

## Assignment-2 (Summer-24) Set-A

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Section: 23

CSE-421

Ans. to the Q. No. 1

I)  $255.255.128.0$

Given,

IPv<sub>4</sub> address is  $107.168.177.108$

Prefix Mask is  $17$ .

So,

Subnet mask in Binary,

$\begin{array}{|c|c|} \hline 1111 & 1111 \\ \hline \end{array} \cdot \begin{array}{|c|c|} \hline 1111 & 1111 \\ \hline \end{array} \cdot \begin{array}{|c|c|} \hline 1000 & 0000 \\ \hline \end{array} \cdot \begin{array}{|c|c|} \hline 0000 & 0000 \\ \hline \end{array}$   
 $\begin{array}{cc} 255 & 255 \\ 128 & 0 \end{array}$

In decimal,

$255.255.128.0$

II)  $107.168.128.0 / 17$

IPv<sub>4</sub> address is  $107.168.177.108$

Subnet Mask is  $255.255.128.0$

Prefix mask is  $17$

So, Network address is,

$107.168.128.0 / 17$

III)

	Total hosts	In $2^n$ format	no of host bit	Subnet Generation	Network Address
LAN-A	$2500+2$ $= 2502$	4096	12	16 (3rd Octate)	107.168.128.0/20
LAN-B	$1200+2$ $= 1202$	2048	11	8 (3rd Octate)	107.168.144.0/21
LAN-C	$1050+2$ $= 1052$	2048	11	8 (3rd Octate)	107.168.152.0/21
WAN-1	$2+2$ $= 4$	4	2	4 (4th Octate)	107.168.160.0/30
WAN-2	$2+2$ $= 4$	4	2	4 (4th Octate)	107.168.160.4/30

Ans. to the Q. No. 2

I) PAT table, unique public port numbers

The ISP router uses its PAT table to map the public destination port in the reply packet to the private ip address and port

which allows it to forward the reply to connect device.

II)

Device A and B use private addresses which are meant for internal network communication. Private addresses are not routable on the internet. The ISP router uses public IP address which is globally unique and routable. It allows multiple private devices to access the internet using PAT.

Ans. to the Q. No. 3

I)

10

Given,

Packet Size 7240 bytes

Header Size 40 bytes

Maximum transmission per fragment 800 bytes

$$\begin{aligned}\text{Packet Size without header} &= 7240 - 40 \\ &= 7200 \text{ bytes}\end{aligned}$$

Each fragment must include its own header.

$$\begin{aligned}\text{So, usable data per fragment} &= 800 - 40 \\ &= 760\end{aligned}$$

$$\begin{aligned}\text{So, Number of fragments} &= 7200 / 760 \\ &= 9.47 \\ &\approx 10\end{aligned}$$

II) 400 bytes

First 9 fragments carry full packets.

$$\text{So, } 9 \times 760 = 6840 \text{ bytes}$$

$$\begin{aligned}\text{Remaining Packets} &= 7200 - 6840 \\ &= 360 \text{ bytes}\end{aligned}$$

So, Last fragment size is  $360 + 40$  (header)

= 400 bytes

III) 665

A full fragment carries 760 bytes.

Before 8th fragment:

$$7 \times 760 = 5320 \text{ bytes}$$

$$\begin{aligned} \text{So, Fragment offset} &= 5320 / 8 \\ &= 665 \end{aligned}$$

IV)

The MF (More fragments) bit is zero for the last fragment because it indicates no more fragments follow, signalling the receiver that this is the final part of the original packet.

Ans. to the Q. No. 4

I) `2001:db8::1:0:0:100`

II) `ff02::1`

III) `2001:0:0:3c10::`

Ans. to the Q.No.5

I)

The problem is single Static route uses only one next hop. If that route fails R1 can not use other paths. For improvement back up static route should be added with higher administrative distance.

Ans. to the Q.No.6

I)

ARP Request uses broadcast. All devices

in the same broadcast domain receive it. Switch forwards ARP broadcast to all ports except the one it came from. Drop frame only if the destination is outside the broadcast domain or already seen. Routers do not forward broadcasts.

So, ARP will be received by A, B, C, D.

forwarded by S2, S1, S3  
drop by R2, C, D, B

II) S2 MAC table:

A	f2	60
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S3 MAC table

A	f2	60
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Ans. to the Q. No. 7



IH, Identifier, flags, fragment offset, Header Checksum, Options and Padding these fields added in IPV6 header that are not present in IPV4.

### Ans. to the Q. No. 8

Distance Vector Routing only knows about its neighbours. Routers periodically exchange routing table with neighbours. But no single router has complete network knowledge. As routing decisions are made locally, this protocol is a decentralized protocol.

Differences between Distance Vector Routing and Link State Routing:

1. DVR knows only neighbours, but LSR knows entire network topology.

. DVR updates periodically, LSR updates whenever changes occur.

3. DVR is slower but LSR is faster.

### Ans. to the Q. No. 9

I) DHCP Requests use broadcast to discover server. But Broadcast can't cross routers by default.

The issue is LAN2 is on a different network and no dhcp relay is configured on the router connecting LAN2 to R1.

Solution is configure dhcp relay.

R2(config)#interface (LAN2 facing interface)

R2(config-if)#ip helper-address (R1 DHCPip)

II) DHCP Request : Client to DHCP server

DHCP Pack : DHCP server to Client

### Ans. to the Q. No. 10

I) AF:CC:FE

II) Mac address is considered flat because it uniquely identifies a device globally without indicating its network location.

### Ans. to the Q. No. 11

Steps:

1. Check if target ip matches router interface
2. If yes, send ARP reply to requester and update ARP table.
3. If no, drop ARP Request.

Destination MAC in ARP request:

FF.FF.FF.FF.FF.FF.FF.FF