

① (a) ① prefix Mask: 12

281 = 11111111 · 11110000 · 00000000 · 00000000,  
255 240

∴ subnet mask = 255.240.0.0

IPv4 address = 200.100.50.25

(819 + 0 = 11001000 · 01100100 · 00110010 · 00011001

IP: 11001000 · 01100100 · 00110010 · 00011001

subnet: 11111111 · 11110000 · 00000000 · 00000000

AND: 11001000 · 01100000 · 00000000 · 00000000

Network Address = 200.96.0.0/12

① Host bit = (32-22) bits = 20 bits

∴ number of hosts =  $2^{20} - 2 = 1048574$

default-gateway = second last usable host

/12 network range:

Increment of 16 in 2nd octet

∴ network range: 200.96.0.0 → 200.111.255.255

Last usable host = ~~200.111~~ 200.111.255.254

second last usable host = 200.111.255.253

∴ Default gateway = 200.111.255.253/12

⑤

$$S3: 1024 + 2 = 1026 = 2^{11} = 8 \text{ (3rd octet)}$$

$$S3(LAN): 250 + 2 = 252 = 2^8 = 8 \text{ (3rd octet)}$$

$$S1: 100 + 2 = 102 = 2^7 = 128 \text{ (4th octet)}$$

$$R-R2: 2 + 2 = 4 = 2^2 = 4 \text{ (4th octet)}$$

$$R-R1: 2 + 2 + 2 = 6 = 2^3 = 8 \text{ (4th octet)}$$

$$R-R2: 2 + 2 = 4 = 2^2 = 4 \text{ (4th octet)}$$

0000100000000000  
8

0000001000000000  
1  
10000000  
00000100

$$S3: 200.06.0.0/21$$

$$S3(LAN): 200.06.8.0/24$$

$$S1: 200.06.0.0/25$$

$$R-R: 200.06.0.128/30$$

$$R-R: 200.06.0.136/29$$

$$R-R: 200.06.0.140/30$$



② (a) Data fragment =  $268 - 20$   
 $= 248$  bytes

Fragment offset of 7th fragment = 186

Data before last fragment =  $186 \times 8$   
 $= 1488$  bytes

Last fragment data = 248 bytes

Total data =  $(1488 + 248)$  bytes

$= 1736$  bytes

∴ Original packet size =  $1736 + 20$

$= 1756$  bytes

② Since all fragments have the same packet size, MTU = 268 bytes

Each fragment carries 248 bytes

offset increase per fragment =  $248 \div 8$   
 $= 31$

1st = 0

2nd =  $0 + 31 = 31$

3rd =  $0 + 31 + 31 = 62$

4th =  $62 + 31 = 93$

5th =  $93 + 31 = 124$

(11) MF=0, since 7th fragment is the last fragment

(6) (1) \* Wrong DHCP network statement: Here ~~the~~ written network 1.1.0.0 255.255.254.0 for R4 but R4 LAN is 1.1.0.0/24, which showing mismatch

\* Incorrect default gateway: Here written default-router 1.1.0.255, but .255 is used for broadcast address. Default-gateway should be a router interface IP. ~~for R4~~

\* DHCP is on another network: DHCP Server 2 is not in R4-LAN. DHCP uses broadcast but Router do not forward broadcast.



① Here DHCP requests from R4-LAN cannot reach DHCP server 2.

To solve it ip helper-address is needed on R4's LAN interface

interface f0  
ip helper-address DHCP-server2 IP

③ @ The topology uses Dijkstra's Shortest Path First algorithm. This is a Link-state routing algorithm used by protocol like

OSPF.

②

Node	Distance
A	0
B	<del>5</del>
C	<del>2</del>
D	<del>4+2=6</del>
E	<del>1</del>
F	<del>4</del>
G	$\infty$
H	$\infty$
I	<del>1+4=5</del>
J	$\infty$
K	$\infty$

∴ Path = A → C → F → B → I  
 cost = 2 + 4 + 5 + 5  
 16

⑤ Link state is better.

Reasons:

- ↳ Has full topology knowledge
- ↳ Uses Dijkstra's algorithm
- ↳ Faster convergence
- ↳ No count to infinity problem
- ↳ More accurate shortest path calculation



② The main cause is the Count-to-  
Infinity problem.

Explanation!

- \* Routers only know neighbor information
- \* Bad news travels slowly
- \* Routing loops occur before convergence
- \* Incremental metric updates cause delay.

④ ① ip route 0.0.0.0 0.0.0.0 1.1.1.2  
1.1.1.2 because Next hop from R4 is R1

① A sign higher AD. Let AD=10  
ip route 0.0.0.0 0.0.0.0 1.1.2.2 10

⑥ ① Network	usable Range
1.1.0.244/29	224 - 231
1.1.0.232/29	232 - 239
1.1.0.240/29	240 - 247
1.1.0.248/29	248 - 255

∴ full continuous range = 224 - 255 = 32 address

1/27 = 32 address

1/27 = 255, 255, 255, 224

static route command:

ip route 1.1.0.224 255.255.255.254 next hop

route summarization:

⑪ Effect of routing table size

- ↳ Reduces
- ↳ Faster Lookup
- ↳ Lower CPU and memory usage
- ↳ improves scalability



- ① Floating static default route  
AD = 50. because:  
↳ Normal static route AD = 1  
↳ AD increased to 50 to make it a backup route  
↳ Used only if a lower AD route fails.

- ⑤ a) IPv6 multicast address start with FF00::/8.  
Multicast replaces broadcast.  
IPv6 does not use broadcast.  
Uses multicast groups instead.  
Hosts join only relevant multicast groups, reduces unnecessary network traffic.

②

⑥ IPv6 uses Extension Headers which carries extra information.

It is inserted between IPv6 header and payload.

If 20 bytes are added!

\* Base IPv6 header remains

\* total packet size increases

\* Header size does not change

~~⑦ DAD is not required in stateful DHCPv6 because,~~

⑥ (a) MAC Address: A2-FE-22-3D-4F

A2 — 1010 0010

1st bit — 1

Locally administered  
Portability:

\* MAC Address move with the device  
\* works on any network interface  
\* Not tied to location



⑥ → ARP request from Host A for Host D

→ switch behavior:

\* Learn MAC of Host A on incoming port

\* ARP is a broadcast

\* Floods frame to all ports except source

→ Only Host A entry may be updated

No change for B, C, D