

Certificates | IPsec

INGI2347: COMPUTER SYSTEM SECURITY (Spring 2014)

Marco Canini

UCL
Université
catholique
de Louvain



Plan for today

Lecture 8

■ Certificates



- Working with certificates

■ VPN

■ IPsec

- Security Association (SA)
- Authentication Header (AH)
- Encapsulated Security Payload (ESP)
- Transport and Tunnel Modes
- Internet Key Exchange (IKE)



Certificates





What is a certificate?

Certificate's goal is to **link** a public key (PK) with its owner

- The pair (PK, owner) is signed by a trusted party (TP)
- The TP is named **Certification Authority** (CA)
- To check the signature, the CA's PK is needed
 - **Root certificate**: the pair (CA's PK, CA) is self-signed
 - The authenticity of the root certificate is fundamental (included in browsers)



Illustration

Lorem ipsum dolor
sit amet, consectetur
adipiscing elit. Fusce
vitae risus ultricies,
dapibus mi ultricies,
suscipit facilisis

**Signature
by Alice**

Certificate

(Alice's PK,
Alice)

**Signature
by TP**

Root Certificate

(TP's PK, TP)

**Signature
by TP**



X.509 Certificates

X.509

- Standard from International Telecommunication Union (ITU), 1988
- Also IETF RFC-2459 (and updates)

Three required fields:

- TBS Certificate (TBS = "To Be Signed")
 - The useful payload of the certificate
- CA signature algorithm
 - Identifier for the crypto algorithm used by the CA to sign this certificate
- CA signature value
 - Signature of the certificate by the CA



X.509 TBS

■ Serial number

- Unique number assigned by the CA to the certificate

■ Issuer field

- Identifies the entity who has signed and issued the certificate

■ Subject

- Identifies the entity associated with the public key
 - O: organization, C: country, OU: organization unit, CN: common name, ST: state, L: city, etc. no IP address



X.509 TBS (Continued)

■ Validity

- Not before
- Not after

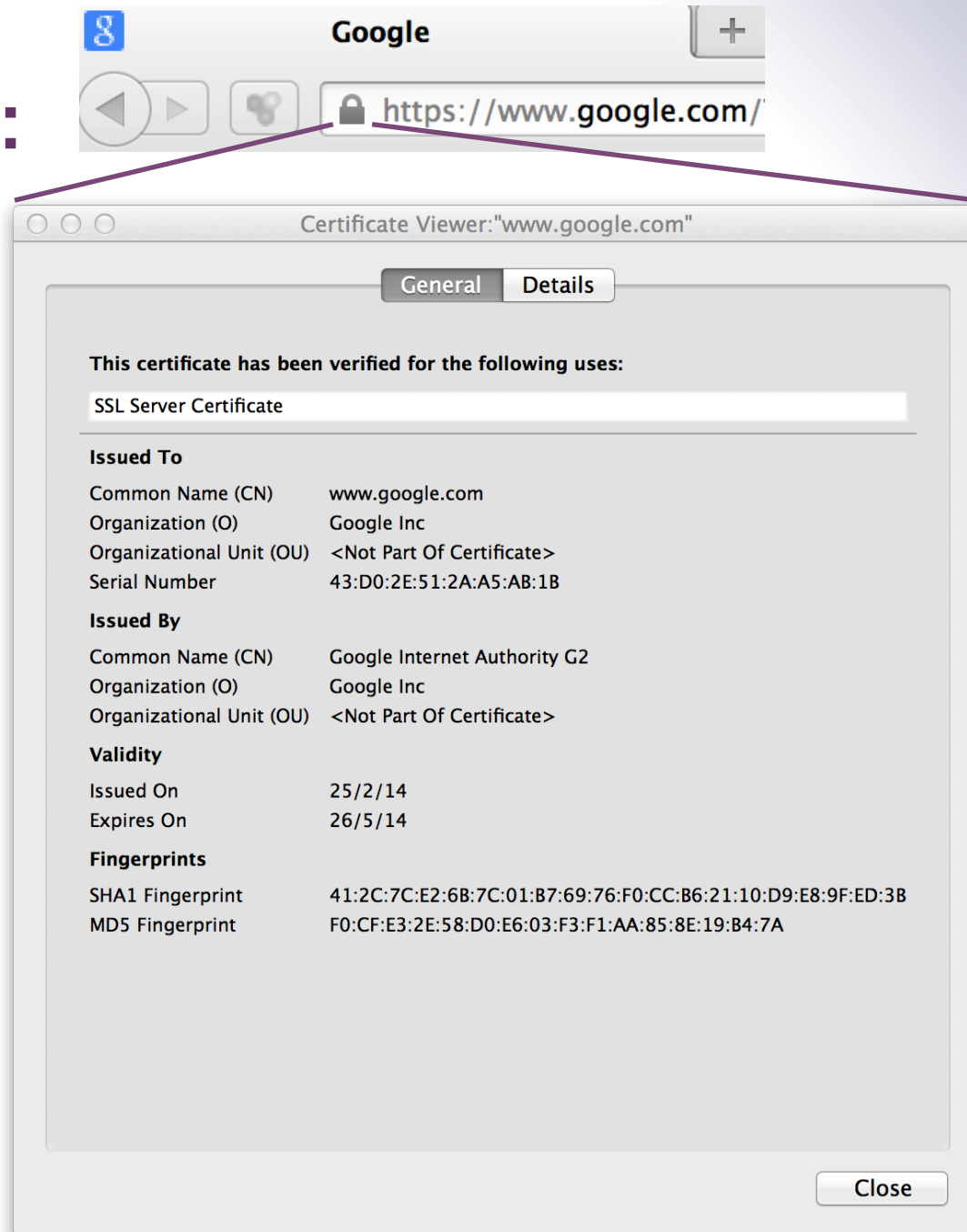
■ Subject Public Key Info

- Public key
- Identifies the algorithm with which the key is used
 - e.g., RSA, DSA, or DH

■ Etc.

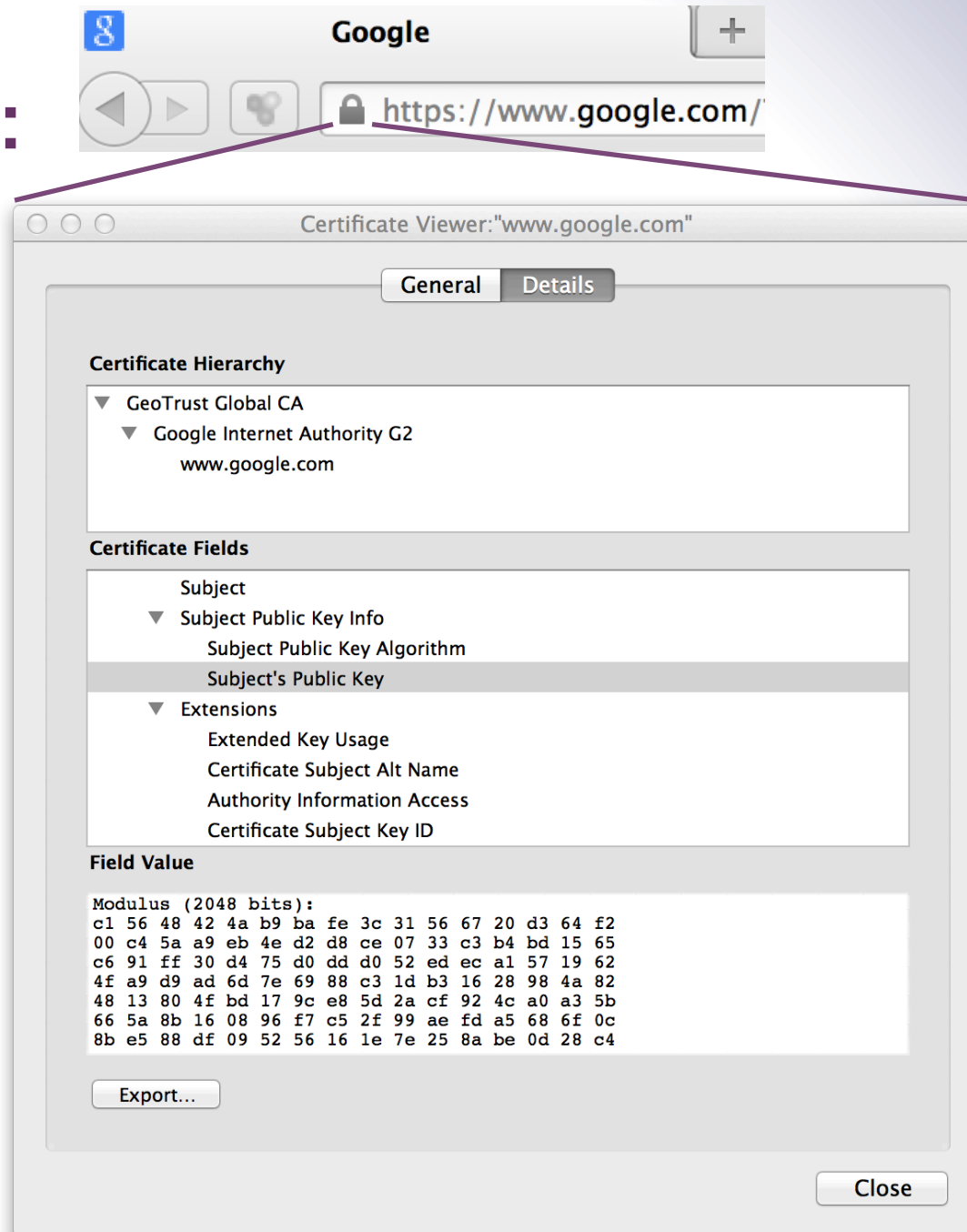


Example:





Example:





Working with Certificates

Certificate Authorities

■ Issuers of certificates found on web servers

CA	Count [%]
GeoTrust	25.19
GoDaddy.com	13.65
Verisign	13.09
Thawte	9.79
Comodo Limited	7.12
Unknown	2.40
DigiCert	2.39
Network Solutions LLC	2.09
Comodo CA Limited	1.77
GlobalSign	1.64

NOTE: GeoTrust, Verisign, and Thawte are the same group

Source: https://secure1.securityspace.com/es/s_survey/data/man.201002/casurvey.html (Feb 2010)

How to obtain a certificate

- Applicant registers with a CA
- CA (physically) authenticates the applicant
- CA asks applicant to generate public/private keys
- CA creates a certificate with the applicant's identity, PK, expiration date, etc., and the CA's signature
- CA provides a copy of its own PK to applicant

Registration Authority (RA)

- CA can delegate the registration of an applicant to the **registration authority** (RA)
- RA does not have CA's private key
- CA trusts the RA to authenticate the applicants
- After applicant is authenticated, applicant generates a pair of keys and sends the public key to the CA to create the certificate
- Technically RA sends a signed Certificate Signing Request (CSR) to the CA



CSR in practice

- Generate a 1024-RSA key-pair
 - `openssl genrsa 1024 > mykey.key`
- Generate a CSR
 - `openssl req -new -key mykey.key -out myreq.csr`
- Verify a CSR
 - `openssl req [-text] [-noout] -verify -in myreq.csr`
- Online checkers
 - <http://support.ecenica.com/ssl-certificates/csr-checker/>
 - <https://ssl-tools.verisign.com/checker/>

Certificate without CA

- Everyone can self-sign a certificate
- Distribute the certificate through an authenticated channel
- Makes sense in enterprise intranet
- Not really for public-facing services
- Rather get a free certificate...





Certificates in practice

■ Generate a certificate

- `openssl x509 -req -in myreq.csr -signkey mykey.key -out mycert.crt`

■ View a certificate

- `openssl x509 -text -in mycert.crt`

■ Verify a certificate

- `openssl verify mycert.crt`

Key escrowing

Keys are held in escrow so that, under certain circumstances, an authorized third party may gain access to those keys

Example:

- A company can provide two key pairs and certificates to each of its employees
 - One for signing | One for encrypting
- CA escrows a copy of the private encryption key
- Only employees can sign, but company can decrypt

Verifying a certificate

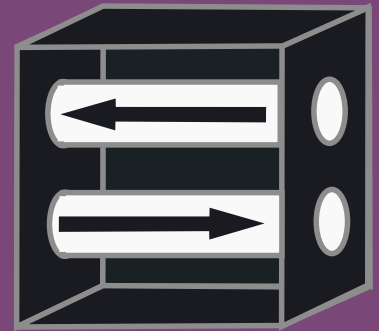
1. Verify the **certification path**

- Performed locally
- Delegated to a server: SCVP
 - Server-based Certificate Validation Protocol

2. Verify the **validity period**

3. Verify that the certificate is **not revoked**

- Performed locally: CRL (certificate revocation lists)
- Delegated to a server: OCSP
 - Online Certificate Status Protocol
 - Supported by all major browsers (enabled by default in Firefox and Safari)



VPN

Virtual Private Network



Virtual Private Network (VPN)

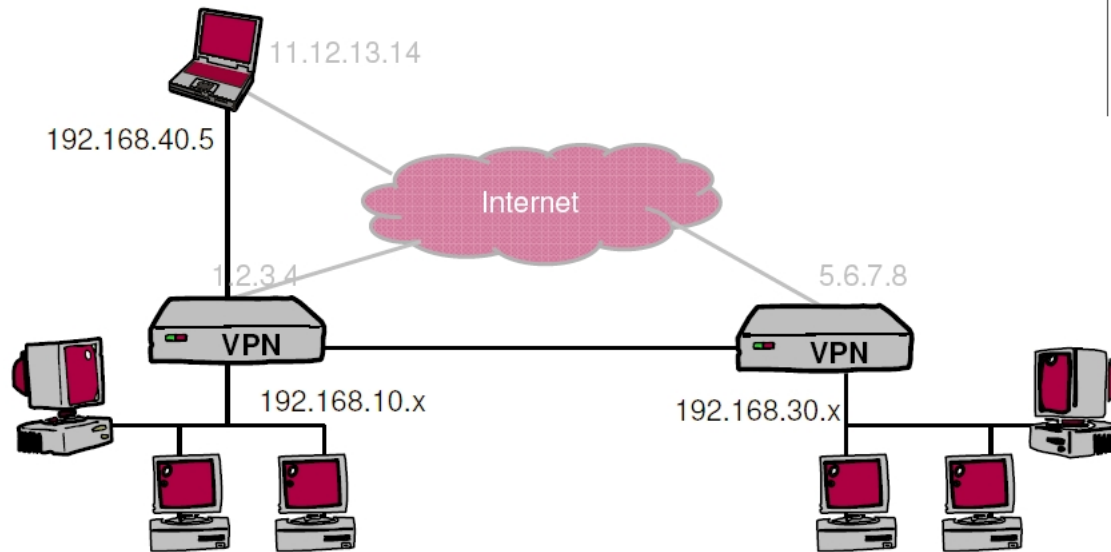
Goal: extend a private network across a public network

■ Scenarios:

- Interconnection of remote sites through the Internet
- Access to a company's network from a laptop connected to Internet



VPN Illustration





VPN basics

- VPN software on **routers** or **PCs** (e.g. laptop)
- Packet **encapsulation** across the Internet
- Encryption of data to guarantee **confidentiality**

Existing VPN Protocols

- Point to Point Tunneling Protocol (PPTP)
 - Microsoft

- Layer 2 Tunneling Protocol (L2TP)
 - IETF
 - Result of merging Cisco's Layer 2 Forwarding (L2F) protocol and Microsoft's PPTP protocol

- IP Security (IPsec)
 - IETF



IPsec

IP Security



IPsec Overview

- Open standard developed by the IETF
 - Public algorithms for confidentiality, authentication, integrity
- Authentication Headers (AH)
 - Provide connectionless **integrity** and origin **authentication** for IP packets
- Encapsulating Security Payloads (ESP)
 - Provide **confidentiality**, data-origin **authentication**, connectionless **integrity**
- Security Associations (SA)
 - Provide algorithms and parameters necessary to AH and/or ESP operations
- Internet Key Exchange (IKE)
 - Key exchange protocol
- Two operation modes: **tunnel, transport**

Security Association (SA)

- End hosts willing to exchange packets securely must first establish a **Security Association (SA)**
- A SA is simply the bundle of algorithms and parameters that is being used to encrypt and authenticate a particular flow in one direction
- SA memorizes algorithms, keys, validity periods, sequence numbers and peer's identity
- In normal bi-directional traffic (like TCP), flows are secured by a pair of SAs



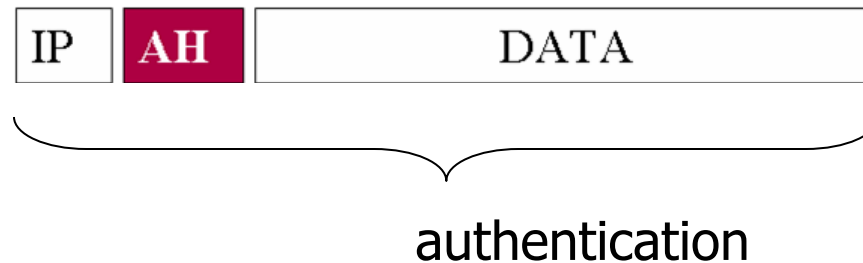
Security Association (SA)

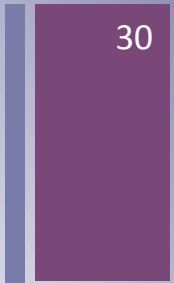
SAs are identified by a Security Parameter Index (SPI)

- The source indicates the SPI on all packets that it sends
- The destination uses the SPI to find the corresponding SA
- The source decides which packets must be processed with which SA
- One SA per destination, per protocol (AH or ESP), per flow

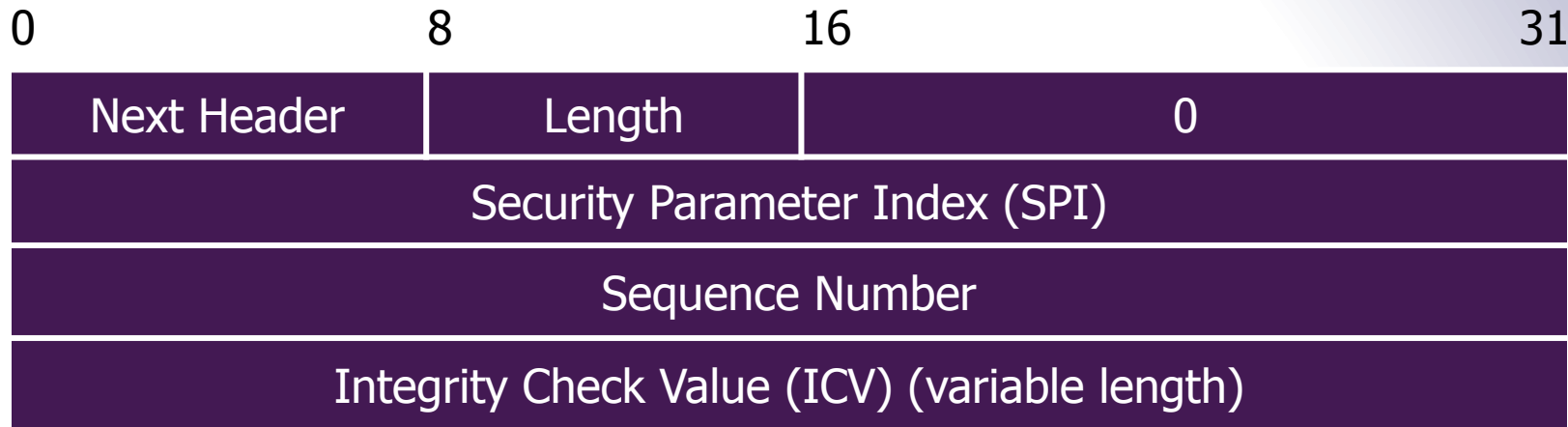
Authentication Header (AH)

- The addition of an authentication header allows verifying the packet's **authenticity** and **integrity**





Authentication Header (AH)



- Next Header: Specifies the encapsulated protocol (ICMP, TCP, UDP,...)
- Length: Size of this Authentication Header in 32-bit units, minus 2 (i.e., - 64 bits)
- Security Parameters Index: Contains a pseudo random value used to identify the security association for this datagram.
- Sequence Number: Monotonically increasing number to avoid replay-attacks.
- Integrity Check Value: Contains keyed-hash value



Authentication Header (AH)

- Authentication is calculated on:
 - Data that follow the AH
 - AH itself (with ICV set to zero)
 - Pseudo IP header
 - Source, destination, protocol, length, version, etc.
- The algorithm to be used to generate the authentication data is negotiated when the SA is created
- Two algorithms **must** be available:
 - HMAC-SHA-96
 - HMAC-MD5-96

Recall HMAC

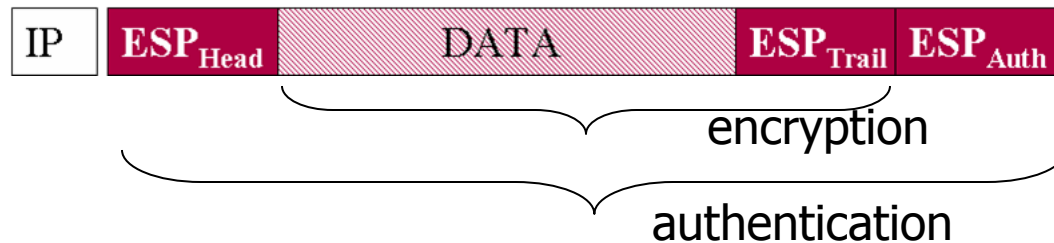
- Most widely used MAC on the Internet
 - Proposed by Bellare, Canetti, Krawczyk in 1996
 - Provably secure
 - Standards: FIPS 198-1, RFC 2104, ISO 9797-2
- Builds a MAC out of a hash function

$$\text{HMAC: } S(k, m) = H\left(k \oplus \text{opad} \parallel H(k \oplus \text{ipad} \parallel m)\right)$$

- Examples:
 - HMAC-SHA256: $H = \text{SHA256}$; output is 256 bits
 - HMAC-SHA1-96: $H = \text{SHA1}$; output truncated to 96 bits

Encapsulated Security Payload (ESP)

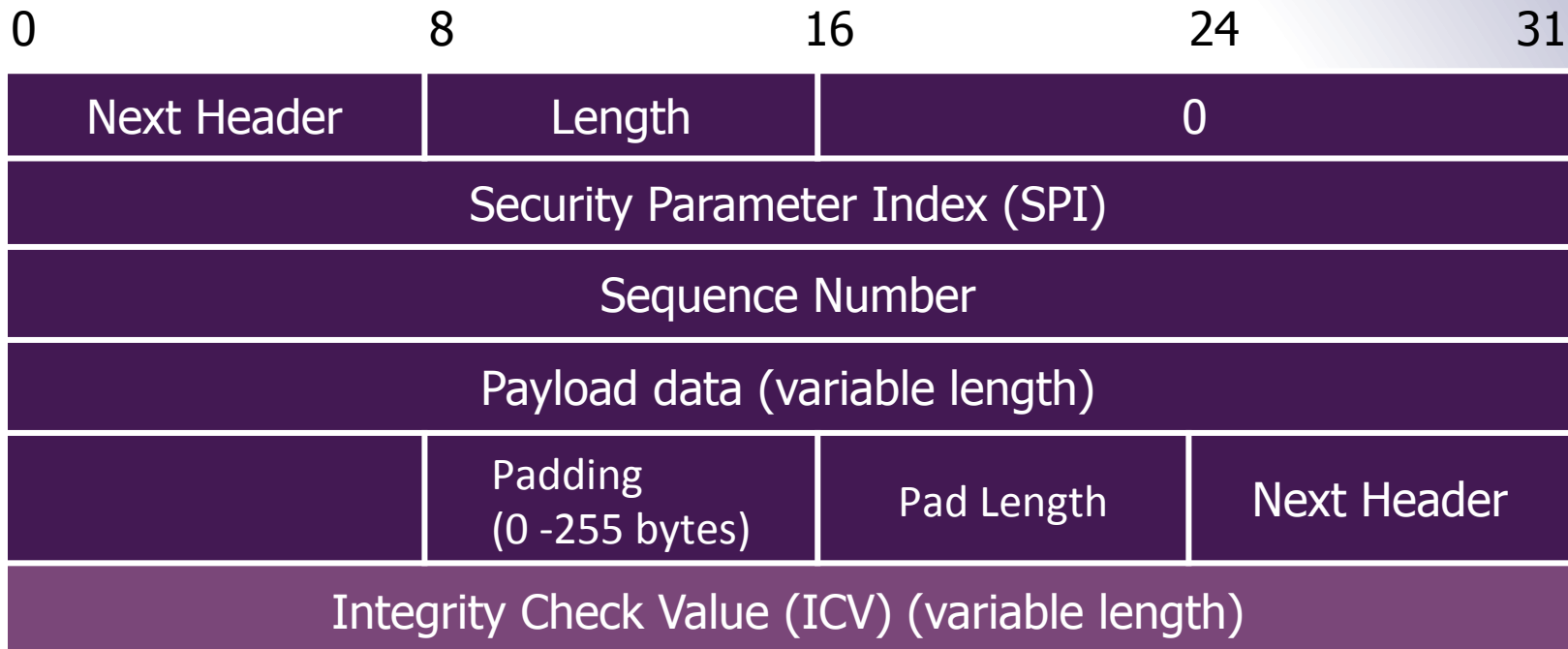
- The ESP header allows packet encryption and authentication



- Encryption is done only on the encapsulated data and the trailer
- Encryption is done neither on the header's fields, nor on the authentication data
- Optional authentication is done on the ESP header and all that follow, but not on the IP header



Encapsulated Security Payload (ESP)



- The mandatory algorithms are:
 - Encryption: DES-CBC, NULL (RFC 2410)
 - Authentication: HMAC-SHA-96 (RFC2404), HMAC-MD5-96 (RFC2403), NULL
- NULL encryption and NULL authentication in the same SA is not allowed

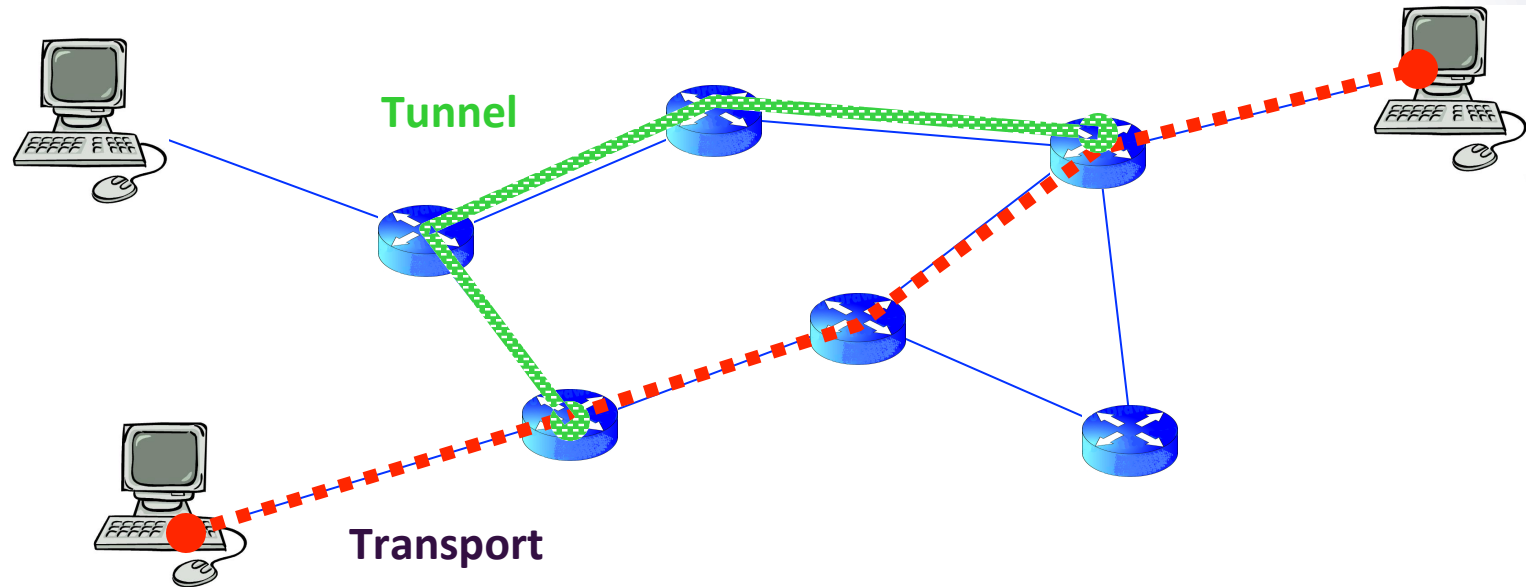


+

Transport and Tunnel Modes

Transport & Tunnel Modes

- **Transport:** only protects the packet's payload
- **Tunnel:** entire packet is encapsulated



Transport & Tunnel Modes

Which packets need to be encrypted/authenticated?

- Each router contains a **Security Policy Database**
- SPD defines which packet needs to be secured
 - According to discriminators: destination address, source address, ...
- **Example:**
 - Secure all HTTP traffic
 - Secure packets sent to remote sites but not to the Internet
 - Secure UDP
 - Secure TCP but not SSL

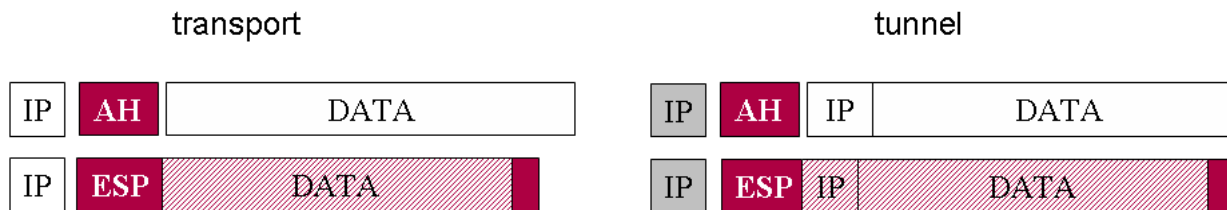
Transport & Tunnel Modes

■ Transport mode:

- Only IP packet **payload** is encrypted and/or authenticated

■ Tunnel mode:

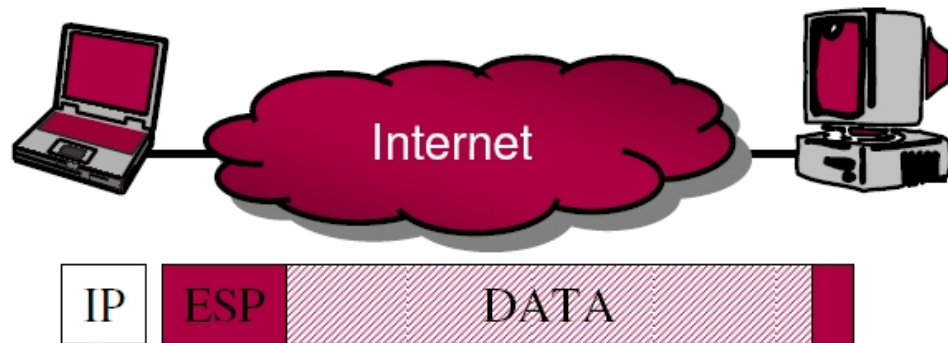
- The entire packet is encapsulated in a new packet





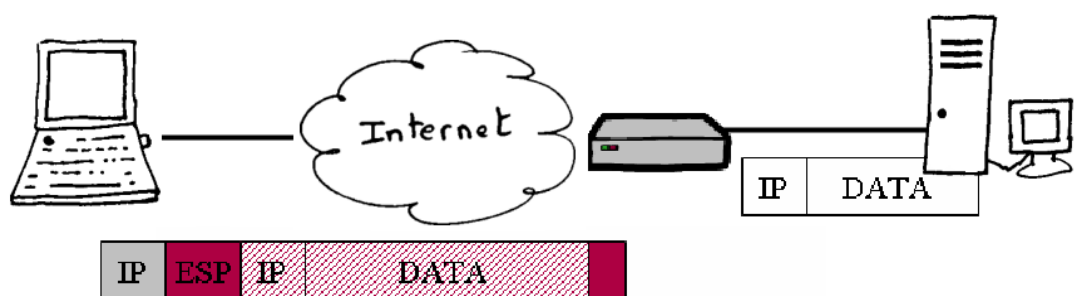
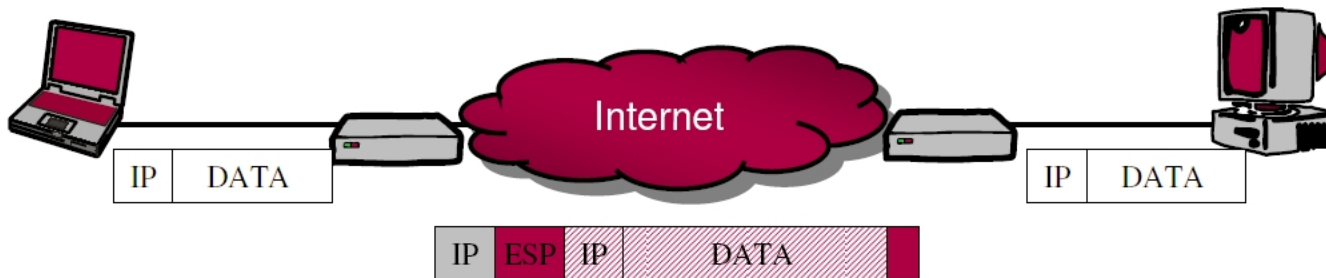
Transport Mode

- Security is done end-to-end



Tunnel Mode

- Security can be added by intermediate routers



Quiz: What is the Applied Mode?

TRANSPORT

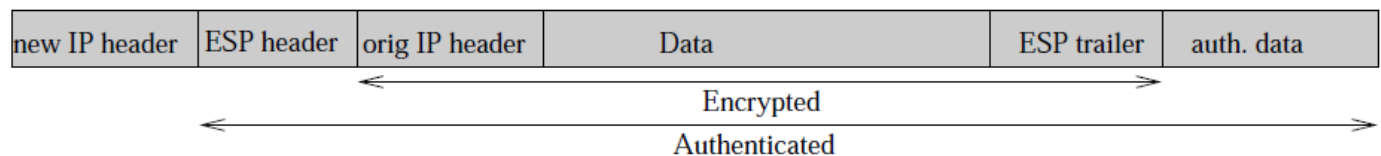
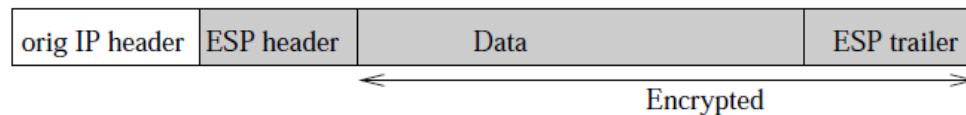
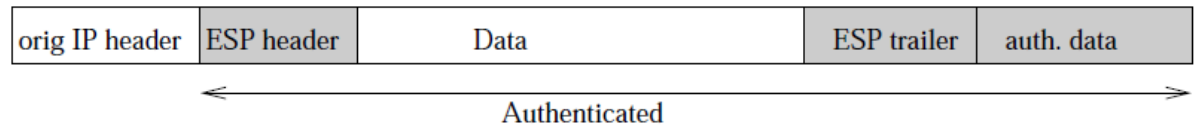
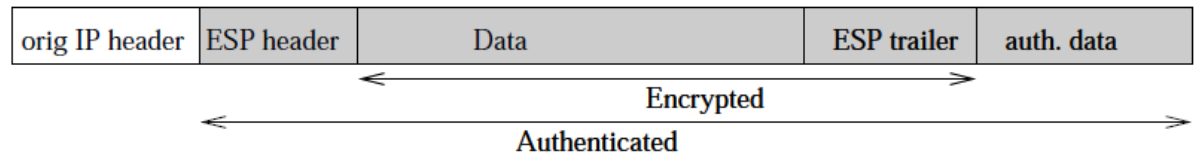
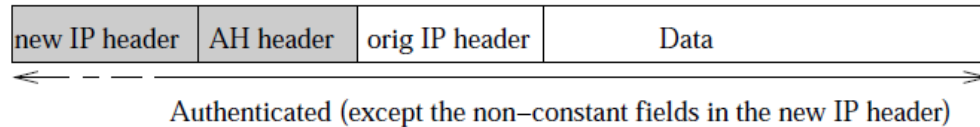
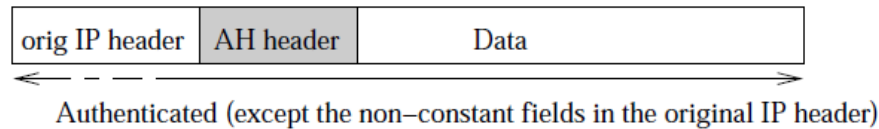
TUNNEL

TRANSPORT

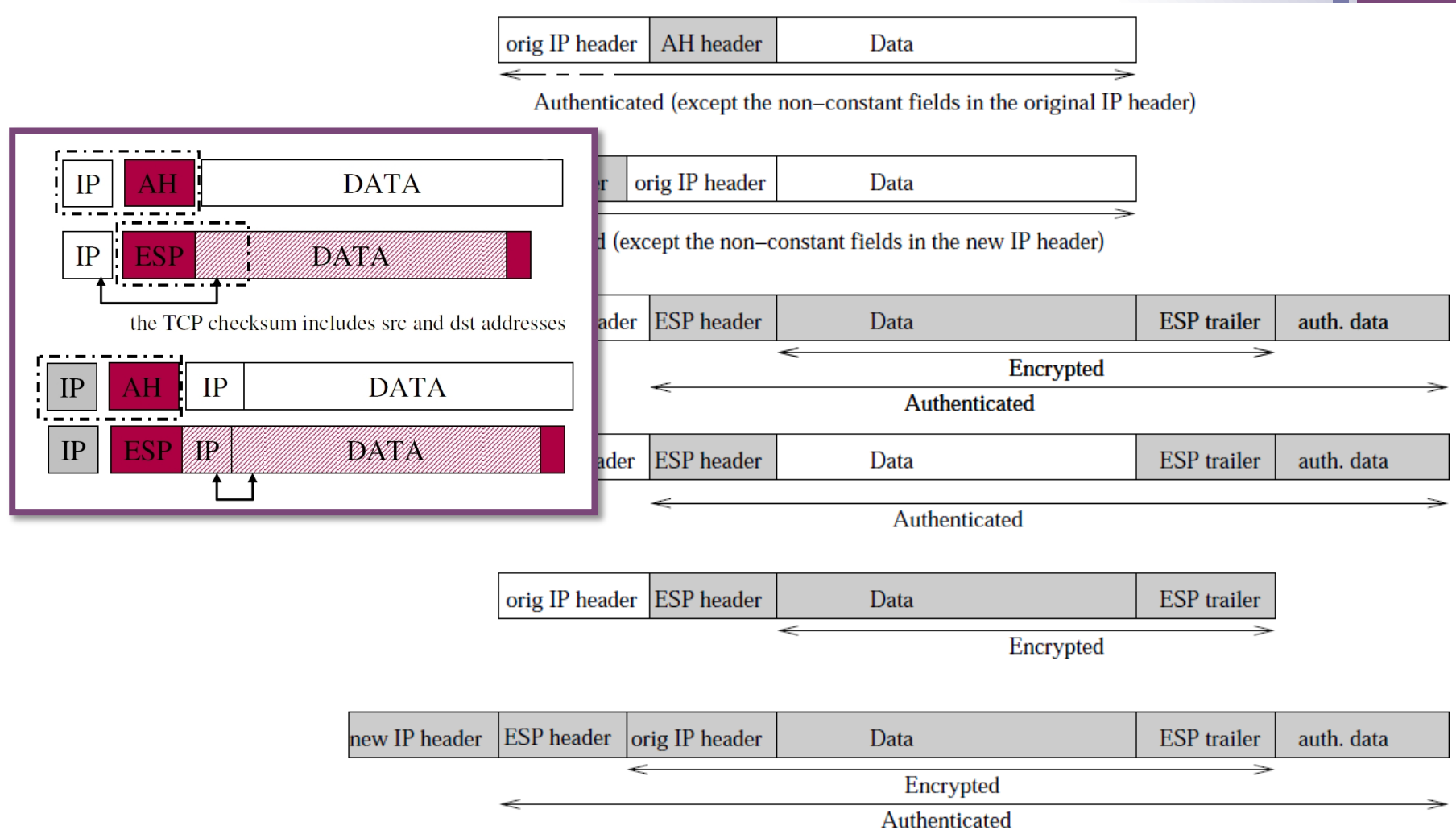
TRANSPORT

TRANSPORT

TUNNEL



IPsec and NAT



IPsec and NAT

- The TCP and UDP checksum calculation includes a pseudo header made of src and dst IP addresses and ports
- When doing NAT, the checksum has to be readjusted every time the source IP address (and port) changes
- This does not work if the payload is encrypted or authenticated
- **NAT-T** mechanism: encapsulate IPsec in UDP to traverse NAT



Internet Key Exchange (IKE)

Internet Key Exchange (IKE)

Internet Key Exchange (IKE) is a protocol used to establish a SA between communicating partners

- IKE's aims:

- Partner authentication
- Key exchange between partners
- Parameters negotiation

- IKE's result is a Security Association (SA)

- SA is identified with a given Security Parameter Index (SPI)

- IKEv1 RFC 2409, 1998

- IKEv2 RFC 4306, 2005, updated RFC 5996, 2010



Two Phases of IKEv1

- Phase 1: set up an SA to protect the negotiations
- Phase 2: set up the SA for an ESP or AH flow

IKEv1: phase 1

■ Authentication

- Pre-shared secrets (PSS)
- Public keys peer-exchanged
- X.509 certificates
 - Require public key of certification authority

■ Key Exchange

■ Generate a shared key using Diffie-Hellman



IKEv1: phase 1

Parameter negotiation

Main Mode

- More negotiation possibilities
 - E.g., DH parameters
- Protects the initiator's identity, and PSS's hash value (if used)

Aggressive Mode

- Faster but less negotiation possibilities
- Reveals the initiator's identity, and PSS's hash value (if used)

Main Mode & Aggressive Mode

main mode:

negotiate crypto
↔

Diffie-Hellman
↔

proof I'm Alice
↔
proof I'm Bob

aggressive mode

$g^a \bmod p$, "Alice"
→

$g^b \bmod p$, proof I'm Bob
←

proof I'm Alice
→

proof can be

- shared secret
- public key
- certificate

Perfect Forward Secrecy (PFS)

A key-agreement protocol has PFS if it ensures that a session key derived from a set of long-term keys will not be compromised if one of the long-term keys is compromised in the future

- Forward secrecy is designed so that the compromise of one message cannot lead to the compromise of others
- Also that there is not a single secret value which can lead to the compromise of multiple messages



IKEv1: phase 2

- Only one mode, called “quick mode”
- Without PFS
 - Keys are periodically refreshed (typically every hour): session keys
 - Session keys are derived from the same secret
 - Stealing this secret compromises all keys
- With PFS
 - A Diffie-Hellman is done for each new session key (slower)
 - Stealing one key does not compromise the previous ones



Glossary

- VPN: Virtual Private Network
- PPTP: Point to Point Tunneling Protocol
- L2TP: Layer 2 Tunneling Protocol
- L2F: Layer 2 Forwarding
- IPsec: IP Security
- AH: Authentication Header
- ESP: Encapsulated Security Payload
- IKE: Internet Key Exchange
- SPI: Security Parameter Index

+

Any questions?

53



Stay tuned



Next time you will learn about

WEP

