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Sec: 22

1) I) 3.12.64.0/19

II) R2 LAN 3.12.64.0/22

SW LAN 3.12.68.0/23

R4 LAN 3.12.76.0/24

WAN Link 3.12.71.0/30

III) Allocated size $2^{22} = 2^{10} = 1024$ addresses

\therefore Wasted: $1024 - 1003$

$= 21$

2) I) ~~no~~ routing table

From the figure, the devices and routers that are directly connected to the respective routers are direct connection.

II) ip route 0.0.0.0 0.0.0.0 192.10.11.1

III) ip route 0.0.0.0 0.0.0.0 50/0/1 10
 $\rightarrow AD$

iv) 40: The AD represents the trustworthiness of the route.

d: Metric, representing cost

v) A directly attached route is better because it avoids the recursive lookup process.

3) i) Total data: $4534 - 42 = 4542$ bytes let

$$80 + 362 = 4542$$

$$= D = 522.576$$

if $D = 520$

$$2) 8 \times 520 = 4160$$

$$\therefore 4542 - 4160 = 382$$

if 9th segment is 362 bytes $80 + 362$ must be \leq Total data

$$\therefore x = 562$$

$$\text{II) Offset} = (4 \times 520) / 8 = 260$$

$$\text{III) } \begin{array}{r} 4562 \\ 2520 \\ \hline \end{array}$$

2973

∴ 9 packets

4) I) Link state algorithms converge once LSPs

are flooded. Here, it takes 3 iterations

$R4 \rightarrow R1 \rightarrow R2 \rightarrow R3$

II) R1 will detect the loss if stops receiving 'Hello' packets from R4 within the dead interval or if the physical layer reports a carrier loss on the interface.

5) I) Longest form: 2001:0db8:12a1:0001:0003
Type: Global unicast, used for generic internet communication.

II) Longest form: 0000:0000:0000:0000:0000:0000:0000:0000
Type: Unspecified address, used by a host as a source address before unique address assigned

8) The client broadcasts the DHCP Request to inform all other DHCP servers that it has accepted an offer from a specific server, allowing those other servers to release the IP addresses they were holding for that client. When renewing, the packet is Unicast directly to the server that granted the lease, rather than broadcast.

9) I) The likely cause is NAT.

II) Rajib should implement Port forwarding.

11) I) Destination MAC for host X:

⇒ Because Host Y is on a different subnet ($10.0.1.x$ vs $10.0.2.x$), Host X sends this frame to its Default gateway (R1).

⇒ Destination MAC: MAC of R1's interface F9

⇒ Host X uses ARP to find MAC address.

II) Switches are 'self-learning' because they build a MAC address table by inspecting the source MAC address of incoming frames on each port. In this scenario, when Host X sends a frame S1 records that MAC 11-11-11-11-11-11 is on port F1.

I) The first time a frame is received

II) The switch should inspect the destination MAC address

I) Destination MAC for host X

Because Host X is on a different subnet

(10.1.1.2 vs 10.1.1.1) Host X is on a different subnet

(10) Forwarding Hosts at a switch

FF:FF:FF:FF:FF:FF to MAC address of the interface F1

Host X uses ARP to get MAC address