CSCI971 Advance Computer Security: Homework #7

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Problem 1

AE-secure $\Leftrightarrow$ semantically secure under CPA and CI.
For the first cipher, assume an attacker who can perform CPA. He intercept the ciphertext $c = E_1(k, m) = (E(k, m), H_1(m))$, He can perform as many as CPA. We assume in CPA attack game. Adversary $\mathcal{A}$ first send m_0, m_0 to challenger $\mathcal{C}$, he get the ciphertext $c = (E(k_0, m_0), H_1(m_0))$. Then $\mathcal{A}$ send m_0, m_1 to $\mathcal{C}$, as E is CPA secure, so key has to be changed. $\mathcal{A}$ get the ciphertext $c = (E(k_1, m_0), H_1(m_0))$ or $c = (E(k_1, m_1), H_1(m_1))$ based on b. Then if $b = 1$, $\mathcal{A}$ can easily differ the plaintext from the tag $H_1(m_b)$. So $Adv_{CPA}(\mathcal{A}, \mathcal{E}) = 1/2$ is not negligible. Cipher1 is not CPA-secure, so it's not AE-secure.
For the second cipher, attacker can intercept the ciphertext $(c, H_2(c))$, so he can learn the mapping model of H_2 function. So in CI attack game, Adversary $\mathcal{A}$ can easily generate an valid ciphertext-tag pair $(c_{atk}, H_2(c_{atk}))$. Then Decryptor $D_2(k, (c_{ack}, H_2(c_{ack}))) \neq \perp$. So $Adv_{CI}(\mathcal{A}, \mathcal{E})$ is not negligible. Cipher1 does not safisfy CI, so it's not AE-secure.

Problem 2

Addition $\mathcal{Z}_6^*$ is a cyclic group.
$\mathcal{Z}_6^* = \{0, 1, 2, 3, 4, 5, 6\}$
0 generate {0},
1 generate {0,1,2,3,4,5}
2 generate {0,2,4}
3 generate {0,3}
4 generate {0,4}
5 generate {5,4,3,2,1,0}
So the generators of $\mathcal{Z}_6^*$ is 0 and 5.

Problem 3

Group under multiplication $\mathcal{Z}_{13}^*$ is a cyclic group.
<0> = {0}
<1> = {1}
<2> = {1 2 4 8 3 6 12 11 9 5 10 7}
<3> = {1 3 9}
<4> = {1 4 3 12 9 10 1 4 3 12 9 10}
<5> = {1 5 12 8}
<6> = {1 6 10 8 9 2 12 7 3 5 4 11}
<7> = {1 7 10 5 9 11 12 6 3 8 4 2}
<8> = {1 8 12 5 }
<9> = {1 9 3 }
<10> = {1 10 9 12 3 4}
<11> = {1 11 4 5 3 7 12 2 9 8 10 6}
<12> = {1 12 1 }
subgroup is 2,4,6,7,11.