

# System Verification

## Tasks:

Note: all tasks execution time is calculated from the actual implemented tasks using GPIOs and the logic analyzer.

Task Name	Periodicity / Deadline (MS)	Execution Time (MS)
Button_1_Monitor	50	0.008
Button_2_Monitor	50	0.008
Periodic Transmitter	100	0.0096
UART Receiver	20	0.017
Load 1 Simulation	10	5
Load 2 Simulation	100	12

## Methods of Verification:

### 1- Using Analytical Method:

#### 1.1- System Hyper period:

It's the Least Common Multiple of all task's periods

$H = \text{LCM}(50, 50, 100, 20, 10, 100) = 100$

#### 1.2- CPU Load

➤  $U = (E_1 + E_2 + E_3 + E_4 + E_5 + E_6) / H$

➤ where E is the Execution time and H is the Hyper period.

➤  $U = (0.008*2 + 0.008*2 + 0.0096 + 0.017*5 + 5*10 + 12) / 100 = 0.621$   
(62.1%)

#### 1.3- System stimulability check using URM and Time Demand Analysis Techniques:

➤  $\sum_{i=1}^n \frac{c_i}{p_i} \leq n(2^{\frac{1}{n}} - 1)$

➤  $L.H.S = \sum_{i=1}^n \frac{c_i}{p_i} = \frac{0.008}{50} + \frac{0.008}{50} + \frac{0.0096}{100} + \frac{0.017}{20} + \frac{5}{10} + \frac{12}{100} = 0.621$

➤  $R.H.S = n(2^{\frac{1}{n}} - 1) = 0.753$

➤  $L.H.S \leq R.H.S$ , so, the system is schedulable.

## 2- Time Demand Analysis:

a- Sort the tasks making the highest priority at the first:

Task Name	Periodicity / Deadline (MS)	Execution Time (MS)
1- Load 1 Simulation	10	5
2- UART Receiver	20	0.017
3- Button_1_Monitor	50	0.008
4- Button_2_Monitor	50	0.008
5- Periodic Transmitter	100	0.0096
6- Load 2 Simulation	100	12

b- Choose the critical instant 0 then:

$$w_1(10) = 5 + 0 = 5 < \text{deadline} \quad w_2(20) = 0.017 + 5 * \frac{20}{10} = 5.017 < \text{deadline}$$

$$w_3(50) = 0.008 + 0.017 * \frac{50}{20} + 5 * \frac{50}{10} = 25.059 < \text{deadline}$$

$$w_4(50) = 0.008 + 0.008 * \frac{50}{50} + 0.017 * \frac{50}{20} + 5 * \frac{50}{10} = 25.067 < \text{deadline}$$

$$w_5(100) = 0.0096 + 0.008 * \frac{100}{50} + 0.008 * \frac{100}{50} + 0.017 * \frac{100}{20} + 5 * \frac{100}{10} = 50.1266 < \text{deadline}$$

$$w_6(100) = 12 + 0.0096 * \frac{100}{100} + 0.008 * \frac{100}{50} + 0.008 * \frac{100}{50} + 0.017 * \frac{100}{20} + 5 * \frac{100}{10} = 62.1266 < \text{deadline}$$

As all Tasks are less Than the deadline. So, the system is schedulable.

## 2- Using SIMSO offline simulator:

Used Scheduler: Fixed priority rate monotonic.

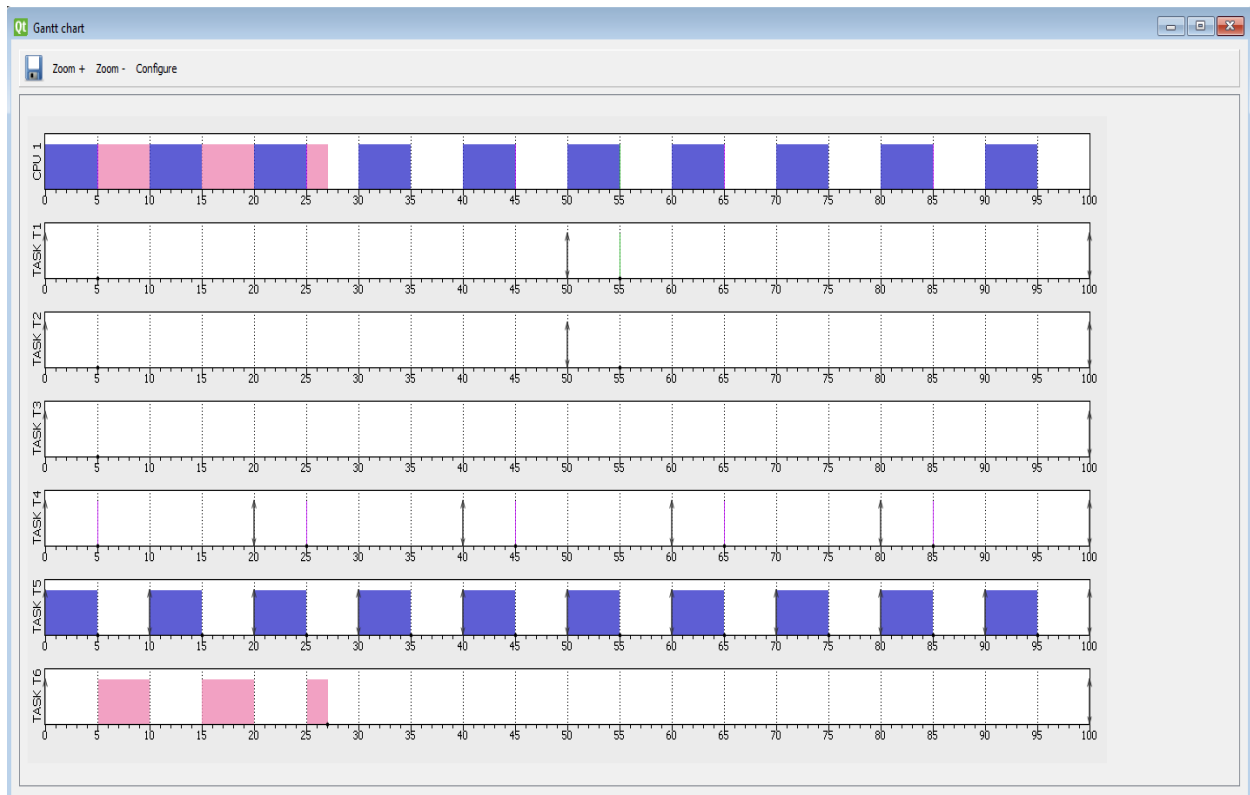
Tasks Simulated:

Qt Model data									
General		Scheduler	Processors	Tasks					
id	Name	Task type	Abort on miss	Act. Date (ms)	Period (ms)	List of Act. dates (ms)	Deadline (ms)	WCET (ms)	
1	TASK T1	Periodic	<input type="checkbox"/> No	0	50	-	50	0.008	
2	TASK T2	Periodic	<input type="checkbox"/> No	0	50	-	50	0.008	
3	TASK T3	Periodic	<input type="checkbox"/> No	0	100	-	100	0.0096	
4	TASK T4	Periodic	<input type="checkbox"/> No	0	20	-	20	0.017	
5	TASK T5	Periodic	<input type="checkbox"/> No	0	10	-	10	5	
6	TASK T6	Periodic	<input type="checkbox"/> No	0	100	-	100	12	

For The CPU load the is the same as the analytical mode

	Total load	Payload	System load
CPU 1	0.6213	0.6213	0.0000
Average	0.6213	0.6213	0.0000

Gantt chart over the Hyper period:



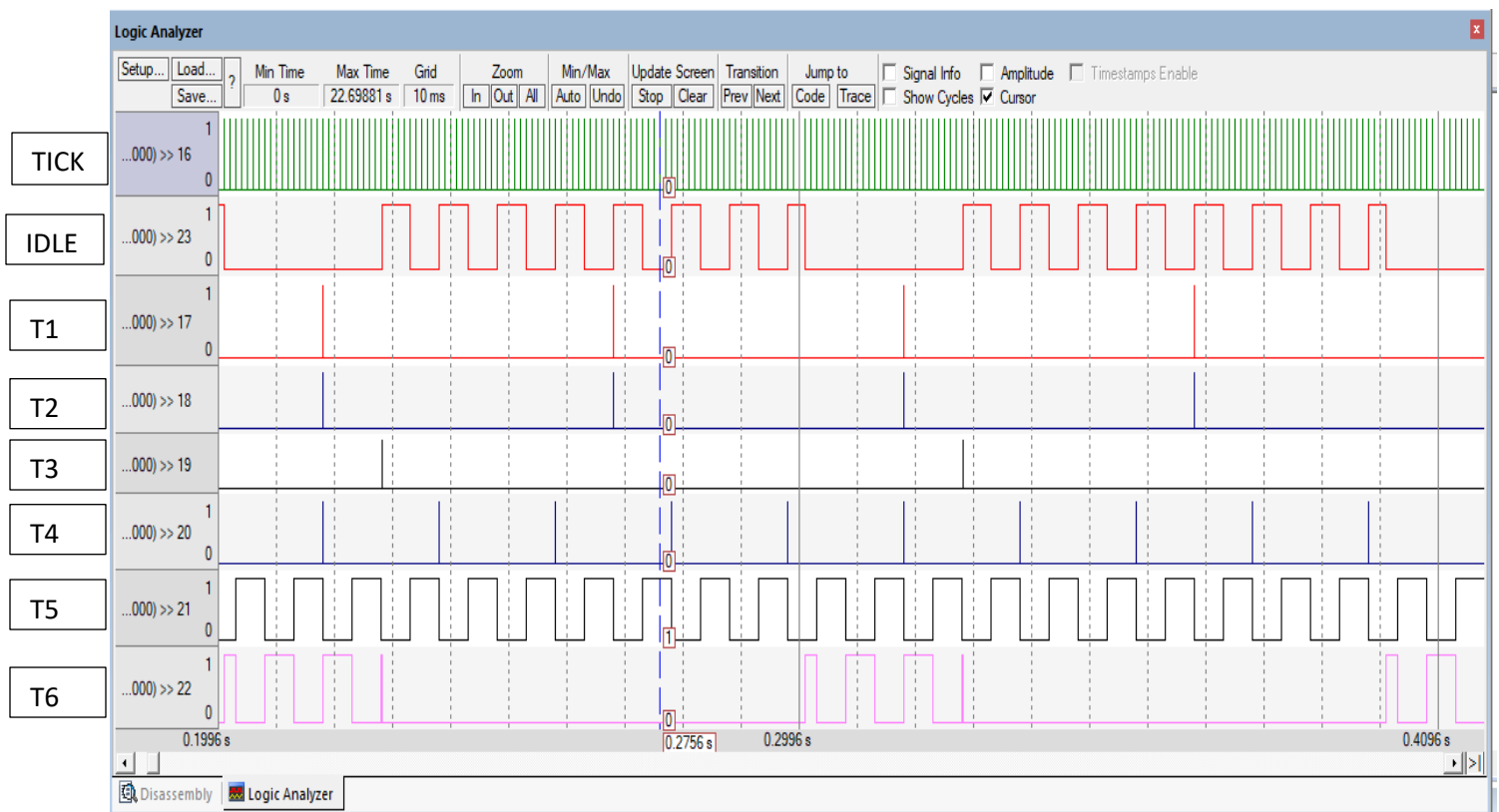
### 3- Using Keil Simulator at Runtime:

1- Calculate the CPU usage time using timer 1 and trace macros

Name	Value	Type
miss	<cannot evaluate>	uchar
deadline_misses[0]	0x00000000	int
deadline_misses[1]	0x00000000	int
deadline_misses[2]	0x00000000	int
deadline_misses[3]	0x00000000	int
deadline_misses[4]	0x00000000	int
deadline_misses[5]	0x00000000	int
total_sys_time	0x0007F0E4	int
cpu_load	63	int
task_1_total	0x000004C2	int
task_2_total	0x000004C9	int
task_3_total	0x00000482	int
task_4_total	0x00000628	int
task_5_total	0x00040013	int
task_6_total	0x0000F59F	int
<Enter expression>		

Note:

- 1- The CPU load is the same as the calculated analytically and the obtained using SIMSO offline simulator.
  - 2- None of The Tasks Miss the Deadline.
2. Using trace macros and GPIOs, plot the execution of all tasks, tick, and the idle task on the logic analyzer:



Note:

- The above chart was using the implemented EDF scheduler. As you can see tasks with closer deadline preempts other tasks.
- The IDLE task never interfered with my other main tasks so my modification in the IDLE task function to make sure it never preempts my main tasks is successful. As I made sure It always offsets the maximum main tasks deadline by a user defined amount So, it never executes unless there are no more tasks in the ready queue.