# CSCI971 Advance Computer Security: Homework #9

Mei Wangzhihui 2019124044

This protocol is like EIgamal encryption mode.

The  $sk \leftarrow k, pk \leftarrow g^k$ .

Alice knows the public key pk and  $F(k,m) = H(m)^k$ , she choose a random  $\rho \leftarrow Z_q$  and sends Bob  $\widehat{m} = H(m) \cdot g^{\rho}$ .

We assume  $v \leftarrow g^{\rho}, \omega \leftarrow pk^{\rho} = g^{\rho k} = v^k$ .

When Bob get the  $\widehat{m}$ , he respond  $res = \widehat{m}^k = H(m)^k \cdot g^{\rho k} = H(m)^k \cdot \omega$  to Alice, as H(m) is random oracle, so Bob cannot know the m from H(m).

Wh en Alice get the res, she knows  $\omega = g^{k\rho}$  so she just get  $H(m)^k = res/(g^{k\rho})$ .

Because It is hard to get k from  $\widehat{m}^k$  as it is a hard problem in number theory. So Alice doesn't know k.

Game 0

fa

 ${\bf Game}\ {\bf 1}$ 

sdfas