> > &inter) {
   int n = g.size();
   dist = g;
   inter.assign(n, vector<int>(n,-1));
   REP(k, n) REP(i, n) REP(j, n) {
     if (dist[i][j] < dist[i][k] + dist[k][j]) {
        dist[i][j] = dist[i][k] + dist[k][j];
        inter[i][j] = k;
     }
   }
}</pre>
```

#### 7.4. Spanning Tree.

7.4.1. Minimum spanning tree (Prim).

```
pair<Weight, Edges> minimumSpanningTree(const Graph &g, int r = 0) {
  int n = g.size();
  Edges T;
  Weight total = 0;

vector<bool> visited(n);
  priority_queue<Edge> Q;
  Q.push( Edge(-1, r, 0) );
  while (!Q.empty()) {
```

7.4.2. Minimum spanning forest (Kruskal).

```
pair<Weight, Edges> minimumSpanningForest(const Graph &g) {
  int n = g.size();
  UnionFind uf(n);
  priority_queue<Edge> Q;
  REP(u, n) EACH(e, g[u]) if (u < e->dst) Q.push(*e);

Weight total = 0;
  Edges F;
  while (F.size() < n-1 && !Q.empty()) {</pre>
```

```
void buildPath(const vector<vector<int> > &inter,int s,int t,vector<int> &path) {
  int u = inter[s][t];
  if (u < 0) path.push_back(s);</pre>
   else buildPath(inter, s, u, path), buildPath(inter, u, s, path);
vector<int> buildPath(const vector< vector<int> > &inter, int s, int t) {
   vector<int> path;
  buildPath(inter, s, t, path);
  path.push_back(t);
   return path;
      Edge e = Q.top(); Q.pop();
      if (visited[e.dst]) continue;
      T.push_back(e);
      total += e.weight;
      visited[e.dst] = true;
      EACH(f, q[e.dst]) if (!visited[f->dst]) Q.push(*f);
   return pair<Weight, Edges>(total, T);
      Edge e = Q.top(); Q.pop();
      if (uf.unionSet(e.src, e.dst)) {
        F.push_back(e);
         total += e.weight;
  return pair<Weight, Edges>(total, F);
```

# 7.4.3. Minimum diameter spanning tree (Cuninghame-Green).

```
Weight minimumDiameterSpanningTree(const Graph &g) {
   int n = g.size();
  Matrix dist(n, Array(n, INF)); // all-pair shortest
  REP(u, n) dist[u][u] = 0;
   REP(u, n) EACH(e, g[u]) dist[e->src][e->dst] = e->weight;
   REP(k, n) REP(i, n) REP(j, n)
      dist[i][j] = min(dist[i][j], dist[i][k]+dist[k][j]);
   Edges E;
   REP(u, n) EACH(e, g[u]) if (e->src < e->dst) E.push_back(*e);
   int m = E.size();
  Weight H = INF;
   vector<Weight> theta(m);
   REP(r, m) {
     int u = E[r].src, v = E[r].dst;
     Weight d = E[r].weight;
     REP(w, n) theta[r] = max(theta[r], min(dist[u][w], dist[v][w]));
```

## 7.4.4. Minimum spanning directed tree (Chu-Liu/Edmonds).

```
void visit (Graph &h, int v, int s, int r,
      vector<int> &no, vector< vector<int> > &comp,
      vector<int> &prev, vector< vector<int> > &next, vector<Weight> &mcost,
      vector<int> &mark, Weight &cost, bool &found) {
   const int n = h.size();
   if (mark[v]) {
      vector<int> temp = no;
      found = true;
      do {
         cost += mcost[v];
         v = prev[v];
         if (v != s) {
            while (comp[v].size() > 0) {
               no[comp[v].back()] = s;
               comp[s].push_back(comp[v].back());
               comp[v].pop_back();
      } while (v != s);
      EACH(\dot{j}, comp[s]) if (\star\dot{j}!= r) EACH(e, h[\star\dot{j}])
         if (no[e->src] != s) e->weight -= mcost[ temp[*j] ];
```

```
H = min(H, d + 2 * theta[r]);
   Weight value = INF;
  REP(r, m) if (2 * theta[r] \le H) {
      int u = E[r].src, v = E[r].dst;
      Weight d = E[r].weight;
      vector< pair<Weight, Weight> > list;
      REP(w, n) list.push_back( make_pair(dist[u][w], dist[v][w]) );
      sort(ALL(list), greater< pair<Weight, Weight> >());
      int p = 0;
      value = min(value, 2 * list[0].first);
      REP(k, n) if (list[p].second < list[k].second)
        value = min(value, d + list[p].second + list[k].first), p = k;
      value = min(value, 2 * list[p].second);
   return n == 1 ? 0 : value;
  mark[v] = true;
  EACH(i, next[v]) if (no[*i] != no[v] && prev[no[*i]] == v)
      if (!mark[no[*i]] || *i == s)
        visit(h, *i, s, r, no, comp, prev, next, mcost, mark, cost, found);
Weight minimumSpanningArborescence(const Graph &g, int r) {
   const int n = g.size();
  Graph h(n);
  REP(u,n) EACH(e,g[u]) h[e->dst].push_back(*e);
  vector<int> no(n);
  vector< vector<int> > comp(n);
  REP(u, n) comp[u].push\_back(no[u] = u);
   for (Weight cost = 0; ;) {
      vector<int> prev(n, -1);
      vector<Weight> mcost(n, INF);
      REP(j, n) if (j!= r) EACH(e, q[j])
         if (no[e->src] != no[j])
```

```
if (e->weight < mcost[ no[j] ])</pre>
              mcost[ no[j] ] = e->weight, prev[ no[j] ] = no[e->src];
     vector< vector<int> > next(n);
     REP(u,n) if (prev[u] \geq= 0)
        next[ prev[u] ].push_back(u);
     bool stop = true;
     vector<int> mark(n);
     REP(u,n) if (u != r && !mark[u] && !comp[u].empty()) {
        bool found = false;
        visit(h, u, u, r, no, comp, prev, next, mcost, mark, cost, found);
        if (found) stop = false;
     if (stop) {
        REP(u,n) if (prev[u] >= 0) cost += mcost[u];
        return cost;
void backward_traverse(int v, int s, int r, matrix &q,
     vector<int> &no, vector< vector<int> > &comp,
     vector<int> &prev, vector<weight> &mcost,
     vector<int> &mark, weight &cost, bool &found) {
   const int n = g.size();
   if (mark[v]) {
     vector<int> temp = no;
     found = true;
     do {
        cost += mcost[v];
        v = prev[v];
        if (v != s) {
           while (comp[v].size() > 0) {
              no[comp[v].back()] = s;
              comp[s].push_back(comp[v].back());
              comp[v].pop_back();
      } while (v != s);
      for (int j = 0; j < n; ++j)
        if (j != r && no[j] == s)
           for (int i = 0; i < n; ++i)</pre>
```

```
if (no[i] != s && q[i][j] < inf)</pre>
                  g[i][j] -= mcost[ temp[j] ];
  mark[v] = true;
   for (int i = 0; i < n; ++i)</pre>
     if (no[i] != no[v] && prev[ no[i] ] == v)
         if (!mark[ no[i] ] || i == s)
            backward_traverse(i, s, r, g,
                  no, comp, prev, mcost, mark, cost, found);
weight minimum spanning arborescence (int r, matrix &g) {
   const int n = g.size();
  vector<int> no(n);
  vector< vector<int> > comp(n);
  for (int i = 0; i < n; ++i) {</pre>
      no[i] = i;
      comp[i].push_back(i);
  weight cost = 0;
   while (1) {
      vector<int> prev(n, -1);
      vector<weight> mcost(n, inf);
      for (int i = 0; i < n; ++i) {</pre>
         for (int j = 0; j < n; ++j) {
            if (j == r) continue;
            if (no[i] != no[j] && g[i][j] < inf) {</pre>
               if (g[i][j] < mcost[ no[j] ]) {</pre>
                  mcost[no[j]] = g[i][j];
                  prev[no[j]] = no[i];
            }
         }
      bool stop = true;
      vector<int> mark(n);
      for (int i = 0; i < n; ++i) {</pre>
         if (i == r || mark[i] || comp[i].size() == 0) continue;
         bool found = false;
         backward_traverse(i, i, r, g,
               no, comp, prev, mcost, mark, cost, found);
         if (found) stop = false;
      if (stop) {
         for (int i = 0; i < n; ++i)
```

```
if (prev[i] >= 0)
    cost += mcost[i];
return cost;
```

7.4.5. Minimum Steiner tree (Dreyfus-Wagner).

```
weight minimum_steiner_tree(const vector<int>& T, const matrix &g) {
 const int n = q.size();
 const int numT = T.size();
 if (numT <= 1) return 0;</pre>
 matrix d(g); // all-pair shortest
 for (int k = 0; k < n; ++k)
  for (int i = 0; i < n; ++i)</pre>
    for (int j = 0; j < n; ++j)
      d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
 weight OPT[(1 << numT)][n];</pre>
 for (int S = 0; S < (1 << numT); ++S)</pre>
  for (int x = 0; x < n; ++x)
    OPT[S][x] = inf;
 for (int p = 0; p < numT; ++p) // trivial case</pre>
  for (int q = 0; q < n; ++q)
   OPT[1 << p][q] = d[T[p]][q];
 for (int S = 1; S < (1 << numT); ++S) { // DP step</pre>
  if (!(S & (S-1))) continue;
  for (int p = 0; p < n; ++p)
    for (int E = 0; E < S; ++E)
     if ((E | S) == S)
       OPT[S][p] = min(OPT[S][p], OPT[E][p] + OPT[S-E][p]);
```

#### 7.5. Flow Cut.

7.5.1. Maximum flow (Edmonds-Karp).

```
#define RESIDUE(s,t) (capacity[s][t]-flow[s][t])
Weight maximumFlow(const Graph &g, int s, int t) {
   int n = g.size();
   Matrix flow(n, Array(n)), capacity(n, Array(n));
   REP(u,n) EACH(e,g[u]) capacity[e->src][e->dst] += e->weight;
```

```
for (int p = 0; p < n; ++p)
    for (int q = 0; q < n; ++q)
      OPT[S][p] = min(OPT[S][p], OPT[S][q] + d[p][q]);
 weight ans = inf;
 for (int S = 0; S < (1 << numT); ++S)</pre>
  for (int q = 0; q < n; ++q)
    ans = min(ans, OPT[S][q] + OPT[((1 << numT)-1)-S][q]);
/*
weight ans = inf;
for (int S = 0; S < (1 << numT); ++S) {
 weight sub = 0;
 for (int P = 0; P < 4; ++P) {
  int E = 0, p = -1;
   for (int k = 0; k < 8; k += 2)
    if (((S >> k) \& 3) == P)
     E += 3 << k, p = T[k];
  sub += (!!E) * OPT[E][p];
 ans = min(ans, sub);
return ans;
*/
```

```
Weight total = 0;
while (1) {
   queue<int> Q; Q.push(s);
   vector<int> prev(n, -1); prev[s] = s;
   while (!Q.empty() && prev[t] < 0) {
     int u = Q.front(); Q.pop();</pre>
```

```
EACH(e,q[u]) if (prev[e->dst] < 0 && RESIDUE(u, e->dst) > 0) {
            prev[e->dst] = u;
            Q.push(e->dst);
      if (prev[t] < 0) return total; // prev[x] == -1 <=> t-side
      Weight inc = INF;
7.5.2. Maximum flow (Dinic).
#define RESIDUE(s,t) (capacity[s][t]-flow[s][t])
Weight augment (const Graph &g, const Matrix &capacity, Matrix &flow,
           const vector<int> &level, vector<bool> &finished,
           int u, int t, Weight cur) {
   if (u == t || cur == 0) return cur;
   if (finished[u]) return 0;
   finished[u] = true;
   EACH(e, g[u]) if (level[e->dst] > level[u]) {
      Weight f = augment(q, capacity, flow, level, finished,
            e->dst, t, min(cur, RESIDUE(u, e->dst)));
      if (f > 0) {
         flow[u][e->dst] += f; flow[e->dst][u] -= f;
         finished[u] = false;
         return f;
     }
   return 0;
Weight maximumFlow(const Graph &g, int s, int t) {
   int n = g.size();
  Matrix flow(n, Array(n)), capacity(n, Array(n)); // adj. matrix
7.5.3. Maximum flow (Goldberg-Tarjan).
#define RESIDUE(s,t) (capacity[s][t]-flow[s][t])
#define GLOBAL_RELABELING() { \
   queue<int> Q; Q.push(t); \
   fill(ALL(d), INF); d[t] = 0; \
   while (!Q.empty()) { \
      int u = Q.front(); Q.pop(); \
```

```
for (int j = t; prev[j] != j; j = prev[j])
        inc = min(inc, RESIDUE(prev[j], j));
      for (int j = t; prev[j] != j; j = prev[j])
         flow[prev[j]][j] += inc, flow[j][prev[j]] -= inc;
     total += inc;
  REP(u,n) EACH(e,g[u]) capacity[e->src][e->dst] += e->weight;
  Weight total = 0;
   for (bool cont = true; cont; ) {
     cont = false;
      vector < int > level(n, -1); level[s] = 0; // make layered network
     queue<int> Q; Q.push(s);
     for (int d = n; !Q.empty() && level[Q.front()] < d; ) {</pre>
         int u = Q.front(); Q.pop();
         if (u == t) d = level[u];
        EACH(e, q[u]) if (RESIDUE(u,e->dst) > 0 && level[e->dst] == -1)
            Q.push(e->dst), level[e->dst] = level[u] + 1;
     vector<bool> finished(n); // make blocking flows
      for (Weight f = 1; f > 0; ) {
        f = augment(g, capacity, flow, level, finished, s, t, INF);
        if (f == 0) break;
        total += f; cont = true;
   return total;
      EACH(e, g[u]) if (RESIDUE(e->dst, u) > 0 && d[u] + 1 < d[e->dst]) \
     Q.push(e->dst), d[e->dst] = d[u] + 1;
  } \
#define PUSH(u, v) { \
  Weight delta = min(excess[u], RESIDUE(u, v)); \
```

```
flow[u][v] += delta; flow[v][w] -= delta; \
    excess[u] -= delta; excess[v] += delta; }

Weight maximumFlow(const Graph &g, int s, int t) {
    int n = g.size(), count = 0;
    Matrix flow(n, Array(n)), capacity(n, Array(n)); // adj. matrix
    REP(u,n) EACH(e,g[u]) capacity[e->src][e->dst] += e->weight;

    vector<Weight> excess(n); excess[s] = INF; // initialize step
    vector<int> d(n);
    GLOBAL_RELABELING();
    vector< queue<int> > B(n); B[ d[s] ].push( s );

for (int b = d[s]; b >= 0; ) {
    if (B[b].empty()) { --b; continue; }
    int v = B[b].front(); B[b].pop();
    if (excess[v] == 0 || v == t) continue;
```

#### 7.5.4. Undirected graph minimum cut across (Nagamochi-Ibaraki/Stoer-Wagner).

```
Weight minimumCut(const Graph &g) {
   int n = g.size();
   vector< vector<Weight> > h(n, vector<Weight>(n)); // make adj. matrix
   REP(u,n) EACH(e,g[u]) h[e->src][e->dst] += e->weight;
   vector<int> V(n); REP(u, n) V[u] = u;

Weight cut = INF;
  for(int m = n; m > 1; m--) {
     vector<Weight> ws(m, 0);
   int u, v;
   Weight w;
   REP(k, m) {
     u = v; v = max_element(ws.begin(), ws.end())-ws.begin();
}
```

# 7.5.5. Thurs Gomory-Hu.

```
#define RESIDUE(s,t) (capacity[s][t]-flow[s][t])
Graph cutTree(const Graph &g) {
   int n = g.size();
   Matrix capacity(n, Array(n)), flow(n, Array(n));
   REP(u,n) EACH(e,g[u]) capacity[e->src][e->dst] += e->weight;
```

```
EACH(e, g[v]) {
    int w = e->dst; // e is the current edge of v
    if (RESIDUE(v,w) > 0 && d[v] == d[w] + 1) { // (w,v) is admissible
        PUSH(v, w);
        if (excess[w] > 0 && w != t) B[d[w]].push( w );
    }
    if (excess[v] == 0) continue;
    d[v] = n;
    EACH(e, g[v]) if (RESIDUE(v, e->dst) > 0)
        d[v] = min(d[v], d[e->dst] + 1);
    if (d[v] < n) B[b = d[v]].push(v);

    if (++count % n == 0) GLOBAL_RELABELING(); // !!HEURISTICS
}
return excess[t];</pre>
```

```
w = ws[v]; ws[v] = -1;
    REP(i, m) if (ws[i] >= 0) ws[i] += h[V[v]][V[i]];
}
REP(i, m) {
    h[V[i]][V[u]] += h[V[i]][V[v]];
    h[V[u]][V[i]] += h[V[v]][V[i]];
}
V.erase(V.begin()+v);
cut = min(cut, w);
}
return cut;
}

vector<int> p(n), prev;
vector<weight> w(n);
for (int s = 1; s < n; ++s) {
    int t = p[s]; // max-flow(s, t)</pre>
```

REP(i,n) REP(j,n) flow[i][j] = 0;

Weight total = 0;

```
while (1) {
         queue<int> Q; Q.push(s);
         prev.assign(n, -1); prev[s] = s;
         while (!Q.empty() && prev[t] < 0) {</pre>
            int u = Q.front(); Q.pop();
            EACH(e,g[u]) if (prev[e->dst] < 0 && RESIDUE(u, e->dst) > 0) {
               prev[e->dst] = u;
               Q.push(e->dst);
         if (prev[t] < 0) goto esc;</pre>
         Weight inc = INF;
         for (int j = t; prev[j] != j; j = prev[j])
            inc = min(inc, RESIDUE(prev[j], j));
         for (int j = t; prev[j] != j; j = prev[j])
            flow[prev[j]][j] += inc, flow[j][prev[j]] -= inc;
         total += inc;
esc:w[s] = total; // make tree
   REP(u, n) if (u != s && prev[u] != -1 && p[u] == t)
```

# 7.5.6. Minimum cost flow (Primal Dual).

```
#define RESIDUE(u,v) (capacity[u][v] - flow[u][v])
#define RCOST(u, v) (cost[u][v] + h[u] - h[v])
pair<Weight, Weight> minimumCostFlow(const Graph &g, int s, int t) {
   const int n = q.size();
  Matrix capacity(n, Array(n)), cost(n, Array(n)), flow(n, Array(n));
  REP (u,n) EACH (e,g[u]) {
     capacity[e->src][e->dst] += e->capacity;
     cost[e->src][e->dst] += e->cost;
   pair<Weight, Weight> total; // (cost, flow)
  vector<Weight> h(n);
   for (Weight F = INF; F > 0; ) { // residual flow
     vector<Weight> d(n, INF); d[s] = 0;
     vector<int> p(n, -1);
     priority_queue<Edge> Q; // "e < f" <=> "e.cost > f.cost"
      for (Q.push(Edge(-2, s)); !Q.empty(); ) {
         Edge e = Q.top(); Q.pop();
         if (p[e.dst] != -1) continue;
         p[e.dst] = e.src;
         EACH(f, q[e.dst]) if (RESIDUE(f->src, f->dst) > 0) {
```

```
p[u] = s;
   if (prev[p[t]] != -1)
     p[s] = p[t], p[t] = s, w[s] = w[t], w[t] = total;
  Graph T(n); // (s, p[s]) is a tree edge of weight w[s]
  REP(s, n) if (s != p[s]) {
     T[s].push_back(Edge(s, p[s], w[s]));
     T[p[s]].push_back( Edge(p[s], s, w[s]) );
   return T;
// Gomory-Hu tree O(n)
Weight maximumFlow(const Graph &T, int u, int t, int p = -1, Weight w = INF) {
  if (u == t) return w;
  Weight d = INF;
  EACH(e, T[u]) if (e->dst != p)
     d = min(d, maximumFlow(T, e->dst, t, u, min(w, e->weight)));
  return d;
            if (d[f->dst] > d[f->src] + RCOST(f->src, f->dst)) {
               d[f->dst] = d[f->src] + RCOST(f->src, f->dst);
              Q.push( Edge(f->src, f->dst, 0, d[f->dst]) );
      if (p[t] == -1) break;
      Weight f = F;
      for (int u = t; u != s; u = p[u])
        f = min(f, RESIDUE(p[u], u));
      for (int u = t; u != s; u = p[u]) {
        total.first += f * cost[p[u]][u];
         flow[p[u]][u] += f; flow[u][p[u]] -= f;
     F -= f;
     total.second += f;
      REP(u, n) h[u] += d[u];
   return total;
```

# 7.6. Matching.

7.6.1. The maximum bipartite graph matching.

7.6.2. Minimum weight maximum bipartite graph matching.

```
#define EVEN(x) (mu[x] == x \mid | (mu[x] != x \&\& phi[mu[x]] != mu[x]))
#define ODD(x) (mu[x] != x && phi[mu[x]] == mu[x] && phi[x] != x)
#define OUTER(x) (mu[x] != x && phi[mu[x]] == mu[x] && phi[x] == x)
int maximumMatching(const Graph &g, Edges &matching) {
   int n = g.size();
  vector<int> mu(n), phi(n), rho(n), scanned(n);
  REP(v, n) mu[v] = phi[v] = rho[v] = v; // (1) initialize
   for (int x = -1; ; ) {
      if (x < 0) { // (2) select even
         for (x = 0; x < n \&\& (scanned[x] || !EVEN(x)); ++x);
         if (x == n) break;
      int y = -1; // (3) select incident
      EACH(e, g[x])
         if (OUTER(e->dst) \mid | (EVEN(e->dst) && rho[e->dst] != rho[x]))
            y = e -> dst;
      if (y == -1) scanned[x] = true, x = -1;
      else if (OUTER(y)) phi[y] = x; // (4) growth
         vector<int> dx(n, -2), dy(n, -2); // (5,6), !TRICK! x % 2 --> x >= 0
         for (int k = 0, w = x; dx[w] < 0; w = k % 2 ? mu[w] : phi[w])
            dx[w] = k++;
         for (int k = 0, w = y; dy[w] < 0; w = k % 2 ? mu[w] : phi[w])
            dy[w] = k++;
```

```
int bipartiteMatching(const Graph& g, int L, Edges& matching) {
   const int n = g.size();
   vector<int> matchTo(n, -1);
   int match = 0;
  REP(u, L) {
      vector<bool> visited(n);
      if (augment(g, u, matchTo, visited)) ++match;
  REP(u, L) if (matchTo[u] >= 0) // make explicit matching
      matching.push_back( Edge(u, matchTo[u]) );
   return match;
         bool vertex_disjoint = true;
         REP(v,n) if (dx[v] >= 0 \&\& dy[v] > 0) vertex_disjoint = false;
         if (vertex_disjoint) { // (5) augment
            REP(v,n) if (dx[v] % 2) mu[phi[v]] = v, mu[v] = phi[v];
            REP(v,n) if (dy[v] % 2) mu[phi[v]] = v, mu[v] = phi[v];
            mu[x] = y; mu[y] = x; x = -1;
            REP(v, n) phi[v] = rho[v] = v, scanned[v] = false;
         } else { // (6) shrink
            int r = x, d = n;
            REP(v,n)
               if (dx[v] >= 0 \&\& dy[v] >= 0 \&\& rho[v] == v \&\& d > dx[v])
                  d = dx[v], r = v;
               if (dx[v] \le d \&\& dx[v] % 2 \&\& rho[phi[v]] != r)
                  phi[phi[v]] = v;
            REP(v,n)
               if (dy[v] <= d && dy[v] % 2 && rho[phi[v]] != r)</pre>
                  phi[phi[v]] = v;
            if (rho[x] != r) phi[x] = y;
            if (rho[y] != r) phi[y] = x;
            REP(v,n) if (dx[rho[v]] \ge 0 \mid | dy[rho[v]] \ge 0) rho[v] = r;
```

#### 7.6.3. Bipartite graph edge coloring.

```
bool augment(vector< vector<int> > &C, int u, int c1, int c2) {
   int v = C[u][c1];
   if (v >= 0) {
      augment(C, v, c2, c1);
      C[v][c2] = u;
      C[u][c1] = -1;
   }
   C[u][c2] = v;
}
int bipartiteColouring(const Graph &g,
      int L, vector< vector<int> > &color) {
   int n = g.size(), deg = 0;
   REP(u, n) deg = max(deg, (int)g[u].size());
   vector< vector<int> > C(n, vector<int> (deg, -1));
```

#### 7.6.4. Maximum Matching.

```
#define EVEN(x) (mu[x] == x || (mu[x] != x && phi[mu[x]] != mu[x]))
#define ODD(x) (mu[x] != x && phi[mu[x]] == mu[x] && phi[x] != x)
#define OUTER(x) (mu[x] != x && phi[mu[x]] == mu[x] && phi[x] == x)
int maximumMatching(const Graph &g, Edges &matching) {
   int n = q.size();
  vector<int> mu(n), phi(n), rho(n), scanned(n);
  REP(v,n) mu[v] = phi[v] = rho[v] = v; // (1) initialize
   for (int x = -1; ; ) {
      if (x < 0) { // (2) select even
         for (x = 0; x < n \&\& (scanned[x] || !EVEN(x)); ++x);
         if (x == n) break;
      int y = -1; // (3) select incident
      EACH(e, g[x])
         if (OUTER(e->dst) \mid \mid (EVEN(e->dst) && rho[e->dst] != rho[x]))
            y = e -> dst;
      if (y == -1) scanned[x] = true, x = -1;
      else if (OUTER(y)) phi[y] = x; // (4) growth
         vector<int> dx(n, -2), dy(n, -2); // (5,6), !TRICK! x \% 2 --> x >= 0
         for (int k = 0, w = x; dx[w] < 0; w = k % 2 ? mu[w] : phi[w])
```

```
return matching.size(); // make explicit matching
REP(u,L) EACH(e,g[u]) {
   int v = e - dst, cu = 0, cv = 0;
   while (C[u][cu] != -1) ++cu;
   while (C[v][cv] != -1) ++cv;
   if (cu != cv) augment(C, v, cu, cv);
   C[u][cu] = v; C[v][cu] = u;
color.assign(n, vector<int>(n)); // make explicit color-table
REP(u, L) REP(c, deg) if (C[u][c] \ge 0)
   color[u][C[u][c]] = c+1;
return deg;
         dx[w] = k++;
      for (int k = 0, w = y; dy[w] < 0; w = k % 2 ? mu[w] : phi[w])
         dy[w] = k++;
      bool vertex_disjoint = true;
      REP(v,n) if (dx[v] >= 0 \&\& dy[v] > 0) vertex_disjoint = false;
      if (vertex_disjoint) { // (5) augment
         REP(v, n) if (dx[v] % 2) mu[phi[v]] = v, mu[v] = phi[v];
         REP(v, n) if (dy[v] % 2) mu[phi[v]] = v, mu[v] = phi[v];
         mu[x] = y; mu[y] = x; x = -1;
         REP(v, n) phi[v] = rho[v] = v, scanned[v] = false;
      } else { // (6) shrink
         int r = x, d = n;
         REP (v, n)
            if (dx[v] >= 0 \&\& dy[v] >= 0 \&\& rho[v] == v \&\& d > dx[v])
               d = dx[v], r = v;
         REP(v,n)
            if (dx[v] \le d \&\& dx[v] % 2 \&\& rho[phi[v]] != r)
               phi[phi[v]] = v;
            if (dy[v] <= d && dy[v] % 2 && rho[phi[v]] != r)</pre>
```

phi[phi[v]] = v;

```
if (rho[x] != r) phi[x] = y;
if (rho[y] != r) phi[y] = x;
REP(v,n) if (dx[rho[v]] >= 0 || dy[rho[v]] >= 0) rho[v] = r;
}
```

# 7.6.5. Minimum weight maximum matching.

```
// adj[i][j] = w(i,j)
Weight adj[N][N];
int size, sorted[N][N], num[N];
void rec(int match[], int p, int 1, Weight w, Weight &opt) {
 int q = p + 1, i;
 if (w >= opt) return;
 for (; q < size; ++q) if (match[q] == -1) break;
 int m = num[q], *list = sorted[q];
 REP(j,m) if (match[i=list[j]] == -1) { // process greedily
  match[q] = i, match[i] = q;
  w += adj[i][q];
  if (1 + 1 < size / 2) rec(match, q, 1+1, w, opt);</pre>
   else opt = min(opt, w);
  w -= adj[i][q];
   match[q] = -1, match[i] = -1;
int gs;
bool compareWith(int i, int j) { return adj[qs][i] < adj[qs][j]; }</pre>
Weight minimumWeightMatching() {
 for (gs = 0; gs < size; ++gs) { // sort adjacent nodes by edge weight
  for (int j = gs+1; j < size; ++j)</pre>
    sorted[gs][num[gs]++] = j;
   sort(sorted[gs], sorted[gs]+num[gs], compareWith);
```

#### 7.7. Tree.

#### 7.7.1. Tree isomorphism.

```
struct Node {
  vector<Node *> child;
};
bool otreeIsomorphism(Node *n, Node *m) {
  if (n->child.size() != m->child.size()) return false;
  REP(i, n->child.size())
```

```
matching.clear();
  REP(u,n) if (u < mu[u]) matching.push_back( Edge(u, mu[u]) );</pre>
  return matching.size(); // make explicit matching
 int match[size]; fill(match, match+size, -1);
 Weight opt = INF;
 rec(match, -1, 0, 0, opt);
 return opt;
Weight best[1<<26];
int minimumWeightMatching() {
 fill(best, best+(1<<size), INF);
 best[0] = 0;
 for (int S = 0; S < (1<<size); ++S) {</pre>
  int i;
  for (i = 0; S & (1 << i); ++i);
  for (int j = i+1; j < size; ++j) {</pre>
   if (S & (1<< j)) continue;
    best[S + (1 << i) + (1 << j)] =
     min(best[S + (1 << i) + (1 << j)], best[S] + adj[i][j]);
 return best[(1<<size)-1];</pre>
  if (!otreeIsomorphism(n->child[i], m->child[i])) return false;
 return true;
```

for (int i = 0; i < n; ++i)

```
struct Node {
 vector<Node *> child:
 vector<int> code;
};
void code(Node *n) {
 int size = 1;
 vector< pair<vector<int>, int> > codes;
 REP(i, n->child.size()) {
  code(n->child[i]);
   codes.push_back( make_pair(n->child[i]->code, i) );
   size += codes[i].first[0];
7.7.2. Off-line least common ancestor (Tarjan).
struct Query {
   int u, v, w;
   Query (int u, int v) : u(u), v(v), w(-1) { }
};
void visit (const Graph &g, int u, int w,
      vector<Query> &qs, vector<int> &color,
      vector<int> &ancestor, UnionFind &uf) {
   ancestor[ uf.root(u) ] = u;
   EACH(e, q[u]) if (e->dst != w) {
      visit(q, e->dst, u, qs, color, ancestor, uf);
      uf.unionSet(e->src, e->dst);
      ancestor[ uf.root(u) ] = u;
7.7.3. Height of the tree.
Weight visit (const Graph &g, Graph& T, int i, int j) {
 if (T[i][j].weight >= 0) return T[i][j].weight;
 T[i][j].weight = g[i][j].weight;
 int u = T[i][j].dst;
 REP(k, T[u].size()) {
  if (T[u][k].dst == i) continue;
  T[i][j].weight = max(T[i][j].weight, visit(g,T,u,k)+g[i][j].weight);
 return T[i][j].weight;
vector<Weight> height(const Graph& g) {
 const int n = q.size();
 Graph T(g); // memoise on tree
```

```
sort(codes.rbegin(), codes.rend()); // !reverse
 n->code.push_back(size);
 for (int i = 0; i < n->child.size(); ++i) {
   swap(n->child[i], n->child[ codes[i].second ]);
  n->code.insert(n->code.end(),
    codes[i].first.begin(), codes[i].first.end());
bool utreeIsomorphism(Node *n, Node *m) {
 code(n); code(m); return n->code == m->code;
  color[u] = 1;
  EACH(q, qs) {
      int w = (q->v == u ? q->u : q->u == u ? q->v : -1);
      if (w \ge 0 \&\& color[w]) q > w = ancestor[uf.root(w)];
void leastCommonAncestor(const Graph &g, int r, vector<Query> &qs) {
  UnionFind uf(q.size());
  vector<int> color(g.size()), ancestor(g.size());
  visit(q, r, -1, qs, color, ancestor, uf);
  for (int j = 0; j < T[i].size(); ++j)</pre>
    T[i][j].weight = -1;
 for (int i = 0; i < n; ++i)</pre>
   for (int j = 0; j < T[i].size(); ++j)</pre>
    if (T[i][j].weight < 0)
     T[i][j].weight = visit(g, T, i, j);
 vector<Weight> ht(n); // gather results
 for (int i = 0; i < n; ++i)
   for (int j = 0; j < T[i].size(); ++j)</pre>
    ht[i] = max(ht[i], T[i][j].weight);
 return ht;
```

#### 7.7.4. Tree diameter.

```
typedef pair<Weight, int> Result;
Result visit(int p, int v, const Graph &g) {
   Result r(0, v);
   EACH(e, g[v]) if (e->dst != p) {
      Result t = visit(v, e->dst, g);
      t.first += e->weight;
   if (r.first < t.first) r = t;
}</pre>
```

#### 7.8. Road Tour.

#### 7.8.1. Directed Euler Path.

```
void visit(Graph& g, int a, vector<int>& path) {
    while (!g[a].empty()) {
        int b = g[a].back().dst;
        g[a].pop_back();
        visit(g, b, path);
    }
    path.push_back(a);
}
bool eulerPath(Graph g, int s, vector<int> &path) {
    int n = g.size(), m = 0;
    vector<int> deg(n);
    REP(u, n) {
        m += g[u].size();
    }
}
```

#### 7.8.2. Undirected Euler Path.

```
void visit(const Graph &g, vector<vector<int> > &adj,int s,vector<int> &path) {
    EACH(e, g[s]) if (adj[e->src][e->dst]) {
        --adj[e->src][e->dst];
        --adj[e->dst][e->src];
        visit(g, adj, e->dst, path);
    }
    path.push_back(s);
}
bool eulerPath(const Graph &g, int s, vector<int> &path) {
    int n = g.size();
    int odd = 0, m = 0;
    REP(i, n) {
        if (g[i].size() % 2 == 1) ++odd;
    }
}
```

```
return r:
Weight diameter (const Graph &g) {
  Result r = visit(-1, 0, g);
  Result t = visit(-1, r.second, g);
   return t.first; // (r.second, t.second) is farthest pair
      EACH(e, q[u]) --deq[e->dst]; // in-deq
      deg[u] += g[u].size(); // out-deg
  int k = n - count(ALL(deg), 0);
  if (k == 0 | | (k == 2 \&\& deg[s] == 1)) {
      path.clear();
      visit(g, s, path);
      reverse (ALL (path));
      return path.size() == m + 1;
   return false:
      m += g[i].size();
  m /= 2;
   if (odd == 0 || (odd == 2 && g[s].size() % 2 == 0)) {
      vector< vector<int> > adj(n, vector<int>(n));
      REP(u, n) EACH(e, g[u]) ++adj[e->src][e->dst];
      path.clear();
      visit(g, adj, s, path);
      reverse (ALL (path));
      return path.size() == m + 1;
   return false:
```

## 7.8.3. Undirected Chinese Postman problem.

```
Weight chinesePostman(const Graph &g) {
  Weight total = 0;
  vector<int> odds;
   REP(u, g.size()) {
      EACH(e, g[u]) total += e->weight;
      if (g[u].size() % 2) odds.push_back(u);
   total /= 2;
   int n = odds.size(), N = 1 << n;</pre>
   Weight w[n][n]; // make odd vertices graph
   REP(u,n) {
      int s = odds[u]; // dijkstra's shortest path
      vector<Weight> dist(g.size(), INF); dist[s] = 0;
      vector<int> prev(q.size(), -2);
      priority_queue<Edge> Q;
      Q.push ( Edge (-1, s, 0) );
      while (!Q.empty()) {
         Edge e = Q.top(); Q.pop();
         if (prev[e.dst] != -2) continue;
         prev[e.dst] = e.src;
```

#### 7.8.4. Shortest Hamiltonian.

```
const int M = 20;
Weight best[1<<M][M];
int prev[1<<M][M];

void buildPath(int S, int i, vector<int> &path) {
   if (!S) return;
   buildPath(S^(1<<i), prev[S][i], path);
   path.push_back(i);
}</pre>
Weight shortestHamiltonPath(Matrix w, int s, vector<int> &path) {
```

#### 7.9. Other.

#### 7.9.1. Linear extension.

```
bool visit(const Graph &g, int v, vector<int> &order, vector<int> &color) {
    color[v] = 1;
```

```
EACH(f,q[e.dst]) {
         if (dist[f->dst] > e.weight+f->weight) {
            dist[f->dst] = e.weight+f->weight;
            Q.push(Edge(f->src, f->dst, e.weight+f->weight));
   REP(v, n) w[u][v] = dist[odds[v]];
Weight best[N]; // DP for general matching
REP(S,N) best[S] = INF;
best[0] = 0;
for (int S = 0; S < N; ++S)
   for (int i = 0; i < n; ++i) if (!(S&(1<<i)))
      for (int j = i+1; j < n; ++j) if (!(S&(1<<j)))
         best[S|(1<<i)|(1<<j)] = min(best[S|(1<<i)|(1<<j)], best[S]+w[i][j]);
return total + best[N-1];
int N = 1 << n;
REP(S,N) REP(i,n) best[S][i] = INF;
best[1 << s][s] = 0;
REP(S,N) REP(i,n) if (S&(1<<i)) REP(j,n)
   if (best[S|(1<<j)][j] > best[S][i] + w[i][j])
      best[S|(1 << j)][j] = best[S][i] + w[i][j],
         prev[S|(1<<j)][j] = i;
int t = min_{element (best[N-1], best[N-1]+n)} - best[N-1];
path.clear(); buildPath(N-1, t, path);
return best[N-1][t];
```

EACH(e, q[v]) {

if (color[e->dst] == 2) continue;

```
if (color[e->dst] == 1) return false;
      if (!visit(q, e->dst, order, color)) return false;
   order.push_back(v); color[v] = 2;
   return true;
bool topologicalSort(const Graph &g, vector<int> &order) {
7.9.2. Graph isomorphism.
typedef vector< vector<int> > Matrix;
typedef pair<int, int> VertexInfo;
#define index second
#define degree first
bool permTest(int k, const Matrix &g, const Matrix &h,
   vector<VertexInfo> &gs, vector<VertexInfo> &hs) {
 const int n = q.size();
 if (k >= n) return true;
 for (int i = k; i < n && hs[i].degree == qs[k].degree; ++i) {</pre>
  swap(hs[i], hs[k]);
  int u = gs[k].index, v = hs[k].index;
   for (int j = 0; j \le k; ++j) {
    if (q[u][qs[j].index] != h[v][hs[j].index]) goto esc;
    if (q[qs[j].index][u] != h[hs[j].index][v]) goto esc;
   if (permTest(k+1, q, h, qs, hs)) return true;
esc:swap(hs[i], hs[k]);
 }
```

#### 8.1. **Base.**

```
struct point{
    double x,y;
    point (double x=0, double y=0): x(x), y(y) {}

    point operator + (point q) { return point(x+q.x,y+q.y); }
    point operator - (point q) { return point(x-q.x,y-q.y); }
    point operator * (double t) { return point(x+t,y*t); }
    point operator / (double t) { return point(x/t,y/t); }

    double operator * (point q) { return x*q.x+y*q.y; }
```

```
int n = q.size();
  vector<int> color(n);
   REP(u, n) if (!color[u] && !visit(q, u, order, color))
      return false;
   reverse (ALL (order));
   return true;
 return false;
bool isomorphism (const Matrix &g, const Matrix &h) {
 const int n = q.size();
 if (h.size() != n) return false;
 vector<VertexInfo> gs(n), hs(n);
 REP(i, n) {
  qs[i].index = hs[i].index = i;
  REP (j, n) {
    gs[i].degree += g[i][j];
    hs[j].degree += h[i][j];
 sort(ALL(gs)); sort(ALL(hs));
 REP(i, n) if (qs[i].degree != hs[i].degree) return false;
 return permTest(0, q, h, qs, hs);
```

#### 8. Geometry

```
double operator % (point q) { return x*q.y-y*q.x; }
int cmp(point q) const {
    if (int t = ::cmp(x,q.x)) return t;
        return ::cmp(y,q.y);
}

bool operator == (point q) const { return cmp(q) == 0; }
bool operator != (point q) const { return cmp(q) != 0; }
```

```
bool operator <(point q) const { return cmp(q) <0; }</pre>
      bool operator <= (point q) const { return cmp(q) <=0; }</pre>
      friend ostream& operator <<(ostream& o, point p) {</pre>
            return o << "(" << p.x << ",.." << p.y << ")";
      static point pivot;
};
point point::pivot;
double abs(point p) { return hypot(p.x,p.y); }
double arg(point p) { return atan2(p.y,p.x); }
8.2. Functions.
// decide se q esta sobre o segmento fechado pr
bool between (point p, point q, point r) {
      return ccw(p,q,r) == 0 && cmp((p-q)*(r-q)) <= 0;
// calcula a distancia do ponto r ao segmento pg
double seq_distance(point p, point q, point r) {
      point A = r-q, B=r-p, C=q-p;
      double a = A*A, b = B*B, c = C*C;
      if (cmp(b,a+c)>=0) return sqrt(a);
      else if(cmp(a,b+c)>=0) return sqrt(b);
      else return fabs(A%B)/sqrt(c);
// classifica o ponto p em relacao ao poligono T
// return: 0 - p esta no exterior
// -1 - p esta na fronteira
// 1 - p esta no interior
int in_poly(point p, polygon &T) {
      double a=0; int N = T.size();
      REP(i,N){
            if (between(T[i],p,T[(i+1)%N])) return -1;
            a += angle(T[i], p, T[(i+1)%N]);
      return cmp(a)!=0;
```

```
typedef vector<point> polygon;
inline int ccw(point p, point q, point r) {
      return cmp((p-r)%(q-r));
inline double angle(point p, point q, point r) {
      point u = p-q, v = r-q;
      return atan2(u%v,u*v);
double dist(point a, point b) {
  point c = a-b;
   return sqrt(c.x*c.x+c.y*c.y);
// comparacao radial
bool radial_lt(point p, point q) {
      point P = p-point::pivot, Q = q-point::pivot;
      double R = P%Q;
      if(cmp(R)) return R>0;
      return cmp (P*P,Q*Q) <0;
// calcula a area orientado do poligono T
// T deve estar orientado positivamente
double poly_area(polygon &T) {
      double s=0; int n=T.size();
      REP(i,n) {
            s += T[i]%T[(i+1)%n];
      return s/2;
// encontra o ponto de intersecao das retas pq e rs
point line_intersect(point p, point q, point r, point s){
      point a = q-p, b = s-r, c = point(p*q,r*s);
      return point (point (a.x,b.x)%c, point (a.y,b.y)%c) / (a%b);
```

#### 8.3. Line intesect.

```
// decide se os segmentos fechados pq e rs tem pontos em comum
bool seg_intersect(point p, point q, point r, point s) {
    point A = q-p, B = s-r, C = r-p, D=s-q;
    int a = cmp(A%C)+2*cmp(A%D);
    int b = cmp(B%C)+2*cmp(B%D);

8.4. Closest points.
```

```
bool compy (point a, point b) {
   return cmp(a.y,b.y) ? cmp(a.y,b.y) < 0 : cmp(a.x,b.x) < 0;
pair<point, point> closest_points_rec(vector<point>& px, vector<point>& py) {
   pair<point, point> ret;
   double d;
   if(px.size() <= 3) {
      double best = 1e10;
      for(int i = 0; i < px.size(); ++i)</pre>
         for(int j = i + 1; j < px.size(); ++j)</pre>
            if(dist(px[i], px[j]) < best) {</pre>
                ret = make_pair(px[i], px[j]);
               best = dist(px[i], px[j]);
      return ret;
   point split = px[(px.size() - 1)/2];
   vector<point> qx, qy, rx, ry;
   for(int i = 0; i < px.size(); ++i)</pre>
      if(px[i] <= split) qx.push_back(px[i]);</pre>
      else rx.push_back(px[i]);
   for(int i = 0; i < py.size(); ++i)</pre>
      if(py[i] <= split) qy.push_back(py[i]);</pre>
      else ry.push_back(py[i]);
   ret = closest_points_rec(qx, qy);
   pair<point, point> rans = closest_points_rec(rx, ry);
   double delta = dist(ret.first, ret.second);
```

```
if (a==3||a==-3||b==3||b==-3) return false;
      int t = (p < r) + (p < s) + (q < r) + (q < s);
      return t!=0 && t!=4;
   if((d = dist(rans.first, rans.second)) < delta) {</pre>
      delta = d;
      ret = rans;
   vector<point> s;
   for(int i = 0; i < py.size(); ++i)</pre>
      if(cmp(abs(py[i].x - split.x), delta) <= 0)</pre>
         s.push_back(py[i]);
   for(int i = 0; i < s.size(); ++i)</pre>
      for(int j = 1; j \le 7 \&\& i + j < s.size(); ++j)
         if((d = dist(s[i], s[i+j])) < delta) {</pre>
            delta = d;
            ret = make_pair(s[i], s[i+j]);
   return ret;
pair<point, point> closest_points(vector<point> pts) {
   if(pts.size() == 1) return make_pair(point(-INF, -INF), point(INF, INF));
   sort(pts.begin(), pts.end());
   for(int i = 0; i + 1 < pts.size(); ++i)</pre>
      if(pts[i] == pts[i+1])
         return make_pair(pts[i], pts[i+1]);
   vector<point> py = pts;
   sort(py.begin(), py.end(), compy);
   return closest_points_rec(pts, py);
```

#### 8.5. Convex Hull.

```
// determina o fecho convexo de um conjunto de pontos no plano
// destroi a lista de pontos T
polygon convex_hull(vector<point> &T) {
      int j=0, k, n=T.size(); polygon U(n);
      point::pivot = *min_element(ALL(T));
      sort(ALL(T), radial_lt);
      for (k=n-2; k>=0 \& \& ccw (T[0], T[n-1], T[k]) ==0; k--);
      reverse((k+1)+ALL(T));
8.6. Spanning Circle.
// encontra o menor circulo que contem todos os pontos dados
typedef pair<point,double> circle;
bool in_circle(circle C, point p) {
      return cmp(abs(p-C.first), C.second) <= 0;</pre>
point circumcenter(point p, point q, point r) {
      point a = p-r, b = q-r, c = point(a*(p+r)/2, b*(q+r)/2);
      return point(c%point(a.y,b.y), point(a.x,b.x)%c)/(a%b);
```

# 8.7. Polygon intersect.

int n = T.size();

circle spanning\_circle(vector<point> &T) {

```
// determina o poligono intersecao dos dois poligonos convexos P e Q
// tanto P quanto Q devem estar orientados positivamente
polygon poly_intersect(polygon &P, polygon &Q) {
   int m = Q.size(), n = P.size();
   int a=0, b=0, aa=0, ba=0, inflag=0;
   polygon R;
   while((aa<n||ba<m) && aa<2*n && ba<2*m) {
      point p1 = P[a], p2 = P[(a+1)*n], q1 = Q[b], q2 = Q[(b+1)*m];
      point A = p2-p1, B=q2-q1;
      int cross = cmp(A*B), ha = ccw(p2,q2,p1), hb = ccw(q2,p2,q1);
      if(cross==0 && ccw(p1,q1,p2)==0 && cmp(A*B)<0) {
         if (between(p1,q1,p2)) R.push_back(q1);</pre>
```

```
REP(i,n){
      // troque o >= por > para manter pontos colineares
      while (j>1&&ccw(U[j-1],U[j-2],T[i])>=0) j--;
      U[j++]=T[i];
U.erase(j+ALL(U));
return U;
random_shuffle(ALL(T));
circle C(point(),-INF);
REP(i,n) if(!in_circle(C,T[i])){
      C = circle(T[i], 0);
      REP(j,i) if(!in_circle(C,T[j])){
            C = circle((T[i]+T[j])/2, abs(T[i]-T[j])/2);
            REP(k,j) if(!in_circle(C,T[k])){
                  point o = circumcenter(T[i],T[j],T[k]);
                  C = circle(o, abs(o-T[k]));
return C;
            if (between (p1, q2, p2)) R.push_back (q2);
            if (between (q1,p1,q2)) R.push_back(p1);
            if (between (q1,p2,q2)) R.push_back (p2);
            if(R.size()<2) return polygon();</pre>
            inflag = 1; break;
      } else if (cross!=0 && seg_intersect(p1,p2,q1,q2)) {
            if(inflag==0) aa = ba = 0;
            R.push_back(line_intersect(p1,p2,q1,q2));
            inflag = (hb>0) ? 1 : -1;
      if(cross==0 && hb<0 && ha<0) return R;
      bool t = cross = 0 && hb == 0 && ha == 0;
```

```
if(t ? (inflag==1) : (cross >=0) ? (ha<=0) : (hb>0)){
    if(inflag==-1) R.push_back(q2);
    ba++; b++; b%=m;
}else{
    if(inflag==1) R.push_back(p2);
    aa++; a++; a%=n;
}
```

```
if(inflag==0){
    if(in_poly(P[0],Q)) return P;
    if(in_poly(Q[0],P)) return Q;
}
R.erase(unique(ALL(R)),R.end());
if(R.size()>1 && R.front() == R.back()) R.pop_back();
return R;
```

#### 9. Useful mathematical facts

9.1. Prime counting function  $(\pi(x))$ . The prime counting function is asymptotic to  $\frac{x}{\log x}$ , by the prime number theorem.

	X	10	$10^{2}$	$10^{3}$	$10^{4}$	$10^{5}$	$10^{6}$	$10^{7}$	$10^{8}$
Ì	$\pi(x)$	4	25	168	1.229	9.592	78.498	664.579	5.761.455

9.2. **Partition function.** The partition function p(x) counts show many ways there are to write the integer x as a sum of integers.

X	36	37	38	39	40	41	42
p(x)	17.977	21.637	26.015	31.185	37.338	44.583	53.174
X	43	44	45	46	47	100	
p(x)	63.261	75.175	89.134	105.558	125.754	190.569.292	

9.3. Catalan numbers. Catalan numbers are defined by the recurrence:

$$C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i}$$

A closed formula for Catalan numbers is:

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1}$$

9.4. Stirling numbers of the first kind. These are the number of permutations of  $I_n$  with exactly k disjoint cycles. They obey the recurrence:

9.5. Stirling numbers of the second kind. These are the number of ways to partition  $I_n$  into exactly k sets. They obey the recurrence:

$${n \brace k} = k {n-1 \brace k} + {n-1 \brace k-1}$$

A "closed" formula for it is:

$${n \brace k} = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^n$$

9.6. **Bell numbers.** These count the number of ways to partition  $I_n$  into subsets. They obey the recurrence:

$$\mathcal{B}_{n+1} = \sum_{k=0}^{n} \binom{n}{k} \mathcal{B}_k$$

X	5	6	7	8	9	10	11	12
$\mathcal{B}_x$	52	203	877	4.140	21.147	115.975	678.570	4.213.597

- 9.7. **Turán's theorem.** No graph with n vertices that is  $K_{r+1}$ -free can have more edges than the Turán graph: A k-partite complete graph with sets of size as equal as possible.
- 9.8. **Generating functions.** A list of generating functions for useful sequences:

$(1,1,1,1,1,1,\ldots)$	$\frac{1}{1-z}$
$(1,-1,1,-1,1,-1,\ldots)$	$\frac{1}{1+z}$
$(1,0,1,0,1,0,\ldots)$	$\frac{1}{1-z^2}$
$(1,0,\ldots,0,1,0,1,0,\ldots,0,1,0,\ldots)$	$\frac{1}{1-z^2}$
$(1, 2, 3, 4, 5, 6, \ldots)$	$\frac{1}{(1-z)^2}$
$(1, \binom{m+1}{m}, \binom{m+2}{m}, \binom{m+3}{m}, \ldots)$	$\frac{1}{(1-z)^{m+1}}$
$(1,c,\binom{c+1}{2},\binom{c+2}{3},\ldots)$	$\frac{1}{(1-z)^c}$
$(1,c,c^2,c^3,\ldots)$	$\frac{1}{1-cz}$
$(0,1,\frac{1}{2},\frac{1}{3},\frac{1}{4},\ldots)$	$\ln \frac{1}{1-z}$

A neat manipulation trick is:

$$\frac{1}{1-z}G(z) = \sum_{n} \sum_{k \le n} g_k z^n$$

9.9. **Polyominoes.** How many free (rotation, reflection), one-sided (rotation) and fixed *n*-ominoes are there?

n	3	4	5	6	7	8	9	10
free	2	5	12	35	108	369	1.285	4.655
one-sided	2	7	18	60	196	704	2.500	9.189
fixed	6	19	63	216	760	2.725	9.910	36.446

9.10. The twelvefold way (from Stanley). How many functions  $f \colon N \to X$  are there?

N	X	Any f	Injective	Surjective
dist.	dist.	$x^n$	$(x)_n$	$x!\binom{n}{x}$
indist.	dist.	$\binom{x+n-1}{n}$	$\binom{x}{n}$	$\binom{n-1}{n-x}$
dist.	indist.	$\binom{n}{1} + \ldots + \binom{n}{x}$	$[n \leq x]$	$\binom{n}{k}$
indist.	indist.	$p_1(n) + \dots p_x(n)$	$[n \leq x]$	$p_x(n)$

Where  $\binom{a}{b} = \frac{1}{b!}(a)_b$  and  $p_x(n)$  is the number of ways to partition the integer n using x summands.

9.11. Common integral substitutions. And finally, a list of common substitutions:

$\int F(\sqrt{ax+b})dx$	$u = \sqrt{ax + b}$	$\frac{2}{a}\int uF(u)du$
$\int F(\sqrt{a^2-x^2})dx$	$x = a\sin u$	$a \int F(a\cos u)\cos u du$
$\int F(\sqrt{x^2+a^2})dx$	$x = a \tan u$	$a \int F(a \sec u) \sec^2 u du$
$\int F(\sqrt{x^2-a^2})dx$	$x = a \sec u$	$a \int F(a \tan u) \sec u \tan u du$
$\int F(e^{ax})dx$	$u = e^{ax}$	$\frac{1}{a} \int \frac{F(u)}{u} du$
$\int F(\ln x)dx$	$u = \ln x$	$\int F(u)e^udu$

9.12. Table of non-trigonometric integrals. Some useful integrals are:

$\frac{1}{a} \arctan \frac{x}{a}$
$\frac{1}{2a} \ln \frac{x-a}{x+a}$
$\frac{1}{2a} \ln \frac{a+x}{a-x}$
$\arcsin \frac{x}{a}$
$\ln\left(u+\sqrt{x^2-a^2}\right)$
$\frac{1}{a}\operatorname{arcsec}\left \frac{u}{a}\right $
$-\frac{1}{a}\ln\left(\frac{a+\sqrt{x^2+a^2}}{x}\right)$
$-\frac{1}{a}\ln\left(\frac{a+\sqrt{a^2-x^2}}{x}\right)$

9.13. **Table of trigonometric integrals.** A list of common and not-so-common trigonometric integrals:

$\int \tan x dx$	$-\ln \cos x $
$\int \cot x dx$	$\ln \sin x $
$\int \sec x dx$	$\ln \sec x + \tan x $
$\int \csc x dx$	$\ln \csc x - \cot x $
$\int \sec^2 x dx$	$\tan x$
$\int \csc^2 x dx$	$\cot x$
$\int \sin^n x dx$	$\frac{-\sin^{n-1}x\cos x}{n} + \frac{n-1}{n}\int \sin^{n-2}x dx$
$\int \cos^n x dx$	$\frac{\cos^{n-1}x\sin x}{n} + \frac{n-1}{n}\int \cos^{n-2}x dx$
$\int \arcsin x dx$	$x \arcsin x + \sqrt{1 - x^2}$
$\int \arccos x dx$	$x \arccos x - \sqrt{1 - x^2}$
$\int \arctan x dx$	$x\arctan x - \frac{1}{2}\ln 1 - x^2 $

#### 10. MISTAKES TO AVOID AND STRATEGY HINTS

# General strategy hints:

- (1) Everyone should have read every problem by the end of the second hour of competition.
- (2) In the last hour, only one question should be attempted at once.
- (3) Always use vectors instead of arrays in the single dimensional case. Use vectors instead of arrays in the bidimensional case if the outermost dimension is not too big. This allows for better debugging.

When thinking (before coding):

- (1) Be organized!
- (2) Don't touch the computer unless the solution is done, including implementation details: avoid thinking too much at the computer. Time spent on paper detailing a solution is time well spent.
- (3) A corollary to the above: **Don't touch the computer if you doubt your idea**.

In case you have no solution ideas:

(1) Talk with member team about one idea for solve. Case he already had read the problem.

In case of Wrong Answer:

(1) Make sure the algorithm is correct (i.e. **sketch a proof of correctness**) as soon as possible. Take your time and check

- it carefully. If you're sure the idea is correct, make sure you don't second-guess it when checking for bugs, even if an implementation bug is nowhere to be found.
- (2) If the code uses vectors (c.f. general strategy hints), add #define \_GLIBCXX\_DEBUG at the very beginning of the source code and submit it again. A runtime error means out-of-bounds array access or other STL misuse.
- (3) Debug on paper, and don't go back to the computer for every bug you find: check the whole solution at least once more after finding each new bug.
- (4) Think of tricky test cases.
- (5) Read the problem again. For every constraint found, check the printed source code.
- (6) If you can't find any bug in five minutes, go to the bathroom. If you still can't find the bug, go to another problem and come back to the wrong solution later. Debugging a single program for too long leads to finding a lot of false bugs and makes it easy to overlook simple mistakes.

#### In case of Runtime Error:

- (1) Check divisions and modulo operations.
- (2) Check array indices (both in declarations and in accesses).
- (3) Check for infinite recursion.

# ACM ICPC TEAM REFERENCE - CONTENTS

# Team Exceptions Universidade Federal do Espírito Santo

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