# CSCI971 Advance Computer Security: Homework #4

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# Problem 1

First, generate differential table for 3-bit Sbox Table 1 and 5-bit Sbox Table 2;

	0	1	2	3	4	5	6	7
0	8 0 0 0 0 0 0	0	0	0	0	0	0	0
1	0	2	0	2	0	2	0	2
2	0	0	2	2	0	0	2	2
3	0	2	2	0	0	2	2	0
4	0	0	0	0	2	2	2	2
5	0	2	0	2	2	0	2	0
6	0	0	2	2	2	2	0	0
7	0	2	2	0	2	0	0	2

Table 1: Differential distribution table for 3-bit Sbox

We set Table 1 as T1, Table 2 as T2:

We can find some extremum value point. T1(1,1), T1(2,2), T1(4,4), T2(1,1), T2(2,2), T2(4,4), T2(8,8), T2(16,16) with the Probability 1/4 to maintain its input differential value.

Because only Sbox can contribute to changes to differential value based on probability, the MixRow and BitRot change the differential value in a fixed mode (with the probibility of 1). We can just take Sbox into concern.

We can enumerate the input differential value 0x01, 0x02, 0x04, 0x08, 0x10 to the Sbox to get the best probability path.

Then we can find when the  $p_0$  of plaintext is 0x20, we can attack as much as round. The plaintext = 0x20 0x00 0

## We assume Sbox keep its differential value.

#### Round1

After entering the Sbox, the  $p_0$  state has the probability of  $2/8 = 2^{-2}$  to maintain 0x01.

Maintain value probability  $P_1 = 1/4 = 2^{-2}$ 

## Round2

We have 0x02 0x10 and 0x01.

Maintain value probability  $P_2 = 1/4 * 1/4 = 2^{-4}$ 

#### Round3

We have  $0x01 \ 0x02 \ 0x04 \ 0x04 \ 0x10 \ 0x02$ .

Maintain value probability  $P_3 = 1/4 * 1/4 * 1/4 * 1/4 = 2^{-8}$ 

#### Round4

We have  $0x04 \ 0x04 \ 0x01 \ 0x08 \ 0x03 => 0x02$ .

Maintain value probability  $P_4 = 1/4 * 1/4 * 1/4 * 1/4 * 1/4 * 1/4 = 2^{-10}$ 

#### Round5

We have  $0x01\ 0x0a => 0x01\ 0x03\ 0x04\ 0x1c => 0x02\ 0x02\ 0x02\ 0x02\ 0x02\ 0x02$ 

Maintain value probability  $P_5 = 1/4 * 1/8 * 1/8 * 1/4 * 1/8 * (1/4)^4 = 2^{-21}$ 

## Round6

We have  $0x10\ 0x02\ 0x04\ 0x02\ 0x01\ 0x10\ x014 => 0x13\ 0x01\ 0x03 => 0x01\ 0x02$ 

Maintain value probability  $P_6 = 1/4 * 1/4 * 1/4 * 1/4 * 1/4 * 1/4 * 1/8 * 1/4 * 1/8 * 1/4 = 2^{-24}$ 

## Round7

We have  $0x09 = >0x07 \ 0x07 = >0x07 \ 0x02 \ 0x19 = >0x19 \ 0x06 = >0x02 \ 0x06 = >0x02 \ 0x01 \ 0x02 \ 0x0a = >0x01 \ 0x09 = >0x0a \ 0x08 = >0x08$ 

Maintain value probability  $P_7 = 1/8 * 1/16 * 1/4 * 1/16 * 1/8 * 1/8 * 1/4 * 1/4 * 1/4 * 1/8 * 1/8 * 1/4 = 2^{-31}$ 

Lefts	State	:		Right	tState	e :	
0×01	0×00	0×00	0×00	0×00	0×00	0×00	0×00
0×00	0×00	0×00	0×00	0×00	0×00	0×00	0×00
0×00	0×00	0×00	0×00	0×00	0×00	0×00	0×00
0×00	0×00	0×00	0×00	0×00	0×00	0×00	0×00
Aftei	r MixF	Row ar	nd BitRot				
0×00	0×00	0×00	0×00	0×00	0×00	0×00	0×00
0×00	0×00	0×00	0×02	0×00	0×00	0×00	0×00
0×00	0×00	0×00	0×00	0×00	0×00	0×00	0×00
0×0	0×00	0×10	0×00	0×00	0×00	0×0	0×00

Figure 1: Round 1

# Round8

Maintain value probability  $P_8 = 1/4 * 1/4 * 1/8 * 1/8 * 1/4 * 1/4 * 1/8 = 2^{-17}$ 

# Round9

Maintain value probability  $P_8 = 1/4*1/16*1/8*1/8*1/16*1/8*1/8*1/4*1/16*1/8*1/8*1/8*1/8*1/8*1/4=2^{-39}$  Accumulate Round1 - Round 9 less than  $2^{-128}$ .

So we get to Round 9

Rounc	l=2:							
LeftS	State			RightState:				
0×00	0×00	0×00	0×00	0×01	0×00	0×00	0×00	
0×00	0×00	0×00	0×02	0×00	0×00	0×00	0×00	
0×00	0×00	0×00	0×00	0×00	0×00	0×00	0×00	
0×00	0×00	0×10	0×00	0×00	0×00	0×00	0×00	
After	MixF	Row ar	nd BitRot					
0×00	0×04	0×00	0×00	0×01	0×00	0×00	0×00	
0×02	0×04	0×00	0×00	0×00	0×00	0×00	0×00	
0×00	0×00	0×00	0×00	0×00	0×00	0×00	0×00	
0×00	0×00	0×00	0×00	0×00	0×00	0×00	0×00	

Figure 2: Round 2

Round=3:						
LeftState:			Right	tState	≘:	
0×01 0×04	0×00	0×00	0×00	0×00	0×00	0×00
0×02 0×04	0×00	0×00	0×00	0×00	0×00	0×02
0×00 0×00	0×00	0×00	0×00	0×00	0×00	0×00
0×00 0×00	0×00	0×00	0×00	0×00	0×10	0×00
After MixF	Row ar	nd BitRot				
0×04 0×00	0×00	0×00	0×00	0×00	0×00	0×00
0×04 0×00	0×00	0×0a	0×00	0×00	0×00	0×02
0×00 0×00	0×00	0×03	0×00	0×00	0×00	0×00
0×00 0×00	0×11	0×00	0×00	0×00	0×10	0×00

Figure 3: Round 3

Round	d=4:						
Lefts	State			Right	tState	≘:	
0×04	0×00	0×00	0×00	0×01	0×04	0×00	0×00
0×04	0×00	0×00	0×08	0×02	0×04	0×00	0×00
0×00	0×00	0×00	0×03	0×00	0×00	0×00	0×00
0×00	0×00	0×01	0×00	0×00	0×00	0×00	0×00
Aftei	· MixF	Row ar	nd BitRot				
0×00	0×00	0×00	0×00	0×01	0×04	0×00	0×00
0×08	0×18	0×00	0×08	0×02	0×04	0×00	0×00
0×03	0×02	0×02	0×02	0×00	0×00	0×00	0×00
0×00	0×02	0×02	0×00	0×00	0×00	0×00	0×00

Figure 4: Round 4

Round=5:						
LeftState	:		Right	State	≘:	
0×01 0×04	0×00	0×00	0×04	0×00	0×00	0×00
0×0a 0×1c	0×00	0×08	0×04	0×00	0×00	0×08
0×03 0×02	0×02	0×02	0×00	0×00	0×00	0×03
0×00 0×02	0×02	0×00	0×00	0×00	0×01	0×00
After Mix	Row ar	nd BitRot				
0×04 0×00	0×01	0×01	0×04	0×00	0×00	0×00
0×14 0×00	0×10	0×0b	0×04	0×00	0×00	0×08
0×00 0×04	0×00	0×03	0×00	0×00	0×00	0×03
0×02 0×02	0×15	0×02	0×00	0×00	0×01	0×00

Figure 5: Round 5

Round=6:					
LeftState:	RightState:				
0×00 0×00 0	×01 0×01	0×01	0×04	0×00	0×00
0×10 0×00 0	×10 0×03	0×0a	0×1c	0×00	0×08
0×00 0×04 0	×00 0×00	0×03	0×02	0×02	0×02
0×02 0×02 0	×14 0×02	0×00	0×02	0×02	0×00
After MixRo	w and BitRot				
0×01 0×06 0	×00 0×00	0×01	0×04	0×00	0×00
0×03 0×05 0	×01 0×01	0×0a	0×1c	0×00	0×08
0×04 0×04 0	×00 0×00	0×03	0×02	0×02	0×02
0×00 0×04 0	×08 0×08	0×00	0×02	0×02	0×00

Figure 6: Round 6

Round=7:						
LeftState	:		Right	State	≘:	
0×00 0×02	0×00	0×00	0×00	0×00	0×01	0×01
0×09 0×19	0×01	0×09	0×10	0×00	0×10	0×03
0×07 0×06	0×02	0×02	0×00	0×04	0×00	0×00
0×00 0×06	0×0a	0×08	0×02	0×02	0×14	0×02
After Mix	Row ar	nd BitRot				
0×02 0×00	0×07	0×03	0×00	0×00	0×01	0×01
0×10 0×00	0×10	0×05	0×10	0×00	0×10	0×03
0×04 0×04	0×00	0×00	0×00	0×04	0×00	0×00
0×0e 0×02	0×04	0×02	0×02	0×02	0×14	0×02

Figure 7: Round 7

9
9
2
8
9
9
2
8
)! )!

Figure 8: Round 8

Round=9:						
LeftStat	e:		Right	State	≘:	
0×02 0×0	6 0×02	0×06	0×02	0×00	0×06	0×02
0×0f 0×1	d 0×07	0×0d	0×00	0×00	0×00	0×06
0×07 0×0	7 0×03	0×02	0×04	0×00	0×00	0×00
0×00 0×0	7 0×0a	0×00	0×0c	0×00	0×10	0×00
After Mi	xRow a	nd BitF	Rot			
0×00 0×0	4 0×07	0×03	0×02	0×00	0×06	0×02
0×10 0×0	a 0×10	0×19	0×00	0×00	0×00	0×06
0×05 0×0	5 0×01	0×01	0×04	0×00	0×00	0×00
0×07 0×0	2 0×07	0×00	0×0c	0×00	0×10	0×00

Figure 9: Round 9

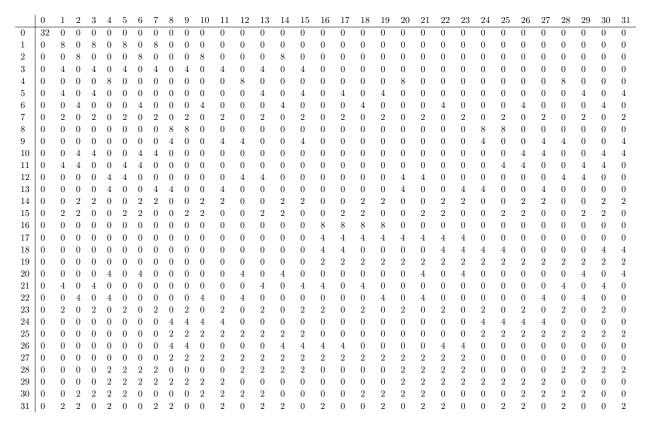


Table 2: Differential distribution table for 5-bit Sbox