

# CSE 260 : Digital Logic Design

## Number Systems and Codes

# Binary Coded Decimal (BCD)

- Decimal numbers are more natural to humans. Binary numbers are natural to computers. Quite expensive to convert between the two.
- If little calculation is involved, we can use some *coding schemes* for decimal numbers.
- One such scheme is **BCD**, also known as the **8421** code.
- Represent each decimal digit as a **4-bit binary code**.



# Binary Coded Decimal (BCD)

Decimal digit	0	1	2	3	4
<b>BCD</b>	<b>0000</b>	<b>0001</b>	<b>0010</b>	<b>0011</b>	<b>0100</b>
Decimal digit	5	6	7	8	9
<b>BCD</b>	<b>0101</b>	<b>0110</b>	<b>0111</b>	<b>1000</b>	<b>1001</b>

- Also known as the **8421** code.
- Represent each decimal digit as a **4-bit binary code**.
- Some codes are unused, eg:  $(1010)_{\text{BCD}}$ ,  $(1011)_{\text{BCD}}$ , ...,  $(1111)_{\text{BCD}}$ . These codes are considered as errors.
- Easy to convert, but arithmetic operations are more complicated.
- Suitable for interfaces such as keypad inputs and digital readouts.

# Binary Coded Decimal (BCD)

Decimal digit <b>BCD</b>	0 <b>0000</b>	1 <b>0001</b>	2 <b>0010</b>	3 <b>0011</b>	4 <b>0100</b>
Decimal digit <b>BCD</b>	5 <b>0101</b>	6 <b>0110</b>	7 <b>0111</b>	8 <b>1000</b>	9 <b>1001</b>

## ■ Examples:

$$(234)_{10} = (0010\ 0011\ 0100)_{\text{BCD}}$$

$$(7093)_{10} = (0111\ 0000\ 1001\ 0011)_{\text{BCD}}$$

$$(1000\ 0110)_{\text{BCD}} = (86)_{10}$$

$$(1001\ 0100\ 0111\ 0010)_{\text{BCD}} = (9472)_{10}$$

Notes: BCD is **not equivalent** to binary.

$$\text{Example: } (234)_{10} = (11101010)_2$$



# Binary Codes

## ■ Other Codes

<b>Decimal Digit</b>	<b>BCD 8421</b>	<b>Excess-3</b>
0	0000	0011
1	0001	0100
2	0010	0101
3	0011	0110
4	0100	0111
5	0101	1000
6	0110	1001
7	0111	1010
8	1000	1011
9	1001	1100
Unused bit combi- nations	1010	0000
	1011	0001
	1100	0010
	1101	1101
	1110	1110
	1111	1111

# Negative Numbers Representation

- There are three common ways of representing signed numbers (positive and negative numbers) for binary numbers:
  - ❖ Sign-and-Magnitude
  - ❖ 1s Complement
  - ❖ 2s Complement

# Sign-and-Magnitude

- Negative numbers are usually written by writing a minus sign in front.

❖ Example:

$$-(12)_{10}, -(1100)_2$$

- In computer memory of fixed width, this sign is usually represented by a bit:

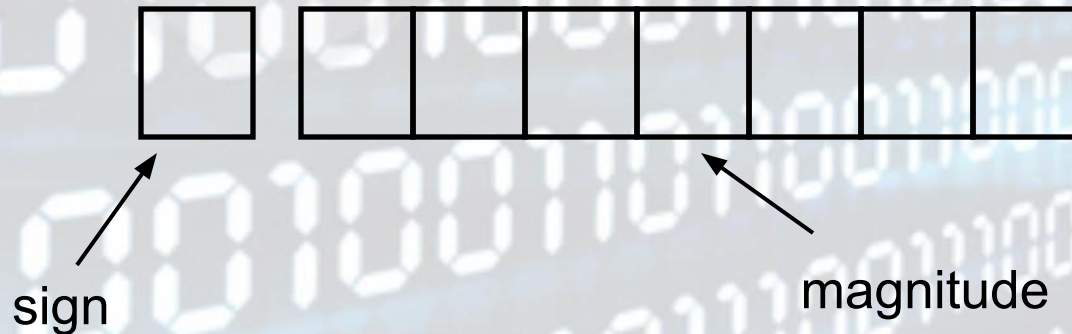
0 for +

1 for -



# Sign-and-Magnitude

- Example: an 8-bit number can have 1-bit sign and 7-bits magnitude.





# Signed magnitude representation

- Examples:

$1101_2 = 13_{10}$  (a 4-bit unsigned number)  
0  $1101 = +13_{10}$  (a positive number in 5-bit signed magnitude)  
1  $1101 = -13_{10}$  (a negative number in 5-bit signed magnitude)

$0100_2 = 4_{10}$  (a 4-bit unsigned number)  
0  $0100 = +4_{10}$  (a positive number in 5-bit signed magnitude)  
1  $0100 = -4_{10}$  (a negative number in 5-bit signed magnitude)

# Sign-and-Magnitude

- Largest Positive Number: 0 1111111  $+(127)_{10}$
- Largest Negative Number: 1 1111111  $-(127)_{10}$
- Zeroes:  
0 0000000  $+(0)_{10}$   
1 0000000  $-(0)_{10}$
- Range:  $-(127)_{10}$  to  $+(127)_{10}$
- Signed numbers needed for negative numbers.
- Representation: Sign-and-magnitude.



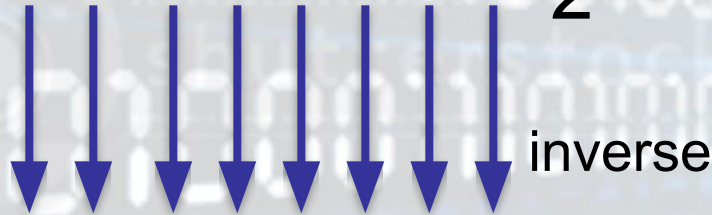
# 1s Complement

- Given a number  $x$  which can be expressed as an  $n$ -bit binary number (*i.e. integer part has  $n$  digits and fraction has  $m$  digit*), its negative value can be obtained in **1s-complement** representation using:

+7	0111	-7	1000
+6	0110	-6	1001
+5	0101	-5	1010
+4	0100	-4	1011
+3	0011	-3	1100
+2	0010	-2	1101
+1	0001	-1	1110
+0	0000	-0	1111

# 1's complement

$-(00001100)_2$



$(11110011)_2$



# 1s Complement

- Essential technique: **invert** all the bits.

Examples: 1s complement of 00000001 =  $(11111110)_{1s}$

1s complement of 01111111 =  $(10000000)_{1s}$

- Range [in 8 bits]:  $-(127)_{10}$  to  $+(127)_{10}$

- **General Formula** =  $-(2^{n-1} - 1)$  to  $+(2^{n-1} - 1)$

- The most significant bit still represents the sign:  
0 = +ve; 1 = -ve.

**Note: Range for n bit no. is  $-(2^{(n-1)}-1)$  to  $(2^{(n-1)}-1)$**

# 1s Complement

- Examples (assuming 8-bit binary numbers):

$$(14)_{10} = (00001110)_2 = (00001110)_{1s}$$

$$-(14)_{10} = -(00001110)_2 = (11110001)_{1s}$$

$$-(80)_{10} = -(?)_2 = (?)_{1s}$$



# 2s Complement

- Given a number  $x$  which can be expressed as an  $n$ -bit (*i.e. integer part has  $n$  digits and fraction has  $m$  digit*) number, its negative number can be obtained in **2s-complement** representation using:

$$-x = 2^n - x$$

Example: With an 8-bit number 00001100, its negative value in 2s complement is thus:

$$\begin{aligned} -(00001100)_2 &= -(12)_{10} \\ &= (2^8 - 12)_{10} \\ &= (244)_{10} \\ &= (11110100)_{2s} \end{aligned}$$

# 2s Complement

- **Method 1:** Essential technique: **invert** all the bits and **add 1**.

Examples:

2s complement of

$$\begin{aligned}(00000001)_{2s} &= (11111110)_{1s} \quad (\text{invert i.e 1's complement}) \\ &= (11111111)_{2s} \quad (\text{add 1})\end{aligned}$$

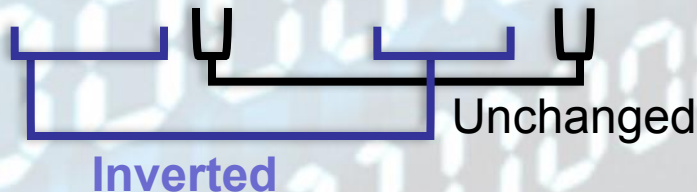
2s complement of

$$\begin{aligned}(01111110)_{2s} &= (10000001)_{1s} \quad (\text{invert i.e 1's complement}) \\ &= (10000010)_{2s} \quad (\text{add 1})\end{aligned}$$

**Official method!**

- **Method 2:** Keep unchanged till 1<sup>st</sup> occurrence of 1 from LSB and invert remaining 1's into 0's and 0's into 1's till MSB

$$(01111110)_{2s} = (10000010)_{2s}$$



**Unofficial method!**



# 2s Complement

- Range in 8 bits:  $-(128)_{10}$  to  $+(127)_{10}$
- General Formula=  $-(2^{n-1})$  to  $+(2^{n-1}-1)$
- The most significant bit still represents the sign:  
0 = +ve; 1 = -ve.

**Note: Range for n bit no. is  $-2^{(n-1)}$  to  $2^{(n-1)}-1$**



# 2s Complement

- Examples (assuming 8-bit binary numbers):

$$(14)_{10} = (00001110)_2 = (00001110)_{2s}$$

$$-(14)_{10} = -(00001110)_2 = (11110010)_{2s}$$

$$-(80)_{10} = -( ? )_2 = ( ? )_{2s}$$

# Comparisons of Sign-and-Magnitude and Complements

- Example: 4-bit signed number (*positive values*)

Value	Sign-and-Magnitude	1s Comp.	2s Comp.
+7	0111	0111	0111
+6	0110	0110	0110
+5	0101	0101	0101
+4	0100	0100	0100
+3	0011	0011	0011
+2	0010	0010	0010
+1	0001	0001	0001
+0	0000	0000	0000

Note: Signed magnitude cannot be used for arithmetic calculations



# Comparisons of Sign-and-Magnitude and Complements

- Example: 4-bit signed number (*negative values*)

Value	Sign-and-Magnitude	1s Comp.	2s Comp.
-0	1000	1111	-
-1	1001	1110	1111
-2	1010	1101	1110
-3	1011	1100	1101
-4	1100	1011	1100
-5	1101	1010	1011
-6	1110	1001	1010
-7	1111	1000	1001
-8	-	-	1000

- Note: Signed magnitude cannot be used for arithmetic calculations



# Exercise:


1. For 2's complement binary numbers, the range of values for 5-bit numbers is
  - a. 0 to 31
  - b. -8 to +7
  - c. -8 to +8
  - d. -15 to +15
  - e. -16 to +15
2. In a 6-bit 2's complement binary number system, what is the decimal value represented by  $(100100)_{2s}$ ?
  - a. -4
  - b. 36
  - c. -36
  - d. -27
  - e. -28

Solution:

1) Following numbers are in 1's complement system. Turn them to their no. negative representation

- A. 1010101
- B. 0111000
- C. 0000001
- D. 00000

2) Now perform 2's complement



**UNSIGNED NO. ARTHMETIC  
OPERATION!**



# Binary Arithmetic Operations for **Unsigned** numbers

## ■ ADDITION

- Like decimal numbers, two numbers can be added by adding each pair of digits together with carry propagation.

$$\begin{array}{r} (11011)_2 \\ + (10011)_2 \\ \hline (101110)_2 \end{array}$$

$$\begin{array}{r} (647)_{10} \\ + (537)_{10} \\ \hline (1184)_{10} \end{array}$$



# Binary Arithmetic Operations for Unsigned Numbers

## ■ SUBTRACTION

- Two numbers can be subtracted by subtracting each pair of digits together with borrowing, where needed.

$$\begin{array}{r} (11001)_2 \\ - (10011)_2 \\ \hline (00110)_2 \end{array} \qquad \begin{array}{r} (627)_{10} \\ - (537)_{10} \\ \hline (090)_{10} \end{array}$$



**Overflow!**

## ▣ Conditions of Overflow Flag :

1. Addition of numbers with same sign,

→  $(+A) + (+B) = '+'$  then  $OF = 0$

$(-)$  "  $OF = 1$

→  $(-A) + (-B) = '-'$  then  $OF = 0$

$(-A) + (-B) = '+'$  "  $OF = 1$

\* 2ଟି same sign -ବା number add କଲାବେଳେ ଓହ୍ଲେ  
sign ନା -ବାସ୍ତବରେ ଗଣନା କଲାବେଳେ  $OF = 1$

that means  $OF = 1$  otherwise  $OF = 0$ .



## 2. Subtraction of numbers with same sign

2. If same sign (+ or -) number subtract (বিভাজন)  
সবসময় overflow হয় না that means  
OF = 0 [  $(+A) - (+B) = \text{no overflow}$ ;  $(-A) - (-B) = \text{no overflow}$  ]

## 3. Addition of numbers with different sign

2. If different sign (+ or -) number (ভাজন) সবসময়  
OF = 0 হয় always overflow হয় না,

$$(-A) + (+B) = \text{no overflow}$$

$$(+A) + (-B) = \text{no overflow}$$

#### 4. Subtraction of numbers with different sign

$$\Rightarrow (+A) - (-B) = A + B$$

- এটা কে simplify করতে  
- আমরা একই সূত্রের form - এ চলে আসি,

- তাহলে  $A + B$  - এর result যদি (+) - তখন

- তাহলে  $OF = 0$  - তার মানে result (-)

- আসে then  $OF = 1$ .

$$\Rightarrow (-A) - (+B) = -A - B = (-A) + (-B)$$

- এটা কে simplify করে - আমরা একই সূত্রের form

- এ চলে আসি, - তাহলে একই same  
sign - এর number add করলে যদি -

sign আসে then  $OF = 0$  - তার মানে

opposite sign আসে তাহলে  $OF = 1$ .





# **ARITHMETIC OPERATIONS ON SIGNED NUMBER**

# 2s Complement Addition/Subtraction

- Algorithm for addition,  $A + B$ :
  1. Perform binary addition on the two numbers.
  2. Ignore the carry out of the MSB (most significant bit).
  3. Check for overflow: Overflow occurs if the 'carry in' and 'carry out' of the MSB are different, or if result is opposite sign of A and B.
- Algorithm for subtraction,  $A - B$ :
$$A - B = A + (-B)$$
  1. Take 2s complement of B by inverting all the bits and adding 1.
  2. Add the 2s complement of B to A.



# 2s Complement Addition/Subtraction

- Examples: 4-bit binary system

+3	0011
+ +4	+ 0100
-----	-----
+7	0111
-----	-----

-1	1111
+ -5	+ 1011
-----	-----
-6	<b>1</b> 1010
-----	-----

+6	0110
+ -3	+ 1101
-----	-----
+3	<b>1</b> 0011
-----	-----

+4	0100
+ -7	+ 1001
-----	-----
-3	1101
-----	-----

- Which of the above is/are overflow(s)?

# 2s Complement Addition/Subtraction

- More examples: 4-bit binary system

-3	1101
+ -6	+ 1010
----	-----
-9	10111
----	-----

+5	0101
+ +6	+ 0110
----	-----
+11	1011
----	-----

- Which of the above is/are overflow(s)?



# 1s Complement Addition/Subtraction

- Algorithm for addition,  $A + B$ :
  1. Perform binary addition on the two numbers.
  2. If there is a carry out of the MSB, **add 1 to the result.**
  3. Check for overflow: Overflow occurs if result is opposite sign of A and B.
- Algorithm for subtraction,  $A - B$ :
$$A - B = A + (-B)$$
  1. Take 1s complement of B by inverting all the bits.
  2. Add the 1s complement of B to A.



# 1s Complement Addition/Subtraction

- Examples: 4-bit binary system

+3	0011
+ +4	+ 0100
-----	-----
+7	0111
-----	-----

+5	0101
+ -5	+ 1010
-----	-----
-0	1111
-----	-----

-2	1101
+ -5	+ 1010
-----	-----
-7	10111
-----	+ 1
	-----
	1000

-3	1100
+ -7	+ 1000
-----	-----
-10	10100
-----	+ 1
	-----
	0101

# Sample Math:

Add -3 with -6 in 6 bits using 2's complement number system and justify whether there is overflow or not.

Solution:

Unsigned 3 = 11

+3 = 011

+3 in 6 bits = 000011

To get -3 in binary, we need to calculate 2's complement.

So, 111100 (1's complement)

+1 (Adding 1)

---

-3 in 6 bits = 111101 (2's complement) ----- (1)

Unsigned 6 = 110

+6 = 0110

+6 in 6 bits = 000110

To get -6 in binary, we need to calculate 2's complement.

So, 111001 (1's complement)

+1 (Adding 1)

---

-6 in 6 bits = 111010 (2's complement) ----- (2)

Now, adding (1) and (2),

111101

111010

---

1110111

Ignoring the carry, final result is 110111

Here, we are adding two negative numbers and the result is also a negative number. Hence, there is no overflow.