



University  
of Glasgow

# Data Storage and Retrieval

## Lecture 5

### Relational Data Model: Enforcing Integrity

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# Overview

- More on the Relational Model
- Integrity
- Next lectures
  - Sets and Relations
  - Relational Algebra



# Remainder: ER -> Schema

- Strong entities
  - build a table with columns for each attribute
- Weak entities
  - build a table with columns for each attribute
  - Add the PK of the owner entity as FK; composite PK
- Relationships
  - 1-1 – FK any side
  - 1-N – FK N side
  - N-M – new table containing FKs; composite PK

# The Relational Model



# The Relational Model

- Recall: In the relational model of a database, all data is represented in terms of tuples, grouped into relations.
- A database organised in terms of the relational model is a relational database.



# A Relation

- A relation may be thought of as a 2D table

STUDENT	name	Student ID	exam1	exam2
	Gedge	891023	12	58
	Kerr	892361	66	90
	Fraser	880123	50	65

- A relation has
  - a **name** - **STUDENT**
  - In the header: an unchanging set of **columns** which are *named* and *typed (domain)* – i.e. the scheme
    - Student(name, StudentID, exam1, exam2)
  - In the body: a time varying set of **tuples**, which are the current set of **records** for the relation



# Degree and Cardinality of a Relation

## STUDENT

name	matric	exam1	exam2
Gedge	891023	12	58
Kerr	892361	66	90
Fraser	880123	50	65

- The relation student has:
  - Degree of 4 (number of attributes/columns)
  - Cardinality of 3 (number of rows/tuples)

**GOTCHA:** Do not confuse this with the cardinality of a relationship type in an E/R diagram



# Characteristics of the Relational Model

- **A relation is a set of tuples**
  - e.g. {<Fraser,880123,50,65>, ....}
- **There are no duplicate tuples**
  - because the tuples form a set
  - this must be checked when
    - a new tuple is added
    - a value is modified
    - a new relation is created as a restriction of an old one
- This implies that **a primary key always exists**
  - Worst case = a key composed of all of the attributes





# Characteristics of the Relational Model (2)

- **The tuples are unordered**
  - a property of sets
  - no meaning is imposed on the relation (we could use any order)
  - a table is only one possible representation of a relation
  - physical storage of a relation must however have an order



# Characteristics of the Relational Model (3)

- **Attributes are also an unordered set**
  - nothing can be inferred from the way we write the set of names or draw a table representation
  - no notion of getting the first attribute, next attribute, etc.
    - no first tuple or next tuple



# Characteristics of the Relational Model (4)

- **All values are atomic**

- no attribute can have a value which is either a set or is decomposable
- So NO MULTI-VALUED attributes

S#	PQ
S1	{(P1,200),(P2,300)}

must  
become

S#	P	Q
S1	P1	200
S1	P2	300

- This simplifies the model
  - called the **First Normal Form Assumption**



# Characteristics of the Relational Model (5)

- Relations represent
  - an entity type and its attributes
  - a relationship
  - a set of values
- Unknown values must be represented
  - replaced by **NULL**



# Null Values

- When adding a new tuple, we might not have all the data we need

name	matric	exam1	exam2
Gedge	891023	12	58
Kerr	892361	66	90
Fraser	880123	50	65
Smith	882854	89	NULL

- The tuple for STUDENT Smith has a NULL value for Exam2
- This is okay if the model allows....but what about if we wanted a NULL value for matric number?
  - Primary keys cannot be allowed to take NULL values

# Integrity



- **Inherent integrity constraints:**
  - must hold for all relational databases
  - typically enforced by DBMS
- **Enterprise constraints:**
  - specific to a particular application



# Integrity Constraints

- Primary key values must be unique
- Primary key values cannot be NULL
- Foreign key values:
  - must exist in the primary key of the referenced relations – we call this **referential integrity**
  - may be NULL (if it is not a mandatory participation)
- NB: In MySQL, foreign keys are constraints that need a unique name.
  - I suggest a notation including both tables: e.g. “fk\_Star\_Studio”





# Enterprise Constraints

- Application dependent
- Examples:
  - specified non-key attributes must not be NULL  
e.g.: all students must have a name, even if the primary key is the student number
  - values of one attribute must be less than values in another attribute  
e.g.: age of parent must be greater than age of child
- SQL standards compliant DBMS platforms allow user-defined constraints to be checked when data is inserted/updated



# Referential Integrity

- Concerns the use of Foreign Keys
- Guarantees that relationships between tuples are coherent
  - Every non-NULL value in a Foreign Key must also exist in the relation for which it is the Primary Key



## (Recap) Foreign Keys: Definition

- A **foreign key** is an attribute (or set of attributes) that *exist in more than one table* and which is the *primary key* for *one* of those tables.
- The foreign key is used to *cross-reference* between tables
- A foreign key is a *referential **constraint*** between two tables, i.e. a value in a foreign key **MUST** exist in the referenced primary key
- A table may have multiple foreign keys and each foreign key may reference a different table
- Foreign keys need *not* have the same attribute name across tables
  - Could be Release(MovieTitle, Year, Length)
  - Names should differ to help readability!



# Enforcing Referential Integrity

- Rules for each foreign key must be enforced independently by the database system
- Improper foreign key/primary key relationships or not enforcing those relationships are often the source of many database and data modeling problems



# Enforcing Referential Integrity

- What happens if we want to update or delete a row of a table?  
(A very likely database operation)



# Referential Integrity: 3 Strategies

- **Restrict** – ban any alterations to a primary key if there are foreign key references to it
- **Cascade** - cascade the effect to all relations in all tables that refer to it
- **Set to NULL** – allow update in original table, set all corresponding FK values to null



# Referential Integrity: Examples

- Schema:
  - Staff(Payroll# : INT, name: VARCHAR(10))
  - Students(matric# : INT, name : VARCHAR(10), adviser : INT)
- A foreign key constraint exists, where Students.adviser refers to Staff.payroll#

STAFF

<u>payroll#</u>	name
8000	Johnson
8450	Gray
8556	Lennon

STUDENTS

<u>matric#</u>	name	adviser
123	Jones	NULL
456	Smith	8556
789	Blair	8450

What is the domain of the the adviser attribute?

Adviser : { 8000, 8450, 8856, NULL } 23



# Referential Integrity: Examples

- What to do if the referent of a foreign key is changed (edited or deleted)?

STAFF

<u>payroll#</u>	name
8000	Johnson
8450	Gray
8556	Lennon

STUDENTS

<u>matric#</u>	name	adviser
123	Jones	NULL
456	Smith	8556
789	Blair	8450

## EXAMPLES:

1. What if Lennon's payroll number changes to 7990?
2. What if Gray's record is deleted?





# Referential Integrity: Example 1

- What if Lennon's payroll number changes to 7990?

STAFF

<u>payroll#</u>	name
8000	Johnson
8450	Gray
<b>8556</b>	<b>Lennon</b>

STUDENTS

<u>matric#</u>	name	adviser
123	Jones	NULL
<b>456</b>	<b>Smith</b>	<b>8556</b>
789	Blair	8450

# Referential Integrity: Example 1

- What if Lennon's payroll number changes to 7990?

## STAFF

<u>payroll#</u>	name
8000	Johnson
8450	Gray
<b>7990</b>	<b>Lennon</b>

## STUDENTS

<u>matric#</u>	name	adviser
123	Jones	NULL
<b>456</b>	<b>Smith</b>	<b>8556</b>
789	Blair	8450

referential integrity violated



# Referential Integrity: Example 1

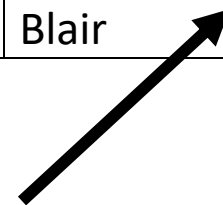
- What if Lennon's payroll number changes to 7990?

STAFF

<u>payroll#</u>	name
8000	Johnson
8450	Gray
<b>7990</b>	<b>Lennon</b>

STUDENTS

<u>matric#</u>	name	adviser
123	Jones	NULL
<b>456</b>	<b>Smith</b>	<b>7990</b>
789	Blair	8450



solution: cascade



# Referential Integrity: Example 2

- What if Gray's record is deleted?

## STAFF

<u>payroll#</u>	name
8000	Johnson
<b>8450</b>	<b>Gray</b>
7990	Lennon

## STUDENTS

<u>matric#</u>	name	adviser
123	Jones	NULL
456	Smith	7990
789	Blair	8450



# Referential Integrity: Example 2

- What if Gray's record is deleted?

STAFF

<u>payroll#</u>	name
8000	Johnson
<b>8450</b>	<b>Gray</b>
7990	Lennon

STUDENTS

<u>matric#</u>	name	adviser
123	Jones	NULL
456	Smith	7990
<b>789</b>	<b>Blair</b>	<b>8450</b>

referential integrity violated



# Referential Integrity: Example 2

- What if Gray's record is deleted?

STAFF

<u>payroll#</u>	name
8000	Johnson
<b>8450</b>	<b>Gray</b>
7990	Lennon

STUDENTS

<u>matric#</u>	name	adviser
123	Jones	NULL
456	Smith	7990
<b>789</b>	<b>Blair</b>	<b>NULL</b>

Solution: Set to NULL



# Referential Integrity

- What if Gray's record is deleted?

Error Number: 1451

Cannot delete or update a parent row: a foreign key constraint fails

STAFF

<u>payroll#</u>	name
8000	Johnson
<b>8450</b>	<b>Gray</b>
7990	Lennon

STUDENTS

<u>matric#</u>	name	adviser
123	Jones	NULL
456	Smith	7990
<b>789</b>	<b>Blair</b>	<b>NULL</b>

Solution: Restrict – raise an error, abort the transaction



# Integrity in MySQL Workbench

- Integrity – controlled by a dialogue box
  - See Foreign Keys Tab of Table Editor





# Steps for Creating a Foreign Key

**Our Schema:** `Movie(Title, Type, Studio, SequelOf)`  
`Release(MovieTitle, Year, Length)`  
where Release Title is a FK reference to Movie Title

1. Check that Movie's Title is a Primary Key
  - And note its datatype e.g. VARCHAR(50)



# Steps for Creating a Foreign Key

**Our Schema:** `Movie(Title, Type, Studio, SequelOf)`  
`Release(MovieTitle, Year, Length)`  
where Release Title is a FK reference to Movie Title

## 2. Create all attributes in Release, including those attributes that will be Foreign Keys

- Make sure that FKs attributes have correct datatype – it MUST be the SAME as the PK attribute you intend to refer to
- Press Apply BEFORE creating the FK constraints



# Steps for Creating a Foreign Key

**Our Schema:** Movie (Title, Type, Studio, Sec  
Release (MovieTitle, Year, Leng  
where Release Title is a FK reference to Movie Title

This name must be globally  
unique in your DB: "Title"  
will not work

## 3. Add a Foreign Key *constraint* for Release

- Name the constraint, e.g. "fk\_release\_movie"
  - Select Movie as the referenced table
  - Select MovieTitle as the (foreign key) column of this table, and Title as the referenced (primary key) column
- 4. Choose a referencing strategy (e.g. CASCADE)

Names...

- (i) the foreign key attribute *name*
- (ii) the referenced table & its PK attribute *name(s)*
- (iii) a *name* for the constraint itself

# Workbench Table Editor: Foreign Keys

2. Make a name for this constraint

Name: Stars

Schema: craigm\_CS1Q

Foreign Key	Referenced Table
fk_actor_studio	`craigm_CS1Q`.`Studio`
<click to edit>	

3. Select the referenced table

Foreign key details 'fk\_actor\_studio'

Column	Referenced Column
<input type="checkbox"/> Name	
<input type="checkbox"/> Address	
<input checked="" type="checkbox"/> Studio	<input checked="" type="checkbox"/> Name
<input type="checkbox"/> Salary	<input type="checkbox"/> Address
	<input type="checkbox"/> Specify Column...

On Update: NO ACTION

On Delete: NO ACTION

Comment:

☐ Skip on SQL generation

Columns

Indexes

Foreign Keys

Triggers

Partitioning

Options

Apply

Revert

1. Select Foreign keys tab

4. Select the foreign key column in this relation, and what attribute it refers to in the referenced table

5. Referencing strategy

Then click 'Apply' to save the table to the DBMS



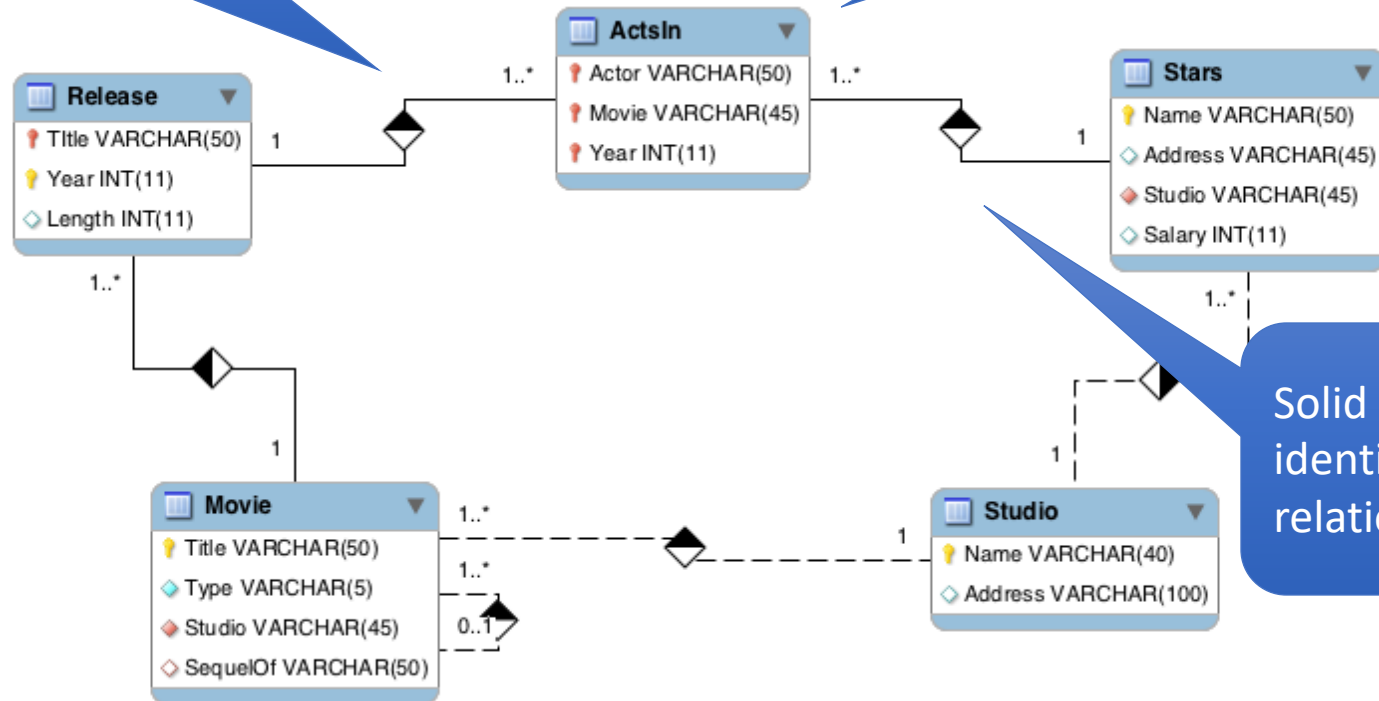
# FK Checklist

- Foreign Keys are difficult! If things go wrong:
  1. Do not panic
  2. *Read* the SQL error message, it will give a hint
  3. Check...
    - ... that you are referencing a PK
    - ... that the datatypes of the PK and FK attributes match
    - ... If you have a composite PK, you must also have a composite FK
  4. *Google* the error message
  5. If all else fails, close the Workbench tabs and try again

# An EER Diagram

Cardinality notation differs

Many-to-many relationships become Relations



Solid lines are identifying relationships

- Use the 'reverse engineer schema' functionality to view the relationships
  - NB: Looks slightly different from the ER diagram?



# Practical Tips

- Create the tables in the same order as the **Translating E-R Diagram to Relational Schema** rules
  - Create Strong entities first
  - Weak entities next
  - When creating a foreign key constraint, the referenced table & column must already exist!
- Which order to create these tables?
  - Movie(Title, Type, Studio, SequelOf) ○ 3
  - Release(Title, Year, Length) ○ 4
  - Stars(Name, Address, Studio, Salary) ○ 2
  - Studio(Name, Address) ○ 1



# Some Tips!

- Follow the stepwise guide – **it works!**
- Write a schema first – then go to access to build the tables
- Add the entities OWN attributes – then decide what FKs to add
- Use “Reverse Engineer Schema” to view the result
- Be careful to select good data types – they must **match** when you go to connect PKs and FKs
- When connecting relationships – think what is the referent PK and what is referring FK