# RW354 Principles of Computer Networking

- Larry L. Peterson and Bruce S. Davie. Computer Networks: A Systems Approach (Fifth Edition). Morgan Kaufmann Publishers. ISBN 1-55860-577-0.
- Behrouz A. Forouzan. Data Communications and Networking. McGraw Hill. ISBN 007-123241-9.
- Alberto Leon-Garcia and Indra Widjaja. Communication Networks: Fundamental Concepts and Key Architectures. McGraw Hill. ISBN 0-07-119848-2.

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# **Expectations**

#### This course IS about

- principles and concepts
- general-purpose computer networks
- the Internet perspective
- network software
- designing and building a system.

#### This course IS NOT about

- a survey of existing protocol standards
- specialized networks (CATV, telephone, . . . )
- the OSI perspective
- network hardware (we do survey)
- network performance using queuing theory models.



# **Perspective**

The expectations that you have of a network depends upon your persepctive

- Network users: network services that user applications need e.g. a guarantee that each message that the application sends will be delivered without error within a certain amount of time.
- Network designers: cost-effective design e.g. network resources are efficiently utilized and fairly allocated to different users.
- Network providers: a system that is easy to administer and manage e.g. faults can easily be isolated and it is easy to charge for usage.



# Characterizing Networks

Communication networks are divided into 2 basic types

- connection oriented: circuit switched
- connectionless: packet switched.

Packet switched networks fall into 3 classes

- WAN: a national/international network
- MAN: a network connecting several LANs
- LAN: a network connecting computers in a building or a campus.

Some networks are deliberately kept small in size.

The Internet is designed to grow to an arbitrarily large size (scale).



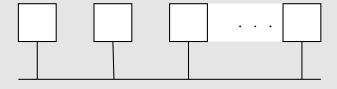
# Connectivity

- network building blocks
  - links: coaxial cable, optical fibre, satellite . . .
  - nodes: routers
- direct links
  - point-to-point: one link connects two nodes



point-to-point network

multiple access: many nodes share a link

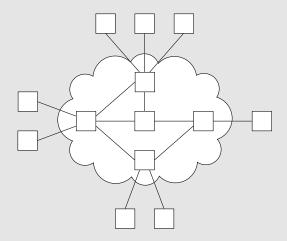


multiple access network

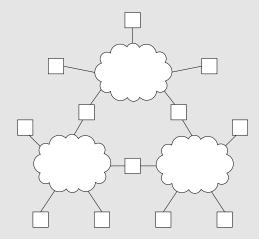


# Indirect connectivity

#### switched networks

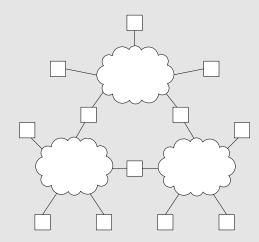


#### internetworks





# Indirect connectivity



- nodes (switches) inside the cloud store & forward packets
- nodes (hosts) outside the cloud support users & applications
- routers (gateways) are connected to two or more networks.



# Switching strategies

#### A network can be defined recursively as

- two or more nodes connected by a physical link, or
- two or more networks connected by one or more nodes.

#### Networks use two switching methods

- circuit switching: dedicated circuits are used to send/receive a bit stream
- packet switching: store-and-forward is used to send/receive messages (packets).



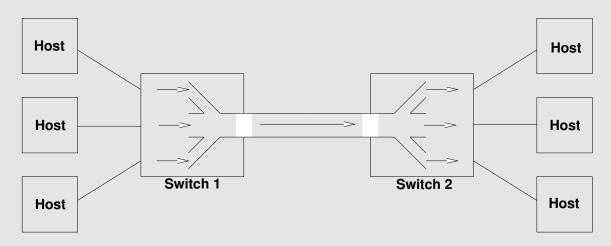
# Addressing & routing

- an address is a byte string that identifies a node; usually unique
- routing is the process of determining how to forward a message towards the destination node based on its address
- there are several types of addresses
  - unicast: node-specific
  - broadcast: all the nodes in the network
  - multicast: some subset of the nodes in the network.



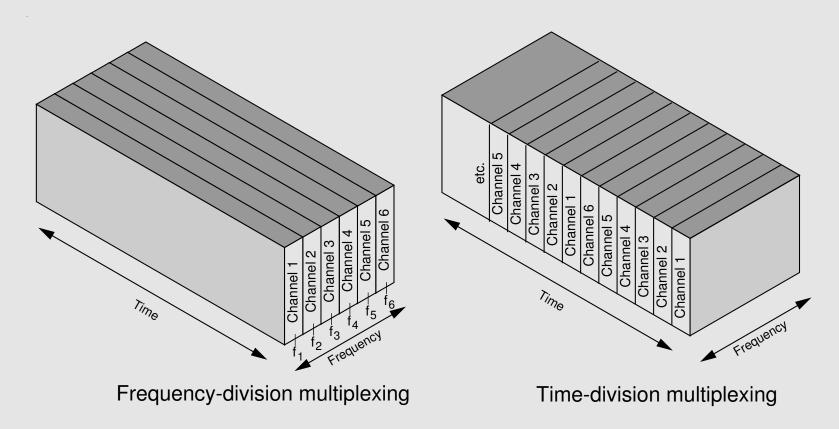
## Cost-effective resource sharing

Networks must share (multiplex) network resources (nodes & links) among multiple users.





# Common multiplexing strategies



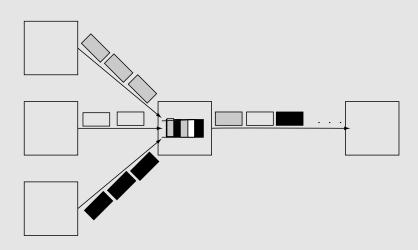
- synchronous time-division multiplexing (STDM)
- frequency-division multiplexing (FDM)

Both STDM and FDM are inefficient.



# Statistical multiplexing

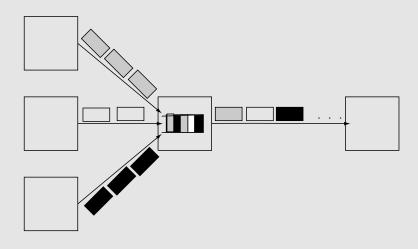
- TDM on demand rather than during a fixed time slot
- the link is rescheduled on a per-packet basis
- packets from different sources are interleaved on the link
- packets that contend for the link are buffered
- the packet queue is usually processed FIFO
- buffer overflow (dropped packets) is called congestion





# Statistical multiplexing

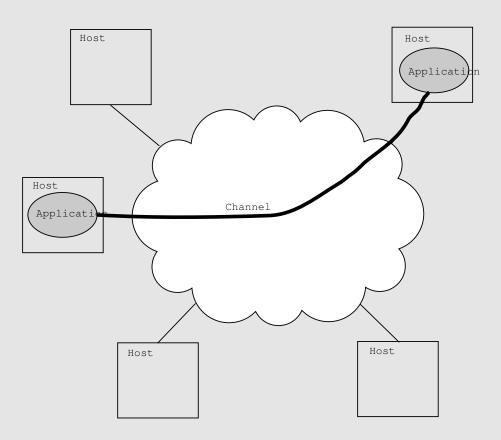
- the packet queue is usually processed FIFO, but not necessarily
  - packets from different flows are serviced in a round robin fashion
  - certain flows receive a certain portion of the link bandwidth
    - Quality of Service





# **Functionality**

Application programs running on hosts connected to the network must be able to communicate in a meaningful way.



The network provides common process-to-process channels. Each channel provides a set of communication services.



## Process-to-process channels

What functionality should the channels provide? Guaranteed/best effort delivery? Delivery in/out of order? Privacy? Constant/variable packet delivery rate?

- request/reply: for file access & digital libraries
- message stream: for video applications
  - video: sequence of frames
  - resolution: 1/4 TV-size image = 352×240 pixels
  - 24-bit color: frame = 352×240×24 / 8 = 247.5KB
  - frame rate: 30 fps = 7500KBps = 60Mbps
  - video on-demand vs video-conferencing



#### Network faults

What can go wrong in the network?

- bit errors, burst errors: rare, error correction
- packet-level errors: usually caused by congestion
- link & node failures
- messages are delayed
- messages are delivered out-of-order
- third parties eavesdrop.

The key problem is to fill in the gap between what applications expect & what the underlying technology provides.



# Performance: bandwidth (throughput)

- the amount of data that can be transmitted per time unit, for example 10Mbps
- link versus end-to-end performance
- notation:  $KB = 2^{10}$  bytes,  $Mbps = 10^{6}$  bits per second
- bandwidth is related to "bit width"

```
1Mbps
(each bit 1 microseconds wide)

2 Mbps
(each bit 0.5 microseconds wide)
```



# Performance: latency (delay)

- latency: the time it takes to send a message from point A to point B
- the round-trip time (RTT): from A to B & back
- the components of latency
  - latency = propagation + transmission + queue
  - propagation = distance / C
  - transmit = size / bandwidth
- the speed of light C
  - $3.0 \times 10^8$  meters/second in a vacuum
  - $2.3 \times 10^8$  meters/second in a cable
  - $2.0 \times 10^8$  meters/second in a fiber



# Performance: latency (delay)

#### **Notes**

- no queuing delays in direct link
- bandwidth is not relevant for the performance of small transfers
- bandwidth is relevant for the performance of large transfers
- process-to-process latency includes software overhead
- software overhead can dominate when distance is small



# Delay-bandwidth product

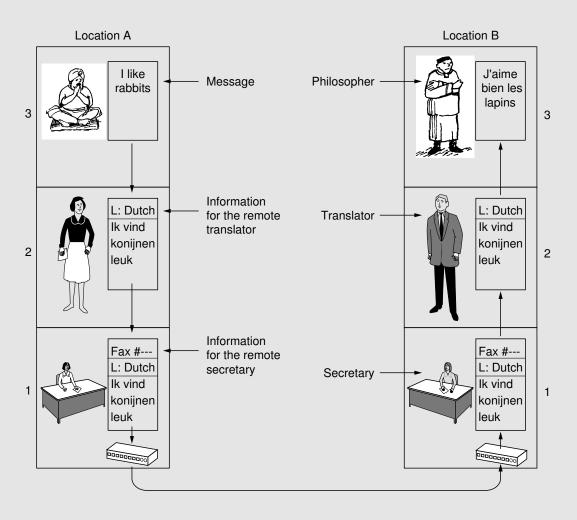
- the relative importance of bandwidth & latency
  - small message (e.g. 1 byte): 1ms vs 100ms RTT dominates 1Mbps vs 100Mbps bandwidth
  - large message (e.g. 25 MB): 1Mbps vs 100Mbps bandwidth dominates 1ms vs 100ms RTT
- delay-bandwidth product: 100ms delay & 45Mbps bandwidth = 560 KB of data in the pipe



- application needs
  - bandwidth requirements: burst size vs peak rate
  - jitter: variance in latency (inter-packet gap)



## Network architecture: Layering



A protocol stack is a list of protocols used by a system, one protocol per layer.



# Network architecture: Layering

- abstractions are used to hide complexity
- abstraction naturally leads to layering

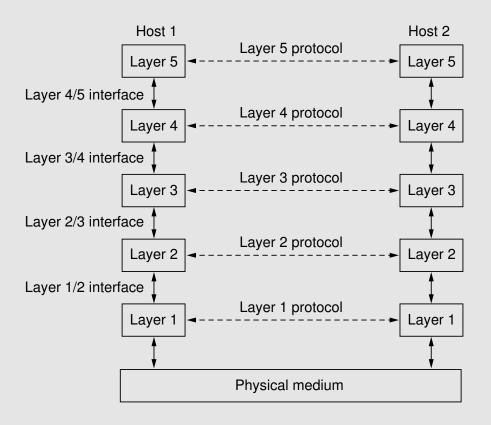
Application Programs		
Process-to-Process Channels		
Host-to-Host Connectivity		
Hardware		

alternative abstractions can be present at each layer

Application Programs		
Request/Reply	Message Stream	
Channel	Channel	
Host-to-Host Connectivity		
Hardware		



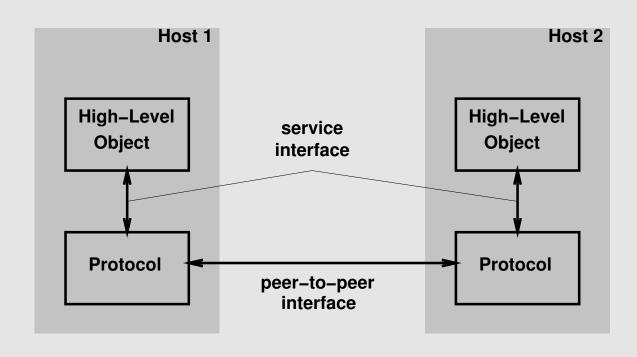
# Network architecture: Layering



- the corresponding layers on different machines are called peers
- an interface is present between each pair of adjacent layers.



- building blocks of a network architecture
- each protocol object has two different interfaces
  - the service interface: defines operations on this protocol
  - the peer-to-peer interface: defines messages exchanged with peer.





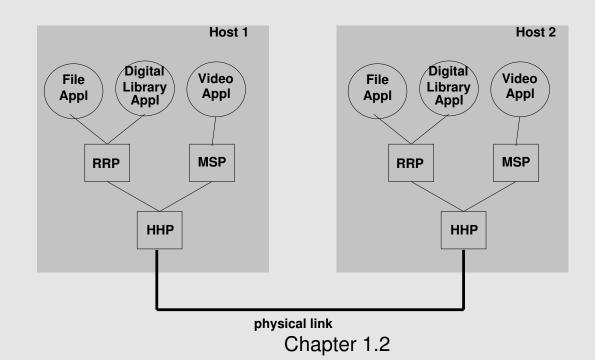
The term "protocol" is overloaded. It implies both

- the specification of the peer-to-peer interface
  - textual, psuedo-code, state transition diagrams, pictures of packet formats
- the module that implements this interface.



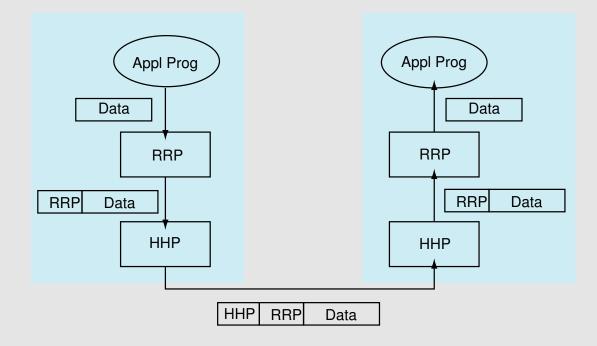
A protocol graph denotes a collection of protocols & their dependencies

- nodes correspond to protocols
- edges correspond to dependencies
- most peer-to-peer communication is indirect
- peer-to-peer communication is direct only at hardware level.





- multiplexing & demultiplexing: the demux key identifies the originating application
- encapsulation (header/body)



The nodes in the network can inspect the HHP header – the payload is not inspected.



#### Standard architecture

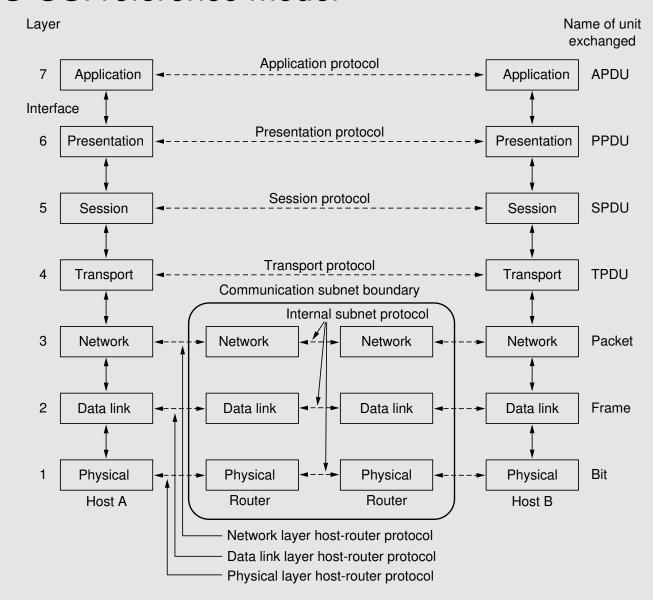
#### Not the first network architecture

- International Standards Organization (ISO)
- Open Systems Interconnect (OSI) Architecture
- International Telecommunications Union (ITU); formerly CCITT
- "X dot" series: X.25, X.400, X.500



#### Standard architecture

#### The ISO OSI reference model





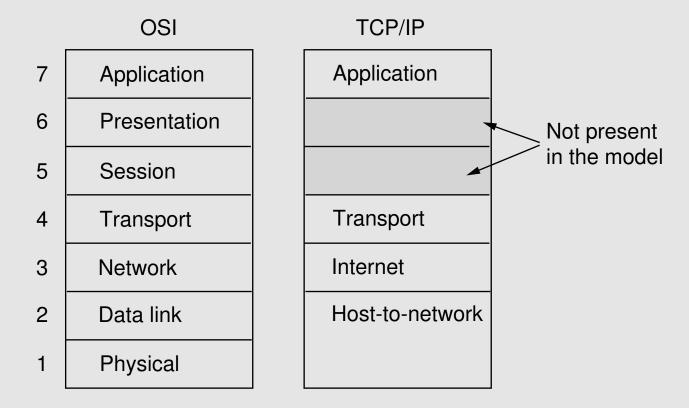
#### Standard architecture

- the physical layer: the transmission of bits on the physical link
- the data link layer: correct transmission of a frame from one node to the next node
- network layer: correct transmission of a packet from source to destination
- the transport layer: correct transmission of a message from source to destination
- the session layer: manages different transport streams that are part of a single application
- the presentation layer: the format of the data exchanged between peers
- the application layer: the application



### Internet architecture

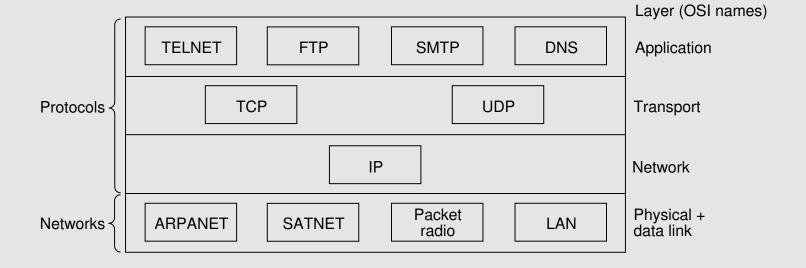
#### The Internet has a 4-layer model





### Internet architecture

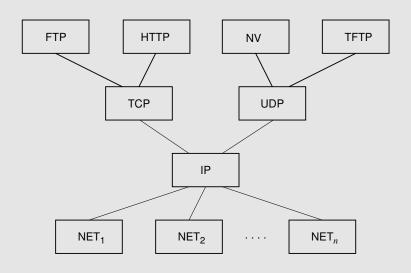
#### The Internet has a 4-layer model





#### Internet architecture

The process of defining the Internet architecture is controlled by the Internet Engineering Task Force (IETF).



- Application vs Application Protocol (FTP, HTTP)
- Features
  - does not imply strict layering
  - hourglass shape IP is the focal point
  - design & implementation go hand-in-hand.

