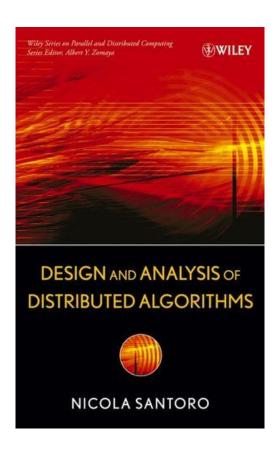
Towards an abstract syntax for distributed algorithms

Presentation by Rúben Guilherme, Paulo Duarte and João Fraga

The book

- Very simple and understandable.
- Very good specification of what a distributed environment is.
- Easy to read pseudo-code for the algorithms which makes it easier for us to create a simple implementation.



The Flooding example - book

PROTOCOL Flooding.

```
    Status Values: S = {INITIATOR, IDLE, DONE};
    S<sub>INIT</sub> = {INITIATOR, IDLE};
    S<sub>TERM</sub> = {DONE}.
```

• Restrictions: Bidirectional Links, Total Reliability, Connectivity, and Unique Initiator.

```
INITIATOR

Spontaneously
begin

send (M) to N(x);
become DONE;
end

IDLE

Receiving (I)
begin

Process (M);
send (M) to N(x) - {sender};
become DONE;
end
```

The Flooding example - syntax part 1

- First the user defines the states of the entities.
- Then, the restrictions this protocol uses.
- Afterwards, the behaviour of the entities during the protocol.

```
section protocol Flooding {
 state = [INITIATOR, IDLE, DONE];
 init = [INITIATOR, IDLE];
 term = [DONE];
 restric = [BidirectionalLinks, TotalReliability, Connectivity, UniqueInitiator];
 behaviour = {
   INITIATOR, spontaneously,
   begin
     send(M) to neighbours();
     become DONE;
    end
    IDLE, Receiving(M),
    begin
     PROCESS(M);
     send(M) to neighbours() - {sender};
     become DONE;
    end
 procedures = {
   PROCESS
   begin
      (* nothing to process *)
```

The Flooding example - syntax part 2

We also decided to define the network on a different section of the code.

```
section network {
  entities = [
    1, INITIATOR;
    2, IDLE;
    3, IDLE;
    4, IDLE;
    5, IDLE;
  edges = [
    1, 2;
    2, 4;
```

AST Implementation

From code To trees

Entity

```
type 'a register = 'a array
type 'a memory = (string, 'a) Hashtbl.t
type state = string
type entity = {
 id : int:
 clock : int;
 register: int register;
 memory : int memory;
 state: state:
```

Capabilities Each entity $x \in \mathcal{E}$ is endowed with local (i.e., private and nonshared) memory M_x . The *capabilities* of x include access (storage and retrieval) to local memory, local processing, and communication (preparation, transmission, and reception of messages). Local memory includes a set of *defined registers* whose values are always initially defined; among them are the *status register* (denoted by *status*(x)) and the *input value register* (denoted by *value*(x)). The register *status*(x) takes values from a finite set of system states x; the examples of such values are "Idle," "Processing," "Waiting,"... and so forth.

In addition, each entity $x \in \mathcal{E}$ has available a local *alarm clock* c_x which it can set and reset (turn off).

An entity can perform only four types of *operations*:

- local storage and processing
- · transmission of messages
- (re)setting of the alarm clock
- changing the value of the status register

External event

type event = MessageReceived of message | AlarmRing | Spontaneous

External Events The behavior of an entity $x \in \mathcal{E}$ is *reactive*: x only responds to external stimuli, which we call *external events* (or just *events*); in the absence of stimuli, x is inert and does nothing. There are three possible external events:

- · arrival of a message
- ringing of the alarm clock
- · spontaneous impulse

Message

```
type message = {
   sender : entity;
   receiver : entity;
   message : string;
   code : int;
}
```

Behaviour

type behaviour = (state * event, statement) Hashtbl.t

Behavior The nature of the action performed by the entity depends on the nature of the event e, as well as on which status the entity is in (i.e., the value of status(x)) when the events occur. Thus the specification will take the form

Status \times Event \longrightarrow Action,

Restrictions

type restriction = MessageOrdering ReciprocalCommunication BidirectionalLink EdgeFailureDetection **EntityFailureDetection** GuaranteedDelivery PartialReliability TotalReliability Connectivity BoundedCommunicationDelays UnitaryCommunicationDelays SynchronizedClocks UniqueInitiator

- *Message Ordering:* In the absence of failure, the messages transmitted by an entity to the same out-neighbor will arrive in the same order they are sent.
- Reciprocal communication: $\forall x \in \mathcal{E}, N_{\text{in}}(x) = N_{\text{out}}(x)$. In other words, if $(x, y) \in \vec{E}$ then also $(y, x) \in \vec{E}$.
- Bidirectional links: $\forall x \in \mathcal{E}$, $N_{\text{in}}(x) = N_{\text{out}}(x)$ and $\lambda_x(x, y) = \lambda_x(y, x)$.
- Edge failure detection: $\forall (x, y) \in \vec{E}$, both x and y will detect whether (x, y) has failed and, following its failure, whether it has been reactivated.
- Entity failure detection: $\forall x \in V$, all in- and out-neighbors of x can detect whether x has failed and, following its failure, whether it has recovered.
- Guaranteed delivery: Any message that is sent will be received with its content uncorrupted.
- Partial reliability: No failures will occur.
- Total reliability: Neither have any failures occurred nor will they occur.
- Connectivity: The communication topology \hat{G} is strongly connected.
- Bounded communication delays: There exists a constant Δ such that, in the absence of failures, the communication delay of any message on any link is at most Δ .
- Synchronized clocks: All local clocks are incremented by one unit simultaneously and the interval of time between successive increments is constant.
- *Unitary communication delays:* In the absence of failures, the communication delay of any message on any link is one unit of time.
- 4. Unique Initiator, that is, only one entity will start.

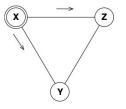
Protocol

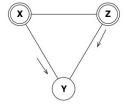
```
type protocol = {
  name : string;
  state : state list;
  init : state list;
  term : state list;
  restric : (restriction, unit) Hashtbl.t;
  behaviour : behaviour;
  procedures : statement list;
}
```

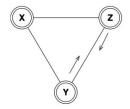
Network

type network = (edge * edge, unit) Hashtbl.t

type edge = entity * entity







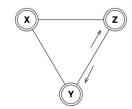


FIGURE 1.3: An execution of Flooding.

Program

```
type program = { protocol : protocol; network : network }
```

General Programming

```
type identifier = string
type uniop = Uneq | Unot
type binop =
    Band
    Bor
    Bxor
    Badd
    Bsub
    Bmul
    Bdiv
    Bmod
    Beq
    Bne
    Blt
    Bgt
    Ble
    Bge
```

```
type constant =
  | Cbool of bool
  | Cint of int
  | Cfloat of float
  | Cstring of string
type expression =
  | Econst of constant
  | Evar of identifier
  | Epair of expression * expression
  | Euniop of uniop * expression
  | Ebinop of binop * expression * expression
type statement =
   Sblock of statement list
   Sexp of expression
  | Satr of identifier * expression
  | Sif of expression * statement * statement
  | Sspontaneously of statement
   Srecieve of event
   Ssend of event
   Sbecome of state
   Sneighbours of entity
```

Overview & future work

```
section protocol Flooding {
 state = [INITIATOR, IDLE, DONE];
 init = [INITIATOR, IDLE];
 term = [DONE];
 restric = [BidirectionalLinks, TotalReliability, Connectivity, UniqueInitiator];
 behaviour = {
   INITIATOR, spontaneously,
    send(M) to neighbours();
    become DONE;
   end
   IDLE, Receiving(M),
   begin
    PROCESS(M);
    send(M) to neighbours() - {sender};
    become DONE;
   end
 procedures = {
   PROCESS
   begin
    (* nothing to process *)
   end
```

```
type 'a register = 'a array
type 'a memory = (string, 'a) Hashtbl.t
type state = string
type entity = {
 id : int;
  clock : int;
  register : int register;
  memory : int memory;
  state : state;
type message = {
 sender : entity;
 reciever : entity;
 message : string;
 code : int;
type event = MessageReceived of message | AlarmRing | Spontaneous
type identifier = string
type uniop = Uneg | Unot
type binop =
  | Band
  Bor
  Bxor
  Badd
  Bsub
  | Bmul
  Bdiv
  I Bmod
  Beq
```