Quantum Computing and Cryptography

Damien, Théo, Matthieu

February 12, 2025

Outline

Outline

Introduction

- Cryptography=TODO
- TODO: secret

Introduction

- Cryptography=TODO
- TODO: secret

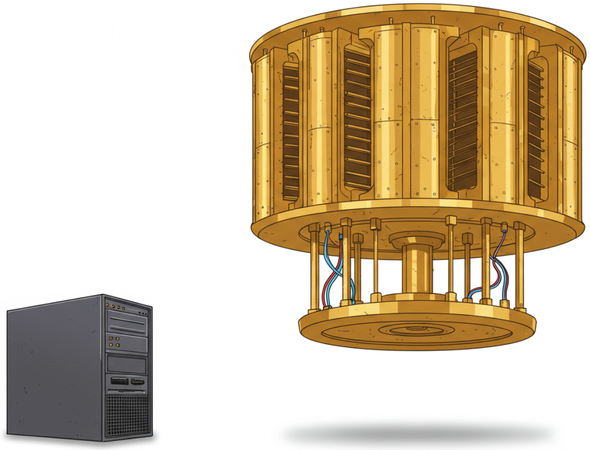
→ Science of secret

Outline

RSA

Outline

Quantum computing



Outline

Qubits

Classical bit

$b \in \{0, 1\}$

Qubits

Classical bit

$b \in \{0, 1\}$

- 0

Qubits

Classical bit

$b \in \{0, 1\}$

- 0
- 1

Qubits

Classical bit

$b \in \{0, 1\}$

- 0
- 1

Quantum bit

$|\psi\rangle \in \mathbb{C}^2$

Classical bit

$b \in \{0, 1\}$

- 0
- 1

Quantum bit

$|\psi\rangle \in \mathbb{C}^2$

- $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$

Qubits

Classical bit

$b \in \{0, 1\}$

- 0
- 1

Quantum bit

$|\psi\rangle \in \mathbb{C}^2$

- $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$
- $|1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$

Qubits

Classical bit

$b \in \{0, 1\}$

- 0
- 1

Quantum bit

$|\psi\rangle \in \mathbb{C}^2$

- $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$
- $|1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$
- $|+\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix}$

Qubits

Classical bit

$b \in \{0, 1\}$

- 0
- 1

Quantum bit

$|\psi\rangle \in \mathbb{C}^2$

- $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$
- $|1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$
- $|+\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix}$
- $|-\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ -1 \end{bmatrix}$

Measurements

Why measuring ?

We cannot read superposition. When we look at a qubit, it collapses to a classical bit.

Measurements

Why measuring ?

We cannot read superposition. When we look at a qubit, it collapses to a classical bit.

What do we get ?

We measure 0 or 1 with a probability that depends on the state of the qubit.

Measurements

Why measuring ?

We cannot read superposition. When we look at a qubit, it collapses to a classical bit.

What do we get ?

We measure 0 or 1 with a probability that depends on the state of the qubit.

- $|0\rangle \rightarrow 0$ (100%)

Measurements

Why measuring ?

We cannot read superposition. When we look at a qubit, it collapses to a classical bit.

What do we get ?

We measure 0 or 1 with a probability that depends on the state of the qubit.

- $|0\rangle \to 0$ (100%)
- $|1\rangle \to 1$ (100%)

Measurements

Why measuring ?

We cannot read superposition. When we look at a qubit, it collapses to a classical bit.

What do we get ?

We measure 0 or 1 with a probability that depends on the state of the qubit.

- $|0\rangle \rightarrow 0$ (100%)
- $|1\rangle \rightarrow 1$ (100%)
- $|+\rangle \rightarrow 0$ (50%), 1 (50%)

Measurements

Why measuring ?

We cannot read superposition. When we look at a qubit, it collapses to a classical bit.

What do we get ?

We measure 0 or 1 with a probability that depends on the state of the qubit.

- $|0\rangle \rightarrow 0$ (100%)
- $|1\rangle \rightarrow 1$ (100%)
- $|+\rangle \rightarrow 0$ (50%), 1 (50%)
- $|-\rangle \rightarrow 0$ (50%), 1 (50%)

NOT gate

X gate

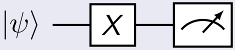
- $X\ket{0} \rightarrow \ket{1}$
- $X\ket{1} \rightarrow \ket{0}$

NOT gate

X gate

- $X|0\rangle \rightarrow |1\rangle$
- $X|1\rangle \rightarrow |0\rangle$

Circuit representation



Hadamard gate

H gate

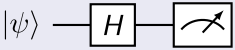
- $H |0\rangle \rightarrow |+\rangle$
- $H |1\rangle \rightarrow |-\rangle$

Hadamard gate

H gate

- $H\,|0\rangle \to |+\rangle$
- $H\,|1\rangle \to |-\rangle$

Circuit representation



Outline

Bernstein-Vazirani Problem

Problem Definition

Given an oracle for a function f:

$$f : \{0,1\}^n \rightarrow \{0,1\}$$

$$f(\mathrm{x}) = \mathrm{x} \cdot \mathrm{s}$$

where s is a secret bit string. Find s with the fewest oracle calls. ($\cdot$ is the bitwise dot product, XOR sum).

Classical Algorithm - Example

Example (n=2)

To find s $= s_0 s_1$:

Requires 2 queries.

Classical Algorithm - Example

Example (n=2)

To find s = $s_0 s_1$:

- Query $f(10)$. Result: s_0.

Requires 2 queries.

Classical Algorithm - Example

Example (n=2)

To find s $= s_0 s_1$:

- Query $f(10)$. Result: s_0.
- Query $f(01)$. Result: s_1.

Requires 2 queries.

Classical Algorithm - General Case

Complexity: $\mathcal{O}(n)$

We need to isolate each bit of s by querying with inputs that have a single '1'. This requires n queries for an n-bit string.

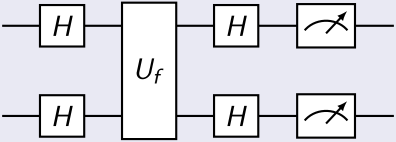
Quantum Algorithm - Overview

Complexity: $\mathcal{O}(1)$

The quantum algorithm can find s with just *one* query. It uses superposition to query all possible inputs simultaneously.

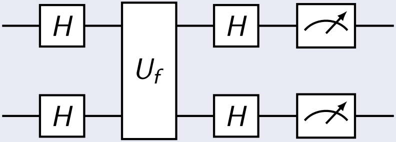
Quantum Algorithm - Circuit

Circuit Diagram



Quantum Algorithm - Circuit

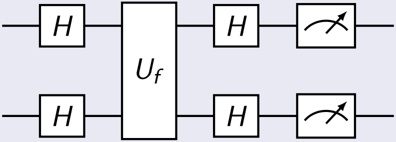
Circuit Diagram



Explanation

Quantum Algorithm - Circuit

Circuit Diagram

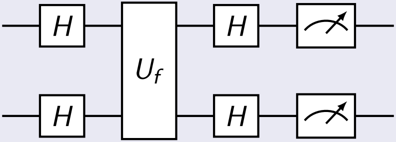


Explanation

- $H^{\otimes n}$: Hadamard gates on all n input qubits (creates superposition).

Quantum Algorithm - Circuit

Circuit Diagram

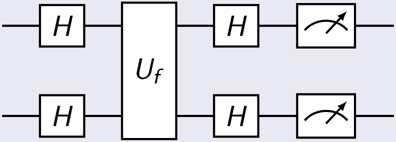


Explanation

- $H^{\otimes n}$: Hadamard gates on all n input qubits (creates superposition).
- U_f : The quantum oracle.

Quantum Algorithm - Circuit

Circuit Diagram



Explanation

- $H^{\otimes n}$: Hadamard gates on all n input qubits (creates superposition).
- U_f : The quantum oracle.
- Final Hadamards and measurement reveal s.

Shor's Algorithm

Complexity Gain

Classical factoring is very slow (roughly $\mathcal{O}(e^{\sqrt[3]{n}})$). Shor's algorithm is much faster (polynomial, $\mathcal{O}((\log N)^3)$).

Shor's Algorithm

Complexity Gain

Classical factoring is very slow (roughly $\mathcal{O}(e^{\sqrt[3]{n}})$). Shor's algorithm is much faster (polynomial, $\mathcal{O}((\log N)^3)$).

Requirements

Shor's Algorithm

Complexity Gain

Classical factoring is very slow (roughly $\mathcal{O}(e^{\sqrt[3]{n}})$). Shor's algorithm is much faster (polynomial, $\mathcal{O}((\log N)^3)$).

Requirements

- Requires a large number of high-quality (low-error) qubits (roughly $2n$ for an n-bit number).

Shor's Algorithm

Complexity Gain

Classical factoring is very slow (roughly $\mathcal{O}(e^{\sqrt[3]{n}})$). Shor's algorithm is much faster (polynomial, $\mathcal{O}((\log N)^3)$).

Requirements

- Requires a large number of high-quality (low-error) qubits (roughly $2n$ for an n-bit number).
- We currently don't have quantum computers large and stable enough to break practical RSA encryption.

Outline

Outline

What is PQ cryptography

- Based on (other) mathematical problems
- Considered unsolvable by a quantum computer

What it is not :

- Cryptography **using** quantum technologies

The problems

- Codes
- Hash functions
- Multivariates polynomials systems
- Isogenies
- Lattices

The problems

- Codes
- Hash functions
- Multivariates polynomials systems
- Isogenies
- Lattices

Why lattices ?

- Well spread
- Good results

Encryption/Key encapsulation	
Crystals-Kyber	Lattices
Signatures	
Crystals-Dilithium	Lattices
Falcon	Lattices
Sphincs+	Hash

Outline

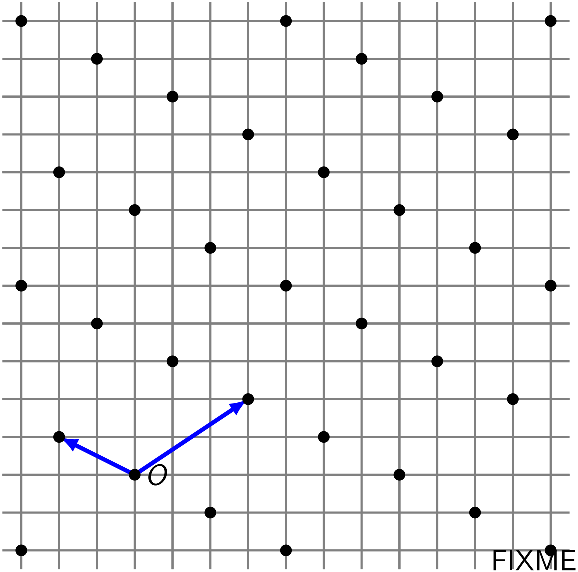
Some definitions

- A discret subgroup of RR^n

Like vector spaces, we have :

- Vectors and matrices
- Linear combination

Example



FIXME : Basis:

Learning with error problem

TODO

(Fully) homomorphic encryption

- Based on lattices (variant of LWE problem)

(Fully) homomorphic encryption

- Based on lattices (variant of LWE problem)
- We can evaluate a circuit (operations) on encrypted data
- Two operations : $+$ and $\cdot$, forms a ring

(Fully) homomorphic encryption

- Based on lattices (variant of LWE problem)
- We can evaluate a circuit (operations) on encrypted data
- Two operations : $+$ and $\cdot$, forms a ring
- We can evaluate (or compile) a function $f : P \times P \rightarrow P$ on encrypted data C_1 and C_2 :

$$eval(f, C_1, C_2) = Enc(f(Dec(C_1), Dec(C_2)))$$

(Fully) homomorphic encryption

- Based on lattices (variant of LWE problem)
- We can evaluate a circuit (operations) on encrypted data
- Two operations : $+$ and $\cdot$, forms a ring
- We can evaluate (or compile) a function $f : P \times P \rightarrow P$ on encrypted data C_1 and C_2 :

$$eval(f, C_1, C_2) = Enc(f(Dec(C_1), Dec(C_2)))$$

- Used to manipulate private data (e.g. Medical data, data science)

Outline

Sizes of the keys and data

TODO	

Not necessarly robust to classical computer

- Example : Supersingular isogenies Diffie-Hellman key exchange

Outline

Conclusion