A Model of Double-Entry Bookkeeping Systems

1 Double-Entry Bookkeeping Space

We can view a double-entry bookkeeping information system as a three dimensional space $S = \{(a, d, m) \mid a \in A, d \in D, m \in M\}$, where,

- A is a finite ordered subset of a predefined set of accounts.
- D is a finite ordered subset of all possible dates in the Gregorian calendar.
- *M* is a finite ordered subset of all possible discrete moments in time, given a precision.

Figure 1 shows a double-entry bookkeeping space sample, where each cube represents a tuple of the form (a,d,m) and is associated with a integer value different from zero. Consider the cubes (4,1,1) and (6,1,1). They share the same moment and date—the only difference are the accounts. We can interpret this pair of cubes as a transaction recorded in the system, with one cube representing the debit and another representing the credit.

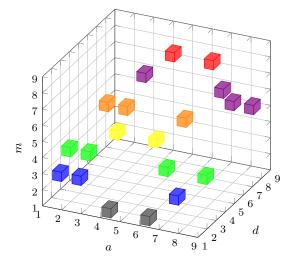


Figure 1: A double-entry bookkeeping space sample.

As, by definition, a double-entry transaction must have debits' sum equals to credits' sum, we can deduce that the values for both cubes (4,1,1) and (6,1,1) are the same, in this example.

More specifically, we can represent a double-entry bookkeeping space as a three dimensional array S:

$$S = (s_{ijk}), where \ 1 \le i \le |A|, \ 1 \le j \le |D|, \ 1 \le k \le |M|, \ s_{ijk} \in \mathbb{Z}$$
 (1)

Satisfying the following property:

$$\forall j, k \left(\sum_{i=1}^{|A|} s_{ijk} = 0 \right) \tag{2}$$

In equations (1) and (2) we assume that debits are represented by positive integers and credits by negative integers.

2 Operations

In the following description of the operations, the notation $\langle x, y \rangle$ means a range specification, and $R^{\langle n \rangle}$ means a set of n range specifications based on the ordered set R.

A double-entry bookkeeping space supports the following operations:

APPEND(S, S') append the space $S' = \left(s'_{ijk}\right)$ to the space S, such that, exists a set M which is associated with both space S and space S'. Let k_1 be the maximum value of k, such that $s'_{ijk} \neq 0$, and let k_2 be the minimum value of k, such that $s'_{ijk} \neq 0$. Given that $k_1 < k_2$, then:

$$\forall i, j \left(s_{ijk} = s'_{ijk} \right), where \ k_2 \leq k \leq |M|.$$

 $\begin{aligned} \mathbf{Projection}(S,A',D^{\langle n\rangle},M^{\langle m\rangle}) \text{ returns another space } S' &= \left(s'_{ijk}\right), \text{ where} \\ A' &\subseteq A, \text{ and } D^{\langle n\rangle} = \{\langle x_1,y_1\rangle,\dots,\langle x_n,y_n\rangle\}, \text{ and } M^{\langle m\rangle} = \{\langle v_1,w_1\rangle,\dots,\langle v_m,w_m\rangle\}, \\ \text{ and for all distinct } \langle x_i,y_i\rangle \text{ and } \langle x_j,y_j\rangle, \, \{x_i,\dots,y_i\} \cap \{x_j,\dots,y_j\} = \emptyset, \text{ and} \\ \text{ for all distinct } \langle v_i,w_i\rangle \text{ and } \langle v_j,w_j\rangle, \, \{v_i,\dots,w_i\} \cap \{v_j,\dots,w_j\} = \emptyset. \end{aligned}$

$$s'_{ijk} = \begin{cases} \sigma_{ijk}, & \text{if } f(j,k) \neq \emptyset \land D_j \in X \land M_k \in V \\ 0, & \text{otherwise,} \end{cases}$$

where:

$$\begin{split} \sigma_{ijk} &= \sum_{p=j}^{j+|x_r,\dots,y_r|-1} \sum_{q=k}^{k+|v_s,\dots,w_s|-1} s_{ipq}, \\ X &= \{x_1,\dots,x_n\}, \\ V &= \{v_1,\dots,v_m\}, \\ x_r &= D_j, \\ v_s &= M_k, \\ f(j,k) &= \{A_i: \sigma_{ijk} \neq 0\} \cap A', \ 1 \leq i \leq |A|. \end{split}$$

SLICE $(S, A', D^{\langle n \rangle}, M^{\langle m \rangle})$ returns another space $S' = \left(s'_{ijk}\right)$, where $A' \subseteq A$, and $D^{\langle n \rangle} = \{\langle x_1, y_1 \rangle, \dots, \langle x_n, y_n \rangle\}$, and $M^{\langle m \rangle} = \{\langle v_1, w_1 \rangle, \dots, \langle v_m, w_m \rangle\}$, such that:

$$s'_{ijk} = \begin{cases} s_{ijk}, & \text{if } f(j,k) \neq \emptyset \land D_j \in \bigcup\limits_{p=1}^n \{x_p, \dots, y_p\} \land M_k \in \bigcup\limits_{q=1}^m \{v_q, \dots, w_q\} \\ 0, & \text{otherwise}, \end{cases}$$

where:

$$f(j,k) = \{A_i : s_{ijk} \neq 0\} \cap A', \ 1 \le i \le |A|.$$

3 Alloy Model