

AI SMPS 2023 Week 5 Algorithms

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Algorithms

$A^*(S)$

```
1  default value of g for every node is  $+\infty$ 
2  parent(S)  $\leftarrow$  null
3  g(S)  $\leftarrow$  0
4  f(S)  $\leftarrow$  g(S) + h(S)
5  OPEN  $\leftarrow$  S : []
6  CLOSED  $\leftarrow$  empty list
7  while OPEN is not empty
8      N  $\leftarrow$  remove node with lowest f value from OPEN
9      add N to CLOSED
10     if GOAL-TEST(N) = TRUE
11         return RECONSTRUCT-PATH(N)
12     for each M in MOVE-GEN(N)
13         if g(M) > g(N) + k(N, M)
14             parent(M)  $\leftarrow$  N
15             g(M)  $\leftarrow$  g(N) + k(N, M)
16             f(M)  $\leftarrow$  g(M) + h(M)
17         if M is in OPEN
18             continue
19         else if M is in CLOSED
20             PROPAGATE-IMPROVEMENT(M)
21         else add M to OPEN      ▷ M is new
22 return empty list
```

PROPAGATE-IMPROVEMENT(M)

```
23 for each X in MOVE-GEN(M)
24     if g(X) > g(M) + k(M, X)
25         parent(X)  $\leftarrow$  M
26         g(X)  $\leftarrow$  g(M) + k(M, X)
27         f(X)  $\leftarrow$  g(X) + h(X)
28     if X is in CLOSED
29         PROPAGATE-IMPROVEMENT(X)
30 return
```

For state space search, BNB maintains the pair (path,g) in the OPEN list, and expands the cheapest path in each iteration.

BNB(S)

```
1  OPEN  $\leftarrow$  (S : [], 0) : []      ▷ use priority queue
2  while OPEN is not empty
3      (path, g)  $\leftarrow$  head OPEN    ▷ (path, cost)
4      N  $\leftarrow$  head path
5      if GOAL-TEST(N) = TRUE
6          return reverse path
7      else newPaths  $\leftarrow$  empty list
8          for each M in MOVE-GEN(N)
9              newPaths  $\leftarrow$  (M : path, g + k(N, M)) : newPaths
10         OPEN  $\leftarrow$  sortg (newPaths ++ tail OPEN)
11 return empty list
```

BNB-2(S)

A variation of BNB, that prevents cyclic expansions like:

$S, \dots, A, \dots, S, \dots, A, \dots,$

GENERALIZED-BNB

```
1  start with all possible solutions
2  repeat
3      refine the least (estimated cost) solution further
4  until the cheapest solution S is fully refined
5  return S
```