CSCI971 Advance Computer Security: Homework #9

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This protocol is like EIgamal encryption mode.

The $sk \leftarrow k, pk \leftarrow g^k$.

Alice knows the public key pk and $F(k,m) = H(m)^k$, she chx ge a random $\rho \leftarrow Z_q$ and sends Bob $\widehat{m} = H(m) \cdot g^{\rho}$.

We assume $v \leftarrow g^{\rho}, \omega \leftarrow pk^{\rho} = g^{\rho k} = v^k$.

When Bob get the \widehat{m} , he respond $res = \widehat{m}^k = H(m)^k \cdot g^{\rho k} = H(m)^k \cdot \omega$ to Alice, as H(m) is random oracle, so Bob cannot know the m from H(m).

Wh en Alice get the res, she knows $\omega = g^{k\rho}$ so she just get $H(m)^k = res/(g^{k\rho})$.

Because It is hard to get k from \widehat{m}^k as it is a hard problem in number theory. So Alice doesn't know k.

Game 0

Barcelona Reasonable.

 ${\bf Game}\ {\bf 1}$

sdfas