

1. (a) (i) address range is $109.64.0.0 - 109.127.255.255$

The block size in the 2nd octet is 64, which corresponds to /10.

subnet mask: $255.192.0.0$ ($/10$)

Network address: $109.64.0.0$

(ii) Host size,

Host bits = $32 - 10 = 22$

Hosts = $2^{22} - 2 = 4194302$

(b)

LAN - 2000 Hosts

Needs $\geq 2000 \rightarrow /21$

subnet size = $2^{11} - 2 = 2046$

11 Host bits

prefix = $32 - 11 = 21$

mask = 255.255.248.0

Network : 109.64.0.0

For LAN - 1460 Hosts, $\rightarrow /21$

$$2^8 - 2 = 2046 \text{ hosts}$$

$$\text{prefix} = 32 - 11 = 21$$

Mask = 255.255.248.0

Network : 109.64.8.0

switch network

needs ~ 4-5 IPs $\rightarrow /29$ (usable)

Network : 109.64.16.0

WAN link

Needs 2 hosts $\rightarrow /30$

Network : 109.64.16.8

$$1S = 11 - 2S = 11 - 2 \times 2 = 7$$

Ques $2, \text{ Q8P8S}$ $7028 + 08 \text{ PSS} = \text{circuit total}$

(a) TTL is an 8-bit field in the IP header that limits how many routers a packet can pass through.

Each router decrements TTL by 1. When TTL reaches 0, the packet is discarded.

so, $\text{TTL} = 104$ means the packet reached computer with 104 hops remaining.

(b) (i) Data per MF=1 packet limit

$$2883 - 35 = 2848 \text{ bytes}$$

Data in 10 such packets $= 10 \times 2848 = 28480$

Data in last packet (MF=0)

$$985 - 35 = 950 \text{ bytes.}$$

$$\text{Total origin} = 28480 + 950 = \boxed{29430 \text{ bytes}}$$

(ii) Data per fragment = 2848 bytes

offset is measured in 8 byte units.

$$\text{offset of 2nd packet} = \frac{2848}{8} = \boxed{356}$$

(iii) if MF flag = 0, this indicates no more fragments follow, so it is the final fragment of the original data.

final fragment of the original data

$$gram, 188 = 88 - 88$$

$$08888 = 8888 \times 01 = \text{first byte of original data}$$

(0 = 7M) first byte of original data

$$\text{last byte} = 88 - 88$$

(C) Reason: PC1's DHCP Discover is a Broadcast packet, and routers do not forward broadcasts by default. Since, PC1 is in subnet 192.168.4.0 and DHCP server is not in a different subnet, the DHCP Discover stops at Router R2 and never reaches R3.

Solution: Configure DHCP relay on the router interface connected to PC1's network using an IP helper address pointing to R3's IP.

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-brugt i et netværk der har 50 til 60 noder (D)

(a)

v) The algorithm used in this topology
topologi er Link-state Shortest Path

Algoritmen (Dijkstrás Algoritmen),

(ii)

order	Destina	shortest path	cost
1	4	5 → 4	1
2	6	5 → 6	2
3	0	5 → 4 → 0	4
4	8	5 → 8	6
5	2	5 → 6 → 2	6

- (b) A router ~~sends~~ ^{detects} a Link state packet to a neighbor only when it detects a change in its Local Link state, such as:
- Link cost change
 - Link failure or recovery
 - New neighbor discovered.
- If no change happens, no LSP is sent, which reduces unnecessary traffic.

(c) Distance vector protocols do not use helo packets because neighbor existence and failure are detected implicitly through regular distance vector updates.

If updates stop arriving, the neighbor is

assumed to be unreachable, (4)

[6]

(a) A MAC address has two parts:

① OUI (first 24 bits) \rightarrow identifies the manufacturer

② NIC (last 24 bits) \rightarrow Uniquely identifies the device.

The LSB of first bytes indicates address type:

$\rightarrow 0 = \text{Unicast}$

$\rightarrow 1 = \text{Multicast}$

(b)

No, Host 0 is in a different network, and does not cross routers, Host A will only learn the MAC address of its default gateway (R1).

The ARP request is broadcast and received by host B, C, R1 and sw1.

→ Host C replies with its MAC address.

→ Host B and R1 ignore it.

→ sw1 floods the request within the LAN

front of flooding is bridge A front network

not known not bros 0AM 2d 2word 142 &

sw1 of flooding the bridge forward

(c) SW1 MAC table:

MAC add	port
A	Fa0/1
B	Fa0/2
D	Fa0/3

SW2:

MAC add	port
A	Fa0/1
D	Fa0/3

When host A sends a packet to host D

→ SW1 knows D's MAC and forward the frame using unicast to SW2

(d)

- sw2 does not know D's MAC, so it floods the frame. (d)
- sw2 learns A's MAC from the incoming frame and later learns D's MAC from the reply. (e)
- [4] (a) LAN C = 100.9.128.128 / 27
- Router C's IP on the Multi-access network is 112.191.63.3
- static router on R-B:
- IP route 100.9.128.255.255.224
112.191.63.3
- no broadcast fix as sw on bfrwd
and two bfrwds pfrwd B

(b) 08.0AM 20 world for each SW2 &

Primary static route on Route A with
AD = 5 using S1/1;

AD = 5 using S1/1;

IP route 100.9.128.255.255.224
S1/105 0ms

If the links fails, using a floating static route with higher AD:

IP route 100.9.128.255.255.255.224-112,

236.224.161.10

(c) The route S 21.1.64.0/26 via 112.225

is not recommended because it causes

recursive routing. It should be

changed to use an exit interface on

directly connected next hop.

ip route add 21.1.64.0 255.255.255.192 soft
via 21.1.64.1

broadcast in fileroutemt trobsans
2299bb0 teo iflum

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- (a) IPv6 do routers do not perform fragmentation, if a packet exceeds the MTU, the router drops the packet and sends an ICMPv6 "Packet too Big" message to the source. The source host performs fragmentation using the fragment Extension header. Router add extra information using IPv6 Extension headers instead of modifying the base header.

We see, address does not support broadcast.

(b) IPv6 does not support broadcast.

Broadcast functionality is achieved using multicast address.

(c) Stateful DHCPv6 does not require DAD.

Because of the DHCPv6 server

assigns unique IPv6 address, duplication

is not possible.

Given information matrix b6 network broadcast

To broadcast message

unicast send

Externally broadcast

b6