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Sec: 2L

I) 3.12.64.0/19

II) R2 LAN 3.12.64.0/22

SW LAN 3.12.68.0/23

R4 LAN 3.12.76.0/24

WAN Linus 3.12.71.0/30

III) Allocated size $2^{22} \cdot 2^{10} = 1024$ addresses

∴ Wasted, $1024 - 1003$

= 21

2) I) ~~no routing topo~~

from the figure, the devices and router that
are directly connected to the respective routers

are direct connection.

II) ip route 0.0.0.0.0.0.0 192.16.11.1

III) ip route 0.0.0.0.0.0.0.0 50/0/1 10

→ AD

iv) 40: The AD represents the trustworthiness of the route.

d: Metric, representing cost

v) A directly attached route is better because it avoids the recursive lookup process.

3) Total data: $4584 - 42 = 4542$ bytes

$$2D + 3C2 = 4542$$

896 → 14 → 14 Data per fragment

$$\rightarrow X = 42$$

$$\therefore D = 522.57 \text{ bytes}$$

$$\therefore D = 520$$

$$2) 8 \times 520 = 4160$$

∴ 4542 - 4160 = 382

If 9th segment is 362 bytes $2D + 3C2$ must be \leq Total data

$$\therefore X = 562$$

$$\therefore \text{Offset} = (4 \times 520) / 8 = 260$$

$$\text{III) } \frac{4562}{2520}$$

29.73

- 9 packets

4) I) Link state algorithms converge once LSPs

are flooded. Here, it takes 3 iterations

$R_4 \rightarrow R_1 \rightarrow R_L \rightarrow R_3$

II) R_1 will detect the loss if stops receiving

'Hello' packets from R_4 within the dead interval

or if the physical layer reports a carrier loss on the

interface.

5) I) Longest form: 2001:0db8:12af:0001:1000:0003

Type: Global unicast, used for generic internet communication.

II) Longest form: 0000:0000:0000:0000:0000:0000:0000:0000

Type: Unspecified address, used by a host as a source address before unique address assigned.

8) The client broadcasts the DHCP Reboot to inform all other DHCP servers that it has accepted an offer from a specific server, allowing those other servers to release the IP addresses they were holding for that client. When renewing, the packet is Unicast directly to the server that granted the lease, rather than broadcast.

- 9) i) The likely cause is NAT.
- ii) Rajib should implement Port forwarding.
- iii) i) Destination MAC for host X:
 - ⇒ Because Host Y is on a different subnet ($10 \cdot 0 \cdot 1 \cdot n$ vs $10 \cdot 0 \cdot 2 \cdot n$), host X sends this frame to its Default gateway (R1).
 - ⇒ Destination MAC: MAC of R1's interface F9
 - ⇒ Host X uses ARP to find MAC address.

II) Switches are 'self-learning' because they build a MAC address table by inspecting the Source MAC address of incoming frames on each port. In this scenario, when Host X sends a frame S1 records that MAC 11-11-11-11-11-11 is on port P1.

TAKE IT DOWN AND IT'S ON

and record host transmission blade digest (II)

X read not SAM matched (II)

forward frame S1 to X read second

switch X reads (e.g. 00-00-00 or 00-1-0-0)

(II) granting HostX right to work

PP arbitration LED to SAM (SAM matched)

switch SAM reads at STA locu X - go to (II)