Flash Storage & Security Design

1. Key Hierarchy & Root of Trust

- **DMS (Device Master Secret)**: one-time provisioned, never leaves device. Burned at provisioning.
- KEK (Key Encryption Key): derived from DMS + boot counter + device UID. Used to wrap DEKs.
- **DEK (Data Encryption Key)**: per-record. Derived via HKDF. Used to encrypt record payloads.
- **KEK_ODD**: session-limited KEK derived after PIN/UV success. Used for records that require user presence/verification.
- **Session tokens**: short-lived, kept in RAM only, never persisted.
- **PIN stretching**: Argon2id (software) → stretched key → input to KEK_ODD derivation.
- All derivations: HKDF-SHA256 (CC310-backed).

2. Record Format

Each record = append-only unit in data segment.

Header

- o Record UUID (random 128b)
- Sequential number (monotonic, prevents replay/reinsertion)
- o Length (payload size)
- Flags (alive/dead, type, etc.)

Body

Encrypted payload (AEAD: AES-256-GCM, 96b nonce, 128b tag)

Trailer

- CRC32 (fast fail detection, optional)
- AEAD authentication tag (mandatory integrity)

```
struct RecordHeader {
 u32 magic = 0x5458524B // "TXRK"
 u8 version = 1
 u8 type // 0=cred, 1=large-blob, 2=meta-note, ...
 u16 header_len
 u128 record_uuid
 u256 rpIdHash // SHA-256(RP ID), for lookup
 u32 seqno
                    // monotonically increasing per record replacement
  u32 created_ts
 u32 last_used_ts
 u16 aead_algo // 1=AES-GCM-256, 2=CH20P1305
 u16 payload_len // plaintext length before AEAD
 u32 crc_header // CRC-32 over header (excl. this field)
struct RecordBody {
 u96 nonce // AEAD nonce; ensure per-(DEK,seqno) uniqueness
     payload[payload_len] // e.g., FIDO2 credential source incl. pubkey
 u32 crc_physical // CRC-32 over full physical record for corruption triage
 u32 commit_mark // 0x8BADF00D when fully written (last word)
```

3. Log-Structured Layout (Flash Partitioning)

Fixed layout for 8 MB serial NOR flash (Append-only, no in-place updates):

- Superblock (metadata) \rightarrow 2 × 128 KB = 256 KB
- **Index region** → 256 KB
- **Data segment** \rightarrow 6 MB (\sim 5000 records @ 1.2 KB avg size)
- **Large blob area** → 1.5 MB (for CTAP2 largeBlob storage)

4. Indexing & Lookup

- Index maps UUID → physical location in data segment.
- Stored in fixed index region.
- Supports fast lookup & GC scan.
- Metadata confidentiality: index stores only hashes of user handles/PII.

5. Anti-Rollback & Replay Protection

- Boot counter (monotonic)
 - Increment on each successful boot.
 - o One-way increment \rightarrow stored in OTP or protected flash.
 - \circ Used in KEK derivation → rollback makes KEK mismatch → old data undecryptable.
- Superblock hash chain
 - Hash = SHA256(prev superblock || boot counter || layout metadata).
 - o Ensures continuity across reboots → tamper-evident.
- Record sequential number
 - o Each UUID increments seq# on update.
 - o Prevents replay of old record with same UUID.
 - Checked during lookup.

6. PIN / User Verification Binding & Rate Limiting

- After PIN/UV success: derive **KEK ODD**.
- Store only a short-lived **ODD session token in RAM**.
- Encrypt UV-bound records with DEKs derived from KEK_ODD → useless without fresh verification.
- Rate limit counters & back-off values
 - Stored as small meta-records in flash.
 - Increment on failed PIN attempts.
 - o Enforce exponential backoff (delay grows per failure).
 - \circ Written with same commit protocol \rightarrow atomic.

7. Garbage Collection

- Append-only writes → dead/stale records accumulate.
- GC process:
 - 1. Select segment for reclaim.
 - 2. Copy live records to new space.
 - 3. Erase segment.
- Trigger when free segments drop below threshold.
- Fragmentation handled automatically (copy compacts records).

8. Crypto Modes, Parameters, Randomness

- **AEAD**: AES-256-GCM (primary).
- **Nonces**: 96-bit, from TRNG (CC310).
- KDFs: HKDF-SHA256 (hardware), Argon2id (software, PIN only).
- Integrity: AEAD tags mandatory; CRC32 optional fast-fail.

9. Power Loss & Crash Consistency

- All mutating ops (writes, updates, deletes) = **two-phase commit**:
 - o **Prepare**: write new record(s) as append.
 - o **Publish**: atomically mark new record as live + old as dead.
- If crash before publish → old record still valid.
- If crash after publish → new record valid, old dead.
- Ensures no half-written corruption.

10. Boot & Provisioning Checklist

Step-by-step burn-in flow:

- 1. **Entropy check**: verify CC310 RNG working.
- 2. **Generate DMS**: 256b random, store in OTP/protected flash.
- 3. **Derive KEK** = HKDF(DMS || UID || boot counter).
- 4. **Provision root keys**: wrap and store securely.
- 5. **Initialize superblocks**: write SB0 + SB1 with clean layout + hash chain base.
- 6. **Initialize index region**: empty.
- 7. **Reserve GC space**: mark one free segment.
- 8. Lock debug interface.
- 9. **PIN provisioning**: store Argon2id-stretched key record.
- 10. **Initialize rate-limit counters** (0 attempts).
- 11. First-boot test: ensure KEK re-derivation, rollback check, GC run.

11. Performance & Wear

- Must always keep ≥1 free segment → GC breathing room.
- If data segment fills (6 MB used):
 - o GC runs to reclaim space.
 - o If no dead space left \rightarrow out-of-storage error \rightarrow reject new record.
- ~5000 records @ 1.2 KB average. Likely sufficient for FIDO2 credentials (tens to hundreds per user, not thousands).
- QSPI-XIP:
 - Quad-SPI Execute-in-Place → allows direct mapped read from flash → speeds lookups & index reads.
 - o No need to copy into RAM for reads.

12. Security Hardening Extras

- Metadata confidentiality: never store raw user handles/PII. Only hashed IDs.
- Side-channel hygiene:
 - Zeroize secrets and buffers.
 - o Use constant-time memcmp, memcpy.
 - o Avoid secret-dependent branching.
 - Keep sensitive ops inside CC310 hardware.
- Firmware-level hardening: disable debug, lock fuses, enforce secure boot.

13. Practical FIDO2 Fields Stored

- Resident keys (RKs).
- User handles (hashed only).
- LargeBlob extension data.
- PIN-stretched key material.
- Rate-limit counters.
- Superblock + index metadata.