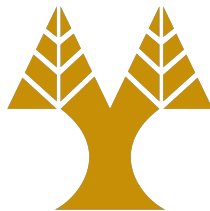


Thesis Dissertation

**WEBFUZZ: IMPLEMENTATION OF A GRAY-BOX  
FUZZING TOOL FOR WEB APPLICATIONS**

**Marcos Antonios Charalambous**

**UNIVERSITY OF CYPRUS**



**COMPUTER SCIENCE DEPARTMENT**

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**webFuzz: Implementation of a Gray-box Fuzzing Tool for Web  
Applications**

**Marcos Antonios Charalambous**

Supervisor

Dr. Elias Athanasopoulos

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# Abstract

Testing software is a common practice for exposing unknown vulnerabilities in security-critical programs that can be exploited with malicious intent. A bug-hunting method that has proven to be very effective is a technique called fuzzing. Specifically, this type of software testing has been in the form of fuzzing of native code, which includes subjecting the program to enormous amounts of unexpected or malformed inputs in an automated fashion. This is done to get a view of their overall robustness to detect and fix critical bugs or possible security loopholes. For instance, a program crash when processing a given input may be a signal for memory-corruption vulnerability.

Although fuzzing significantly evolved in analysing native code, web applications, invariably, have received limited attention, so far. This thesis explores the technique of grey box fuzzing of web applications and the construction of a fuzzing tool that will automate the process of discovering bugs in web applications.

We design, implement and evaluate webFuzz, which is the first gray-box fuzzer for web applications. webFuzz leverages instrumentation for successfully detecting reflective Cross-site Scripting (XSS) vulnerabilities faster than other black-box fuzzers. The functionality of webFuzz is demonstrated using WordPress and Drupal.

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# Chapter 1

## Introduction

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A highly automated testing technique that covers numerous boundary cases using invalid data (from files, network protocols, API calls, and other targets) as application input to better ensure the absence of exploitable vulnerabilities. The name comes from modem applications' tendency to fail due to random input caused by line noise on fuzzy telephone lines.

A highly automated testing technique that covers numerous boundary cases using invalid data (from files, network protocols, API calls, and other targets) as application input to better ensure the absence of exploitable vulnerabilities. The name comes from modem applications' tendency to fail due to random input caused by line noise on fuzzy telephone lines.

Numerous fuzzers have been developed in the past few years that try to optimize the fuzzing process by proposing various methodologies [7, 8, 20, 21, 35, 42, 46]. For instance, most of the fuzzers take advantage of instrumentation on the source or binary level. That is, inserting code to the program in order to receive feedback when a code block gets triggered and try to adjust the generated inputs to improve code coverage. Others utilize concolic/symbolic execution in order to extract useful information about the program and use that information for improving the input generation process [20, 21, 46]. However, all

these fuzzers are currently targeted towards finding vulnerabilities in native code, while web applications have received limited attention.

mutation-based fuzzer, might actively see the code paths executed in the target and make adjustments accordingly, which is very smart. EFS and AFL do exactly this

## **1.1 Motivation**

## **1.2 Related Work**

## **1.3 Contributions**

# Chapter 2

## Background

### Contents

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In this section we provide background information, to give a detailed understanding about various key points regarding this thesis. First, we define what a Cross-Site Scripting bug is in web applications, and elaborate on an example regarding this vulnerability. Then, we briefly discuss what fuzzing is and the various categories that constitute it and continue on how instrumentation helps when used during gray box fuzzing. After, we continue discussing the concept of concurrency in Python and we close with the containerization of services using Docker.

### 2.1 Web Application Bugs

The internet has been growing exponentially since its inception. Although there are over 1 billion pages currently on-line, writing a web application that it is secure from any available vulnerability, can be extremely hard. Every significant web application, especially large-scale that are composed with thousands of lines of code, have bugs in them.

Even the simplest ones can be the root of irreparable damage when they are exploited by attackers with malicious intentions. In fact, web application vulnerabilities account for the majority of the vulnerabilities reported in the Common Vulnerabilities and Exposures database [?]. The OWASP Top 10 represents a broad consensus about the most critical security risks to web applications [?]. One of the most pressing security problems on the Internet, according to the aforementioned list, is Cross-Site Scripting, also known as XSS.

XSS flaws occur whenever an application includes untrusted data in its web page responses without validating or escaping them first. In other words, the web application accepts input from the user and then attempts to display it without filtering for HTML tags or script code, such as JavaScript. As a result these untrusted data can get executed which can in turn hijack the browser, deface the web site, redirect the user to dangerous sites and many other attacks. Some XSS types include Reflected(aka Non-Persistent or Type II), Stored (Persistent or Type I) and DOM-based(Type-0).

Reflected XSS [?] vulnerabilities arise when data is copied from a request and echoed into the application's immediate response. This way, scripting language code included within a request can be dynamically executed. In the case of Stored XSS vulnerabilities, the malicious payload is first permanently stored in storage such a database residing on a server, and is only later outputted by an unsuspecting query. Examples might be Web forums or blog comments. Currently webFuzz focuses in detecting bugs that can lead to Reflected Cross-Site Scripting, which is among the most common of XSS attacks.

It is imperative that we understand what an RXSS (reflected XSS) bug typically looks like, in order to grasp the thesis' perspective. Most of the time RXSS is caused due to a failure to sanitise the user input. For instance, let us assume that we have a simple login page with two input fields: the username and password. The login page also displays appropriate error messages back to the user if the login fails. An implementation of this in PHP could look something like Listing 2.1.

---

```
1 <?php
2 $username=$_POST['username'];
3 $pwd=$_POST['password'];
4 if (search_username($username)) {
5     if (match_username_password($username, $pwd)) {
6         // do normal login work
```

```

7      } else {
8          echo 'Wrong_Password';
9      }
10 } else {
11     echo 'Error' . $username . 'was_not_found.';
12 }
13 ?>

```

---

Listing 2.1: Vulnerable login form.

The code above is faulty for two reasons. First, letting the user know that the username exists can help an attacker guess a set of correct credentials much faster, since only the password is left to find. But this design choice is not linked with Cross-Site Scripting. The source of the bug is on line 11 where the error message "the \$username was not found" is displayed. Because \$username is a tainted variable that has not been sanitized, an attacker can inject malicious payload in this field which will freely be interpreted by the HTML parser according to whatever its content is. Exploit: A victim is fooled into submitting a form located in an attacker controlled website. This malicious form is designed to trigger the vulnerability found in the above login form. As soon as the form is submitted the vulnerable login page is opened with the XSS script executed in it. If the victim now tries to login, the XSS script can easily send the credentials to the attacker as well.

Defeating XSS attacks is similar to defending against other types of code injection. The input must be sanitized. User input containing HTTP code needs to be escaped or encoded in order to avoid its execution. Also, systemwide measures such as Content Security Policy(CSP) [?] may be set as well to eliminate or mitigate XSS attacks. Nevertheless, flaws such as buffer overflows or cross-site scripting issues comprise a majority of security incidents, and malicious hackers abuse them on a daily basis.

## 2.2 Fuzzing

A promising technique for discovering unknown vulnerabilities in programs and web applications proven to be very effective, is a technique called fuzzing [?]. With this quality

assurance technique, software is exercised using a vast amount of anomalous inputs for inferring if any of them introduces security-related side effects. A fuzzer, which is the tool that can automate the aforementioned stress-testing process, can be categorized in relation to its awareness of the program structure as black-, white-, or gray-box [?].

A black-box fuzzer treats the program as a black box and is unaware of internal program structure. It conducts its test on the target through external interfaces and produces random inputs using no information of the target's underlying structure. Hence, black-box fuzzers are only able to scratch the surface usually and expose "shallow" bugs. [?] A white-box fuzzer infers source code knowledge, such as source code auditing, to reveal flaws in the software. It leverages program analysis to systematically increase code coverage or to reach certain critical program locations. Program analysis can be based on either static or dynamic analysis, or their combination [?]. They may also leverage symbolic execution in order to derive what inputs cause each part of a program to execute [?]. Therefore, they can be very effective at exposing bugs that hide deep in the program. By studying the application code, you may be able to detect optional or proprietary features, which should be tested as well. A fuzzer is considered gray-box when it leverages instrumentation rather than program analysis to glean information about the coverage of a generated input from the program it tries to fuzz. In this thesis we explore gray-box fuzzing, which is a combination of both the white-box and black-box approaches since it uses the internals of the software to assist in generating better test cases.

## **2.3 Instrumentation**

Typically, a fuzzer is considered more effective if it achieves a higher degree of code coverage. This can be explained by the fact that to be able to trigger any given bug, the fuzzer must first execute the code where the bug lies, so increasing code coverage increases the chances of executing unsafe pieces of code where bugs may reside. As we mentioned in the previous section, using instrumentation may be the key yield a higher code coverage percentage. Currently available fuzzers for web applications act in a blackbox fashion [16](FIX THIS REFERENCE FROM PAPER); they just brute force the target with URLs

that embed known web-attack payloads, with little or no information about the underlying structure of the target.

In contrast, webFuzz firstly instruments a web application by adding code that tracks all control flows triggered by an input and notifies the fuzzer, accordingly. Notifications can be embedded in the web application's HTTP response using custom headers or can be outputted to a shared file or memory region. On the other hand, the fuzzer starts sending requests to the target and analyses the responses in order to realize any interesting requests that would later help to improve the code coverage and as a result, trigger vulnerabilities nested deep in the web application's code.

We instrument web applications for delivering feedback once they are fuzzed. As opposed to native applications, where several options exist for instrumenting their source or binary representation, we decide to instrument web applications by modifying the Abstract Syntax Tree (AST) of PHP files and then reverting it back to source code form. This in turn provides us feedback on the basic blocks that are visited during analysis. For altering the AST of PHP files, PHP-Parser [?] is used. Instrumentation performed by webFuzz on our targeted web application is similar to how AFL instruments binaries but adapted to work in web applications. An elaborated approach of the instrumenting functionality provided by webFuzz is out of the scopes of this thesis.

## **2.4 Concurrency**

Concurrency is defined as working on multiple things at the same time [?]. However, in Python this does not mean that they work in parallel, since only one core of the CPU is active at any given time. Instead, each task takes turns in occupying the core and executing their code. When a task is interrupted, the state of each task is stored, so it can be restarted right from the point where it left off. Concurrency is aimed to speed up the overall performance of input/output (I/O) bound problems, whose performance can be slowed down dramatically when they are obliged to wait often for I/O operation from some external resource. An example of such resource are requests on the internet or any kind of network traffic that can take several orders of magnitude longer than CPU



instructions. An illustration of the above can be seen at Figure 2.1:

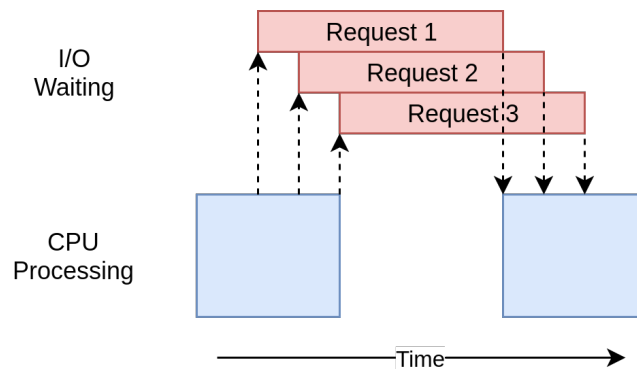


Figure 2.1: Requests over the internet processed in concurrent fashion [?].

More specifically in Python, concurrency can be expressed either through the Threading or AsyncIO(short for Asynchronous Input Output) [?] modules. Due to the infamous Global Interpreter Lock (GIL) [?] Python has, both AsyncIO and Threading are single-threaded, single-process design. Thus, there was no clear advantage of using the latter so AsyncIO was opted instead. Not to mention the complexity using threads and making the program thread-safe is added. In a few words, GIL makes sure there is only one thread running at any given time, thus making the use of multiple cores/processors with threads infeasible. In the Python community, there is a general rule of thumb when it comes to I/O-bound problems; “Use asyncio when you can, threading when you must”. More info on the AsyncIO module and its use in the webFuzz implementation can be found in Section 4

## 2.5 Docker

Docker containers provides developers the commodity of creating software locally with the knowledge that it will run identically regardless of the host environment [?]. Containers are an encapsulation of an application with its dependencies that share resources with the host OS, unlike Virtual Machines. During the evaluation, which can be seen in detailed in Section 5, a docker-compose YAML file was created to allow multiple containers to be initiated and managed at once, with a set of predefined configuration.

# Chapter 3

## Architecture

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### 3.1 Modelling of Mining Procedure

creating a Session object allows requests to do some fancy networking tricks and really speed things up.

### **3.1.1 Defining Mining Environment**

### **3.1.2 Strategies as State Machines**

## **3.2 Selfish Mining**

## **3.3 Stubborn Selfish Mining**

### **3.3.1 Lead**

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### **3.4.1 Safe-Lead**

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### 4.1 Selfish Miner

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### 5.1 Methodology

### 5.2 Dominant Strategies

### 5.3 Revenue and Comparison with Stubborn Strategies

### 5.4 Fairness of Blockchain

### 5.5 Risk Safety Property

## Chapter 6

### Future Work

examples of concurrency in this article run only on a single CPU or core in your computer. The reasons for this have to do with the current design of CPython and something called the Global Interpreter Lock, or GIL. SAY ABOUT aimultiprocessing... Hold out on adding concurrency until you have a known performance issue and then determine which type of concurrency you need. As Donald Knuth has said, “Premature optimization is the root of all evil (or at least most of it) in programming.” Parallelism consists of performing multiple operations at the same time. Multiprocessing is a means to effect parallelism, and it entails spreading tasks over a computer’s central processing units (CPUs, or cores).

# Chapter 7

## Conclusion

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### 7.1 Conclusion

### 7.2 Future Work

As Donald Knuth has said, “Premature optimization is the root of all evil (or at least most of it) in programming.”

# Appendix A



## **Appendix B**