INGI2347 : EXERCISES* LAB SESSION 4

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Exercise 1: Amplification Attack (5 min)

Steve Pirate noticed that a DNS zone transfer request (used for example to provide to a slave DNS server the information it has to know regarding the zone(s) it serves) is always 27 bytes long, and the DNS server dns.thingummy.com response is 745 bytes. Let us assume that the communications are based on UDP. Steve Pirate is able to use source address spoofing when it connects to the dns.thingummy.com.

- 1. How can S. Pirate carry out a denial of service attack against the server poor.victim.com?
- 2. If Steve Pirate has a bandwidth of 256 Kbps and is able to use 100% of it for his attack, how much of the bandwidth belonging to poor.victim.com will be used by the attack?

Exercise 2: Network Traffic Analysis (5 min)

Simon Osterwald, a computer engineer working for a Swiss bank, is responsible for the proper working and security of the employer's network. One morning, during a routine control, he notices suspicious traffic on the server that manages the database of the customers bank accounts. Below is a portion of this traffic.

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02:01:27 arp reply 100.20.94.1 is-at 0:3:47:48:5c:ee
02:01:27 arp reply 100.20.94.1 is-at 0:3:47:48:5c:ee
02:01:28 arp reply 100.20.94.1 is-at 0:3:47:48:5c:ee 02:01:28 arp reply 100.20.94.1 is-at 0:3:47:48:5c:ee
02:01:28 arp reply 100.20.94.1 is-at 0:3:47:48:5c:ee
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02:01:28 arp reply 100.20.94.1 is-at 0:3:47:48:5c:ee
02:01:28 arp reply 100.20.94.1 is-at 0:3:47:48:5c:ee
02:01:29 arp reply 100.20.94.1 is-at 0:3:47:48:5c:ee
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^{*}A part of these exercises comes from the book "Computer System Security". The reproduction and distribution of these exercises or a part of them are thus forbidden.

- 1. Knowing that 100.20.94.1 is the gateway's address, give a short interpretation of this traffic.
- 2. Explain how you can protect against ARP spoofing.

Exercise 3: ARP/DNS Spoofing (10 min)

Let us consider a local area network (LAN) composed of two workstations and separated from the outside by a router (gateway) as represented in Figure 1. The workstations are configured to use the DNS server at 128.178.33.38, outside the LAN and do not use an internal DNS cache. Finally, let us consider two HTTP servers outside the LAN, www.site.ch and www.fakesite.ch. The objective of this exercise is to propose an attack based on *DNS spoofing*, such that when station1 user (victim) tries to reach the site www.site.ch, he will transparently end on the site www.fakesite.ch. The attack will be carried out from station2.

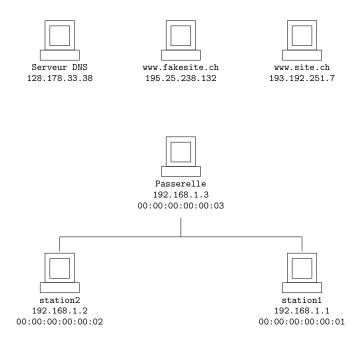


Figure 1: Architecture of the network to be attacked by ARP spoofing

When a workstation wants to communicate with the outside of the LAN, it uses, as destination MAC address the MAC address of the gateway. The gateway receives the packet and retransmits it towards its destination (which is outside the LAN); the destination address in the IP packet remains unchanged. For now, let us suppose that none of the computers on the LAN (including the gateway) know the MAC addresses of the other computers and that the ARP protocol is used to obtain the MAC addresses.

- 1. The user on computer station1 executes the command ping 192.168.1.2. Write down the ARP and IP messages exchanged on the LAN as well as the destination and source address (MAC, and IP when relevant) of these until the ping packet can actually be sent.
- 2. The same question for ping 128.178.33.38.
- 3. The same question for ping www.site.ch.
- 4. We suppose that station2 has launched an ARP spoofing attack to impersonate the gateway and to replace www.site.ch by www.fakesite.ch. Draw in Figure 1 the paths taken by the

packets transiting on the LAN when station1 executes the command ping www.site.ch (do not draw the ARP requests and replies).

Exercise 4: DHCP Vulnerabilities (5 min)

DHCP schematically works in the following way: when a computer needs an IP address to connect to a network, it sends a DHCPDISCOVER request to all computers of the network located on the same segment. A DHCP server is then responsible to listen to the network in order to allocate IP addresses to the different computers. When it receives a request from a computer, it returns a DHCPOFFER packet containing an IP address along with a validity period, as well as other information, such as for example, the gateway's IP address, or those of the DNS servers. These communications are unfortunately not secure.

- 1. Describe a simple attack that allows a pirate to prevent the clients from obtaining IP Addresses.
- 2. How is it possible to extend the previous attack for the pirate to be able to intercept the outbound LAN communications of a targeted computer?

Exercise 5: TCP session Hijacking (10 min)

A pirate (Philippe) is sniffing a telnet connection between Gildas and Pascal. He forges a TCP packet to insert the \n echo HACKED \n command into the data flow. The last exchanges before the insertion is shown below. Complete Figure 2 with the inserted packet and the following packets.

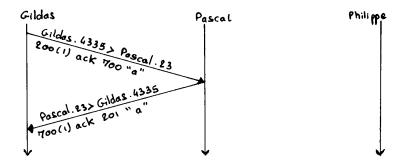


Figure 2: Packets exchanged between Gildas and Pascal before Philippe hijacks the session