

# Language Design for Everyone

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September, 29th 2023  
Lowell, MA, USA

# Programming Languages (PLs)

**As the theory and practice of PL advances:**

## **New PLs:**

**TypeScript (Microsoft), Hack (Facebook), Go (Google), Dart (Google), Swift (Apple), Rust (Mozilla), Julia, Elixir, and so on.**

## **New Domain-specific Languages (DSLs):**

**Halide. MIT (with help from Stanford, Google, Adobe.) for GPU  
Turi. (formerly GraphLab) CMU/UoW acquired by Apple, for ML  
GraphIt for Graph analytics.  
and so on**

# This Talk: Some of My Work on

*Empowering people to quickly deploy their own programming languages and domain-specific languages*

## Flexibility:

- Lang-n-play
- Lang-n-change
- Lang-n-send

## Safety:

- Lang-n-check
- Lang-SQL (or Lang-n-query)
- Lang-n-prove

# Lang-n-play

## (Languages as first-class)

C-like array access `a[10]`

Java-like array access `a[10]`

**access anything, unsafe**

**checks length, safe**

```
addCArrays(Language lan):
```

```
  return lan + { evaluator of C-like _[_] }
```

```
addJavaArrays(Language lan):
```

```
  return lan + { evaluator of Java-like _[_] }
```

```
mylan := {syntax + evaluator }
```

```
if trustworthy(code):
```

```
  then addCArrays(mylan) exec code
```

```
  else addJavaArrays(mylan) exec code
```

Matteo Cimini. *A Calculus for Multi-language Operational Semantics*. VSTTE 2021.

Matteo Cimini. *On the Effectiveness of Higher-Order Logic Programming in Language-Oriented Programming*. FLOPS 2020.

Matteo Cimini. *Languages as First-Class Citizens (Vision Paper)*. SLE 2018.

# Lang-n-send (Processes that send languages)

Server of language definitions

request  
array operations



send syntax + semantics  
of array operations

client



defines *mylan*  
defines *code*

Server: replicate[send(*arraysLan*)]  
parallelWith  
Client: receive(Server,*arrays*) ->  
(*mylan + arrays*) exec *code*

# Lang-n-change (Compiling languages to languages)

```
f(float x) { truncate(x) }
```

```
int x := 3;
```

```
f(x);
```

**without subtyping: reject**  
**with subtyping: accept**



```
type_check(part_of_prg):
```

```
  if part_of_prg is function type var body arg
```

```
  if type == type_check(arg) then accept
```

```
  else reject
```

**must become**

```
type_check_with_subtype(part_of_prg):
```

```
  if part_of_prg is function type var body arg
```

```
  type2 := type_check_with_subtype(arg)
```

```
  if subtype(type2, type) then accept
```

```
  else reject
```

# Lang-n-change (Compiling languages to languages)

Lang-n-change can express:

```
retrieve type_checker,  
compile 'type1 == type2' into:  
  'newVar1 := type1;  
   newVar2 := type2;  
   if subtype(newVar1, newVar2) then accept  
   else reject'
```

Matteo Cimini, Benjamin Mourad.

*Language Transformations in the Classroom.* EXPRESS/SOS 2021.

Benjamin Mourad, Matteo Cimini.

*A Declarative Gradualizer with Language Transformations.* IFL 2020.

Benjamin Mourad, Matteo Cimini. *System Description: Lang-n-Change - A Tool for Transforming Languages.* FLOPS 2020.

Benjamin Mourad, M. Cimini. *A Calculus for Language Transformations.* SOFSEM 2020.

# Lang-n-change (Compiling languages to languages)



add subtyping  
add object-oriented programming  
add pattern-matching  
...

Matteo Cimini, Benjamin Mourad.

*Language Transformations in the Classroom.* EXPRESS/SOS 2021.

Benjamin Mourad, Matteo Cimini.

*A Declarative Gradualizer with Language Transformations.* IFL 2020.

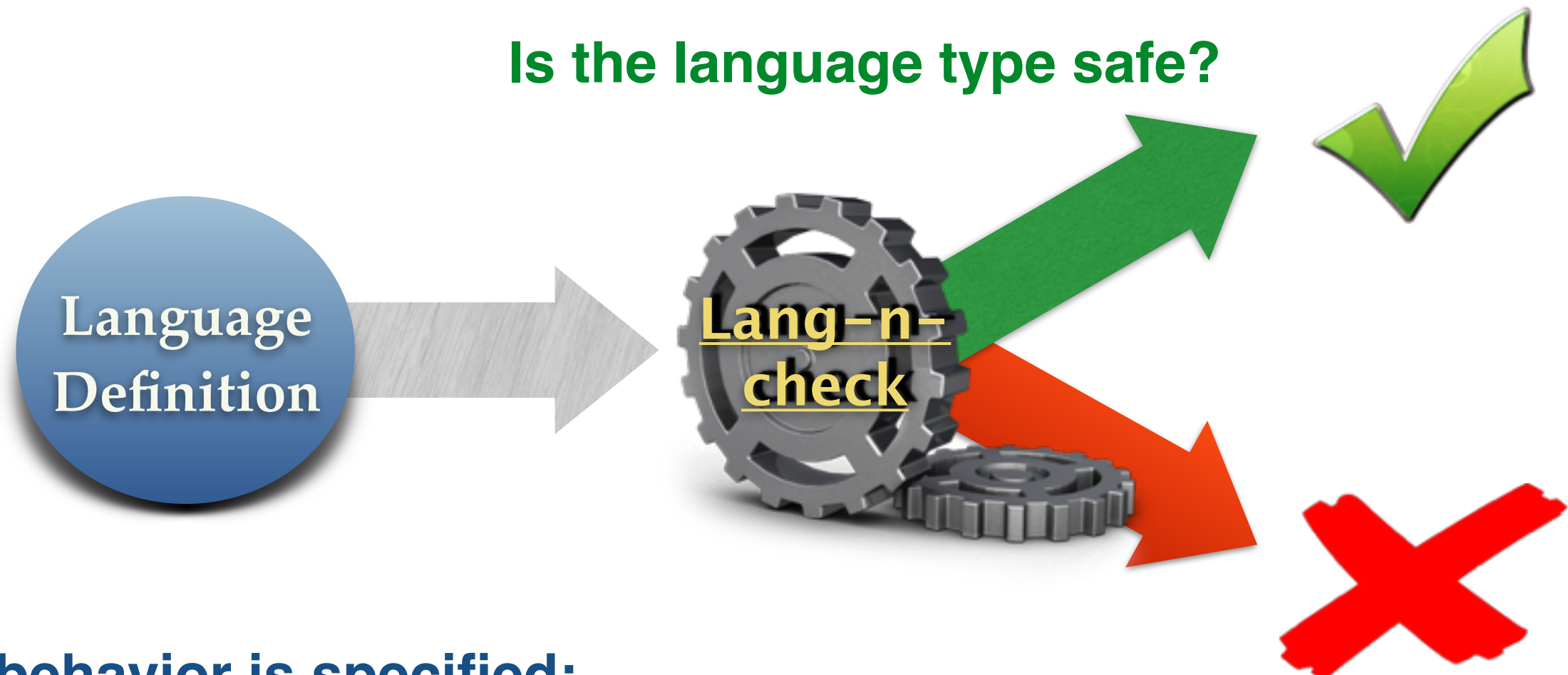
Benjamin Mourad, Matteo Cimini. *System Description: Lang-n-Change - A Tool for Transforming Languages.* FLOPS 2020.

Benjamin Mourad, M. Cimini. *A Calculus for Language Transformations.* SOFSEM 2020.



# Safety: Lang-n-check

(A type checker for language definitions)



**all behavior is specified:**

**reject** if `eval( n div 0 )` is not specified

**all behavior is consistent:**

**reject** `int_to_string(4/2) ==> int_to_string(4) / int_to_string(2)`

Matteo Cimini, Dale Miller, Jeremy G. Siek.

*Extrinsically Typed Operational Semantics for Functional Languages. SLE 2022.*

Matteo Cimini. *Early Experience in Teaching the Basics of Functional Language Design with a Language Type Checker. TFP 2019.*

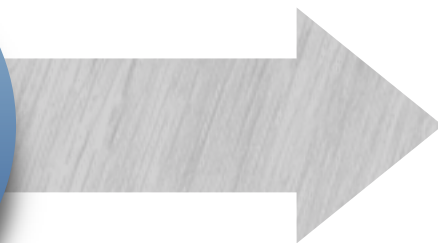
# Language Testing

small questions



what constructors bind names?  
what polymorphism?  
what are the canonical forms?  
...

Language  
Definition



Language  
Analysis



canonical forms

value	type
<i>true</i>	<i>bool</i>
<i>false</i>	<i>bool</i>
<i>function</i>	<i>arrow</i>
$[a_1, \dots, a_n]$	<i>list</i>

*helpful for debugging*

# Lang-SQL

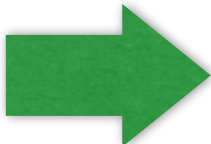
(A query language for querying languages)

*“Interrogating languages should be akin to interrogating databases”*

- **languages-as-databases.** *Stores languages as database tables.*
- **Lang-SQL,** *tailored for interrogating languages*

## canonical forms

```
SELECT GET-OPNAME(term) AS value,  
       GET-OPNAME(EXPRESSION(args)) AS type  
FROM Value, rule  
WHERE predname = type_check AND role = CONCL  
AND GET-OPNAME(term) = GET-OPNAME(OUTPUT-TYPE(args))
```



value	type
<i>true</i>	<i>bool</i>
<i>false</i>	<i>bool</i>
<i>fun</i>	<i>arrow</i>
<i>[...]</i>	<i>list</i>

Matteo Cimini. *Testing Languages with a Languages-as-Databases Approach.* TAP 2023.

Matteo Cimini. *A Declarative Validator for GSOS Languages.* PLACES 2023.

Matteo Cimini. *A Query Language for Language Analysis.* SEFM 2022.

# Lang-n-prove: a proof language for language proofs

1. Theorem *progress* :  $\forall e, T. \text{Main} : \vdash e : T \Rightarrow e \text{ is a result } \vee \exists e'. e \longrightarrow e'$ .
2. Proof
3. *induction on Main.*
4. *apply inductive hypothesis on the typing of the arg. ( $e$  of  $(\text{head } e)$  and  $(\text{fst } e)$ ).*
5. *case analysis on the progress of that argument.*
  - (a)  $\text{head } e$ 
    6. *apply canonical-form-list.*
    7. *appeal to existence of a rule for nil.*
    8. *appeal to existence of a rule for cons*
    9. *proof continues.*
  - (b)  $\text{fst } e$ 
    6. *apply canonical-form- $\times$ . (product)*
    7. *appeal to existence of a rule for  $\langle v, v \rangle$ .*
    8. *proof continues.*

## Lang-n-prove can express:

if  $e$  handles type  $t$  then

    apply canonical-form\_ $\underline{t}$ .

    for each  $v$  in valuesOf( $t$ ): QED “rule exists”.

Matteo Cimini.

*Towards the Complexity Analysis of Programming Language Proof Methods. ICTAC 2023.*

Matteo Cimini. *Lang-n-Prove: A DSL for Language Proofs. SLE 2022.*

Thankful to NSF for grant *Language-agnostic proofs*. PI: Matteo Cimini.

# Conclusion: Some of My Work on

*Empowering people to quickly deploy their own programming languages and domain-specific languages*

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**Thank you!**