

Homework 1. Part 2**Due: Monday, January 15, 2018 before 8am EST.****Problem 1 [DPV] 6.2 – Hotel stops with minimum penalty.**

(a) Define the entries of your table in words. E.g., $T(i)$ or $T(i, j)$ is

Let $p(i)$ be the minimum penalty on a_1, a_2, \dots, a_i where a_i is included.

(b) State recurrence for entries of table in terms of smaller subproblems.

$$p(i) = \min_j \{p(j) + (200 - (a_i - a_j))^2\} \text{ where } j < i$$

(c) Write pseudocode for your algorithm to solve this problem.

```
Input : an array  $a$  of size  $N$   
Output: the minimum penalty by taking the optimal path of hotels  
initialize array  $p$  of size  $N$  to 0;  
for  $i \in 1 \rightarrow N$  do  
     $min\_j = i - 1$  ;  
    for  $j \in 0 \rightarrow i - 2$  do  
        if  $p[j] + (200 - (a[i] - a[j]))^2 < p[min\_j] + (200 - (a[i] - a[min\_j]))^2$  then  
             $min\_j = j$   
        end  
    end  
     $p[i] = p[min\_j] + (200 - (a[i] - a[min\_j]))^2$   
end  
return last element in array  $p$ 
```

(d) Analyze the running time of your algorithm.

Since there are two for loops (one looping through N items and the other going back to find the minimum value), it will run in $O(n^2)$

Problem 2 [DPV] 6.3 – Yuckdonald's

(a) Define the entries of your table in words. E.g., $T(i)$ or $T(i, j)$ is

Let $z(i)$ be the maximum profit in $\begin{cases} m_1, m_2, \dots, m_i \\ p_1, p_2, \dots, p_i \end{cases}$ where m_i and p_i are included.

(b) State recurrence for entries of table in terms of smaller subproblems.

$$z(i) = p_i + \max_j \{z_j : (m_i - m_j) < k\}$$

(c) Write pseudocode for your algorithm to solve this problem.

```

Input : an array  $m$  and array  $p$  of size  $N$  and minimum distance  $k$ 
Output: the maximum profit gathered from optimal store placement

initialize array  $z$  of size  $N$  to 0;
for  $i \in 1 \rightarrow N$  do
     $max\_j = 0$  ;
    for  $j \in 0 \rightarrow i$  do
        if  $p[j] > p[min\_j]$  then
             $max\_j = j$ 
        end
    end
    if  $max\_j \neq None$  and  $(m[i] - m[j]) \geq k$  then
         $z[i] = p[i] + p[max\_j]$ 
    else
         $z[i] = p[i]$ 
    end
end
 $max\_profit = 0$  for  $i \in 1 \rightarrow N$  do
    if  $z[i] > z[max\_profit]$  then
         $max\_profit = i$ 
    end
end
return  $z[max\_profit]$ 

```

(d) Analyze the running time of your algorithm.

Since there are two for loops (one looping through N items and the other going back to find the minimum value), it will run in $O(n^2)$

Problem 3 [DPV] 6.17 – Coin changing (unlimited supply of each denomination) Practice Only

(a) Define the entries of your table in words. E.g., $T(i)$ or $T(i, j)$ is

(b) State recurrence for entries of table in terms of smaller subproblems.

(c) Write pseudocode for your algorithm to solve this problem.

(d) Analyze the running time of your algorithm.

Problem 4 [DPV] 6.18 – Coin changing (use each denomination at most once) Practice Only

(a) Define the entries of your table in words. E.g., $T(i)$ or $T(i, j)$ is

(b) State recurrence for entries of table in terms of smaller subproblems.

(c) Write pseudocode for your algorithm to solve this problem.

(d) Analyze the running time of your algorithm.