Scaling Abstraction Refinement via Pruning

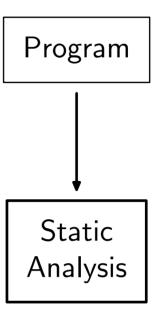
PLDI - San Jose, CA

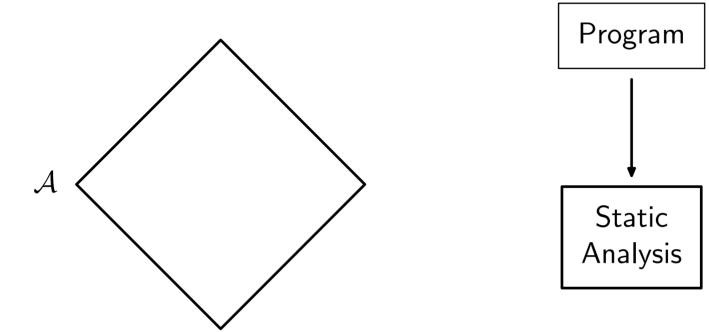
June 8, 2011

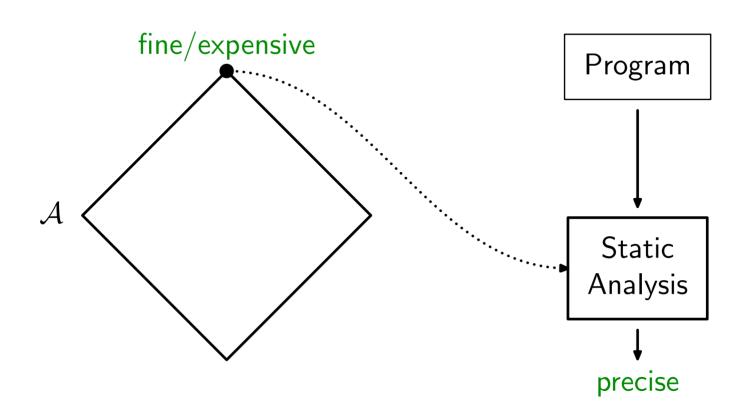
Percy Liang
UC Berkeley

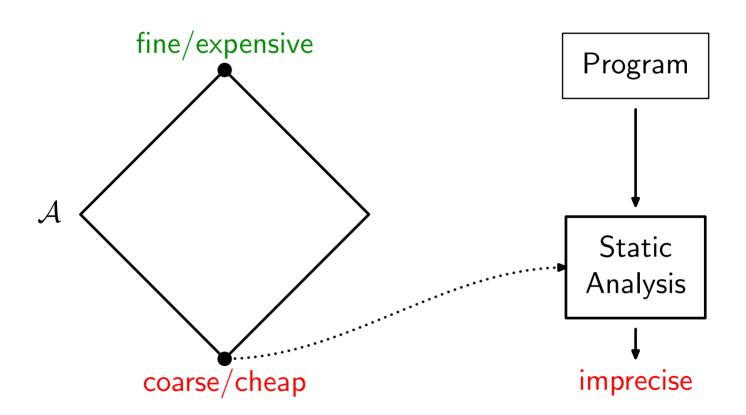
Mayur Naik Intel Labs Berkeley

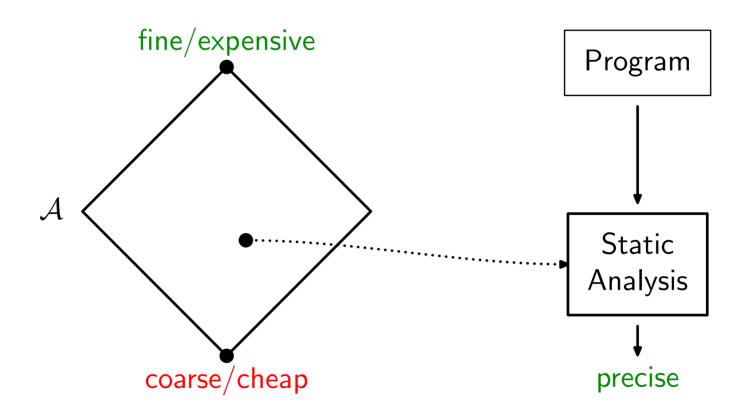
Program





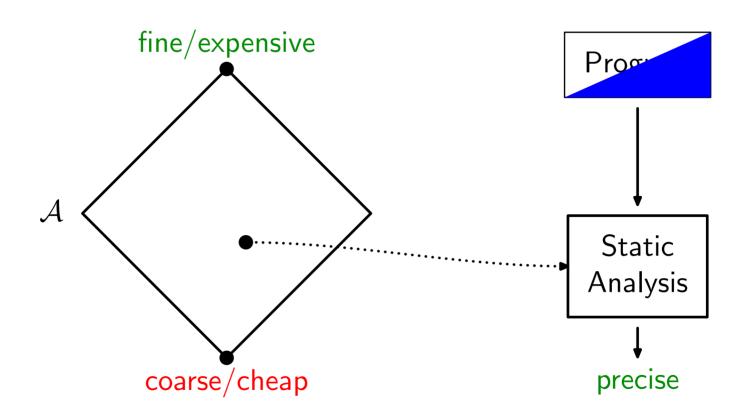






selected refinement

[Heintze & Tardieu 2001] [Guyer & Lin 2003] [Sridharan et al. 2005] [Zheng & Rugina 2008] [Liang et al. 2011]



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[Heintze & Tardieu 2001] [Guyer & Lin 2003] [Sridharan et al. 2005] [Zheng & Rugina 2008] [Liang et al. 2011]



Query: do x and y alias? (no)

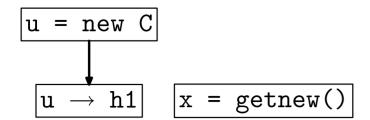
Query: do x and y alias? (no)

$$u = new C$$

$$v = new C$$

Query: do x and y alias? (no)

0-CFA:

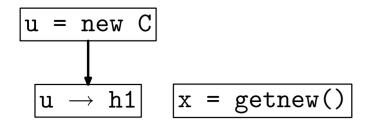


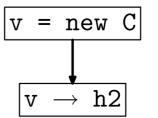
$$v = new C$$

y = getnew()

Query: do x and y alias? (no)

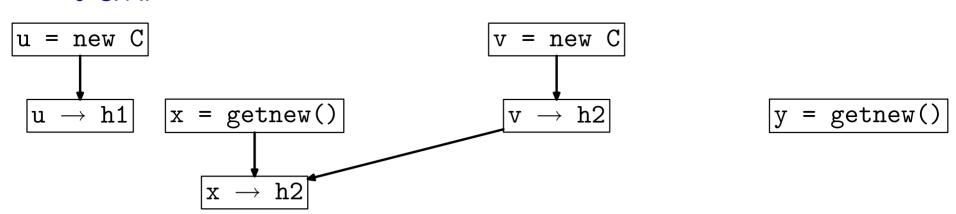
0-CFA:



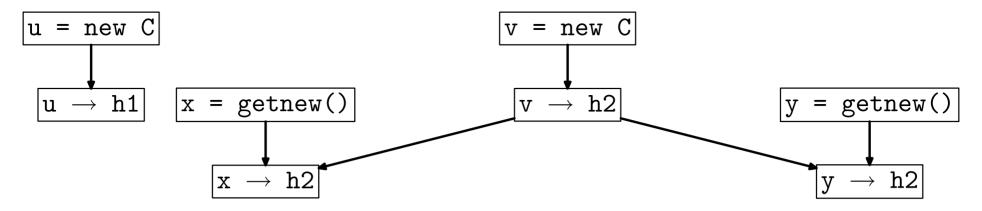


y = getnew()

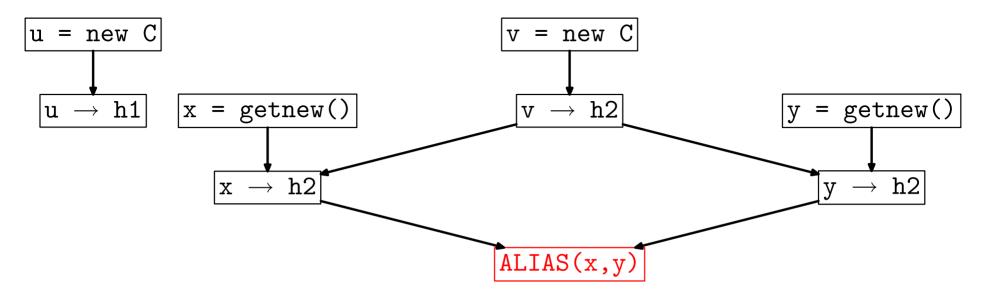
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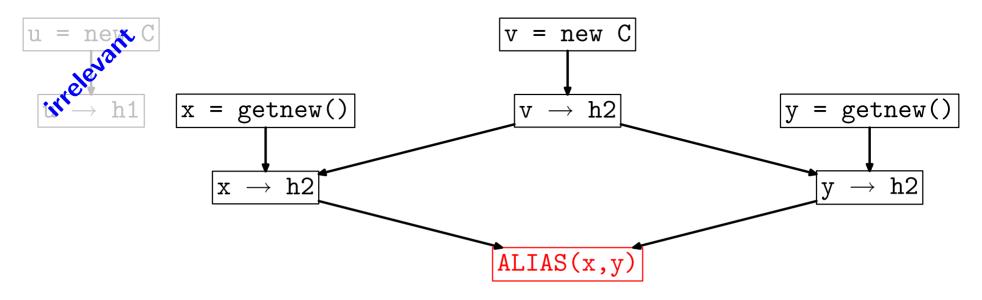


Query: do x and y alias? (no)



```
getnew() {
h1:     u = new C
h2:     v = new C
     return v
}
i1: x = getnew()
i2: y = getnew()
```

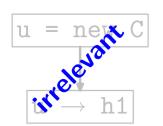
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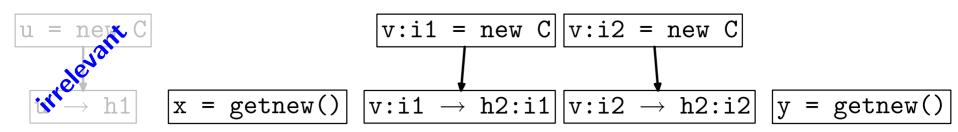
1-CFA on pruned:



$$x = getnew()$$

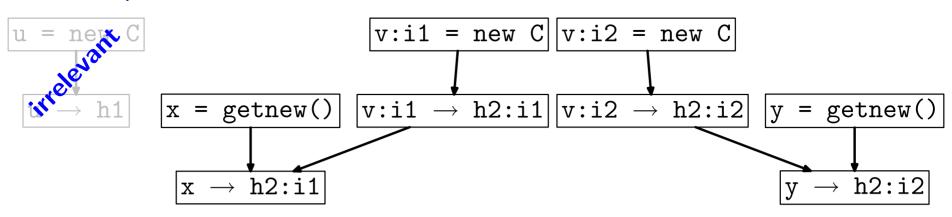
Query: do x and y alias? (no)

1-CFA on pruned:



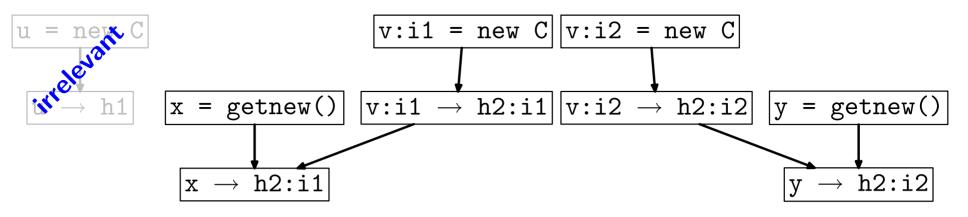
Query: do x and y alias? (no)

1-CFA on pruned:



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1-CFA on pruned:



(not aliasing - query proven)

Input tuples A_0

$$v = new C \cdots$$

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Query tuple x_{Q}

Input tuples A_0

$$v = new C \cdots$$

Query tuple $x_{\rm Q}$

Datalog rules

$$v_2 \rightarrow h \Leftarrow v_2 = v_1$$
 , $v_1 \rightarrow h$

. . .

Input tuples A_0

$$v = new C \cdots$$

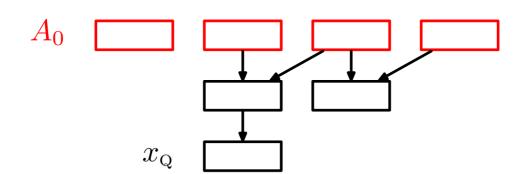
Query tuple x_0

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Prune/prove operator **P**

$$A_0 \xrightarrow{\mathbf{P}}$$
 subset of A_0 used to derive $x_{\mathbf{Q}}$



Input tuples A_0

$$v = new C \cdots$$

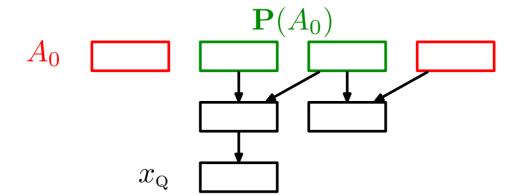
Query tuple x_0

Datalog rules

$$\begin{bmatrix} v_2 & o & h \end{bmatrix} \Leftarrow \begin{bmatrix} v_2 & = & v_1 \end{bmatrix}$$
 , $\begin{bmatrix} v_1 & o & h \end{bmatrix}$

Prune/prove operator **P**

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Input tuples A_0

$$v = new C \cdots$$

Query tuple $x_{\rm O}$

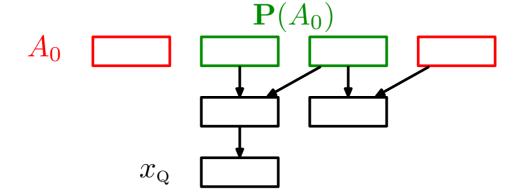
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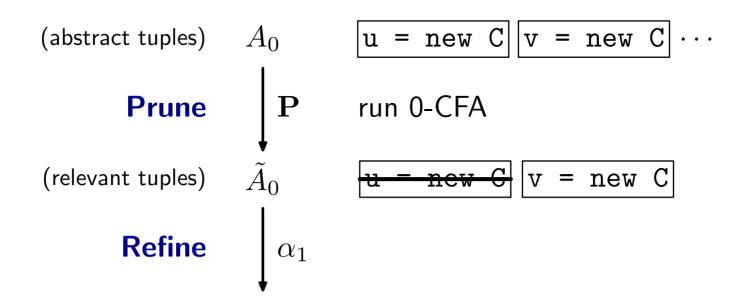
$$A_0 \xrightarrow{\mathbf{P}}$$
 subset of A_0 used to derive $x_{\mathbf{Q}}$

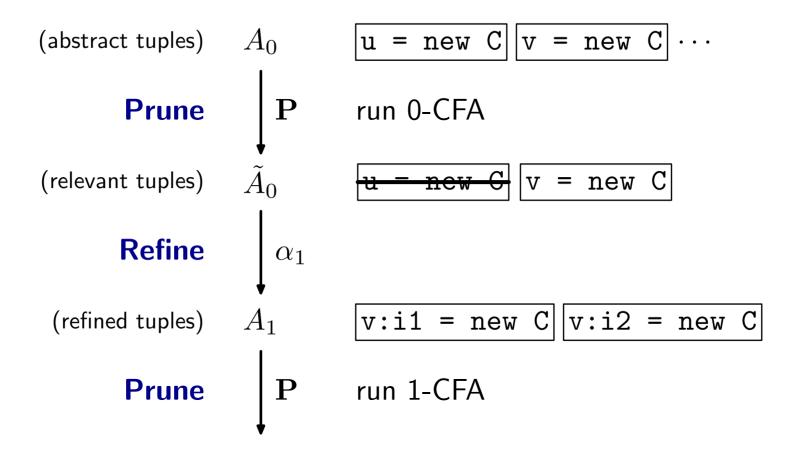


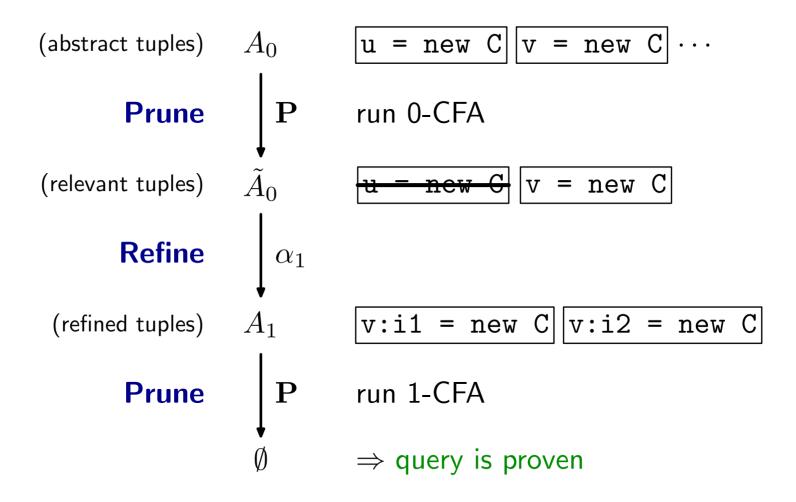
Query proven
$$\Leftrightarrow$$
 $\mathbf{P}(A_0) = \emptyset$

(abstract tuples) A_0 $[u = new C][v = new C] \cdots$

(abstract tuples)
$$A_0$$
 $\boxed{\mathbf{u} = \operatorname{new} \mathbf{C}} \boxed{\mathbf{v} = \operatorname{new} \mathbf{C}} \cdots$ Prune $\boxed{\mathbf{P}}$ run 0-CFA







Prune-Refine Algorithm

Input:

Sequence of abstractions: $\alpha_0 \leq \alpha_1 \leq \alpha_2 \leq \cdots$

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For $t = 0, 1, 2, \dots$:

Prune: $\tilde{A}_t = \mathbf{P}(A_t)$. If $\tilde{A}_t = \emptyset$: return proven.

Refine: $A_{t+1} = \alpha_{t+1}(\tilde{A}_t)$.

Input:

Sequence of abstractions: $\alpha_0 \leq \alpha_1 \leq \alpha_2 \leq \cdots$ A_0 , initial set of tuples

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Main Result:

Prune-Refine after t iterations

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Full Analysis on α_t

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Full Analysis on α_t

fast

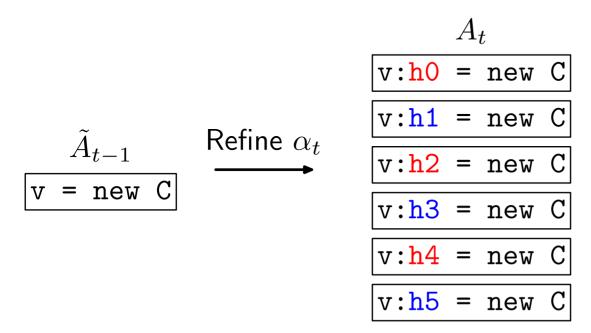
Rest of Talk

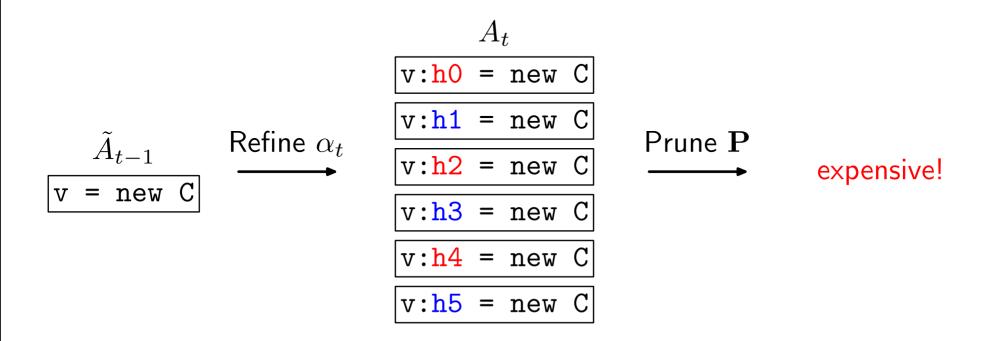
Pre-Pruning Extension

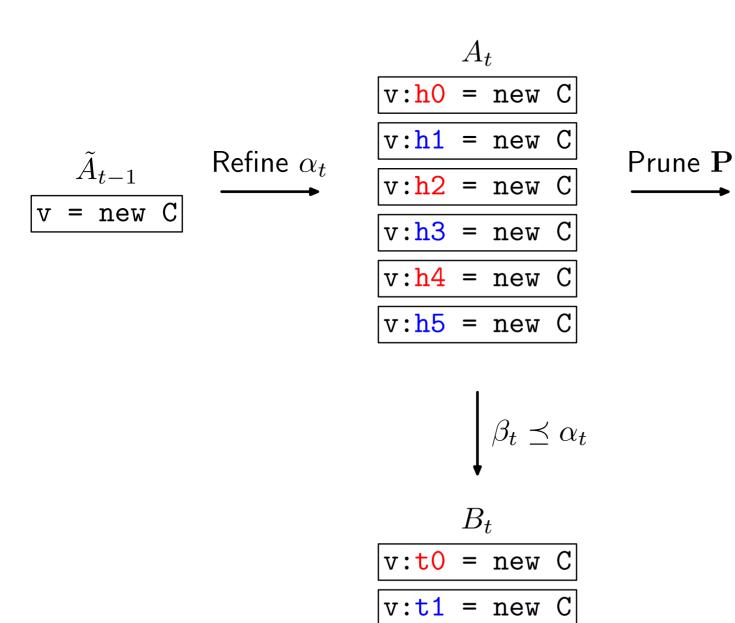
Experiments

$$\widetilde{A}_{t-1}$$

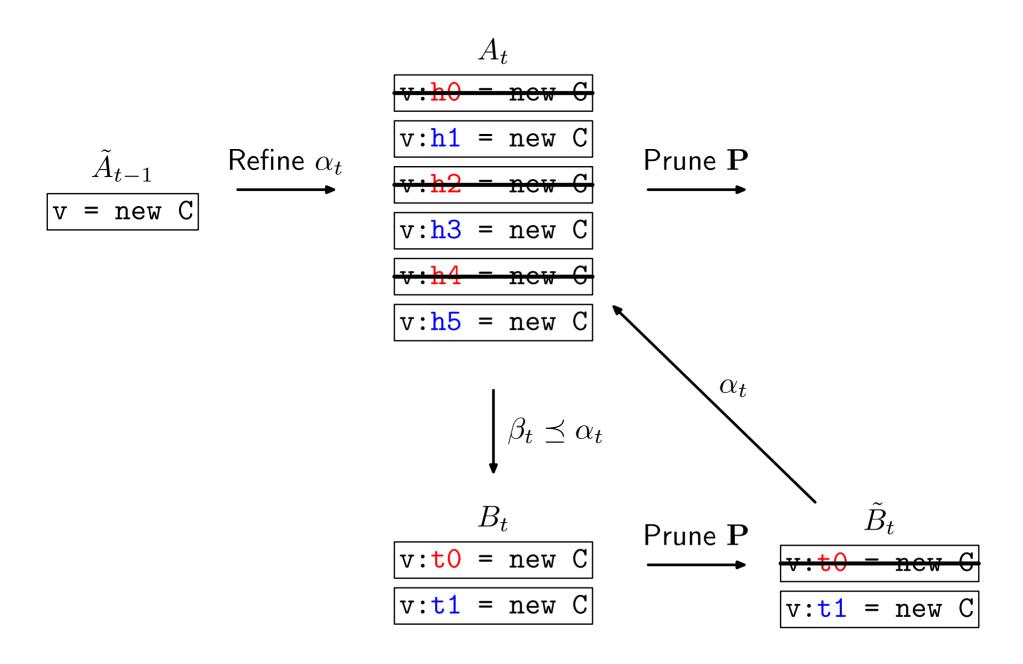
$$v = \text{new } C$$

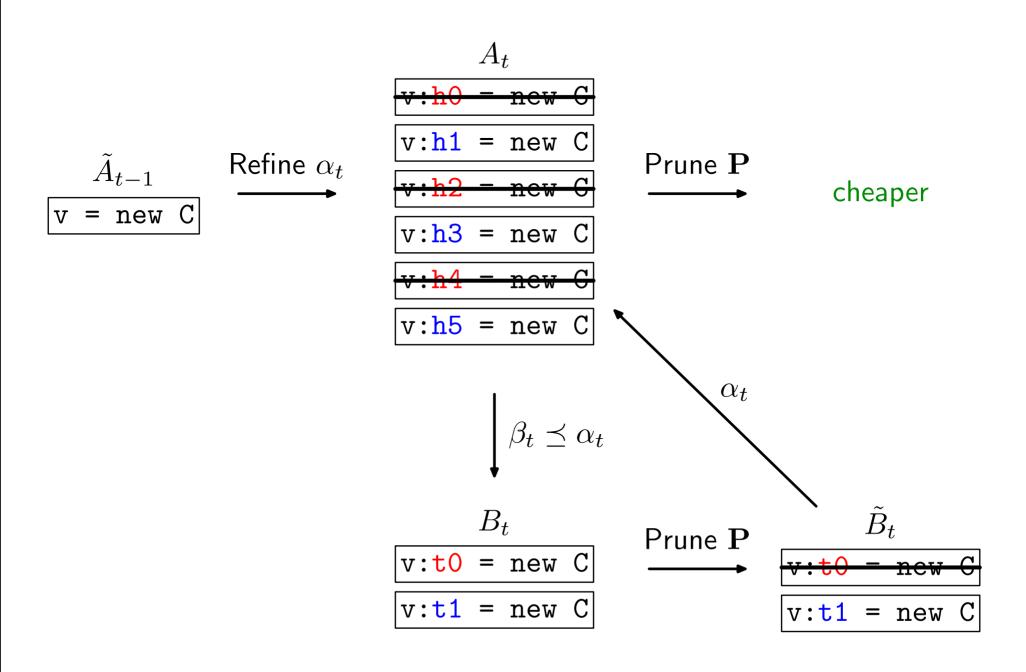






$$\begin{array}{c} A_t \\ \hline v: h0 = \text{new C} \\ \hline v: h1 = \text{new C} \\ \hline v: h2 = \text{new C} \\ \hline v: h3 = \text{new C} \\ \hline v: h4 = \text{new C} \\ \hline v: h5 = \text{new C} \\ \hline \end{array}$$

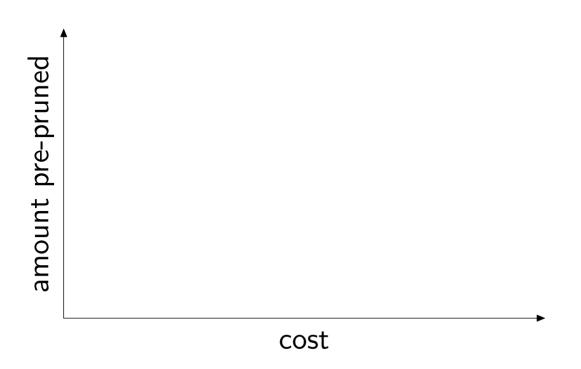




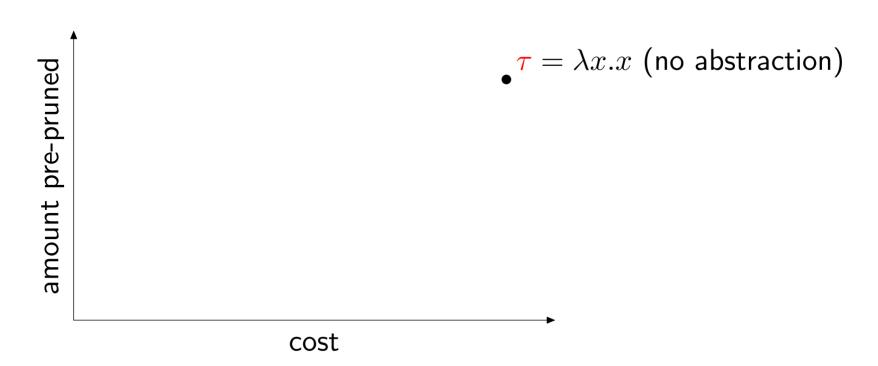
Choose abstraction τ ;

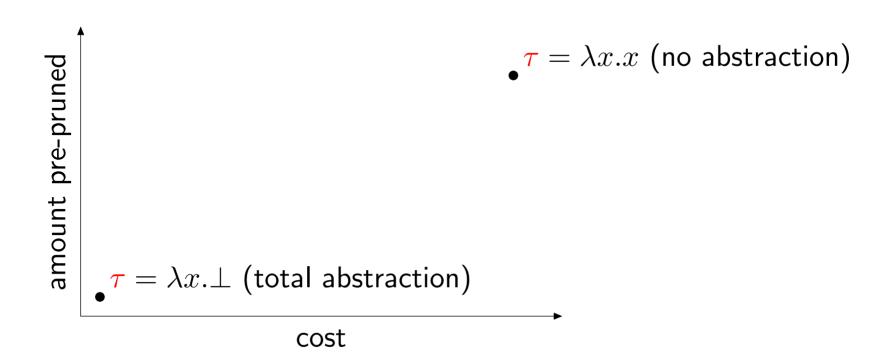
(main)
$$\alpha_0 \preceq \alpha_1 \preceq \alpha_2 \preceq \cdots$$
 $\gamma_1 \qquad \gamma_1 \qquad \gamma_1$
(auxiliary) $\beta_0 \preceq \beta_1 \preceq \beta_2 \preceq \cdots$

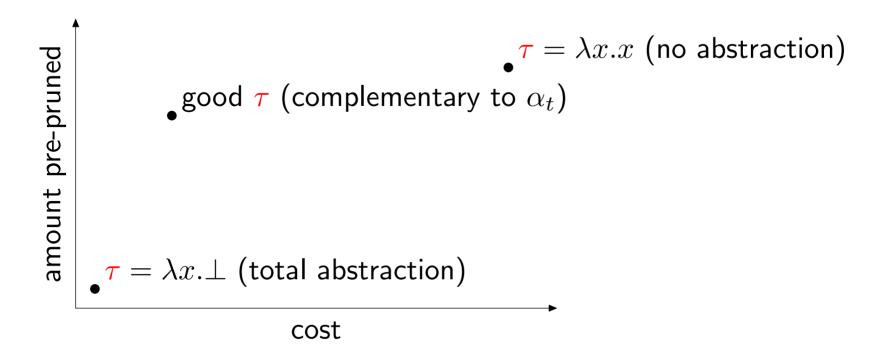
(main)
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k-limited: $\alpha_k = \text{take length } k \text{ prefix}$ $v:h5:h8 = \text{new C} \xrightarrow{\alpha_1} v:h5 = \text{new C}$

k-limited: α_k = take length k prefix $v:h5:h8 = new C \longrightarrow v:h5 = new C$ Type-based: τ = replace alloc. sites with types [Smaragdakis et al. 2011]

$$v:h5:h8 = new C \xrightarrow{\mathcal{T}} v:t1:t0 = new C$$

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We use $\tau = \text{type}$ of containing class $\times \text{type}$ of allocation site

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```

```
We use \tau = \text{type} of containing class \times type of allocation site class C1 {
    h1: x = \text{new C2}
```

k-limited: $\alpha_k = \text{take length } k \text{ prefix}$ $v:h5:h8 = \text{new C} \xrightarrow{\alpha_1} v:h5 = \text{new C}$

Type-based: τ = replace alloc. sites with types [Smaragdakis et al. 2011]

$$v:h5:h8 = new C \longrightarrow v:t1:t0 = new C$$

We use $\tau = \mathsf{type}$ of containing class $\times \mathsf{type}$ of allocation site

class C1 {
h1:
$$x = \text{new C2}$$
} $t = \text{new C2}$

$$t = \text{h1} = t =$$

k-limited: α_k = take length k prefix $v:h5:h8 = new C \longrightarrow v:h5 = new C$

Type-based: τ = replace alloc. sites with types [Smaragdakis et al. 2011]

$$\boxed{\text{v:h5:h8 = new C}} \xrightarrow{\tau} \boxed{\text{v:t1:t0 = new C}}$$

We use $\tau = \mathsf{type}$ of containing class $\times \mathsf{type}$ of allocation site

class C1 {
h1:
$$x = \text{new C2}$$
} $(C1,C2)$

Composed:
$$\beta_1 = \alpha_1 \circ \tau$$

$$v:h5:h8 = new C \longrightarrow v:t1 = new C$$

Rest of Talk

Pre-Pruning Extension

Experiments

Clients (based on flow-insensitive k-object-sensitivity):

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Benchmarks:

	description	# bytecodes	# alloc. sites
elevator	discrete event simulation program	39K	637
hedc	web crawler	151K	1,494
weblech	website downloading and mirroring tool	230K	2,545
lusearch	text indexing and search tool	267K	2,822
avrora	AVR microcontroller simulation/analysis framework	312K	4,823

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Details:

64-bit IBM J9VM 1.6, Chord with bddbddb Datalog solver Terminate a process if it exceeds 8GB of memory

Curbing the Exponential Growth

Methods:

no pruning

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- × selected refinement [Liang et al. 2011]

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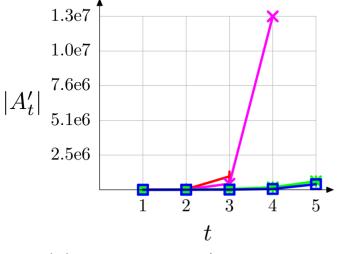
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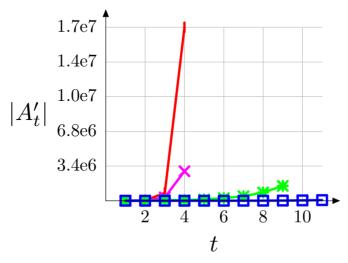
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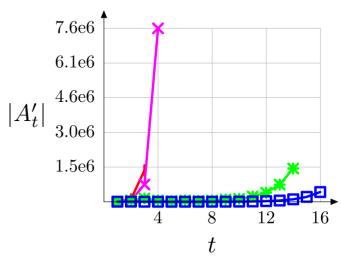
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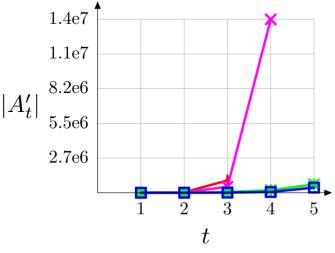
(b) DOWNCAST/lusearch



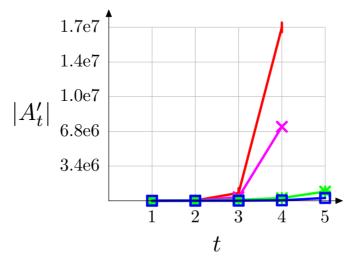
(c) DOWNCAST/avrora

Methods:

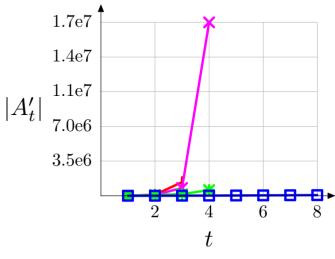
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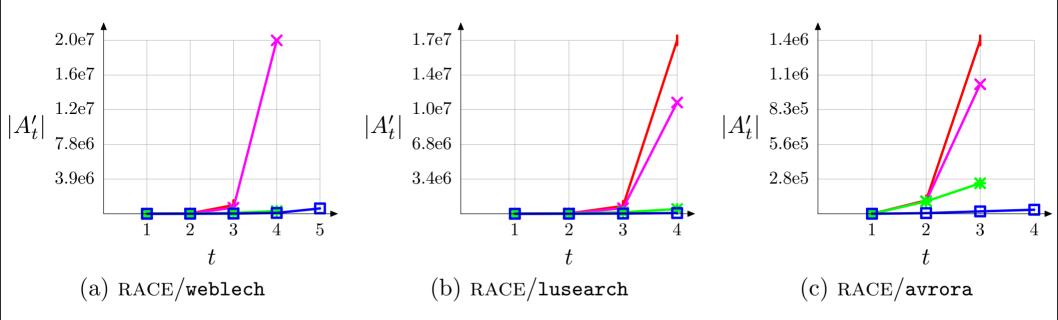
(b) MONOSITE/lusearch



(c) MONOSITE/avrora

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In each iteration, what fraction of tuples are kept?

 A_t

100%

$$A_t \xrightarrow{\text{Pre-Prune}} A'_t$$

$$100\% \qquad 22\%$$

$$A_t$$
 Pre-Prune A'_t Prune P \tilde{A}_t Refine α_{t+1} A_{t+1} 100% 22% 18% 272%

In each iteration, what fraction of tuples are kept?

$$A_{t} \xrightarrow{\text{Pre-Prune}} A'_{t} \xrightarrow{\text{Prune P}} \tilde{A}_{t} \xrightarrow{\text{Refine } \alpha_{t+1}} A_{t+1}$$

$$100\% \qquad 22\% \qquad 18\% \qquad 272\%$$

(numbers averaged across all clients/benchmarks/iterations)

In each iteration, what fraction of tuples are kept?

Take Away: Pruning (especially pre-pruning) helps a lot to maintain tractability

Impact on Queries Proven

How many queries remain unproven?

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client/benchmark $\setminus k$	1	2	3	4	5
DOWNCAST/elevator	0	-	-	_	_
$ ext{DOWNCAST}/ ext{hedc}$	10	8	3	2	2
${ t DOWNCAST}/{ t weblech}$	24	14	6	6	_
DOWNCAST/lusearch	36	14	6	5	5
DOWNCAST/avrora	12	10	6	6	6
MONOSITE/elevator	1	1	1	1	1
MONOSITE/hedc	164	149	149	149	-
MONOSITE/weblech	273	258	252	252	-
MONOSITE/lusearch	593	454	447	447	-
MONOSITE/avrora	288	278	272	_	-
$\mathrm{RACE}/\mathtt{elevator}$	475	440	437	437	437
$\mathrm{RACE}/\mathtt{hedc}$	23,033	22,043	21,966	_	-
$\mathrm{RACE} / \mathtt{weblech}$	7,286	4,742	4,669	_	-
$\mathrm{RACE}/\mathtt{lusearch}$	33,845	23,509	16,957	_	-
RACE/avrora	62,060	61,807	61,734	-	-

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Take Away: By using Prune-Refine, able to prove two additional queries

• Goal: scale up static analyses

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• Contribution: new general pruning framework

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• Key Idea: use coarse abstraction to remove irrelevant tuples

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- Theoretical Result: pruning is correct

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- Contribution: new general pruning framework
- Key Idea: use coarse abstraction to remove irrelevant tuples
- Theoretical Result: pruning is correct
- Empirical Result: enable much finer abstractions

- Goal: scale up static analyses
- Contribution: new general pruning framework
- Key Idea: use coarse abstraction to remove irrelevant tuples
- Theoretical Result: pruning is correct
- Empirical Result: enable much finer abstractions

http://code.google.com/p/jchord

Thank you!