

i Network bits = 17 ; Host bits = 32-17
= 15 bits

$$\therefore \text{Total Host} = 2^{15} \times 2 \\ = 32768$$

ii Subnet mask = 255.255.128.0

3rd octet block size = 128 [0-127]

Network address = 32.99.0.0

Broadcast address = 32.99.127.255

iii

Network	Required host	Host bits	Subnet	Network address
LAN A	2000	11	121	32.99.0.0
LAN B	1024	10	122	32.99.8.0
LAN C	3	2	123	32.99.12.0
WAN 1	2	2	130	32.99.12.8
WAN 2	2	2	131	32.99.12.12

I To allow external users to initiate FTP access (port 21) to an internal server static NAT with port forwarding is required.

Public IP 203.0.113.10:21

Internal 10.0.0.50:21

II It is using NAT because: One public IP

- Multiple internal users borrowing simultaneously
- Port number used to differentiate sessions

$$3. \quad \text{III} \quad n = \frac{6397}{1624} = 3.99 \approx 4$$

$$\text{II, } 6397 (16293) = 1525$$

$$\text{IV} \quad \text{fragment offset of last fragment} = \frac{16293}{8} = \frac{6876}{8} = 809$$

[Correction 8 byte unit]

V Reassembly of destination -

uses source IP, destination IP, identification field.

- Fragments are labeled with fragment offset.

- MF flag tells when last fragment arrives

- Packet is reassembled only when all fragments arrive.

V Significance of MF Flag:-

MF = 1 → more fragment coming

MF = 0 → last fragment.

Q5 R₁, R₂ and ISP router will converge faster because they run a Link state Protocol.

- Routers build a complete topology map.
- Changes are flooded immediately.
- No Periodic full-table updates.

Distance Vector

- Track Neighbors implicitly
- Learn via:
 - Routing table updates
 - Incoming Interface
- No explicit neighbour table

Link state

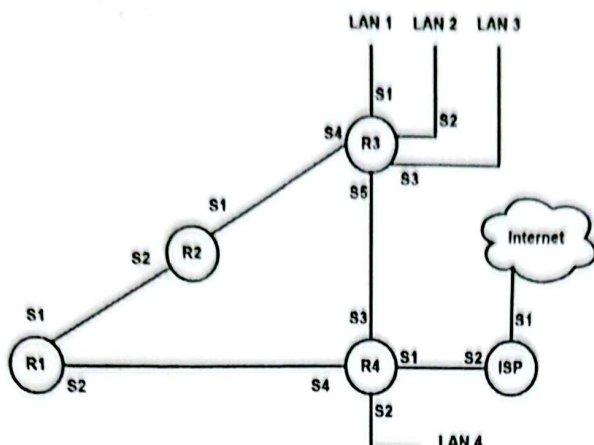
- Track neighbours explicitly
- Use:
 - Hello Packets
 - Neighbor table
 - Link state Database (LSDB)

Q7 I Source MAC = PC A's MAC

Destination MAC = FF:FF:FF:FF:FF:FF (Broadcast)

II

Q4



Device	Interface	IP	Network
R1	S1	.101	192.168.10.100/30
	S2	.225	192.168.11.224/30
R2	S1	.193	192.168.12.192/30
	S2	.102	192.168.10.100/30
R3	S1	.1	192.168.96.0/24
	S2	.1	192.168.64.0/25
	S3	.1	192.168.80.0/26
	S4	.194	192.168.12.192/30
	S5	.97	192.168.10.96/30
R4	S1	.1	192.168.9.0/30
	S2	.1	192.168.72.0/27
	S3	.98	192.168.10.96/30
	S4	.226	192.168.11.224/30
ISP	S1	.1	210.1.1.0/24
	S2	.2	192.168.9.0/30

Given the following topology where R1-R4 denote routers and the respective IP table of the topology.

- Configure a directly attached default static route in R2 with AD = 10.
- Configure a backup route of the above default static route using the next hop IP address.

Q5 Referring to the Q4's topology, R1, R2 and R3 are running Distance Vector protocol and R1, R4 and the ISP router are running Link State Protocol. Determine which routers will converge faster and why. Also, state which routers will keep track of their neighbors and how.

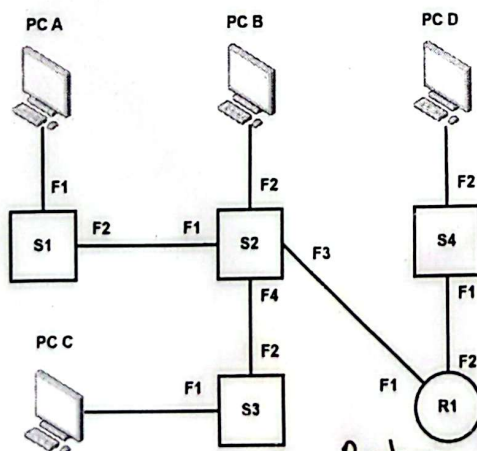
Q6 Write the shortened version of the following IPv6 addresses:

- 2001:0db8:85a3:0000:0000:8a2e:0370:7334 *2001:db8::8a2e:370:7334*
- 2607:0000:0000:0805:0000:0000:0000:200e *2607::0:0:805::200e*
- 3ffe:1900:4545:1003:1200:a0f8:fe21:67cf *3ffe:1900:4545::1200:a0f8:fe21:67cf*
(no compression possible)

END OF SECTION B

[CO2] SECTION C [Answer ANY THREE out of FIVE in this section] - 18 MARKS

Q7



Refer to the figure, PC A sends an ARP requests for PC D.

I. State the source and destination MAC addresses in the ARP request packet.

II. State what will router R1 do with the packet and why.

III. After receiving the ARP reply, state what is the first action made by PC A.

Router R1 will drop the ARP packet. Because: ARP is a Layer 2 Broadcast. Routers do not forward broadcast packets.

III PC A: Stores MAC-IP mapping in ARP cache. Sends the data frame to PC D.

I MAC-Source = PC A's MAC address

MAC-Destination = FF:FF:FF:FF:FF:FF (Broadcast)

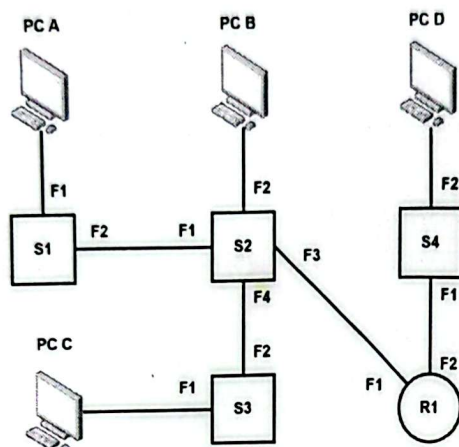
Q8 Consider a situation where someone is involved in malicious activities in our USIS system. The system administrator detected the source socket address associated with these actions. Despite this, they could not pinpoint the specific user due to the absence of end-to-end traceability. **Discuss and evaluate** the potential factors that could have contributed to the loss of this traceability. 6

Q9 In a DHCP-enabled network, each IP address is assigned to a device for a specific period, known as the lease time. Upon expiration of this lease, the device must renew its IP address to maintain network connectivity. 2 + 4

- I. List the steps involved in the IP address **renewal** process with DHCP.
- II. Additionally, **discuss** potential reasons a device might fail to renew its IP address. **Mention** which actions the device takes if it cannot renew the IP address.

Q10 Explain the function of the **loopback address** in network communications. Additionally, provide the loopback address used in IPv4 and IPv6. 4 + 2

Q11



Refer to the figure given, PC A sends a packet to PC C, at this stage all switch tables contain information about all devices shown. *No worries!!* 4 + 2

- I. State the actions that the switch S2 will take after it receives the packet.
- II. Explain, using the above scenario, why we call switches to be self learning.

END OF SECTION C

THE END

Why did the computer network go to therapy?

It had too many unresolved IP issues.

Q11 I Switch S2 Packet Handling:

- Reads Source MAC
- updates MAC table
- Looks up destination MAC
- Forwards packet via correct port (No flooding since table is complete)

- II • Learn MAC's from incoming frames
- No manual configuration needed
- Automatically updates table

Ab It is used to send and verify ~~the~~ test the proper functioning of the networking stack on a local device without sending packets to the external network. It allows a computer to send network packets to itself, which helps in troubleshooting diagnostics, and testing network applications. Packet sends to the loopback address never leave the host and are immediately routed.

back internally, IPv4: 127.0.0.1 IPv6::1