CS5413 - Assignment: DFS and etc.

The input graph is to be given in the form of adjacency matrix as follows. Note that the input should in a file. Note that each edge is to be assigned with a weight ("1" is just to show a format).

```
f
           b
                         d
                                                   h
                                                          i
                                                                j
     a
                  С
                               е
                                            g
                                            0
                                                   0
     0
            1
                  0
                         0
                               0
                                      1
                                                          0
                                                                0
a
b
     0
            0
                  1
                         0
                               0
                                      0
                                            1
                                                   0
                                                          0
                                                                0
     0
            0
                  0
                         1
                               0
                                      0
                                            0
                                                   1
                                                          0
                                                                0
С
     0
           0
                  0
                         0
                                      0
                                            0
                                                   1
d
                               1
                                                          1
                                                                1
     0
           0
                  0
                         0
                               0
                                      0
                                            0
                                                   0
                                                          0
                                                                1
е
f
     0
            0
                         0
                               0
                                      0
                                                          0
                                                                0
                  0
            0
                         0
                               0
                                      0
                                            0
                                                   1
                                                          0
                                                                0
g
     1
     0
                  1
                                      0
                                            0
                                                   0
                                                          0
h
           0
                         0
                               0
                                                                0
i
     0
           0
                  0
                         0
                               0
                                      0
                                            0
                                                   0
                                                          0
                                                                1
j
     0
            0
                  0
                         0
                               1
                                      0
                                            0
                                                   0
                                                          0
                                                                0
```

```
(a) Implement the following algorithm (70%). DFS(G)
```

```
\begin{array}{lll} \textbf{1 for each vertex } u \in V[G] \\ 2 & \textbf{do } color[u] \leftarrow WHITE \\ 3 & \pi[u] \leftarrow NIL \\ 4 \text{ time } \leftarrow 0 \\ 5 \text{ for each vertex } u \in V[G] \\ 6 & \textbf{do if } color[u] = WHITE \\ 7 & \textbf{then } DFS\text{-}VISIT(u) \end{array}
```

```
DFS-VISIT(u)

1 color[u] \leftarrow GRAY

2 time \leftarrow time+1

3 d[u] \leftarrow time

4 for each v \in Adj[u]

5 do if color[v] = WHITE

6 then \pi[v] \leftarrow u

7 DFS-VISIT(v)
```

8
$$color[u] \leftarrow BLACK$$

9 $f/u/ \leftarrow time \leftarrow time + 1$

Output: an adjacency matrix with discovery time and finishing time on each vertex.

(b) Then implement the following algorithm. (20%)

TOPOLOGICAL-SORT(G)

- 1 call DFS(G) to compute finishing times f/v for each vertex v
- 2 as each vertex is finished, insert it onto the front of a linked list
- 3 return the linked list of vertices

Output: a sorted list of vertices with the finishing time in decreasing order.

(c) Given directed graph G = (V, E) (10%), a **strongly connected component (SCC)** of G is a maximal set of vertices $C \subseteq V$ such that for all $u, v \in C$, both $u \leadsto^b v$ and $v \leadsto^b v$. Implement an algorithm in the book to find the SCCs in the graph.

Output: a set of sub-adjacency graphs as resulting SCC's. No birding edges need be indicated.