

Quantum Computing and Cryptography

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- 1 Intro
- 2 RSA
- 3 Quantum computing
 - Introduction to the quantum world
 - Quantum algorithms
- 4 Post-Quantum cryptography
 - Intro to PQ cryptography
 - Lattice cryptography
 - Limits of PQ cryptography
- 5 Conclusion

Outline

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- Cryptography=TODO
- TODO: secret

Introduction

- Cryptography=TODO
 - TODO: secret
- Science of secret

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Classical bit

$$b \in \{0, 1\}$$

Classical bit

$$b \in \{0, 1\}$$

- 0

Classical bit

$b \in \{0, 1\}$

- 0
- 1

Classical bit

$$b \in \{0, 1\}$$

- 0
- 1

Quantum bit

$$|\psi\rangle \in \mathbb{C}^2$$

Classical bit

$$b \in \{0, 1\}$$

- 0
- 1

Quantum bit

$$|\psi\rangle \in \mathbb{C}^2$$

- $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$

Classical bit

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Quantum bit

$$|\psi\rangle \in \mathbb{C}^2$$

- $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$
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$$|\psi\rangle \in \mathbb{C}^2$$

- $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$
- $|1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$

- $|+\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix}$

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Quantum bit

$$|\psi\rangle \in \mathbb{C}^2$$

- $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$
- $|1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$

- $|+\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix}$
- $|-\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ -1 \end{bmatrix}$

- $\langle 0| \Rightarrow 0$ (100%)

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- $\langle 1| \Rightarrow 1$ (100%)

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- $\langle 1| \Rightarrow 1$ (100%)
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- $\langle 0 | \Rightarrow 0$ (100%)
- $\langle 1 | \Rightarrow 1$ (100%)
- $\langle + | \Rightarrow 0$ (50%), 1 (50%)
- $\langle - | \Rightarrow 0$ (50%), 1 (50%)

- Gate X

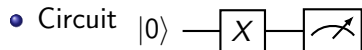
- Gate X
 - $X|0\rangle \rightarrow |1\rangle$

- Gate X

- $X|0\rangle \rightarrow |1\rangle$
- $X|1\rangle \rightarrow |0\rangle$

- Gate X

- $X|0\rangle \rightarrow |1\rangle$
- $X|1\rangle \rightarrow |0\rangle$



- Gate H

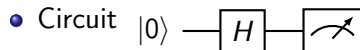
- Gate H
 - $H|0\rangle \rightarrow |+\rangle$

- Gate H

- $H|0\rangle \rightarrow |+\rangle$
- $H|1\rangle \rightarrow |-\rangle$

- Gate H

- $H|0\rangle \rightarrow |+\rangle$
- $H|1\rangle \rightarrow |-\rangle$



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Problème B.V

Given the oracle of a function f :

$$f : \{0, 1\}^n \rightarrow \{0, 1\} \quad f(x) = x \cdot s$$

Find s in the few request possible.

with $n = 2$ try :

- $f(10) = s_0$

2 requests.

with $n = 2$ try :

- $f(10) = s_0$
- $f(01) = s_1$

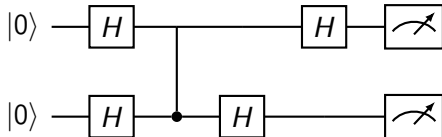
2 requests.

in general : $\mathcal{O}(n) \rightarrow$ Try every x that contains one bit to 1. At each query, we get the value of that bit in s

$\mathcal{O}(1) \rightarrow$ Just try every x at the same time.

Not only the x with only one bit at one but every possible x .

Algo Quantique - Slide 2



- Gain de complexité : $\mathcal{O}(e^b) \rightarrow \mathcal{O}(b)$

- Gain de complexité : $\mathcal{O}(e^b) \rightarrow \mathcal{O}(b)$
- combien de qubit il faut

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- combien de cubit on as

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What is PQ cryptography

- Based on (other) mathematical problems
- Considered unsolvable by a quantum computer

What it is not :

- Cryptography **using** quantum technologies

The problems

- Codes
- Hash functions
- Multivariate polynomials systems
- Isogenies
- Lattices

The problems

- Codes
- Hash functions
- Multivariate polynomials systems
- Isogenies
- **Lattices**

Why lattices ?

- Well spread
- Good results

Encryption/Key encapsulation	
Crystals-Kyber	Lattices
Signatures	
Crystals-Dilithium	Lattices
Falcon	Lattices
Sphincs+	Hash

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What is a lattice ?

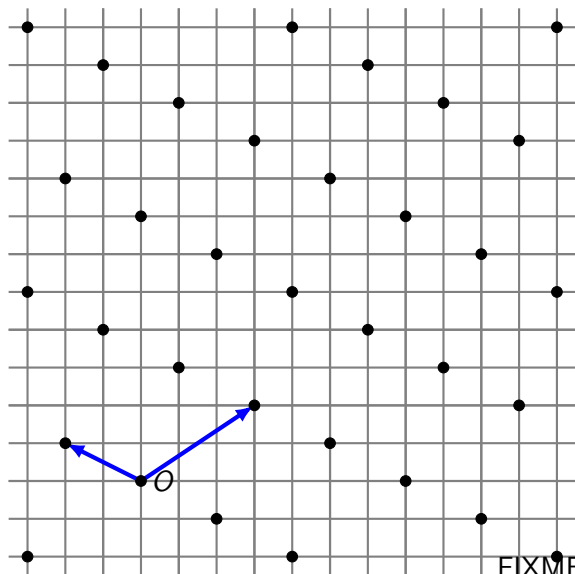
- A discrete subgroup of \mathbb{R}^n

abc akpqzsfkeokf akoczckdp kpzqkfs pkzapqfkkpkde pzkpkd czqks qkp
kfsdkvoesd, okpze kswkw k

Like vector spaces, we have :

- Vectors and matrices
- Linear combination

What is a lattice ? (cont'd)



FIXME : Basis:

TODO

(Fully) homomorphic encryption

TODO

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Size

TODO	
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Not necessarily robust to classical computer

- Example : Supersingular isogenies Diffie-Hellman key exchange

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