## COMP 360 - Winter 2016 - Sample Final Exam

- 1. (10 points) Prove that the following problem belongs to P: Given a graph G, we want to know whether G has an independent set of size 100.
- 2. (10 Points) Prove that the following problem belongs to PSPACE: Given a graph G and an integer k, we want to know whether the number of independent sets in G is equal to k.
- 3. (10 Points) Show that the following problem is NP-complete:
  - Input: An undirected graph G and an edge e.
  - Question: Does G have a Hamiltonian cycle that passes through the edge e.
- 4. (10 points) In the MAX-CUT problem, given an undirected graph G we want to partition the vertices of G into two parts (A, B) such that the number of edges between A and B is maximized. Prove that the following is a  $\frac{1}{2}$ -factor approximation algorithm for this problem:
  - Let  $v_1, \ldots, v_n$  be all the vertices of G.
  - Initially set  $A = B = \emptyset$ .
  - For  $i = 1, \ldots, n$  do
  - IF  $v_i$  has more neighbours in A than in B THEN
  - add  $v_i$  to B
  - Else
  - add  $v_i$  to A
  - EndFor
- 5. (10 points) Prove that the following algorithm is a 2-factor approximation algorithm for the minimum vertex cover problem:
  - While there is still an edge e left in G:
  - Delete all the two endpoints of e from G
  - EndWhile
  - Output the set of the deleted vertices
- 6. (10 points) Prove that the following algorithm is a  $\frac{1}{2}$ -factor approximation algorithm for the MAX-SAT problem: Given a CNF  $\phi$  on n variables  $x_1, \ldots, x_n$ :
  - For  $i = 1, \ldots, n$  do
  - IF  $x_i$  appears in more clauses than  $\overline{x_i}$  THEN
  - Set  $x_i = T$

- Else
- Set  $x_i = F$
- Remove all True clauses from  $\phi$  and remove  $x_i$  and  $\overline{x_i}$  from all the other clauses
- EndFor
- 7. (10 points) A *kite* is a graph on an even number of vertices, say 2k, in which k of the vertices form a clique and the remaining k vertices are connected in a tail that consists of a path joined to one of the vertices of the clique. Prove that KITE problem defined as in the following is NP-complete.
  - Input: An undirected graph G, and a positive integer k.
  - Question: Does G contain a kite on 2k vertices as a subgraph?
- 8. Consider a graph G = (V, E). The chromatic number of G is the minimum number of colors required to color the vertices of G properly. Let  $\mathcal{I}$  be the set of all independent sets in G (Note that every element in  $\mathcal{I}$  is a set).
  - (a) (10 Points) Prove that the solution to the following linear program provides a lower-bound for the chromatic number of G.

$$\begin{array}{ll} \min & \sum_{I \in \mathcal{I}} x_I \\ \text{s.t.} & \sum_{I:v \in I} x_I \geq 1 \\ & x_I \geq 0 \end{array} \qquad \forall v \in V$$

- (b) (10 Points) Write the dual of the above linear program.
- (c) (10 Points) Prove that every clique in G provides a solution to the dual linear program.