Homework #1

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Information Theory

Problem 2.3): Let $\mathbb{P}^n$ be the set of all n-dimensional probability vectors, with elements $\vec{p} \in \mathbb{P}^n$ defined as $\vec{p} = (p_1, p_2, \ldots, p_i, \ldots, p_n)$ for $i \in \mathbb{Z}^+ \ni i \leq n$. By the definition of a probability space, we must have

$$\vec{p} \cdot \vec{1} = \sum_{i=1}^{n} \{p_i\} = 1 \quad , \quad \forall \vec{p} \in \mathbb{P}^n, \tag{2.3-1}$$

where vector $\vec{1}$ is defined as $\vec{1} = (q_1, q_2, \ldots, q_k, \ldots, q_n) \in \mathbb{Z}^n$ with $q_k = 1, \forall k \in [1, n]$ (*where the interval* $[1, n]$ *is defined such that* $[1, n] \subseteq \mathbb{Z}^+$). Futhermore, the definition of a probability space also requires that, for any $\vec{p} \in \mathbb{P}^n$, the elements of $\vec{p}$ (*the* $p_i \in \vec{p}$ *such that* $i \in \mathbb{Z}^+ \ni i \leq n$) satisfy the condition

$$p_i \geq 0 \tag{2.3-2}$$

for all $i \in \mathbb{Z}^+ \ni i \leq n$.

The expression in 2.3-1 guarantees that the p_i of any $\vec{p} \in \mathbb{P}^n$ satisfy the bound $0 \leq p_i \leq 1$ where $i \in \mathbb{Z}^+ \ni i \leq n$. Therefore, the relation

$$p_i \log_2 [p_i] \geq 0 \tag{2.3-3}$$

holds for all p_i of any $\vec{p} \in \mathbb{P}^n$. Moreover, for the cases where $p_i = 0$ or $p_i = 1$, it is clear that the expression in 2.3-3 reduces to equality. Specifically, the realation $p_i \log_2 [p_i]$ becomes

$$p_i \log_2 [p_i] = 0 \tag{2.3-4}$$

for the case where $p_i = 0$ or $p_i = 1$. Moreover, the relation in 2.3-4 also represents the **smallest** possible value/result for the expression $p_i \log_2 [p_i]$. That is to say, that when $p_i = 0$ or $p_i = 1$, then $p_i \log_2 [p_i]$ is at a minimum.

The result in 2.3-1, makes it is clear that *only* **ONE** p_i in each $\vec{p} \in \mathbb{P}^n$ may have the value $p_i = 1$; therefore the probability vectors $\vec{p} \in \mathbb{P}^n$ which result in a minimum value for $p_i \log_2 [p_i]$ all have exactly one non-zero element with the non-zero element having a value of one. This implies that there are only n such probability vectors, $\vec{p}^{\star}$ within any $\mathbb{P}^n$. Furthermore, the value of $H(X) = \sum_{i=1}^{n} \{p_i \log_2 [p_i]\}$ for any such $\vec{p}^{\star}$ is also zero.

Problem 2.4 a): Recall the chain-rule for conditional entropies of X given Y,

$$H(X \mid Y) = H(X, Y) - H(Y) \tag{2.4-1}$$

We apply the expression in 2.4-1 to the case of $g(X)$ given X to obtain

$$H(g(X) \mid X) = H(g(X), X) - H(X) \tag{0.0.1}$$

by rearranging the expression in the previous result as follows,

$$H(g(X), X) = H(X) + H(g(X) \mid X) \tag{2.4-2}$$

we obtain the desired result.

Problem 2.4 b): For any given value of X, we automatically know $g(X)$. Therefore, the expression for $H(g(X), X)$ in 2.4-2 becomes

$$H(g(X), X) = H(X)$$

which is the desired result.

Problem 2.4 c): Recalling the expression for the conditional entropy chain rule in 2.4-1 and using it for the case where $X = X$ and $Y = g(X)$ yields the result

$$H(X \mid g(X)) = H(X, g(X)) - H(g(X))$$

rearranging the above expression yields

$$H(X, g(X)) = H(g(X)) + H(X \mid g(X)) \tag{2.4-3}$$

we obtain the desired result.

Problem 2.4 d): For any arbitrary function, $g(X)$, of a random variable X, the entropy $H(X \mid g(X))$ satisfies the condition

$$H(X \mid g(X)) \geq 0 \tag{2.4-4}$$

for the case where $g(X)$ is one-to-one, the relation in 2.4-4 simplifies to

$$H\left(X \mid g\left(X\right)\right) = 0$$

Applying the relation in 2.4-4 to the expression in 2.4-3 yields

$$\begin{aligned}
H\left(X, g\left(X\right)\right) &= H\left(g\left(X\right)\right) + H\left(X \mid g\left(X\right)\right) \\
&\geq H\left(g\left(X\right)\right) + H\left(X \mid g\left(X\right)\right) - H\left(X \mid g\left(X\right)\right) \\
&\geq H\left(g\left(X\right)\right)
\end{aligned}$$