Problem 8): For a public key encryption scheme $\prod = (\mathsf{Gen}, \mathsf{Enc}, \mathsf{Dec})$, we define **CPA** security according to the probability obtaining a secure result, as defined in the privacy experiment $\mathsf{PubK}^{\mathsf{LR\text{-}cpa}}_{\mathcal{A},\prod}$. This experiment goes as follows

The LR-orcale experiment $\mathsf{PubK}^{\mathsf{LR\text{-}cpa}}_{\mathcal{A},\prod}(n)$

1. $\mathsf{Gen}(1^n)$ *is run to obtain keys* (pk, sk).
2. *A uniform bit* $b \in \{0,1\}$ *is chosen.*
3. *The adversary* $\mathcal{A}$ *is given input* pk *and oracle access to* $\mathsf{LR}_{pk,b}(\cdot,\cdot)$.
4. *The adversary* $\mathcal{A}$ *outputs a bit* b'.
5. *The adversary* $\mathcal{A}$ *is defined to be* 1 *if* $b' = b$, *and* 0 *otherwise. If* $\mathsf{PubK}^{\mathsf{LR\text{-}cpa}}_{\mathcal{A},\prod}(n) = 1$, *we say that* $\mathcal{A}$ **succeeds**.

Using this definition for the experiment $\mathsf{PubK}^{\mathsf{LR\text{-}cpa}}_{\mathcal{A},\prod}$, we say that the encryption scheme $\prod$ is secure if the probability of $\mathcal{A}$ succeeding during $\mathsf{PubK}^{\mathsf{LR\text{-}cpa}}_{\mathcal{A},\prod}$