

Assignment - 2

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Section - 23 (CSE421)

(a)	First IP = 109.64.0.1 Last IP = 109.127.255.253	Network Address = 109.64.0.0/16 Broadcast Address = 109.127.255.255	(b)
			(0000) N/A (0011) N/A

∴ Network address = 109.64.0.0/16 Ans.

2nd octet range 64 → 127 Ans.

$$\therefore \text{block size} = 127 - 64 + 1 = 64$$

∴ Subnet = 255.192.0.0 (16) Net = 192 (b) L

② Total bits = 32 Ans.

Net Host = 10 Ans.

∴ Host bits = 32 - 10 = 22 Ans.

Total hosts = $2^{22} = 4194304$ Ans.

usable = $4194304 - 2 = 4194302$ Ans.

A- Subnetting

(b)

	Network address	Prefix mask	Usable host
LAN (2000)	109.64.0.0	/21	2046
LAN (1400)	109.64.8.0	/21	2046
Switch (4 routers)	109.64.16.0	/29	6
WAN Link	109.64.16.8	/30	2

$$P_2 = 1 + P_2 - FS = 0.5 \text{ (old)}$$

- 2] (a) TTL = 104 (01) denotes remaining Time to live (hop count) of the ICMP packet.

It means the packet has passed through several routers (24 hops if the initial TTL was 128).

TTL is used to prevent infinite routing loops by discarding packets when the value reaches zero.

$$\text{SOP} = \text{I-SOP} = 3/256$$

(uA)

(b) (i) Data per MF = 1 fragment transport + nof (111)

$$2883 - 35 = 2848 \text{ bytes. } O = \text{first 3M } \rightarrow$$

→ Data of last fragment

$$\text{no frag} 85 - 35 = 6 \times 950 \text{ byte}$$

∴ flushed pd after last fragment from last window
∴ original data size = $(10 \times 2848) + 950$
= 29430 bytes

(ii) offset field of the 2nd packet

offset unit = 8 bytes

first fragment data = 2848 bytes

$$\text{so } \text{initial offset} = \frac{2848}{18} = 356$$

(ps10.1.8.81.5.01 browser side note)

(iii) Last fragment \rightarrow I = 7M and last \rightarrow ①

↳ MF flag = 0. offset bytes = 28 - 8885

(c) Reason:

PC1's DHCP Discover is a broadcast packet and routers do not forward broadcasts by default. Since PC1 is in network 192.168.4.0/24 and the DHCP server (R3) is in a different subnet, the DHCP Discover is dropped at Router R2.

Solution: Set MF flag = 1 in transport part

Configure DHCP Relay (IP helper address) on the router interface connected to PC1's LAN (R2 interface toward 192.168.4.0/24).

3] If no modification at 92.1 then return A (d)

i. Algorithm: Dijkstra's Algorithm (Link-State Routing Algorithm)

Want to return shortest path to node 92.1 out ←

ii. Shortest path from Node 5 to first five nodes

triggers for node 92.1 out to (LTT) sys (excluding itself)

order	Node	shortest cost	shortest Path
1	4	infinity	5 → 4
2	6	2	5 → 6
3	0	4	5 → 4 → 0
4	2	6	5 → 6 → 2
5	1	8	5 → 4 → 0 → 1

most important to implement shortest path ←

• shortest path

(using message on the log off)

(b) A router sends its LSP to neighbour only if;
either (1) it has not received any LSP from the neighbour or
it has received an LSP with a higher sequence number than the stored one.

→ The LSP is new

(new)

→ The LSP has a higher sequence number than the stored one

one with the stored one

→ The age (TTL) of the LSP has not expired

old fib	new fib	9601	22610
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(c) In the distance vector algorithm;

J → E

S

J

L

→ Router do not need neighbor discovery

O → P → E

P

O

C

→ They periodically exchange routing tables

with directly connected neighbors;

→ Neighbor reachability is inferred from routing updates.

(Hello packets are unnecessary)

Q1 (a) Link : $100 \cdot 9 \cdot 128 \cdot 128/27$ SV9E (A) LC

Router C through multi-access network ($112 \cdot 191 \cdot 63 \cdot 0/29$)

Router C on that network = $112 \cdot 191 \cdot 63 \cdot 3$

∴ Static route on Router B

• feeding link

$100 \cdot 9 \cdot 128 \cdot 128 \quad 255 \cdot 255 \cdot 255 \cdot 224 \quad 112 \cdot 191 \cdot 63 \cdot 3$

(b) AD = 5

exit interface : 31/1

∴ ip route $100 \cdot 9 \cdot 128 \cdot 128 \quad 255 \cdot 255 \cdot 255 \cdot 224 \quad 31/1 \quad 5$

(c) Next hop only static route cause recursive

will lookup ; use fully specified static route.

5) (a) IPv6 routers drop some oversized packets

and notify the source; fragmentation is done by the sender using fragment extension headers.

→ if a packet larger than MTU, the router drops the packet.

→ The router sends an ICMPv6 "packet too big"

message to source

E = DA (d)

(b) IPv6 replaces broadcast with multicast.

since IPv6 doesn't have true broadcast it functions

achieved using multicast addresses like;

FF02::1 → All nodes on same local link

FF02::2 → All routers on same link

(c) ~~for stateful DA~~ DHCPv6 doesn't require DAD because because the server guarantees unique address.

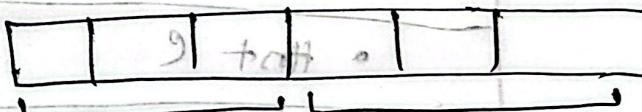
Reason:

- Router advertises a no reply to
- DHCP assigns addresses centrally
- Server ensures address uniqueness
- DAD is unnecessary.

Q 6] (a) Two part of MAC address.

1. OUI (organizationally unique identifier)

2. NIC → last 24 bits.



0 → unic平 (individual)

1 → Multicast (Group)

(b) No. (Host A will never learn the Mac of host C)

Reason;

- Host C is on a different network
- ARP works only within the local broadcast domain

ARP request sent by Host A for host C

The ARP request is a Layer-2 broadcast, so it is

received by, number 3AM to try out (A) 12

Device	Action
• Host B	Recognise IP → sends ARP reply
• Host C	Recognise IP → sends ARP reply
• Router R1 (fa0/1 interface)	IP doesn't match - Ignore

(Invisible) transmit → 0

(Answer) transmit → 1

(c) SW1 unicasts the frame since D is known, while SW2 floods it due to an empty MAC table; (~~No flooding at SW1~~)

After delivery, SW2 learns both A and D's MAC addresses.

Updated MAC Address Table:

SW1 - MAC Address Table

MAC address	port
Host A	Fa0/1
Host D	Fa0/3

(No change)

SW2 MAC address Table

MAC address	port
Host A	Fa0/1
Host D	Fa0/3