

MATH 145

Personal Notes

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Chapter 1: Fundamentals of Set Theory

*note: this assumes knowledge of MATH 147 Chap 1 & 2.

THE SIX FUNDAMENTAL AXIOMS

EXISTENCE + EXTENSIONALITY

EXISTENCE

The Axiom of Existence states that there exists a set with no elements; for this set A , A satisfies $A \in X$ for all sets X .

EXTENSIONALITY

The Axiom of Extensionality states that if two sets have the same elements, they are equal; ie if every element in X is in Y and v.v., $X = Y$.

THE UNIQUENESS OF \emptyset

These first 2 axioms can be used to prove the empty set, \emptyset , is unique.

Proof. By AoExi, \emptyset exists, so we only need prove its uniqueness.

Now, suppose $\exists X_1, X_2$ such that X_1 & X_2 have no elements.

Then, every element of X_1 is in X_2 and v.v., since they both do not have any elements.
 $\therefore X_1 = X_2$ by AoExt, and we are done.

COMPREHENSION

The Axiom Schema of Comprehension states that if $P(X)$ is a property of a set X ($P(X)$ is a statement), then for any set A , there exists a set B such that $X \in B$ if and only if $X \in A$, & $\forall X P(X)$ is true;

ie

$$\forall P(X), A : \exists B = \{X \in A : P(X)\}.$$

*note: A, B are sets of sets!

*AoC is referred to as a "schema" as it actually is a large collection of axioms (1 for every possible $P(X)$).

Furthermore, we can show this set B is unique.

Proof. The AoC shows B exists, so we need only prove it is unique.

Suppose $\exists B_1, B_2$ such that $B_1 = \{X \in A : P(X)\}$ and $B_2 = \{X \in A : P(X)\}$ for some $A, P(X)$.

Clearly, if $X \in B_1$, $X \in B_2$ also, and v.v.

By the AoExt, this implies $B_1 = B_2$.

PAIR, UNION & THE POWER SET

We need to define other axioms to give us sets other than \emptyset . (The first 3 only tell us \emptyset exists!)

*again, C is a set of sets!

PAIR

The Axiom of Pair states that, given any sets A & B , there exists a set C whose elements are exactly A and B ; ie $\forall X : (X=A \vee X=B) \Rightarrow X \in C$.

We can also show C is unique, using the AoExt. Denote $\{A, B\}$ as the set containing A & B , and the shorthand $\{\{A\}\}$ for the set $\{\{A\}\}$.

In combination with the previous 3 axioms, we can create 2 new sets:

① Let $A=B=\emptyset$. Then, the AoP tells us

$\{\emptyset, \emptyset\}$, or $\{\{\emptyset\}\}$, exists.

② Now, let $A=\emptyset$ & $B=\{\emptyset\}$. Applying AoP once again, we can also deduce $\{\emptyset, \{\emptyset\}\}$ exists.

UNION

The Axiom of Union states that, for any set S , there exists a set U such that, for any set X , $X \in U$ if and only if $X \in A$ for some set $A \in S$; ie

$$\forall S : \exists U \Rightarrow [\forall X : X \in U \Leftrightarrow \exists A \in S : X \in A].$$

Again, using the AoExt, we can show U is unique. Denote $\cup S$ to represent the union of elements in S , and $A \cup B$ as shorthand for $\cup \{A, B\}$.

We can use the AoP to create larger sets from smaller ones.

POWER SET

The Axiom of Power Set states for any set A , there exists a set P such that for any set X , $X \in P$ if and only if $X \subseteq A$;

ie $\forall A : \exists P \ni [\forall X : X \in P \Leftrightarrow X \subseteq A]$.

* P can also be written as $\mathcal{P}(A)$.

We can also show $P(X)$ is unique from the preceding axioms. Denote $\mathcal{P}(A)$ as the power set of A .

Examples:

$$① P(\emptyset) = \emptyset$$

$$② P(\{\emptyset\}) = \{\emptyset, \{\emptyset\}\}$$

$$③ P(\{x_1, x_2\}) = \{\emptyset, \{x_1\}, \{x_2\}, \{x_1, x_2\}\}.$$

*note:
both \emptyset & A are always elements of $\mathcal{P}(A)$.

OTHER SET CONSTRUCTIONS

SET DIFFERENCE

1: We do not need any axioms to describe the intersection of two sets.

2: Given two sets $A \& B$, we use the AoC, with $P(X) = "X \in B"$. We then define the output set as the intersection between $A \& B$; ie $A \cap B = \{x \in A : x \in B\}$.

1: Similarly, we can define $A \setminus B$ by letting $P(X) = "X \notin B"$ in the AoC, resulting in

$$A \setminus B = \{x \in A \mid x \notin B\}.$$

CONSTRUCTION OF \mathbb{N}

1: We can use the sets we have defined to assign each natural number (and zero) to a set.

2: In particular, we would like $n \in \mathbb{N}$ to correspond to a set with n elements.

ZERO

1: Since \emptyset is the only set with no elements, we therefore conclude that $0 = \emptyset$.

ONE

1: Although there are many sets with one element, we define 1 as the set $\{\emptyset\}$; ie $1 = \{\emptyset\} = \{\{\emptyset\}\}$.

THE SUCCESSOR, & DEFINING SUCCESSIVE ELEMENTS

1: The "successor" of any set x is defined to be $x \cup \{x\}$.

2: We can use this to define a method to obtain the next natural number; given any $n \in \mathbb{N}$, we define $n+1 = n \cup \{n\}$.

ORDERED PAIRS

1: Given any two sets $X \& Y$, we define the ordered pair (X, Y) to be $(X, Y) = \{\{X\}, \{X, Y\}\}$.

2: We can use the AoExt to show that (X, Y) is unique; ie if $X = X'$ & $Y = Y'$, then $(X, Y) = (X', Y')$, and vice versa.

Proof Case 1: $X = Y$.

Then, $(X, Y) = \{\{X\}\} = \{\{X\}\}$.

Similarly, $(X', Y') = \{\{X'\}\} = \{\{X\}\}$. (proved).

Case 2: $X \neq Y$.

Then, $\{X, Y\} \neq \{X, Y'\}$, as $\{X, Y\} \neq \{X, Y'\}$.

Hence, $\{\{X\}\} \neq \{\{X'\}\}$, so $X = X'$.

So $\{\{X', Y'\}\} = \{\{X, Y'\}\} = \{\{X, Y\}\}$; $\therefore Y = Y'$. (proved).

3: We can use this definition of an ordered pair to create ordered n -tuples for any n :

$(x_1, x_2, \dots, x_n) = ((x_1, x_2, \dots, x_{n-1}), x_n)$.
(Where a 1-tuple is simply the set $\{x_1\}$.)

INDUCTIVE SETS

1: A set I is inductive if

- ① $0 \in I$; and
- ② If $n \in I$, then $S(n) = n+1 \in I$.

THE AXIOM OF INFINITY

1: The Axiom of Infinity states that an inductive set exists.

2: We can then define \mathbb{N} using AoI & AoC, resulting in

$$\mathbb{N} = \{x \in I : x \in I \ \forall I\},$$

where I is the inductive set defined by AoI, and I is an inductive set.

3: We can also show \mathbb{N} is unique from the AoExt.

4: Next, we can show \mathbb{N} is inductive.

Proof: Since $0 \in I$, where I is an inductive set, $0 \in \mathbb{N}$ also, based on its definition.

Then, suppose $n \in \mathbb{N}$. By defn, this means $n \in I$ for all inductive sets I . For each I , $(n+1)$ must also be it (by defn). Hence $(n+1) \in \mathbb{N}$ also.

This is enough to prove the claim. \square

CARTESIAN PRODUCTS

1: Given any two sets $X \& Y$, we define $X \times Y$ to be

$$X \times Y = \{w \in P(P(X \cup Y)) : w = (x, y), x \in X, y \in Y\}.$$

Proof: We know $X \times Y = \{w = (x, y) \in Z : x \in X, y \in Y\}$. But we need to find a set Z that contains (x, y) !

First, $(x, y) = \{\{x\}, \{x, y\}\}$.

Observe $\{\{x\}, \{x, y\}\} \subseteq (X \cup Y)$; so $\{\{x\}, \{x, y\}\} \in P(X \cup Y)$.

Similarly, $\{\{y\}, \{x, y\}\} \subseteq P(X \cup Y)$; hence,

$$(x, y) \in P(P(X \cup Y)) \quad \forall x \in X, y \in Y.$$

So, our set $Z = P(P(X \cup Y))$.

METHOD OF ORDERING \mathbb{N}

1: Let $m, n \in \mathbb{N}$. We say $m < n$ (ie m is less than n) if $m < n$. * note: this is an "order relation" on a set.

Chapter 2: Relations & Functions

RELATIONS

Given two sets X & Y , a binary relation from X to Y is a subset of $X \times Y$.

More generally, a set R is called a relation if all the elements of R are ordered pairs.

TERMINOLOGY

DOMAIN OF R

The domain of a binary relation R , or $\text{dom}(R)$, is the set of all x for which $(x,y) \in R$, for some y .

RANGE OF R

The range of a binary relation R , or $\text{ran}(R)$, is the set of all y for which $(x,y) \in R$, for some x .

FIELD OF R

The field of a binary relation R , or $\text{field}(R)$, is defined to be $\text{field}(R) = \text{dom}(R) \cup \text{ran}(R)$.

IMAGE OF A SET UNDER R

The image of a set A under a binary relation R is the set

$$R(A) = \{b \in \text{ran}(R) : (a,b) \in R, a \in A\}.$$

INVERSE IMAGE OF A SET UNDER R

The inverse image of a set B under a binary relation R is the set

$$R^{-1}(B) = \{a \in \text{dom}(R) : (a,b) \in R, b \in B\}.$$

INVERSE OF R

The inverse relation R^{-1} of a binary relation R is defined by $R^{-1} = \{z \in \text{ran}(R) \times \text{dom}(R) : z = (b,a), (a,b) \in R\}$.

COMPOSITION OF R_1 & R_2

Given two relations R_1 & R_2 , the composition of R_1 & R_2 , or $R_2 \circ R_1$, is defined by

$$R_2 \circ R_1 = \{z \in \text{dom}(R_1) \times \text{ran}(R_2) : z \in (a,c), \text{ where } \exists b \ni (a,b) \in R_1 \wedge (b,c) \in R_2\}.$$

FUNCTIONS

A function f is a relation such that if $(a,b_1), (a,b_2) \in f$, then $b_1 = b_2$; ie for any first coordinate in f , there is only one ordered pair with that first coordinate.

For any two sets A & B , f is a function from A to B if each $a \in A$ is related to exactly one element $b \in B$. We use the notation $f: A \rightarrow B$ to represent this.

INJECTIVITY OF f (1-1)

A function f is injective (ie 1-1) if $\forall x_1, x_2 \in \text{dom}(f), x_1 \neq x_2 : f(x_1) \neq f(x_2)$.

SURJECTIVITY OF f (ONTO)

A function f is surjective (ie onto) if $\text{ran}(f) = B$; ie, $\forall b \in B, \exists a \in A : f(a) = b$.

BIJECTIVITY OF f (1-1 & ONTO)

A function f is bijective if it is both injective & surjective.

INVERTIBILITY OF f

A function f is invertible if f^{-1} exists.

We can prove f is injective if and only if f is invertible.

EQUIVALENCE RELATIONS

A relation R on a set A is:

- ① reflexive if $\forall a \in A, aRa$;
- ② symmetric if $\forall a, b \in A, aRb \Leftrightarrow bRa$;
- ③ transitive if $\forall a, b, c \in A, aRb \wedge bRc \Leftrightarrow aRc$.

\therefore A relation is an equivalence relation if it is reflexive, symmetric & transitive.

EQUIVALENCE CLASSES

Let E be an equivalence relation on a set A . Given an element $a \in A$, the equivalence class of a modulo E is the set

$$[a]_E = \{x \in A : aEx\}.$$

\therefore We can prove that:

- 1) aEb if and only if $[a]_E = [b]_E$; and
- 2) $\neg aEb$ if and only if $[a]_E \cap [b]_E = \emptyset$.

(Proof: MATH 147 WAI Q2a)

PARTITIONS

Given any set A , a partition P of A is a collection of non-empty sets that satisfy the following:

$$\text{① } P_1 \cap P_2 = \emptyset \quad \forall P_1, P_2 \in P; \quad \&$$

$$\text{② } \bigcup P = A.$$

\therefore Every equivalence relation E on A gives a partition of A . (Proof from above.)

$$\rightarrow \text{Denote } A/E = \{[a]_E : a \in A\}.$$

\therefore We can prove A/E is a partition of A .

Proof. $[a]_E \neq \emptyset$ since $a \in [a]$. ($a \in Ea$).

Then $\forall a, b \in A, [a]_E \cap [b]_E = \emptyset$; ie they are disjoint. ①

Next, $\forall a \in A, a \in [a]$. (since E is reflexive).

Hence $\bigcup (A/E) = A$. ②

① & ② are sufficient to prove the claim. \square

\therefore Then, let P be a partition on a set A .

Let the relation E be such that if

$\exists P \in P$ such that $a_1 \in P$ & $a_2 \in P$, then $a_1 E a_2$.

We can prove E is an equivalence relation.

Proof:

① For any $a \in A$, $\exists P \in P$ containing a , as P is a partition. Hence $a \in E a$; thus E is reflexive.

② For any $a_1, a_2 \in A$ such that $a_1 E a_2$, it must be that $\exists P \in P \ni a_1 \in P$ & $a_2 \in P$. Thus $a_2 \in P$ and $a_1 \in P$, implying $a_1 E a_2$, also. Hence E is symmetric.

③ Let $a_1, a_2, a_3 \in A$ such that $a_1 E a_2$ and $a_2 E a_3$. Suppose $a_1 \in P_1$ & $a_3 \in P_2$, where $P_1, P_2 \in P$ & $P_1 \neq P_2$. By definition, $a_2 \in P_1$ and $a_2 \in P_2$. However, this implies $P_1 \cap P_2 \neq \emptyset$, which is a contradiction as $P_1 \neq P_2$. Thus $P_1 = P_2$, implying $a_3 \in P_1$, and so $a_1 E a_3$. Hence E is transitive.

④ As E is symmetric, transitive & reflexive, E is an equivalence relation — and so we are done. \square

SET OF REPRESENTATIVES

Let E be an equivalence relation on a set A . A set X is a set of representatives for E if X contains exactly one element of each equivalence class; ie $\forall [a] \in A/E, X \cap [a] = \{a\}$ for some $a \in [a]$.

ORDER RELATIONS

PARTIAL ORDERING

A relation R is "antisymmetric" if aRb & bRa implies $a=b$.

A relation " \leq " on a set A is an order relation on A if it is reflexive, antisymmetric and transitive.

We also call " \leq " a partial ordering on A .

eg the identity relation R on A , where

$$R = \{(a, a) : a \in A\}.$$

- R is reflexive;

- R is transitive; $(aRb, bRc \Rightarrow aRc)$ & $aRa \Rightarrow aRc$)

- R is antisymmetric. $(aRb, bRa \Rightarrow a=b)$

A relation R on a set A is "asymmetric" if aRb and bRa cannot be simultaneously true.

For any partial ordering " \leq " on A , we can define a strict ordering " $<$ " from it by declaring that $\forall a, b \in A, a < b$ if $a \leq b$ but $a \neq b$.

TOTAL/LINEAR ORDERING

For any partial ordering " \leq " on A , we say

$a, b \in A$ are comparable if either $a \leq b$ or $b \leq a$.

A total/linear ordering is a partial ordering in which every pair of elements are comparable.

CHAIN

For any partial ordering " \leq " on A , a set $C \subseteq A$ is called a chain if every pair of elements in C are comparable.

LEAST/GREATEST ELEMENT

Let A be a set with a partial ordering " \leq ". Let $B \subseteq A$. Then, an element $b \in B$ is a least element of B if $b \leq b' \quad \forall b' \in B$.

Similarly, b is a greatest element of B if $b' \leq b \quad \forall b' \in B$.

MINIMAL / MAXIMAL ELEMENT

Let A be a set with partial ordering " \leq ".

Let $B \subseteq A$. Then, an element $b \in B$ is a minimal element of B if there are no "smaller" elements of B ;

ie if $b' \leq b$ for some $b' \in B$, then $b = b'$.

Similarly, b is a maximal element of B if there are no "larger" elements of B ;

ie if $b \leq b'$ for some $b' \in B$, then $b = b'$.

Note: a least element is always minimal, but a minimal element may not be a least element.

eg for $a, b \in \mathbb{Z}^+$: let $a \leq b$ if $\frac{b}{a} \in \mathbb{Z}^+$

Then 1 is a least element, as it divides every tve integer.

However, $\mathbb{Z}^+ \setminus \{1\}$ no longer has a least element; there is no other integer that divides both 2 & 3.

But there are infinitely many minimal elements; every prime p is minimal, since it is not divisible by any other tve integers other than 1 and itself.

BOUNDS ON SETS

Suppose A is a set with order relation " \leq ".

Then, $a \in A$ is a lower bound on $B \subseteq A$ if $a \leq b \quad \forall b \in B$.

$b \in A$ is an upper bound on B .

Similarly, $b \in A$ is the greatest lower bound / infimum if it is the greatest possible lower bound.

Then, an element $a \in A$ is the greatest upper bound / supremum if it is the lowest possible upper bound.

Similarly, an element $b \in A$ is the lowest upper bound / supremum if it is the lowest possible upper bound.

Chapter 3: Fundamentals of Set Theory II

THE AXIOM OF CHOICE

B₁: Let C be a collection of sets, where $C \neq \emptyset$. Then, the Cartesian product $\prod_{c \in C} C$ is the set of functions $\alpha: C \rightarrow \bigcup C$, with the property that $\forall c \in C, \alpha(c) \in C$.

* note: this defn works for both finite and infinite collections of sets.

B₂: The Axiom of Choice states every Cartesian product of any non-empty collection of sets is non-empty.

ZORN'S LEMMA

B₁: ZL states that for any partially ordered set A , with order relation \leq , if every chain in A has an upper bound in A , then A has a maximal element.

B₂: We can prove Zorn's Lemma is logically equivalent to the Axiom of Choice.

WELL-ORDERING THEOREM

B₁: A set A with order relation " \leq " is well-ordered if every non-empty subset of A has a least element.

B₂: Then, the WOT states every non-empty set A has a well-ordering; ie there exists an order relation \leq on A such that A is well-ordered (with respect to \leq .)

B₃: Again, WOP is also logically equivalent to the Axiom of Choice.

CARDINALITY

E1: Two sets A & B have the same cardinality if there exists a bijection $f: A \rightarrow B$, and write $|A| = |B|$.

E2: Then:

- ① $|A| = |A| \quad \forall A$;
- ② $|A| = |B| \Leftrightarrow |B| = |A| \quad \forall A, B$; &
- ③ $|A| = |B| \wedge |B| = |C| \Rightarrow |A| = |C| \quad \forall A, B, C$.

*shown in
A & Q2.

"≤" ON CARDINALITY

E1: We say $|A| \leq |B|$ if there is an injective function $f: A \rightarrow B$.

E2: Then, we can prove for any sets A_1, B, A : if $A_1 \subseteq B \subseteq A$, then $|A_1| \leq |A|$ implies $|B| \leq |A|$.

Proof. Since $|A_1| = |A|$, $f: A_1 \rightarrow A$ is a bijection.

Then, let $\{A_n\}$, $\{B_n\}$ be such that $A_0 = A$ & $B_0 = B$, and $A_{n+1} = f(A_n)$ & $B_{n+1} = f(B_n) \quad \forall n \in \mathbb{N}$.

We can use induction to prove $A_{n+1} \subseteq A_n \quad \forall n \in \mathbb{N}$.

Next, set $C_n = A_n \setminus B_n$ for some n .

Let $C = \bigcup_{n=0}^{\infty} C_n$. Then, if $a \in f(C_n)$, then $a = f(c)$ for some $c \in C_n$. By defn, $c \in A_n$ & $c \notin B_n$,

implying $f(c) \in f(A_n) = A_{n+1}$.

Similarly, $f(c) \notin f(B_n)$, as if $f(c) = f(b)$ for some $b \in B_n$, then necessarily $c = b$ (as f is injective) and so $c \in B_n$, which is a contradiction.

We now know $f(c) \in A_{n+1} \setminus B_{n+1} = C_{n+1}$; hence $f(C_n) \subseteq C_{n+1}$. On the other hand, if $a \in C_{n+1}$, then $a \in A_{n+1} \wedge a \notin B_{n+1}$, and so $a = f(a')$ for some $a' \in A_n$. Since $a \notin B_{n+1}$, we also know $a' \notin B_n$. Thus $a' \in C_n$, so that $a = f(a') \in f(C_n)$.

Therefore $C_{n+1} \subseteq f(C_n)$, which implies $C_{n+1} = f(C_n)$.

Subsequently, if $a \in f(C)$, then $a = f(c)$ for some $c \in C$.

Then $c \in C_n$ for some $n \in \mathbb{N}$, implying $a = f(c) \in f(C_n) = C_{n+1}$, proving $a \in \bigcup_{n=1}^{\infty} C_n$.

Conversely, if $a \in \bigcup_{n=1}^{\infty} C_n$, then $a \in C_n$ for some $n \in \mathbb{N}$.

In particular, since $C_n = f(C_{n-1})$, $a \in f(C_{n-1})$; thus

$a = f(c)$ for some $c \in C_{n-1}$, implying $a \in f(C)$.

Therefore, since $f(C)$ and $\bigcup_{n=1}^{\infty} C_n$ have the same elements,

by extensionality $f(C) = \bigcup_{n=1}^{\infty} C_n$.

From here, we can show " \leq " behaves like an order relation; ie

- ① $\forall A, B, C: |A| \leq |B| \wedge |B| \leq |C| \Rightarrow |A| \leq |C|$; ie \leq is "transitive".
- ② (Cantor-Schröder-Bernstein Theorem)
 $\forall A, B: \text{if } |A| \leq |B| \wedge |B| \leq |A|, \text{ then } |A| = |B|$. ? "antisymmetric".

② (Cantor-Schröder-Bernstein Theorem)

$\forall A, B: \text{if } |A| \leq |B| \wedge |B| \leq |A|, \text{ then } |A| = |B|$. ? "antisymmetric".

Proof. All the results of ① can be proved

by noticing (gof) is injective if both f & g are injective.

We now prove ②.

Suppose we have injective functions $f: A \rightarrow B$ and $g: B \rightarrow A$, for some sets A & B . Then $(gof): A \rightarrow A$ is also injective.

Then, let $X = g(B)$ and $Y = g(f(A))$.

Clearly $X \subseteq A$, and since $f(A) \subseteq B$ we get that $Y = g(f(A)) \subseteq g(B) = X$.

Additionally, since (gof) is injective, it is also a bijective function from A to $(gof)(A) = Y$. Hence $|Y| = |A|$.

Therefore, $Y \subseteq X \subseteq A$. By the previous prof, $|Y| = |A|$ implies $|X| = |A|$.

But since $X = g(B)$, and g is injective, we must have $g: B \rightarrow X$ is a bijection.

So $|X| = |B|$, implying $|A| = |B|$, which we wanted to show. \blacksquare

Finally, let $D = A \setminus C$. Define $g: A \rightarrow B$ by

$$g(x) = \begin{cases} f(x), & x \in C \\ x, & x \in D \end{cases}$$

If $x \in C$, then $f(x) \in C_n$ for some $n \geq 1$,

implying $f(x) \in A_n$.

Moreover, since $A_0 \supseteq A_1 \supseteq A_2 \dots$, we can deduce

$f(x) \in A_1$. We also know $A_1 \subseteq B$ by construction, consequently, $f(x) \in B$.

Similarly, if $x \in D$, then $x \notin C$. This implies $x \notin C_0$,

and hence $x \notin A \setminus B$, so $\neg x \in B$ (ie $x \in B$.)

Therefore, $\forall x \in A: g(x) \in B$. (So $g: A \rightarrow B$.)

Next, suppose $x_1, x_2 \in A$ such that $g(x_1) = g(x_2)$.

If $x_1, x_2 \in C$, then $f(x_1) = f(x_2)$, implying $x_1 = x_2$ since f is injective.

If $x_1, x_2 \in D$, then clearly $x_1 = x_2$ also.

If $x_1 \in C$ & $x_2 \in D$, then $f(x_1) = x_2$; however $f(x_1) \in C$

whilst $x_2 \in D$, so this is a contradiction.

A similar argument also shows $x_1 \in D$ & $x_2 \in C$ leads to a contradiction too.

Therefore, $x_1 = x_2$, proving g is injective.

Similarly, let $b \in B$ be arbitrary.

If $b \in f(C)$, then $\exists c \in C$ such that $f(c) = b$, so $g(c) = f(c) = b$.

If $b \notin f(C)$, then either $b \in C_0$ or $b \in D$. If $b \in D$, $g(c) = b$ trivially.

and if $b \in C_0$, then $b \in A \setminus B$, contradicting the assumption that $b \in B$.

Therefore $\forall b \in B, \exists c \in C$ such that $g(c) = b$. This is sufficient to show g is surjective.

Since g is both injective & surjective, it is also bijective.

Therefore $|A| = |B|$, and we are done. \blacksquare

E4: We write $|A| < |B|$ to signify $|A| \leq |B|$ but $|A| \neq |B|$.

E5: Note that for all sets $B \neq \emptyset$, $|\emptyset| < |B|$.

Proof. Note that the empty relation from \emptyset to B is both vacuously a function and injective, but not surjective (and so not bijective.) This is sufficient to prove the result. \blacksquare

FINITE / INFINITE SETS

A set A is finite if it has the same cardinality as some $n \in \mathbb{N}$.

In this case, we write $|A|=n$.

We can prove $\forall n \in \mathbb{N}$, there is no injective mapping from n to $X \subset \mathbb{N}$.

Proof. Suppose such a mapping exists.

Then by WOP, there exists a least element $n \in \mathbb{N}$ which satisfies this property.

Clearly $n \neq 0$, since there are no proper subsets of 0. Then, either $(n-1) \in X$ or $(n-1) \notin X$.

If $(n-1) \notin X$, then $X \subseteq (n-1)$.

Subsequently, if $f: n \rightarrow X$ is injective, we can define a new injective mapping $g: (n-1) \rightarrow X \setminus \{(n-1)\}$ by $g(b) = f(a) \quad \forall b \in (n-1)$. * g is the "restriction" of f to the set $n-1$.

Then g is injective as f is, and g maps $(n-1)$ to a proper subset of $(n-1)$, since $X \subseteq (n-1)$.

This contradicts the minimality of n .

On the other hand, if $(n-1) \in X$, we know $(n-1)=f(b)$ for some unique $b \in n$ (since f is injective.)

Then, let $g: (n-1) \rightarrow X \setminus \{(n-1)\}$ by

$$g(i) = \begin{cases} f(i), & i \neq k \\ f(n-1), & i = k. \end{cases}$$

Again, g is injective, and it maps $(n-1)$ to a proper subset of itself, which contradicts the minimality of n .

Thus no such n exists, proving our claim. \square

UNIQUENESS OF CARDINALITY

For any finite set A , if $|A|=m$ & $|A|=n$, then $m=n$.

Proof. If $m \neq n$, then either $m < n$ or $n < m$. However, by the above proof, both cases lead to contradiction.

Hence necessarily $m=n$. \square

\mathbb{N} IS INFINITE

We can prove \mathbb{N} is infinite, ie not finite.

Proof. Consider $d: \mathbb{N} \rightarrow \mathbb{N}$ by $d(n)=2n$.

Clearly, d is a bijection, but the set of even numbers is a proper subset of all naturals!

Hence \mathbb{N} cannot be finite, as it would lead to a contradiction otherwise. \square

SUBSETS OF A FINITE SET ARE ALSO FINITE

Let A be a finite set, and $B \subseteq A$. Then $|B| \leq |A|$, and B is finite also.

Proof. First, let the "inclusion mapping" $i: B \rightarrow A$ given by $i(b) = b \quad \forall b \in B$.

Clearly this mapping is injective, showing $|B| \leq |A|$.

Then, let $|A|=n$ for some $n \in \mathbb{N}$.

If $n=0$, $|B|=0$ necessarily, implying $B=\emptyset$.

We consider what happens if $n \geq 1$.

First, let $A = \{a_0, a_1, \dots, a_{n-1}\}$.

Then, if $B=\emptyset$, we are done. Otherwise,

there is a least index i for which $a_i \in B$

Call this element b_0 .

Again, if $B=\{b_0\}$, we are done. Otherwise,

$\exists i > i$ such that $a_i \in B$, and call this element b_1 .

Then, if $B=\{b_0, b_1\}$, we are done.

Otherwise, we repeat this procedure until we get B .

Since A is finite, this procedure cannot go on indefinitely; thus $\exists m \in \mathbb{N}$ such that $B=\{b_0, b_1, \dots, b_{m-1}\}$.

This is sufficient to prove B is finite. \square

THE IMAGE OF A FINITE SET UNDER A FUNCTION IS ALWAYS FINITE

Let A & B be sets, such that A is finite. Suppose $f: A \rightarrow B$ is a function.

Then $f(A)$ is a finite subset of B , & $|f(A)| \leq |A|$.

Proof. Assume $|A| \geq 1$, since the proof is trivial if $A=\emptyset$.

Then $A=\{a_0, a_1, \dots, a_{n-1}\}$. So $f(A)=\{f(a_0), f(a_1), \dots, f(a_{n-1})\}$.

Subsequently, let $b_0=f(a_0)$, and then find the least index $i > 0$ for which $f(a_i) \neq f(a_0)$ (if it exists) and set $b_1=f(a_i)$.

Again, if $\{b_0, b_1\} \neq f(A)$, then we continue this process until it does, since we know $0 < i_1, i_2, \dots$ cannot surpass $n-1$.

Hence $f(A)=\{b_0, b_1, \dots, b_{m-1}\}$ for some $m \in \mathbb{N}$, proving

the former claim. *

Then, let the function $g: f(A) \rightarrow A$ by $g(b_0)=a_0$ & $g(b_k)=a_k \quad \forall k > 0$.

Clearly g is injective; therefore $|f(A)| \leq |A|$. \square

THE UNION OF A FINITE COLLECTION OF SETS IS ALSO FINITE

Let A & B be finite sets. Then $|A \cup B| \leq |A|+|B|$, with $|A \cup B| = |A|+|B|$ iff $A \cap B = \emptyset$.

Proof. Let $|A|=m$ and $|B|=n$, ie $A=\{a_0, a_1, \dots, a_{m-1}\}$ & $B=\{b_0, b_1, \dots, b_{n-1}\}$.

Then $A \cup B = \{a_0, a_1, \dots, a_{m-1}, b_0, b_1, \dots, b_{n-1}\}$, implying

$(A \cup B)$ has at most $m+n$ elements.

However, since some elements might be repeated, $|A \cup B|$ might be less than $(m+n)$; hence $|A \cup B| \leq (m+n) = |A|+|B|$. \square

For any finite collection of sets S , $\bigcup S$ is also finite.

Proof. If $S=\emptyset$, $\bigcup S=\emptyset$, establishing our base case.

Then, assume for some $S=\{A_0, A_1, \dots, A_{n-1}\}$ the claim holds.

Subsequently, $A_n \cup (\bigcup S)$ is also finite by the above proof, implying the claim also holds for $S=\{A_0, A_1, \dots, A_n\}$.

By induction, this is sufficient to prove the claim. *

THE POWER SET OF A FINITE SET IS ALSO FINITE

For any finite set A , $P(A)$ is also finite.

Proof. If $A=\emptyset$ (ie $|A|=0$), then $P(A)=\{\emptyset\}$, which is finite.

Then, assume for some set $A=\{a_0, a_1, \dots, a_{n-1}\}$ for which $|P(A)|$ is finite.

Subsequently, let $A' = A \cup \{a_n\}$.

Then $P(A') = B \cup C$, where $B = \{B \in P(A') : a_n \in B\}$ & $C = \{C \in P(A') : a_n \notin C\}$. Clearly $B \cap C = \emptyset$.

Subsequently, observe $|B| = |C| = |P(A')|$.

Since $|P(A')|$ is finite, necessarily $|B|$ & $|C|$ also are.

Thus $|P(A')| = |B| + |C|$ is also finite.

By induction, this is sufficient to prove the claim. \square

COUNTABLE SETS

\exists_1 A set A is "countable" if $|A| = |\mathbb{N}|$.

\exists_2 A set A is "at most countable" if $|A| \leq |\mathbb{N}|$. this implies they are either
1) finite; or
2) countable.

\exists_3 For any countable set A , there exists a bijection $f: \mathbb{N} \rightarrow A$; ie we can write out the elements of A as an infinite sequence;

$$A = \{a_0, a_1, a_2, \dots\}$$

EVERY SUBSET OF A COUNTABLE SET IS AT MOST COUNTABLE

\exists_1 Let $B \subseteq A$. If A is countable, then B is either finite or countable.

Proof: If B is finite we have our conclusion, so suppose B is infinite.

Let k_0 be the smallest index k_0 such that $a_{k_0} \in B$, and set $b_0 = a_{k_0}$.

Then, let k_1 be the smallest index larger than $a_{k_0} \in B$, setting $b_1 = a_{k_1}$.

Subsequently, we continue this process for each $i \in \mathbb{N}$, as since B is infinite,

$$B \setminus \{b_0, b_1, \dots, b_{i-1}\} \neq \emptyset.$$

Hence $B = \{b_0, b_1, \dots\} = \{a_{k_0}, a_{k_1}, \dots\}$ so B is a subsequence of A , and we are done. \star

THE CARTESIAN PRODUCT OF COUNTABLE SETS IS COUNTABLE

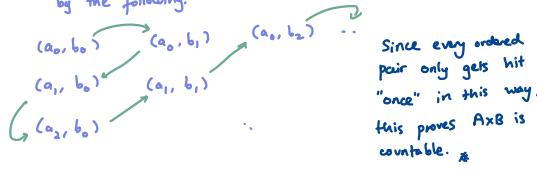
\exists_1 Let A and B be countable sets. Then necessarily $A \times B$ is also countable.

Proof: Since $|A| = |\mathbb{N}| = |B|$, we can represent them as infinite sequences;

$$\text{ie } A = \{a_0, a_1, a_2, \dots\} \text{ and } B = \{b_0, b_1, b_2, \dots\}$$

$$\text{Then } (a_i, b_j) \in (A \times B) \quad \forall i, j \in \mathbb{N}.$$

We can then define a bijection from \mathbb{N} to $A \times B$ by the following:



\exists_2 Let A_0, A_1, \dots, A_n be finitely many countable sets. Then $\prod_{i=0}^n A_i$ is also countable.

Proof: When $n=1$, this is the result above.

Suppose the claim is true for some $n=N$.

Then clearly

$$\prod_{i=0}^{N+1} A_i = \underbrace{\prod_{i=0}^N A_i}_{\text{countable}} \times \underbrace{A_{N+1}}_{\text{countable}}.$$

Hence $\prod_{i=0}^{N+1} A_i$ is also countable, proving the claim for $n=N+1$.

(By induction, this is sufficient to prove the claim.) \star

\mathbb{Z} IS COUNTABLE

\exists_1 We can show that $|\mathbb{Z}| = |\mathbb{N}|$.

Proof: Let $a_k = \begin{cases} \frac{k}{2}, & k \text{ is even} \\ (\frac{k+1}{2})/2, & k \text{ is odd} \end{cases}$

Then clearly $\{a_k\} = \{0, -1, 1, -2, 2, \dots\} = \mathbb{Z}$.

So each integer is hit exactly once, so this is a bijection from \mathbb{N} to \mathbb{Z} . \star

\mathbb{Q} IS COUNTABLE

\exists_1 We can similarly show $|\mathbb{Q}| = |\mathbb{N}|$.

Proof: $\forall q \in \mathbb{Q}$, $q = \frac{a}{b}$, where $a, b \in \mathbb{Z}$ & $b \neq 0$.

Hence every rational number can be associated to an ordered pair (a, b) of integers.

So $|\mathbb{Q}| \leq |\mathbb{Z} \times \mathbb{Z}|$, and since $\mathbb{Z} \times \mathbb{Z}$ is countable,

\mathbb{Q} must be at most countable.

However $\mathbb{Z} \subseteq \mathbb{Q}$, so $|\mathbb{Z}| \leq |\mathbb{Q}|$.

Since $|\mathbb{Z}| = |\mathbb{N}| = |\mathbb{Z} \times \mathbb{Z}|$, by the Cantor-Bernstein

Theorem $|\mathbb{Q}| = |\mathbb{N}|$, as required. \square

THE UNION OF COUNTABLE SETS IS COUNTABLE

\exists_1 Let A and B be countable sets. Then $A \cup B$ is also countable.

Proof: Again, we can write

$$A = \{a_0, a_1, a_2, \dots\} \text{ and } B = \{b_0, b_1, b_2, \dots\}.$$

Then, let the sequence $\{c_n\}$ by $c_{2k} = a_k$ & $c_{2k+1} = b_k \forall k \in \mathbb{N}$.

Let $C = \{c_n\}$. It follows that since $A \subseteq A \cup B$ and $A \cup B \subseteq C$, thus

$$|A| \leq |A \cup B| \leq |A| + |B|.$$

However since $|A| = |\mathbb{N}| = |A| + |B|$, by the Cantor-Bernstein-Schröder theorem we must have that $|A \cup B| = |\mathbb{N}|$, and we are done. \star

\exists_2 Let A_0, A_1, \dots, A_n be a finite collection of countable sets.

Then $\bigcup_{i=0}^n A_i$ is also countable.

Proof: Again, we can use induction.

If $n=1$, this is just the above result.

Suppose the claim is true for $n=N$.

$$\text{So } \bigcup_{i=0}^{N+1} A_i = \underbrace{\bigcup_{i=0}^N A_i}_{\text{countable}} \cup \underbrace{A_{N+1}}_{\text{countable}}$$

Hence $\bigcup_{i=0}^{N+1} A_i$ is also countable, and so the claim is true for $n=N+1$.

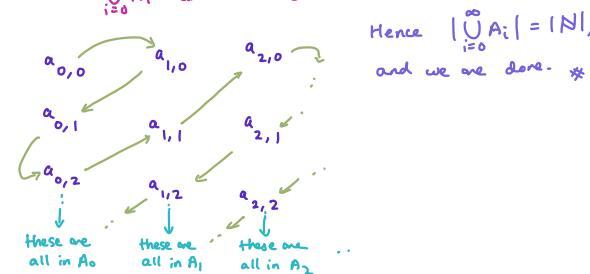
It follows by induction that the claim is true $\forall n \in \mathbb{N}$. \star

\exists_3 Let A_0, A_1, A_2, \dots be a countable collection of countable sets.

Then $\bigcup_{i=0}^{\infty} A_i$ is also countable.

Proof: Let $A_i = \{a_{i,0}, a_{i,1}, a_{i,2}, \dots\}$ for each $i \in \mathbb{N}$.

Then we can define a bijection between the elements of $\bigcup_{i=0}^{\infty} A_i$ and \mathbb{N} by



Hence $|\bigcup_{i=0}^{\infty} A_i| = |\mathbb{N}|$, and we are done. \star

UNCOUNTABLE SETS

\mathbb{R} IS UNCOUNTABLE

Q: First, we prove that the open interval $(0,1)$ has the same cardinality as \mathbb{R} ; ie there exists a bijection $f: (0,1) \rightarrow \mathbb{R}$.

Proof. Let $f(x) = \frac{1-2x}{x(x-1)}$. We claim $f(x)$ is a bijection.

Suppose $f(a) = f(b)$ for some $a, b \in (0,1)$, $a \neq b$.

$$\text{Then } \frac{1-2a}{a(a-1)} = \frac{1-2b}{b(b-1)}.$$

$$\Rightarrow 0 = (a-b)(a+b-2ab-1).$$

Since $a \neq b$ we get that $a+b-2ab-1 = 0$.

This implies $ab = a+b-ab-1 = -(a-1)(b-1)$.

However $ab > 0$ and $-(a-1)(b-1) < 0$, since $a, b \in (0,1)$; thus this is a contradiction; hence f is injective.

Then, suppose $f(x) = r$.

$$\text{This leads to } r = \frac{1-2x}{x^2-x}.$$

$$\Rightarrow rx^2 + (2-r)x - 1 = 0.$$

$$\therefore x = \frac{(r-2) \pm \sqrt{(r-2)^2 - 4(r)(-1)}}{2r} = \frac{(r-2) \pm \sqrt{r^2+4}}{2r}.$$

$$\text{We claim } 0 < \frac{(r-2) + \sqrt{r^2+4}}{2r} < 1 \quad \forall r \in \mathbb{R} \setminus \{0\}.$$

Proof: if $r > 0$, then $(r-2) + \sqrt{r^2+4} > 0$.

$$\text{Hence } \frac{r-2 + \sqrt{r^2+4}}{2r} > 0.$$

We also can infer $(r-2) + \sqrt{r^2+4} < 2r$,

$$\text{so } \frac{r-2 + \sqrt{r^2+4}}{2r} < 1.$$

There is a similar proof for when $r < 0$.

Therefore for any $r \in \mathbb{R}$, there exists an $x \in (0,1)$ such that $f(x) = r$, so $f(x)$ is surjective (and hence bijective). \blacksquare

Q: We can now prove \mathbb{R} is uncountable.

Proof. We can prove \mathbb{R} is uncountable if we show $(0,1)$ is uncountable, so we will do so.

Suppose \mathbb{R} is countable. Then we can write out all the real numbers in a list:

$$0.b_{11}b_{12}b_{13}b_{14}\dots$$

Subsequently, we take the diagonal "digits" of each number (circled in purple), and construct a number $r = 0.c_1c_2c_3\dots$

$$0.b_{21}b_{22}b_{23}b_{24}\dots$$

by $c_i = \begin{cases} 4, & b_{ii} \neq 4 \\ 5, & b_{ii} = 4 \end{cases}$

$$0.b_{31}b_{32}b_{33}b_{34}\dots$$

Then, by definition, $r \neq r_i \quad \forall i \in \mathbb{N}$, as r differs from r_i in the i th decimal place.

So $|R| > |\mathbb{N}|$, showing \mathbb{R} is uncountable. \blacksquare

$P(\mathbb{N})$ IS ALSO UNCOUNTABLE, AND $|P(\mathbb{N})| = |\mathbb{R}|$

Q: We can also show $|P(\mathbb{N})| = |\mathbb{R}|$, proving $P(\mathbb{N})$ is also uncountable.

Proof. again, we can just prove $P(\mathbb{N})$ has the same cardinality as the interval $(0,1)$.

Then, let $f: (0,1) \rightarrow P(\mathbb{N})$ by the following:
Given a $r \in (0,1)$, with decimal expansion $r = 0.b_1b_2b_3\dots$, let $f(r) = \{10^{n-1}b_n : n \in \mathbb{N}^+\}$.
eg $r = \frac{1}{3}$, $f(r) = \{3, 30, 300, \dots\}$

We claim f is injective.

Proof. Suppose $f(r) = f(s)$, for $r, s \in (0,1)$.

Let $r = 0.b_1b_2b_3\dots$ and $s = 0.c_1c_2c_3\dots$

Then for any $i \in \mathbb{N}^+$, $10^{i-1}b_i \in R$, so $10^{i-1}c_i \in S$ also.

However:

1) if $b_i \neq 0$, then $10^{i-1}b_i$ is the unique integer between 10^{i-1} and $9 \cdot 10^{i-1}$ belonging to R , and a similar statement holds for S .

$\therefore 10^{i-1}b_i$ must be the unique integer between 10^{i-1} & $9 \cdot 10^{i-1}$ as well, proving $10^{i-1}b_i = 10^{i-1}c_i$, so $b_i = c_i$.

2) if $b_i = 0$, then $0 \in R$, and so there is no integer between 10^{i-1} & $9 \cdot 10^{i-1}$ in R .

Then since $R = S$, the same holds true for S , thus $c_i = 0$ as well.

Since r and s agree in every decimal place, $r = s$; hence f is injective.

Next, let $g: P(\mathbb{N}) \rightarrow (0,1)$ such that

$$\forall A \in \mathbb{N}. \quad g(A) = 0.a_1a_2a_3\dots$$

$$\text{where } a_i = \begin{cases} 4, & i \in A \\ 5, & i \notin A \end{cases}$$

Then $a_i \in (0,1)$ with an unique decimal expansion.

Suppose $A, B \in \mathbb{N}$ such that $g(A) = g(B)$.

$$\Rightarrow 0.a_1a_2a_3\dots = 0.b_1b_2b_3\dots$$

Then necessarily $a_i = b_i \quad \forall i \in \mathbb{N}$, implying that if $i \in A$, then $i \in B$, and vice versa.

Hence $A \subseteq B$ & $B \subseteq A$, proving that $A = B$, so g is injective.

Since there exist injective functions $f: (0,1) \rightarrow P(\mathbb{N})$ and $g: P(\mathbb{N}) \rightarrow (0,1)$, consequently $|P(\mathbb{N})| = |(0,1)|$,

and so $|R| = |P(\mathbb{N})|$. \blacksquare

Chapter 4: Algebraic Structures

BINARY OPERATION

- E₁: Let S be a set. A binary operation on S is a function from $S \times S$ to S .
- E₂: If $*: S \times S \rightarrow S$ is a binary operation, we write $a * b$ instead of $*(a, b)$ as the output of $*$.
- E₃: Examples of binary operations:
① $+, -, \times$ on \mathbb{Z} or \mathbb{R} (NOT $\div, \frac{1}{\cdot}$)
② \div on $\mathbb{R}^* = \mathbb{R} \setminus \{0\}$
③ \cup and \cap on $P(A)$

ASSOCIATIVITY

- E₁: Let $*$ be a binary operation on S . Then $*$ is "associative" if $\forall a, b, c \in S$, we have $(a * b) * c = a * (b * c)$.
- E₂: We can show if a_1, a_2, \dots, a_n are arbitrary elements in S , with $n \geq 1$, then $a_1 * a_2 * \dots * a_n$ is well-defined, regardless of the choice of bracketing in the product.
- Proof: We let the "standard product" of a_1, a_2, \dots, a_n , denoted as $\langle a_1, a_2, a_3, \dots, a_n \rangle$, to be defined recursively by $\langle a_1 \rangle = a_1$, $\langle a_1, a_2 \rangle = a_1 * a_2$, and $\langle a_1, a_2, \dots, a_n \rangle = \langle a_1, a_2, \dots, a_{n-1} \rangle * a_n \ \forall n \geq 3$. (So $\langle a_1, a_2, a_3 \rangle = (a_1 * a_2) * a_3$; $\langle a_1, a_2, a_3, a_4 \rangle = ((a_1 * a_2) * a_3) * a_4$; etc.)

We claim every product of the n elements is equal to the standard product.

If $n=1$ or $n=2$, the choice of bracketing; thus the claim holds trivially in these cases.

Then, assume the claim is true for $1, 2, \dots, n$ elements, and suppose we are given $n+1$ elements of S ,

e.g. a_1, a_2, \dots, a_{n+1} .

Subsequently, any product of these elements can be expressed in the form $b * c$, where b is the product of some elements a_1, a_2, \dots, a_k with $1 \leq k \leq n$, and c is the product of the remaining elements a_{k+1}, \dots, a_{n+1} .

But since $b = \langle a_1, a_2, \dots, a_k \rangle$ and $c = \langle a_{k+1}, a_{k+2}, \dots, a_{n+1} \rangle$, we can express $b * c$ by

$$\begin{aligned} b * c &= \langle a_1, a_2, \dots, a_k \rangle * (\langle a_{k+1}, a_{k+2}, \dots, a_n \rangle * a_{n+1}) \\ &= (\langle a_1, a_2, \dots, a_k \rangle * \langle a_{k+1}, a_{k+2}, \dots, a_n \rangle) * a_{n+1} \\ &= \langle a_1, a_2, \dots, a_k, a_{k+1}, \dots, a_n, a_{n+1} \rangle \quad (\text{by induction hypothesis}). \end{aligned}$$

Hence the claim is also true for $n+1$.

Therefore, by induction, the claim is true $\forall n \in \mathbb{N}$. \square

COMMUTATIVITY

- E₁: Let $*$ be a binary operation on S . Then $*$ is "commutative" if $\forall a, b \in S$, we have $a * b = b * a$.

UNITY / IDENTITY

- E₁: Let $*$ be a binary operation on S . Then an element $e \in S$ is an "identity" for $*$ if $a * e = e * a = a \ \forall a \in S$.

CLOSED UNDER *

- E₁: For any binary operation $*$ on a set S , we say that S is "closed under $*$ " if $\forall a, b \in S : a * b \in S$.

MONOID

- E₁: Let S be a set with binary operation $*$. Then S is a "monoid" if $*$ is both associative, and there exists an identity element $e \in S$ with respect to $*$.
* we use the notation " ab " = $a * b$, and " a^n " = $\underbrace{axax\dots x a}_{n \text{ times}}$, with $a^0 = e$, where e is the identity element wrt $*$.

- E₂: We can prove the identity element of any monoid is unique.

Proof: let S be a monoid which has two identity elements, i.e. e_1 and e_2 .

Then $\forall a \in S$, $ae_1 = e_1 a = a$ and $ae_2 = e_2 a = a$. So if $a = e_2$, then $e_2 e_1 = e_1 e_2 = e_2$, and if $a = e_1$, then $e_2 e_1 = e_1 e_2 = e_1$. Thus $e_1 = e_2$, proving uniqueness. \square

UNIT & INVERSE

- E₁: For any monoid S , an element $a \in S$ is called a "unit" of S if there exists some $b \in S$ for which $ab = ba = e$, where e is the identity element of S .

- E₂: In this case, we call b the "inverse" of a .

* we usually use the notation " a^{-1} " or " $-a$ " to denote the inverse of a , and $a^{-m} = (a^{-1})^m$.

- E₃: Once again, we can show the inverse of any unit in a monoid S is unique.

Proof: let S be a monoid, and $a \in S$ be a unit.

Suppose b_1 & b_2 are both inverses of a .

so that $ab_1 = b_1 a = e$ and $ab_2 = b_2 a = e$.

Hence $b_1 = b_1 e = b_1 (ab_2) = (b_1 a) b_2 = eb_2 = b_2$,

proving $b_1 = b_2$, giving the desired uniqueness. \square

GROUP

B1 A set G is called a "group" if G is a monoid, and every element of G is a unit.

B2 Hence, a set G with a binary operation is called a group if:

- ① it is associative, ie $(ab)c = a(bc) \quad \forall a, b, c \in G$;
- ② there exists an identity, ie $\exists e \in G$ such that $ae = ea = e \quad \forall a \in G$; &
- ③ every element of G has an inverse, ie $\forall a \in G : \exists b \in G$ such that $ab = ba = e$.

ABELIAN GROUP

B1 A group G is further considered "abelian" if the operation $*$ on G is also commutative.

B2 Hence, any abelian group G satisfies the 3 properties above, in addition to:

- ④ it is commutative; ie $\forall a, b \in G : ab = ba$.

EXTRACTING A GROUP FROM A MONOID

B1 Let M be a monoid, and M^* be the set of all units of M . Then M^* is a group.

Proof. Let $a, b \in M^*$ be arbitrary.

Then by defn., a^{-1} & b^{-1} exist such that $aa^{-1} = e$ & $bb^{-1} = e$.

$$\begin{aligned} \text{Hence } (ab)(b^{-1}a^{-1}) &= a(bb^{-1})a^{-1} \\ &= aea^{-1} \\ &= aa^{-1} \\ &= e. \end{aligned}$$

We can similarly show $(b^{-1}a^{-1})(ab) = e$ also.

This implies ab is a unit, with inverse $b^{-1}a^{-1}$.

Hence M^* is "closed" under the binary operation on M .

Subsequently, M^* is also associative & has an identity element, as M has those properties. Furthermore, since all $a \in M^*$ has an inverse, it implies M^* is a group, and we are done. \square

MANIPULATING EXPONENTS IN

A GROUP

B1 Let G be a group, and $g \in G$.

$$\text{Then } \forall n, m \in \mathbb{Z}, \quad g^{n+m} = g^n \cdot g^m.$$

Proof. If $n=0$, $g^{n+m} = g^m$ trivially.

Assume the claim is true for some $n \in \mathbb{N}$ and all $m \in \mathbb{N}$.

$$\begin{aligned} \text{Then } g^{n+1} \cdot g^m &= (g \cdot g^n) \cdot g^m \\ &= g \cdot (g^n \cdot g^m) \\ &= g \cdot g^{n+m}. \end{aligned}$$

If $n, m > 0$, then the RHS is the result of multiplying $n+m$ copies of g by one more g , which is $g^{(n+m)+1}$ by definition.

If $n, m < -1$, then the above is equal to $g \cdot g^{-1}$, which is equal to $e = g^0 = g^{(n+m)+1}$.

Lastly, if $n, m \leq -2$, then the RHS is multiplying $|n+m|-1$ copies of g^{-1} by one copy of g , which results in $|n+m|-1$ copies of g^{-1} , which can again be written as $g^{(n+m)+1}$.

A similar argument holds if $n \in \mathbb{Z}$ and $m \in \mathbb{N}$.

Hence the claim is true $\forall m, n \in \mathbb{Z}$. by induction, so we are done. \square

B2 Similarly, for any group G , if $g \in G$, then $(g^m)^n = g^{mn} \quad \forall m, n \in \mathbb{Z}$.

Proof. If $n=0$, then $(g^m)^0 = (g^m)^0 = e = g^0$, so the claim is true.

Then, if the claim is true for some $n \in \mathbb{N}$,

$$\begin{aligned} \text{it implies } (g^m)^{n+1} &= (g^m)^n \cdot (g^m)^1 \\ &= g^{mn} \cdot g^m \\ &= g^{m(n+1)}. \end{aligned}$$

So the claim is true $\forall m, n \in \mathbb{Z}$, $n \in \mathbb{N}$.

However, if $n \in \mathbb{Z}$, then $(g^m)^n = (g^m)^{-l}$ for some $l \in \mathbb{N}$. But since for any $h \in G$, $h^{-l} = (h^l)^{-1}$, thus, $(g^m)^{-l} = [(g^m)^l]^{-1} = (g^{ml})^{-1} = g^{-ml} = g^{m(-l)} = g^{mn}$.

So the claim is true $\forall m, n \in \mathbb{Z}$, and we are done. \square

B3 Lastly, if $g, h \in G$ commute, ie that $gh = hg$, where G is a group, then $(gh)^n = g^n h^n \quad \forall n \in \mathbb{Z}$.

Proof. If $n=0$, then $(gh)^0 = e = ee = g^0 h^0$, so the claim is true for this value of n .

Then, assume the claim is true for some $n \in \mathbb{N}$.

$$\begin{aligned} \text{It follows that } (gh)^{n+1} &= (gh)^n \cdot (gh)^1 \\ &= (g^n h^n) \cdot gh \\ &= g^n (h^n) h \\ &= g^n (gh^n) h \\ &= g^{n+1} h^{n+1}, \end{aligned}$$

so the claim is true $\forall n \in \mathbb{N}$.

Lastly, if $n = -l$ for some $l \in \mathbb{N}$,

$$\begin{aligned} (gh)^n &= (gh)^{-l} = [(gh)^l]^{-1} = (g^l h^l)^{-1} = (h^l)^{-1} (g^l)^{-1} \\ &= h^{-l} g^{-l} = h^n g^n = g^n h^n \quad (\text{since } g, h \text{ commute}). \end{aligned}$$

Hence the claim is true $\forall n \in \mathbb{Z}$, so we are done. \square

CYCLIC GROUPS

SUBGROUP

\exists Let G be a group, and $g \in G$ be an element. Then

$\langle g \rangle = \{g^k : k \in \mathbb{Z}\}$ is a group in G , called the "subgroup of G generated by g ".

Proof. We know $\forall g^m, g^n \in \langle g \rangle$, $g^m \cdot g^n = g^{m+n} \in \langle g \rangle$, so $\langle g \rangle$ is closed under the group operation on G .

Thus, the operation on G restricts to a binary operation on $\langle g \rangle$, and so it is associative by default.

Furthermore, $e = g^0 \in \langle g \rangle$, so an identity exists, and for a given g^m there exists its inverse $g^{-m} \in \langle g \rangle$ as $g^m \cdot g^{-m} = g^0 = e$.

Hence $\langle g \rangle$ is a group. \square

CYCLIC GROUP

\exists A group G is called a "cyclic group" if $G = \langle g \rangle$ for some $g \in G$.

\exists_2 In this case, we say that g is a "generator" of G .

THE SET OF INTEGERS MODULO n

\exists_1 For any $n \in \mathbb{Z}^+$, the "set of integers modulo n ", denoted by " $\mathbb{Z}/n\mathbb{Z}$ ", is the set of equivalence classes of \mathbb{Z} under the relation "congruence modulo n ", which is defined by: $a \equiv b \pmod{n}$ if $\frac{a-b}{n} \in \mathbb{Z}$.
* this is also written as $n \mid a-b$.

\exists_2 We can prove that for any $n \in \mathbb{Z}$, $\mathbb{Z}/n\mathbb{Z}$ is a cyclic group.

Proof. Let $[a], [b] \in \mathbb{Z}/n\mathbb{Z}$.

We can define an "addition" operation by

$$[a] + [b] = [a+b].$$

why? \rightarrow assume $[a]=[a']$ and $[b]=[b']$.

Then $a-a' = kn$ and $b-b' = ln$ for some

$$k, l \in \mathbb{Z}.$$

$$\text{Hence } (a+b) - (a'+b') = (k+l)n.$$

$$\text{so } a+b \equiv a'+b' \pmod{n}.$$

Thus $[a+b] = [a'+b']$, so this operation is well defined.

Next, $+$ is trivially associative, with an identity $[0]$ and any element $[a] \in \mathbb{Z}/n\mathbb{Z}$ has an inverse $[-a]$. Moreover, $+$ is commutative, so $\mathbb{Z}/n\mathbb{Z}$ is an abelian group under this operation.

Lastly, since $[a] = a \cdot [1] \quad \forall a \in \{0, 1, \dots, n-1\}$, $\mathbb{Z}/n\mathbb{Z}$ is cyclic, with generator $[1]$. \square

ORDER OF AN ELEMENT

\exists Let G be a group. Then, for any $g \in G$, the "order" of g , denoted by $\text{o}(g)$, is the smallest integer $n \geq 1$ such that $g^n = e$, if such an integer exists.

If $g^n \neq e \quad \forall n \geq 1$, then we say $\text{o}(g) = \infty$.

ORDER OF THE INVERSE OF AN ELEMENT

\exists Let G be a group. Then, for any $g \in G$, $\text{o}(g^{-1}) = \text{o}(g)$.

Proof. Case 1: $\text{o}(g) = \infty$.

Then $\forall k \geq 1$, $g^k \neq e$.

So $(g^{-1})^k \neq e$ also, so $\text{o}(g^{-1}) = \infty$.

Case 2: $\text{o}(g) = n, \quad n \in \mathbb{N}$.

Then n is the smallest true integer such that $g^n = e$.

So $(g^{-1})^n = e$.

If $\exists m < n$ such that $(g^{-1})^m = e$, taking inverses of both sides implies that $g^m = e$, a contradiction.

So $\text{o}(g^{-1}) = \text{o}(g) = n$. \square

ORDER OF ELEMENTS IN A FINITE GROUP

\exists If G is a finite group (ie it only has finitely many elements), then every element of G has finite order.

Proof. Let G be a finite group.

Then for any $g \in G$, g^0, g^1, g^2, \dots cannot contain infinitely many elements, so $\exists m, n \in \mathbb{N}$ such that $g^m = g^n$.

This implies $g^{m-n} = e$, so g has finite order. \square

CONNECTING ORDER TO SUBGROUPS

\exists_1 Let G be a group, and $g \in G$ such that $\text{o}(g) = n, \quad n \in \mathbb{N}$. Then $g^k = g^m$ iff $k \equiv m \pmod{n}$, and $g^k = e$ iff $n \mid k$.

Proof. Suppose $k \equiv m \pmod{n}$.

So $k-m = ln, \quad l \in \mathbb{Z}$.

$$\text{So } g^{k-m} = g^{ln} = (g^n)^l = e, \text{ implying that } g^k = g^m.$$

Then, suppose $g^k = g^m$.

$$\text{So } g^{k-m} = e.$$

Assume $k-m = nr$. This implies

$$e = g^{nr} = e^n g^r = g^r, \text{ forcing } r=0.$$

So $n \mid (k-m)$, implying that $k \equiv m \pmod{n}$.

If $m=0$, then $g^k = g^0 = e$ iff $k=0$, i.e. iff $n \mid k$. \square

\exists_2 Let G be a group, and let $g \in G$ with $\text{o}(g) = n$.

Then $\langle g \rangle = \{e, g, g^2, \dots, g^{n-1}\}$, where $e, g, g^2, \dots, g^{n-1}$ are all distinct.

Proof. Clearly $\{e, g, g^2, \dots, g^{n-1}\} \subseteq \langle g \rangle$.

Then, suppose $k \in \mathbb{Z}$, and $k \neq nq+r$.

Since $g^k = g^r$, hence $g^k \in \{e, g, g^2, \dots, g^{n-1}\}$, so $\langle g \rangle \subseteq \{e, g, g^2, \dots, g^{n-1}\}$.

Subsequently, each element of $\{e, g, g^2, \dots, g^{n-1}\}$ are distinct,

because if $g^k = g^m$ for some k, m where $0 \leq m < k \leq n-1$, then $n \mid (k-m)$; this is impossible, so $k=m$,

proving the elements of this set are all distinct. \square

SUBGROUPS

\exists_1 Let G be a group. Then $H \subseteq G$ is called a "subgroup" of G if H itself is a group (with respect to the binary operation defined on G).

* note that $\{e\}$ is a subgroup of G , called the "trivial" subgroup of G .

SUBGROUP TEST

\exists_2 Let G be a group, and $H \subseteq G$, $H \neq \emptyset$.

Then H is a subgroup of G iff

$\forall a, b \in H$, we have $ab^{-1} \in H$.

Proof. First, suppose H is a subgroup of G .

Then $\forall a, b \in H$, $b^{-1} \in H$ necessarily.

and since H is closed w.r.t the operation on G , we get that $ab^{-1} \in H$, as desired. *

Then, suppose $\forall a, b \in H$, we have $ab^{-1} \in H$. Since $H \neq \emptyset$, we can deduce that $\forall a \in H$,

$a \cdot a^{-1} = e \in H$, so H has an identity element.

Then, taking any $b \in H$ and $a \in e$, we get that $ab^{-1} = e \cdot b^{-1} = b^{-1} \in H$, so each element of H is a unit.

Lastly, since the binary operation on G is associative, necessarily the binary operation on H is also associative.

Therefore H is a group, and since $H \subseteq G$, it follows that H is a subgroup of G . \square

CENTER

\exists_1 For any group G , the "center" of G , denoted as $Z(G)$, is defined to be the set

$$Z(G) = \{z \in G \mid zg = gz \ \forall g \in G\}.$$

\exists_2 We can show $Z(G)$ is an abelian subgroup of G .

Proof. Since $e \in Z(G)$, $\therefore Z(G) \neq \emptyset$.

Then, suppose $a, b \in Z(G)$. Let $g \in G$ be arbitrary.

Since $bg = gb$ (as $b \in Z(G)$), therefore

$$b^{-1}(bg) = b^{-1}gb \Rightarrow g = b^{-1}gb,$$

and hence $b^{-1} \in Z(G)$.

It follows that

$$(ab^{-1})g = a(b^{-1}g) = a(gb^{-1})$$

$$= (ga)b^{-1} = (ga)b^{-1} = g(ab^{-1}),$$

and so $ab^{-1} \in Z(G)$, proving $Z(G)$ is a subgroup of G .

Lastly, since $ab = ba \ \forall a, b \in G$ (by defn),

$Z(G)$ is also abelian, so we are done. \square

\exists_3 Note that $Z(G) = G$ iff G is abelian. (So $Z(G)$ is a "measure" of how commutative the group is.)

CONJUGATE

\exists_1 For any group G , the "conjugate" of H in G by some $g \in G$ to be the set

$$gHg^{-1} = \{ghg^{-1} : h \in H\}.$$

\exists_2 Once again, we can verify that gHg^{-1} is a subgroup of G .

Proof. Since $H \neq \emptyset$, $\therefore gHg^{-1} \neq \emptyset$ also.

Then, for any $a, b \in gHg^{-1}$, by defn

$$a = gh_1g^{-1} \quad \& \quad b = gh_2g^{-1} \text{ for some } h_1, h_2 \in H.$$

$$\text{Hence } ab^{-1} = (gh_1g^{-1})(gh_2g^{-1})^{-1}$$

$$= (gh_1g^{-1})(gh_2^{-1}g^{-1})$$

$$= g(h_1h_2^{-1})g^{-1},$$

and since $h_1h_2^{-1} \in H$ (by defn), since H is a subgroup),

therefore gHg^{-1} is also a subgroup. \square

NORMAL

\exists_1 Let G be a group, and H be a subgroup of G .

Then H is a "normal subgroup" if all the

conjugate subgroups of H are equal to H .

SYMMETRIC GROUPS

\exists_1 For any $n \in \mathbb{N}$, the "symmetric group of degree n " is the set of all bijections $f: \{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$, and we denote the set by S_n .

\exists_2 We can represent an element $\sigma \in S_n$ as a sort of matrix:

$$\text{ie } \begin{pmatrix} 1 & 2 & 3 & \dots & n \\ a_1 & a_2 & a_3 & \dots & a_n \end{pmatrix} \xleftarrow{\text{"input"} \quad \text{"output", where } a_i = f(i).}$$

\exists_3 For any $n \in \mathbb{N}$, $|S_n| = n!$.

Proof. There are n choices for a_1 .

$(n-1)$ choices for a_2 , (since f is injective) and so on, until we get that there is 1 choice for a_n .

So the number of permutations of a_1, a_2, \dots, a_n is $n(n-1) \cdots (1) = n!$. *

\exists_4 For $n \geq 3$, S_n is not abelian.

Proof. Let $\sigma = \begin{pmatrix} 1 & 2 & 3 & \dots & n \\ 2 & 1 & 3 & \dots & n \end{pmatrix}$ and $\tau = \begin{pmatrix} 1 & 2 & 3 & \dots & n \\ 3 & 1 & 2 & \dots & n \end{pmatrix}$.

$$\text{Then } (\sigma \circ \tau) = \begin{pmatrix} 1 & 2 & 3 & \dots & n \\ 3 & 2 & 1 & \dots & n \end{pmatrix}, \text{ but } (\tau \circ \sigma) = \begin{pmatrix} 1 & 2 & 3 & \dots & n \\ 3 & 1 & 2 & \dots & n \end{pmatrix},$$

so $(\sigma \circ \tau) \neq (\tau \circ \sigma)$, so S_n is not abelian. \square

HOMOMORPHISMS

Q: Let G_1 and G_2 be groups.

Then, a function $\phi: G_1 \rightarrow G_2$ is a "homomorphism" if $\forall a, b \in G_1$, we have

$$\phi(ab) = \phi(a) \cdot \phi(b),$$

where " \cdot " denotes the binary operation on G_2 .

Q: Examples:

① The "trivial" homomorphism

$$\phi(a) = e_{G_2} \quad \forall a \in G_1$$

② "Reduction modulo n" map

$$\phi: \mathbb{Z} \rightarrow \mathbb{Z}/n\mathbb{Z} \text{ by}$$

$$\phi(m) = [m] \quad \forall m \in \mathbb{Z}$$

③ The "exponent" map

$$\phi: \mathbb{Z} \rightarrow \langle g \rangle, \quad g \in G$$

$$\phi(m) = g^m \quad \forall m \in \mathbb{Z}$$

HOMOMORPHISMS "PRESERVE"

THE IDENTITY

Q: Let $\phi: G_1 \rightarrow G_2$ be a group homomorphism.

$$\text{Then } \phi(e_{G_1}) = e_{G_2}.$$

Proof. By defn.

$$\phi(e_{G_1}) = \phi(e_{G_1} \cdot e_{G_1}) = \phi(e_{G_1}) \cdot \phi(e_{G_1}).$$

$$\text{Hence } [\phi(e_{G_1})]^{-1} [\phi(e_{G_1})] = [\phi(e_{G_1})]^{-1} [\phi(e_{G_1}) \cdot \phi(e_{G_1})]$$

$$\Rightarrow e_{G_2} = \phi(e_{G_1}), \text{ as required.} \blacksquare$$

HOMOMORPHISMS "PRESERVE"

INVERSES

Q: Let $\phi: G_1 \rightarrow G_2$ be a group homomorphism.

$$\text{Then } \forall g \in G_1, \phi(g^{-1}) = [\phi(g)]^{-1}.$$

Proof. For each $g \in G_1$, observe that

$$\phi(g) \cdot \phi(g^{-1}) = \phi(g \cdot g^{-1}) = \phi(e_{G_1}) = e_{G_2}.$$

$$\text{Hence } [\phi(g)]^{-1} \cdot \phi(g) = [\phi(g)]^{-1} \cdot e_{G_2},$$

$$\text{or that } \phi(g^{-1}) = [\phi(g)]^{-1}. \blacksquare$$

HOMOMORPHISMS PRESERVE

POWERS

Q: Let $\phi: G_1 \rightarrow G_2$ be a group homomorphism.

$$\text{Then } \forall g \in G_1, \forall k \in \mathbb{Z}, \text{ we have } \phi(g^k) = [\phi(g)]^k.$$

Proof. We first prove this **Kernel** by induction.

$$\text{when } k=0, \phi(g^0) = \phi(e_{G_1}) = e_{G_2} = [\phi(g)]^0,$$

establishing our base case.

Then assuming the claim holds for some $k \in \mathbb{N}$,

note that $\phi(g^{k+1}) = \phi(g^k \cdot g) = \phi(g^k) \cdot \phi(g)$

$$= [\phi(g)]^k \cdot \phi(g) = [\phi(g)]^{k+1},$$

so the claim holds $\forall k \in \mathbb{N}$.

Subsequently, if $k=-m$ for some $m \in \mathbb{Z}^+$,

$$\text{then } \phi(g^k) = \phi(g^{-m}) = \phi((g^{-1})^m)$$

$$= [\phi(g^{-1})]^m = [[\phi(g)]^{-1}]^m = [\phi(g)]^{-m},$$

and so this is sufficient to prove the claim holds $\forall k \in \mathbb{Z}$. \blacksquare

THE COMPOSITION OF HOMOMORPHISMS

IS A HOMOMORPHISM

Q: Let $\phi: G_1 \rightarrow G_2$ and $\psi: G_2 \rightarrow G_3$ be group homomorphisms.

Then $(\psi \circ \phi): G_1 \rightarrow G_3$ is also a group homomorphism.

Proof. By defn.

$$\phi(a_1 \cdot a_2) = \phi(a_1) \cdot_2 \phi(a_2) \quad \forall a_1, a_2 \in G_1,$$

$$\text{and } \psi(b_1 \cdot_2 b_2) = \psi(b_1) \cdot_3 \psi(b_2) \quad \forall b_1, b_2 \in G_2,$$

where \cdot_1, \cdot_2 & \cdot_3 are the binary operations on G_1, G_2 , and G_3 respectively.

Hence, if we let $b_1 = \phi(a_1)$ and $b_2 = \phi(a_2)$, then

$$\psi(\phi(a_1) \cdot_2 \phi(a_2)) = \psi(\phi(a_1)) \cdot_3 \psi(\phi(a_2))$$

$$\Rightarrow \psi(\phi(a_1 \cdot a_2)) = \psi(\phi(a_1)) \cdot_3 \psi(\phi(a_2))$$

$$\Rightarrow (\psi \circ \phi)(a_1 \cdot a_2) = (\psi \circ \phi)(a_1) \cdot_3 (\psi \circ \phi)(a_2),$$

verifying that $(\psi \circ \phi)$ is a group homomorphism as well. \blacksquare

KERNEL

Q: Let $\phi: G_1 \rightarrow G_2$ be a group homomorphism.

Then the "kernel" of ϕ , denoted " $\ker \phi$ ",

is defined to be the set

$$\ker \phi = \{g \in G_1 \mid \phi(g) = e_{G_2}\}.$$

We can prove that $\ker \phi$ is a subgroup of G_1 .

Proof. We apply the Subgroup Test.

First, $\ker \phi$ is non-empty $\because \phi(e_{G_1}) = e_{G_2}$, so $e_{G_1} \in \ker \phi$.

Then, if $a, b \in \ker \phi$, then

$$\phi(ab^{-1}) = \phi(a)\phi(b^{-1}) = \phi(a)[\phi(b)]^{-1} = (e_{G_2})[e_{G_2}]^{-1} = e_{G_2},$$

and so $ab^{-1} \in \ker \phi$ also, verifying that $\ker \phi$ is a subgroup of G_1 . \blacksquare

Moreover, we can also show that ϕ is injective if and only if $\ker \phi = \{e_{G_1}\}$.

Proof. Assume ϕ is injective.

Let $g \in \ker \phi$ be arbitrary.

Then $\phi(g) = e_{G_2} = \phi(e_{G_1})$, so by injectivity

of ϕ , necessarily $g = e_{G_1}$, proving $\ker \phi = \{e_{G_1}\}$.

Then, assume $\ker \phi = \{e_{G_1}\}$.

Suppose there exists $a, b \in G_1$, such that $\phi(a) = \phi(b)$.

Hence $\phi(a) \cdot \phi(b^{-1}) = \phi(b) \cdot \phi(b^{-1})$

$$\Rightarrow \phi(ab^{-1}) = e_{G_2}, \text{ implying that } ab^{-1} \in \ker \phi.$$

Thus $ab^{-1} = e_{G_1}$, implying that $a = b$.

This is sufficient to prove that ϕ is injective. \blacksquare

COSETS

Q: Let G be a group, and H be a subgroup of G .

Then let the relation " \sim " on G be such that

$\forall a, b \in G$, we have $a \sim b$ if and only if $ab^{-1} \in H$.

Q: We can prove that " \sim " is an equivalence relation.

Proof. Since for any $a \in G$, $\phi(a \cdot a^{-1}) = \phi(a) = e_{G_2}$,

$a \sim a$; hence \sim is reflexive.

Then, for any $a, b \in G$, if $a \sim b$, it implies that $ab^{-1} \in H$; since H is closed under taking inverses, we get that $(ab^{-1})^{-1} = ba^{-1} \in H$ also, saying that $b \sim a$.

So \sim is also symmetric.

Lastly, if $a \sim b$ & $b \sim c$ for some $a, b, c \in G$, it implies that $ab^{-1} \in H$ and $bc^{-1} \in H$.

But since H is closed under multiplication, hence $(ab^{-1})(bc^{-1}) = ac^{-1} \in H$ also, implying that $a \sim c$.

So \sim is transitive (and hence an equivalence relation). \blacksquare

Q: Then, for any element $a \in G$, the "right coset of H generated by a " is defined to be the set

$$Ha = \{ha \mid h \in H\},$$

and is the equivalence class of a with respect to " \sim ".

Proof. $[a] = \{g \in G \mid g \sim a\}$

$$= \{g \in G \mid ga^{-1} \in H\}$$

$$= \{g \in G \mid ga^{-1} = h \text{ for some } h \in H\}$$

$$= \{g \in G \mid g = ha \text{ for some } h \in H\}$$

$$= Ha. \blacksquare$$

Q: Similarly, we can define an equivalence relation \sim_L on G such that $\forall a, b \in G$, we have that $a \sim_L b$ if and only if $b^{-1}a \in H$.

Then, the "left coset of H generated by a " for some $a \in G$ is the set

$$aH = \{ah \mid h \in H\},$$

and is an equivalence class of \sim_L .

Q: Note that if G is abelian, then $aH = Ha$ for any $a \in G$ and subgroup H .

QUOTIENT GROUPS

NORMAL SUBGROUP

\exists_1 Let G be a group. Then, a subgroup $H \triangleleft G$ is called a "normal subgroup" of G if $gHg^{-1} = \{ghg^{-1} : h \in H\} = H \quad \forall g \in G$.

\exists_2 We use the notation " $H \triangleleft G$ " to denote that H is a normal subgroup of G .

MULTIPLICATION OF RIGHT COSETS

\exists_1 Let G be a group. Then, for every subgroup $H \triangleleft G$, the formula $(Ha)(Hb) = H(ab)$ gives a well-defined multiplication of right cosets if and only if $H \triangleleft G$.

Proof. First, assume $(Ha)(Hb) = Hab$ for all right cosets of H in G .

Let $g \in G$ and $h \in H$ be arbitrary.

Then clearly $Hh = H$.

So $(Hg)(Hh) = (Hg)(H)$, or $Hgh = Hg$.

By defn., this implies $(gh)(ge)^{-1} = ghg^{-1} \in H \quad \forall g \in G$. Letting $g^{-1} = g$, we see that $g^{-1}hg \in H$ also, and hence $gHg^{-1} \subseteq H$ and $g^{-1}Hg \subseteq H$, or $HgHg^{-1}$. Thus $H = gHg^{-1}$, implying $H \triangleleft G$. \blacksquare

Conversely, assume $H \triangleleft G$.

Let $a, b, a_1, b_1 \in G$ be such that $Ha = Ha_1$ and $Hb = Hb_1$. Then by definition, $aa_1^{-1} \in H$ and $bb_1^{-1} \in H$.

Next, observe that

$$abb_1^{-1}a_1^{-1} = a(bb_1^{-1})a_1^{-1} = a(bb_1^{-1})a^{-1}(aa_1^{-1}).$$

Since $aHa^{-1} = H$ and $bb_1^{-1} \in H$, so $a(bb_1^{-1})a^{-1} \in aHa^{-1} = H$.

But since we assumed $aa_1^{-1} \in H$, thus $(a(bb_1^{-1})a^{-1})(aa_1^{-1}) \in H$

(by closure of multiplication), so $abb_1^{-1}a_1^{-1} \in H$.

But since $abb_1^{-1}a_1^{-1} = (ab)(b_1^{-1}a_1^{-1}) = (ab)(a_1b_1)^{-1} \in H$, it follows that $(Ha)(Hb) = (Ha_1)(Hb_1)$, and so multiplication of right cosets is well-defined.

QUOTIENT GROUP

\exists_1 Let G be a group, and $H \triangleleft G$. Then, the "quotient group" of G by H

is the set G/H of right cosets of H under the operation $(Ha)(Hb) = Hab$.

\exists_2 We can verify that G/H is a group.

Proof. By the above, we know the binary operation on G/H is well-defined.

Then, notice that for any cosets $Ha, Hb \in G/H$:

$$(Ha)(Hb)(Hc) = (Hab)(Hc) = H(ab)c = Ha(bc) \\ = (Ha)(Hbc) = (Ha)(Hb)(Hc),$$

so the operation is associative.

Then the identity of G/H is the coset $H = He$, and the inverse of the coset Ha is the coset Ha^{-1} ,

so G/H is indeed a group. \blacksquare

QUOTIENT MAPPING

\exists_1 Let G be a group, and $H \triangleleft G$. Then the "quotient mapping" is

the function $\phi: G \rightarrow G/H$ by $\phi(g) = Hg$.

\exists_2 We can show ϕ is surjective, and a group homomorphism.

Proof. By defn., $\phi(ab) = Hab = (Ha)(Hb) = \phi(a)\phi(b)$, so ϕ is a homomorphism.

Then for any coset $Ha \in G/H$, $\phi(a) = Ha$, so ϕ is also surjective. \blacksquare

THE QUOTIENT GROUP IS ABELIAN IF THE ORIGINAL GROUP IS ABELIAN

\exists_1 Let G be a group, and $H \triangleleft G$. Suppose G is abelian. Then G/H is also abelian.

Proof. If G is abelian, then $Hab = Hba$.

So $(Ha)(Hb) = Hab = Hba = (Hb)(Ha)$, implying that G/H is also abelian. \blacksquare

THE QUOTIENT GROUP IS CYCLIC IF THE ORIGINAL GROUP IS CYCLIC

\exists_1 Let G be a group, and $H \triangleleft G$. Suppose G is cyclic. Then G/H is also cyclic.

Proof. If $G = \langle g \rangle$ for some $g \in G$, then every element of G is in the form g^k for some $k \in \mathbb{Z}$. So, given any $Ha \in G/H$, we know $a = g^k$ for some $k \in \mathbb{Z}$, and so $Ha = Hg^k = \phi(g^k) = [\phi(g)]^k = (Hg)^k$, where ϕ is the quotient homomorphism.

Hence $G/H = \langle Hg \rangle$, so it is also cyclic. \blacksquare

ALL SUBGROUPS OF AN ABELIAN GROUP ARE ALSO NORMAL

\exists_1 Let G be an abelian group. Then all subgroups $H \triangleleft G$ are normal.

Proof. Let $g \in G$ be arbitrary.

Then $gHg^{-1} = \{ghg^{-1} : h \in H\} = \{h : h \in H\} = H$, so H is normal. \blacksquare

EXAMPLE OF A QUOTIENT GROUP: \mathbb{Q}/\mathbb{Z}

\exists_1 Consider the group \mathbb{Q} , with addition as its binary operation. Then \mathbb{Z} is a subgroup of \mathbb{Q} under addition, and since

\mathbb{Q} is abelian, \mathbb{Z} is a normal group.

Hence the quotient group \mathbb{Q}/\mathbb{Z} exists.

\exists_2 Note the elements of \mathbb{Q}/\mathbb{Z} are of the form $\mathbb{Z} + q$, where $q \in \mathbb{Q}$.

In fact, we can further deduce each element of \mathbb{Q}/\mathbb{Z} is uniquely represented by a coset of the form $\mathbb{Z} + s$, where $0 \leq s < 1$.

TORSION

\exists_1 Let G be a group.

Then G is also a "torsion" if every element of G has finite order.

LAGRANGE'S THEOREM

INDEX

Let G be a group, and H be a subgroup of G . Then, the "index" of H in G , denoted by $|G:H|$, is equal to the number of distinct left cosets of H in G . *this can be finite or infinite.

* $|G:H|$ is also equal to the number of distinct left cosets of H .

LAGRANGE'S THEOREM

Let G be a finite group, and let H be a subgroup of G .

$$\text{Then } |G:H| = \frac{|G|}{|H|}.$$

Proof: Since the right cosets of H in G forms a partition of G , let H_1, H_2, \dots, H_n denote the collection of distinct right cosets of H in G .

$$\text{Then, by definition, } \bigcup_{i=1}^n H_i = G,$$

and $H_i \cap H_j = \emptyset$ if $i \neq j$.

Next, since $f: H \rightarrow H_1$ by $f(h) = h_1$ is a bijection, consequently $|H| = |H_1|$.

$$\text{Hence } |H_1| = |H_2| = \dots = |H_n|.$$

$$\text{It follows that } |G| = |H_1| + |H_2| + \dots + |H_n| = n|H|$$

$$\text{and so } n = |G:H| = \frac{|G|}{|H|}. \quad \square$$

$\text{o}(g)$ DIVIDES $|G|$

Let G be a finite group.

Then, for any $g \in G$, $\text{o}(g)$ divides $|G|$.

Proof: Note $H = \langle g \rangle$ is a subgroup of G , and that $|H| = \text{o}(g)$.

Since $|H|$ divides $|G|$, thus

$\text{o}(g)$ divides $|G|$ also. \square

$$|G|=n \Rightarrow g^n=e$$

Let G be a finite group, with $|G|=n$.

Then $\forall g \in G$, we have $g^n=e$.

Proof: Let $g \in G$ be arbitrary.

Then, necessarily $\text{o}(g)$ divides n , so $n = kl$, where $k = \text{o}(g)$, and $l \in \mathbb{Z}$.

Hence $g^n = g^{kl} = (g^k)^l = e^l = e$, as needed. \square

EVERY GROUP WITH A PRIME NUMBER OF ELEMENTS IS CYCLIC

Suppose G is a group with $|G|=p$, where p is a prime number.

Then for any non-identity element $g \in G$, $G = \langle g \rangle$; ie G is cyclic.

Proof: Since $p \geq 2$, there exists a non-identity element $g \in G$.

Let $H = \langle g \rangle$. Then $|H| > 1$, since $\text{o}(g) > 1$.

But since $|H|$ must divide $|G|$ by Lagrange's Theorem, and $|G|=p$ only has 1 and p as positive divisors, it follows that $|H|=p=|G|$, so $G=H=\langle g \rangle$, showing that G is cyclic. \square

RINGS

E1: A ring is a set R equipped with two binary operations (usually denoted by addition and multiplication) that satisfy the following:

① R is an abelian group wrt the binary operation "+", with identity "0";

② R is a monoid wrt the binary operation "·", with identity "1"; and

③ left & right distributive laws hold; ie

$$\forall a, b, c \in R, \quad a(b+c) = ab+ac \quad \& \quad (ab)c = a(bc)$$

E2: If multiplication in R is commutative, ie $ab=ba \quad \forall a, b \in R$, then we call R a "commutative ring".

$\forall a, b, c \in R:$

- ① $(ab)+c = a+(bc)$
- ② $\exists 0 \in R$ such that $a+0=0+a=a$
- ③ $\exists (-a) \in R$ such that $a+(-a)=(-a)+a=0$
- ④ $a+b=b+a$.
- ⑤ $(ab)c = a(bc)$
- ⑥ $\exists 1 \in R$ such that $a \cdot 1 = 1 \cdot a = a$
- ⑦ $(a+b)c = ac+bc$
- ⑧ $a(c+b) = ac+ab$

PROPERTIES OF RINGS

UNIQUENESS OF IDENTITIES AND INVERSES

E1: Let R be an arbitrary ring. Then:

- ① the additive and multiplicative inverses of R are unique; and
- ② for any $a \in R$, its additive inverse is unique, and usually denoted $-a$.

* these follow from previous theorems.

ADDITIONAL IDENTITY IN MULTIPLICATION

E1: Let R be a ring, with additive identity 0 .

$$\text{Then } \forall a \in R, \quad a \cdot 0 = 0 \cdot a = 0.$$

Proof: Since 0 is the additive identity,

$$\text{hence } 0+0=0.$$

So, by distributivity,

$$a \cdot 0 = a(0+0) = a \cdot 0 + a \cdot 0.$$

Adding $-a \cdot 0$ to both sides yields

$$0 = a \cdot 0, \text{ as needed.}$$

(The proof that $0 \cdot a = 0$ is similar.)

ADDITIONAL INVERSES AND MULTIPLICATION

E1: Let R be a ring, and $a, b \in R$ be arbitrary.

Then,

$$\textcircled{1} \quad (-a)b = a(-b) = -(ab); \text{ and}$$

$$\textcircled{2} \quad (-a)(-b) = ab.$$

Proof: By distributivity,

$$(-a)b + ab = (-a+b)b = 0b = 0,$$

and by commutativity of addition $ab+(-a)b=0$ also.

This implies $(-a)b$ is the additive inverse of ab ,

$$\text{and so } (-a)b = -(ab).$$

The proof that $a(-b) = -(ab)$ is similar. *

Then, note $(-a)(-b) = -(-a(-b)) = -(-(ab)) = ab$,

as ab is the unique additive inverse of $-ab$. ■

CHARACTERISTIC

E1: Let R be a ring. Then, the "characteristic of R ", denoted as "char R ", is the order of the multiplicative identity 1 in the group R under addition, if this order is finite.

If it is not, we denote $\text{char } R = 0$.

CHARACTERISTICS AND MULTIPLICATION

E1: Let R be a ring with $\text{char } R \neq 0$.

$$\text{Then } k \cdot r = \underbrace{r+r+\dots+r+r}_{k \text{ times}} = 0 \quad \forall r \in R$$

if and only if $n | k$ (ie n divides k).

Proof: Let $k \in \mathbb{Z}$ such that $n | k$, and let $r \in R$ be arbitrary.

Then $k = mn$ for some $m \in \mathbb{Z}$.

$$\text{Note } k \cdot r = (mn) \cdot r = (mn \cdot 1) \cdot r \text{ by distributivity.}$$

Furthermore, $n \cdot 1 = 0$, so $(mn) \cdot 1 = m \cdot (n \cdot 1) = m \cdot 0 = 0$;

therefore $k \cdot r = 0 \cdot r = 0$, as needed. *

Next, suppose $k \in \mathbb{Z}$ such that $k \cdot r = 0 \quad \forall r \in R$.

In particular, $k \cdot 1 = 0$, so k must be a multiple of the order of 1 in the group R under addition.

So necessarily $n | k$ (see the previous theorems.) ■

E2: Let R be a ring with $\text{char } R = 0$.

$$\text{Then } k \cdot r = 0 \quad \forall r \in R \text{ if and only if } k=0.$$

Proof: If $k=0$, then $k \cdot r = 0 \cdot r = 0 \quad \forall r \in R$.

Conversely, if $k \cdot r = 0 \quad \forall r \in R$, then $k \cdot 1 = 0$, and since 1 has infinite order, we can deduce that this occurs only when $k=0$. ■

ENDOMORPHISMS

E1: Let G be an abelian group, with the group operation denoted by addition.

Then the "set of endomorphisms" of G ,

denoted by $\text{End}(G)$, is the set of all

group homomorphisms $\phi: G \rightarrow G$.

E2: We can define an addition on $\text{End}(G)$ by

$$(\phi+\psi)(g) = \phi(g) + \psi(g) \quad \forall \phi, \psi \in \text{End}(G).$$

E3: We can also define a multiplication on $\text{End}(G)$ by

$$(\phi\psi)(g) = (\phi \circ \psi)(g) \quad \forall \phi, \psi \in \text{End}(G).$$

SUBRINGS

B₁ Let R be a ring. Then a subset $S \subseteq R$ is a "subring" if the "+" and "x" operations on R restrict to binary operations on S , and if S is a ring with respect to these restricted operations from R .

B₂ We also insist $l_R = l_S$; ie the multiplicative identity of the rings R and S must be the same.

Why? → they may not agree.

eg $R = \mathbb{Z}/6\mathbb{Z}$, $S = \{[0], [2], [4]\}$

$l_R = [1]$, $l_S = [4]$, $l_R \neq l_S$.

SUBRING TEST

B Let R be a ring, and $S \subseteq R$ with $S \neq \emptyset$.

Then S is a subring if and only if:

① $l_R \in S$, where l_R is the multiplicative identity for R ;

② $a - b \in S \quad \forall a, b \in S$; and

③ $ab \in S \quad \forall a, b \in S$.

Proof: First, assume $S \subseteq R$ satisfies the three conditions above.

Then by ②, S is a subgroup of R with respect to addition.

Then by ③, S is closed under multiplication, and associativity of "x" follows from the fact that it is true for R .

By ①, S has a multiplicative identity l_S , and by the uniqueness of the identity of a monoid we know that $l_R = l_S$.

Finally, the distributive laws hold in S because they also hold in R .

Hence S must be a ring, and so is a subring of R . \blacksquare

Conversely, assume S is a subring of R .

Then since S is a subgroup of R wrt "+", it follows from the Subgroup Test that $a - b \in S \quad \forall a, b \in S$.

Similarly, since S is closed under the multiplication on R , we know $ab \in S \quad \forall a, b \in S$ by definition.

Finally, since by definition $l_R = l_S \in S$, we can see all three conditions are satisfied. \blacksquare

CENTRE OF A RING

B₁ Let R be a ring. Then the "centre" of R , denoted by $Z(R)$, is defined to be

$$Z(R) = \{z \in R : zr = rz \quad \forall r \in R\}.$$

B₂ We can prove $Z(R)$ is also a ring.

Proof: Note $Z(R) \neq \emptyset$, as $l \in Z(R)$.
(So ① in the Subring Test is satisfied.)

Then, for any $r \in R$ and $a, b \in Z(R)$:

$$\begin{aligned} (a - b)r &= (a + (-b))r = ar + (-b)r = ar + (-br) \\ &= ra - rb = ra + r(-b) = r(a - b), \\ \text{so } (a - b) &\in Z(R), \text{ satisfying ② in the test.} \\ \text{Lastly, } (ab)r &= a(br) = a(-b) = (ar)b = (ra)b = r(ab), \\ \text{so } ab &\in Z(R) \text{ also, satisfying ③ in the test.} \end{aligned}$$

Hence $Z(R)$ must be a ring. \blacksquare

B₃ Note if R is commutative, then $Z(R) = R$.

RING HOMOMORPHISMS

B Let R and S be rings. Then, a function $\phi: R \rightarrow S$ is a "ring homomorphism" if:

- ① $\phi(a+b) = \phi(a) + \phi(b) \quad \forall a, b \in R$;
- ② $\phi(ab) = \phi(a)\phi(b) \quad \forall a, b \in R$; and
- ③ $\phi(l_R) = l_S$.

PROPERTIES OF HOMOMORPHISMS

B Suppose $\phi: R_1 \rightarrow R_2$ is a ring homomorphism. Let $r \in R_1$ be arbitrary. Then:

- ① $\phi(0) = 0$;
- ② $\phi(-r) = -\phi(r)$;
- ③ $\phi(kr) = k\phi(r) \quad \forall k \in \mathbb{Z}$;
- ④ $\phi(r^n) = [\phi(r)]^n \quad \forall n \in \mathbb{N}$; and
- ⑤ $\phi(r^{-1}) = [\phi(r)]^{-1} \quad \forall n \in \mathbb{N}$ if r^{-1} exists.

IDEALS AND QUOTIENT RINGS

RELATIONSHIP BETWEEN $\phi: G \rightarrow G_1$ AND $\ker \phi$

$\textcircled{1}$: Let $\phi: G \rightarrow G_1$ be a group homomorphism. Then $\ker \phi$ is a normal subgroup of G .

Proof: Let $K = \ker \phi$.

Then we know K is a subgroup of G .

Let $g \in G$ be arbitrary.

Suppose $h = gkg^{-1}$ for some $k \in K$.

Observe that since

$$\begin{aligned}\phi(h) &= \phi(gkg^{-1}) = \phi(g)\phi(k)\phi(g)^{-1} \\ &= \phi(g)(e)\phi(g)^{-1} \\ &= e,\end{aligned}$$

it follows that $h \in \ker \phi = K$, and so $gkg^{-1} \in K$.

If we let g^{-1} take the place of g , it shows that $g^{-1}kg \in K \quad \forall g \in G$, so $K \trianglelefteq G$.

It follows that $K = gKg^{-1}$, and since this holds $\forall g \in G$, we thus have that K is normal in G . \blacksquare

$\textcircled{2}$: Let H be a normal subgroup of G .

Then there exists a group homomorphism

$\phi: G \rightarrow G_1$ such that $H = \ker \phi$.

Proof: Let $G_1 = G/H$, and consider the quotient homomorphism $q: G \rightarrow G/H$ given by

$$q(g) = Hg \quad \forall g \in G.$$

Observe that $g \in \ker q$ if and only if $Hg = He$,

which occurs if and only if $g^{-1} \in H$,

which occurs if and only if $g \in H$.

Hence $H = \ker q$, and we are done. \blacksquare

KERNEL OF ABELIAN GROUPS

$\textcircled{1}$: Let R and S be rings, and $\phi: R \rightarrow S$ be a ring homomorphism.

Then, the "kernel of ϕ ", denoted by " $\ker \phi$ ", is defined to be the set $\ker \phi = \{r \in R : \phi(r) = 0_S\}$.

$\textcircled{2}$: Note that by construction, $\ker \phi$ is an additive subgroup of R .

IDEAL OF A RING

$\textcircled{1}$: Let R be a ring. Then, a subset $I \subseteq R$ is called an ideal of R if

(1) I is a subgroup of the additive group R ; and

(2) I "absorbs" multiplication; ie $\forall r \in I, a \in R$, we have that $ar, ra \in I$.

$\textcircled{2}$: Example: the "zero ideal" of R ; $\{0\}$.

First, note $\{0\}$ is the trivial subgroup of R under addition.

Then, for any $a \in R$, certainly $a \cdot 0 = 0 \cdot a = 0 \in \{0\}$,

so $\{0\}$ absorbs multiplication as well.

So $\{0\}$ is an ideal of R .

PRINCIPAL IDEAL GENERATED BY a

$\textcircled{1}$: Let R be a commutative ring, and $a \in R$ be arbitrary.

Then the "principal ideal generated by a " is the set

$$aR = Ra = \{s \in R : s = ar \text{ for some } r \in R\}.$$

QUOTIENT RING

$\textcircled{1}$: Let R be a ring, and let I be an ideal of R .

Then the set of right cosets R/I can be given the structure of a ring, with

- i) addition given by $(I+a) + (I+b) = I + (a+b)$; and
- ii) multiplication given by $(I+a)(I+b) = I + ab$,

for any $I+a, I+b \in R/I$.

Proof: First, since I is an additive subgroup of R , and R is abelian (under $+$), then by definition R/I is well defined under addition.

Then, suppose $\exists a', b' \in R$ such that $I+a = I+a'$ and $I+b = I+b'$.

This implies $a-a' \in I$ and $b-b' \in I$.

Subsequently, note that

$$ab - a'b' = ab - a'b + a'b - a'b' = (a-a')b + a'(b-b').$$

Since $a-a' \in I$, and I absorbs multiplication, necessarily $(a-a')b \in I$. Similarly, since $b-b' \in I$,

we must have that $a'(b-b') \in I$.

Finally, since I is closed under addition, $(a-a')b + a'(b-b') \in I$, and so $ab - a'b' \in I$.

This shows that $I+ab = I+a'b'$, proving that multiplication is well-defined.

Certainly, $I+I$ is an identity wrt multiplication, since $(I+a)(I+b) = I+a+b = I+a \quad \forall (I+a) \in R/I$.

Also, multiplication in R/I is associative,

since

$$\begin{aligned}((I+a)(I+b))(I+c) &= (I+ab)(I+c) = (I+(ab)c) \\ &= (I+a(bc)) = (I+a)(I+bc) = (I+a)((I+b)(I+c))\end{aligned}$$

for any $I+a, I+b, I+c \in R/I$.

We can similarly show R/I also satisfies the distributive laws.

RELATIONSHIP BETWEEN $\phi: R \rightarrow R_1$ AND $\ker \phi$

$\textcircled{1}$: Let $\phi: R \rightarrow R_1$ be a ring homomorphism.

Then $\ker \phi$ is an ideal of R .

(Proof in previous section)

$\textcircled{2}$: Let I be an ideal of R . Then there exists a ring homomorphism $\phi: R \rightarrow R_1$ such that

$$\ker \phi = I.$$

Proof: Let $R_1 = R/I$, and consider the quotient mapping

$$q: R \rightarrow R_1 \text{ by } q(a) = I+a \quad \forall a \in R.$$

Then q is a ring homomorphism since R/I is a ring.

But since $a \in \ker q$ if and only if $q(a) = I+a = I+0$,

which holds if and only if $a \in I$,

we must have that $I = \ker q$, and we are done. \blacksquare

Chapter 5: Elementary Number Theory

INTEGRAL DOMAINS

ZERO DIVISOR

Let R be a commutative ring. Then, an element $a \in R$ is called a "zero divisor" if there exists a $b \in R$, $b \neq 0$ such that $ab = 0$.

INTEGRAL DOMAIN

Let R be a commutative ring, where $R \neq \{0\}$. Then R is an "integral domain" if 0 is the only zero divisor; i.e. if $ab = 0$, then either $a=0$ or $b=0$.

Example: the ring \mathbb{Z} .

CANCELLATION PROPERTY OF AN INTEGRAL DOMAIN

Let R be a commutative ring. Then, we can show that $R \neq \{0\}$ is an integral domain if and only if

$\forall a, b, c \in R$, if $ab = ac$ and $a \neq 0$, then $b=c$.

Proof. First, suppose R is an integral domain. Let $a, b, c \in R$ such that $ab = ac$ and $a \neq 0$.

Then $ab - ac = 0$, so $a(b - c) = 0$.

But since R is an integral domain, and $a \neq 0$, by assumption, we must have that $b - c = 0$, and so $b = c$.

Conversely, assume the "cancellation" property holds.

Let $a, b \in R$ such that $ab = 0$.

If $a \neq 0$ the claim follows trivially, so assume $a \neq 0$. Then $ab = 0 = a \cdot 0$ and $a \neq 0$, so necessarily $b = 0$. \blacksquare

FIELD

Let F be a ring. Then, F is a "field" if

- ① it is commutative; and
- ② every non-zero element of F has a multiplicative inverse in F .

In other words, for each $a \in F$, $\exists a^{-1} \in F$ such that $aa^{-1} = 1$.

Examples: \mathbb{Q} and \mathbb{R} .

EVERY SUBRING OF A FIELD IS AN INTEGRAL DOMAIN

Let F be a field. Then every subring R of F is an integral domain of F .
Proof. Since multiplication is defined the same as in F , and multiplication in F is commutative, we know R is a commutative ring.

Then, suppose we have $a, b \in R$ such that $ab = 0$ in R .

Necessarily, this equation also holds in F . Next, if $a \neq 0$, then a^{-1} exists in F by definition,

and so $a^{-1}ab = b = a^{-1} \cdot 0 = 0$, and so $b = 0$.

This proves if $ab = 0$ in R , then $a = 0$ or $b = 0$, and so R is an integral domain. \blacksquare

* this also proves F is an integral domain!

GAUSSIAN INTEGERS

The ring of "Gaussian integers", denoted by $\mathbb{Z}[i]$, is defined to be the set

$$\mathbb{Z}[i] = \{a+bi : a, b \in \mathbb{Z}\}$$

with addition given by

$$(a+bi) + (c+di) = (a+c) + (b+d)i$$

and multiplication given by

$$(a+bi)(c+di) = (ac-bd) + (ad+bc)i.$$

* i is such that $i^2 = -1$.

NORM

Let $a+bi \in \mathbb{Z}[i]$ be arbitrary. Then,

the "norm" of $a+bi$, denoted by

$$N(a+bi)$$

$$N(a+bi) = a^2 + b^2.$$

PROPERTIES OF INTEGRAL DOMAINS

CHAR R IS ZERO OR PRIME

Let R be an integral domain. Then either $\text{char } R = 0$ or $\text{char } R = p$, where p is prime.

Proof. Suppose $\text{char } R \neq 0$ and $\text{char } R$ is not prime.

Then either $\text{char } R = 1$ or $\text{char } R$ is composite.

If $\text{char } R = 1$, then $R = \{0\}$, and so cannot be an integral domain.

If $\text{char } R$ is composite, then $\text{char } R = ab$, where $1 < a, b < n$. Then $r = a \cdot 1_R$ and $s = b \cdot 1_R$ are non-zero, but $rs = (ab) \cdot 1_R = 0$, which is a contradiction.

Hence necessarily if R is an integral domain, then $\text{char } R = 0$ or $\text{char } R$ is prime. \blacksquare

EVERY FINITE INTEGRAL DOMAIN IS A FIELD

Let R be a finite integral domain. Then R is also a field.

Proof. Let $n = |R|$, where $n \in \mathbb{N}$.

Let $a \in R$ be arbitrary, with $a \neq 0$.

Let the multiplication map $\phi: R \rightarrow R$ by

$$\phi(r) = ar \quad \forall r \in R.$$

Note ϕ is injective: if $\phi(r) = \phi(s)$, then $ar = as$, and by the cancellation property necessarily $r = s$.

Then, since R is finite, ϕ must also be surjective.

By injectivity, $|\phi(R)| = n$, and since $\phi(R) \subseteq R$ and $|R| = n$, necessarily $R = \phi(R)$.

Subsequently, by surjectivity, there exists a $b \in R$ such that $\phi(b) = 1$,

which says that $ab = ba = 1$.

Thus each $a \in R$ has a multiplicative inverse $b \in R$, and so

R must be a field. \blacksquare

DIVISIBILITY

Let R be an integral domain, and let $a, b \in R$ be arbitrary. Then, we say " a divides b " if there exists a $c \in R$ such that $b = ac$, and write $a | b$.

* we could write $a + b$ if a does not divide b .

THE DIVISIBILITY RELATION IS REFLEXIVE

Let R be an integral domain. Then $a | a$ $\forall a \in R$.

Proof. This follows from the fact that $a = 1 \cdot a$ $\forall a \in R$. \square

THE DIVISIBILITY RELATION IS TRANSITIVE

Let R be an integral domain. Then $\forall a, b, c \in R$, if $a | b$ and $b | c$, necessarily $a | c$.

Proof. Since $a | b$ and $b | c$, thus $ak = b$ and $c = bl$ for some $k, l \in R$.

$$\begin{aligned} \text{Thus } c &= bl \\ &= (ak)l \\ &= a(kl). \end{aligned}$$

Since $(kl) \in R$, this tells us $a | c$, so we are done.

$$a | b, b | c \Rightarrow a | (bx+cy)$$

Let R be an integral domain. Then $\forall a, b, c \in R$, if $a | b$ and $a | c$, necessarily $a | (bx+cy)$ $\forall x, y \in R$.

Proof. Since $a | b$ and $a | c$, so $ak = b$ and $al = c$ for some $k, l \in R$.

$$\begin{aligned} \text{Thus for any } x, y \in R, \\ bx+cy &= (ak)x + (al)y \\ &= a(kx+ly) \end{aligned}$$

and so $a | (bx+cy)$, as required. \square

ASSOCIATE

We can define an equivalence relation " \sim " on any ring R by stating that $a \sim b$ if $a | b$ and $b | a$.

why? $\rightarrow A4Q3$.

We say a and b are "associate" in R if $a \sim b$.

UNIT OF A RING

Let R be a ring. Then, an element $r \in R$ is called a "unit" of R if r has a multiplicative inverse in R .

We denote R^* to be the set of all units of R .

$$a \sim b \Leftrightarrow a = ub$$

Let R be an integral domain. Then, given any $a, b \in R$, we have $a \sim b$ if and only if $a = ub$, where $u \in R^*$.

Proof. First, suppose $a \sim b$ in R .

Then $a | b$ and $b | a$, so there exists $k, l \in R$ such that $b = ak$ and $a = bl$.

$$\text{Hence } b = ak = (bl)k = b(lk).$$

If $b = 0$, then $a = 0 \cdot l = 0$, and so $a = b = l \cdot b$, where $l \in R^*$.

On the other hand, if $b \neq 0$, then $b \cdot l = b \cdot (lk)$,

and by the cancellation property necessarily $lk = 1$.

So $l \in R^*$, and the proof follows.

Conversely, suppose $a = ub$ for some $u \in R^*$.

Then $b | a$ follows immediately.

But note $u^{-1}a = u^{-1}ub = b$, and so $a | b$ also.

Thus $a \sim b$, completing the proof. \square

DIVISION WITH REMAINDER

DIVISION WITH REMAINDER IN \mathbb{Z}

💡 Let $a, b \in \mathbb{Z}$ with $b > 0$. Then there exist unique integers q and r with $0 \leq r < b$, such that $a = bq + r$.

Proof. First, assume $a > 0$.

Let $S = \{n \in \mathbb{N} : n = a - bq, q \in \mathbb{Z}\}$.

Note $S \neq \emptyset$, since $a = a - b(0)$, so $a \in S$.

Thus, by WOP, S has a least element r .

Then by construction $r = a - bq$, so $a = bq + r$.

Moreover, since both

i) $r \geq 0$; and

ii) $r < b$; since if $r \geq b$, then $r - b > 0$, and so $r - b = a - bq - b = a - b(q+1)$, implying $r - b < r$, contradicting the minimality of r ,

we get that $0 \leq r < b$, and we are done. *

Conversely, if $a < 0$, then $-a > 0$, so by the above there exists $q_0, r_0 \in \mathbb{Z}$ such that $-a = bq_0 + r_0$ where $0 \leq r_0 < b$.

Then,

① if $r_0 = 0$, then $a = b(-q_0)$, so $q = -q_0$ and $r = 0$ satisfy the theorem's conditions; and

② if $r_0 \neq 0$, then $a = b(-q_0) - r_0 = b(-q_0 - 1) + (b - r_0)$, and so $q = -q_0 - 1$ and $r = b - r_0$ satisfy the conditions of the theorem. *

Hence, we have shown q, r exist $\forall a \in \mathbb{Z}$, and so we only need to prove uniqueness.

Suppose $\exists q', r' \in \mathbb{Z}$ such that $a = bq' + r'$ and $0 \leq r' < b$.

Then $bq + r = bq' + r'$, and so

$$r - r' = b(q - q').$$

If $q = q'$, necessarily $r = r'$, completing the uniqueness proof.

If $q \neq q'$, then by taking absolute values of both sides:

$$|r' - r| = |b||q - q'| \geq b.$$

But since $r < b$ and $r' < b$, this is impossible, so we must get that $q = q'$, and so $r = r'$, proving q and r are unique. ■

💡 Note that

① q is known as the "quotient" of the division; and

② r is known as the "remainder" of the division.

GREATEST COMMON DIVISOR

💡 Let R be an integral domain, and $a, b \in R$ be arbitrary, with $a \neq 0$ and $b \neq 0$.

Then, an element $d \in R$ is called a "greatest common divisor", or "gcd", of a and b if

① $d | a$ and $d | b$; and

② If $e \in R$ is another common divisor of a and b , so that $e | a$ and $e | b$, then necessarily $e | d$.

💡 Note the greatest common divisor of a and b may not be unique!

DIVISION WITH REMAINDER IN $\mathbb{Z}[i]$

💡 Let $\alpha, \beta \in \mathbb{Z}[i]$, with $\beta \neq 0$. Then there exists $\gamma, \delta \in \mathbb{Z}[i]$ such that $\alpha = \beta\gamma + \delta$, with $0 \leq N(\delta) < N(\beta)$.

Proof. Let $\alpha = abi$ and $\beta = cdi$, for some $a, b, c, d \in \mathbb{Z}$, with $c \neq 0$ and $d \neq 0$.

$$\text{Then } \frac{\alpha}{\beta} = \frac{abi}{cdi} = \frac{abi}{cdi} \cdot \frac{c-di}{c-di} = \frac{(ac+bd)i + (bc-ad)}{c^2+d^2} i$$

$$\text{with } r = \frac{ac+bd}{c^2+d^2} \in \mathbb{Q} \text{ and } s = \frac{bc-ad}{c^2+d^2} \in \mathbb{Q}.$$

Subsequently, choose $m, n \in \mathbb{Z}$ such that $|m - r| \leq \frac{1}{2}$ and $|n - s| \leq \frac{1}{2}$.

Let $\gamma = m + ni \in \mathbb{Z}[i]$ and $\delta = \alpha - \beta\gamma$.

Note $\delta \in \mathbb{Z}[i]$, as "+" and " \times " are closed in the ring $\mathbb{Z}[i]$.

Then, certainly $\alpha = \beta\gamma + \delta$.

Moreover, note that

$$\begin{aligned} \delta &= \alpha - \beta\gamma = \beta\left(\frac{\alpha}{\beta}\right) - \beta\gamma \\ &= \beta(r+si) - \beta(m+ni) \\ &= \beta((r-m) + (s-n)i), \end{aligned}$$

$$\text{so } N(\delta) = |\delta|^2 = |\beta((r-m) + (s-n)i)|^2 = |\beta|^2 ((r-m)^2 + (s-n)^2) = N(\beta)((r-m)^2 + (s-n)^2).$$

But since $|r-m| \leq \frac{1}{2}$ and $|s-n| \leq \frac{1}{2}$, consequently $(r-m)^2 \leq \frac{1}{4}$ and $(s-n)^2 \leq \frac{1}{4}$, and so $(r-m)^2 + (s-n)^2 \leq \frac{1}{2}$.

Hence $N(\delta) \leq N(\beta) \frac{1}{2} < N(\beta)$, proving the other condition of the theorem. ■

DIVISION ALGORITHM & DIVISOR FUNCTION

💡 Let R be an integral domain. Then, we say R has a "division algorithm" if there exists a function $d: R \setminus \{0\} \rightarrow \mathbb{N}$ such that $\forall a, b \in R$ with $b \neq 0$, there exist $q, r \in R$ such that $a = bq + r$, with $d(r) < d(b)$, or $r = 0$.

💡 In this case, we call d the "divisor function" of R .

RELATION BETWEEN \sim AND GCD'S

💡 Let R be an integral domain, and $a, b \in R \setminus \{0\}$ be arbitrary. Then,

① If d_1 and d_2 are both greatest common divisors

of a and b , then $d_1 \sim d_2$; and

② If d_1 is a greatest common divisor of a and b and $d_2 \in R$ is such that $d_2 \sim d_1$, then necessarily d_2 is a gcd of a and b .

Proof. First, suppose d_1 and d_2 are both gcds of a and b .

Then since d_1 is a common divisor of a and b , and d_2 is the greatest common divisor of a and b , necessarily $d_1 | d_2$.

Symmetrically, since d_1 is a gcd of a and b and d_2 is a common divisor of a and b , we must also get that $d_2 | d_1$.

Consequently since $d_1 | d_2$ and $d_2 | d_1$, so $d_1 \sim d_2$, proving ①.

Then, assume d_1 is a gcd of a and b , and $d_2 \sim d_1$.

So $d_2 | d_1$ and $d_1 | d_2$.

But since $d_1 | a$ and $d_1 | b$, the transitivity of divisibility tells us that $d_2 | a$ and $d_2 | b$.

So d_2 is a common divisor of a and b .

Furthermore, for any common divisor e of a and b , we know $e | d_1$, since d_1 is the gcd of a and b .

But since $d_1 | d_2$, transitivity of divisibility

once again yields that $e | d_2$, proving that d_2 is indeed a gcd of a and b .

proving ②. ■

THE EUCLIDEAN ALGORITHM

Let R be an integral domain with a division algorithm, and $a, b \in R$ be arbitrary.

Then, we can use the Euclidean algorithm to efficiently find the gcd of a and b .

In other words, the Euclidean algorithm shows that given any $a, b \in R \setminus \{0\}$, $\gcd(a, b)$ always exists.

FUNDAMENTAL IDEA

The fundamental idea that makes the Euclidean algorithm work is as follows:

Let R be an integral domain, and suppose there exist $a, b, q, r \in R$ such that $a = bq + r$.

Then, an element $d \in R$ is a gcd of a and b if and only if it is a gcd of b and r ,

$$\text{i.e. } \gcd(a, b) \sim \gcd(b, r).$$

Proof: First, suppose d is a gcd of a and b . This implies $d \mid a$ & $d \mid b$.

So $d \mid (a \cdot 1 + b \cdot (-q))$, telling us that $d \mid r$, so d is a common divisor of b and r .

Then, suppose e is a common divisor of b and r . Thus $e \mid b$ and $e \mid r$.

Hence $e \mid b \cdot q + r \cdot 1$, or $e \mid a$, showing that e is a common divisor of a and b . But since d is a gcd of a and b , it follows that necessarily $e \mid d$, and so d is also a gcd of b and r .

A similar proof can be used to prove the backward argument. ■

PROCEDURE

Let R be an integral domain with a division algorithm, and suppose D denotes the divisor function in this case.

Then, given some $a, b \in R$ with $a \neq 0$ and $b \neq 0$, we can calculate the gcd of a and b as follows:

① We first carry out a division with remainder to get that

$$a = bq_0 + r_1,$$

where $q_0, r_1 \in R$ and either $r_1 = 0$ or $D(r_1) < D(b)$.

② Then, if $r_1 = 0$, it implies $b \mid a$, and consequently that $b = \gcd(a, b)$.

③ Otherwise, if $r_1 \neq 0$, then we can use the lemma to the left to deduce that $\gcd(a, b) \sim \gcd(b, r_1)$.

④ So, we can carry out another division with remainder of b by r_1 to get that

$$b = r_1 q_1 + r_2,$$

where $q_1, r_2 \in R$ and either $r_2 = 0$ or $D(r_2) < D(r_1)$.

⑤ Again, if $r_2 = 0$, then $r_1 = \gcd(b, r_1)$, implying that $r_1 = \gcd(a, b)$ as well.

⑥ Otherwise, we can infer $\gcd(b, r_1) \sim \gcd(r_1, r_2)$, and so we are reduced to calculate a gcd of r_1 and r_2 .

⑦ We can continue this process to obtain a succession of divisions with remainder:

$$a = bq_0 + r_1$$

$$b = r_1 q_1 + r_2$$

$$r_1 = r_2 q_2 + r_3$$

$$\vdots$$

$$r_{n-2} = r_{n-1} q_{n-1} + r_n$$

$$r_{n-1} = r_n q_n + 0.$$

where r_n is the last non-zero remainder obtained.

Then, since $D(b) > D(r_1) > \dots$, we must eventually get a remainder of 0, as this is a strictly decreasing sequence of natural numbers.

⑧ But since $\gcd(a, b) \sim \gcd(b, r_1) \sim \dots \sim \gcd(r_{n-1}, r_n)$, and $r_n \mid r_{n-1}$,

we can consequently infer $r_n = \gcd(r_n, r_{n-1})$, and so

$r_n = \gcd(a, b)$ also.

EXAMPLE: $\gcd(1009, 33)$

We can use the Euclidean algorithm to find the gcd of 1009 and 33.

First, we perform divisions with remainder:

$$1009 = 33 \cdot 30 + 19$$

$$33 = 19 \cdot 1 + 14$$

$$19 = 14 \cdot 1 + 5$$

$$14 = 5 \cdot 2 + 4$$

$$5 = 4 \cdot 1 + 1$$

$$4 = 1 \cdot 4 + 0$$

So the last non-zero remainder is 1,

implying $1 = \gcd(1009, 33)$.

THE EXTENDED EUCLIDEAN ALGORITHM

The Euclidean algorithm can be used to find elements $x, y \in \mathbb{R}$ such that $ax + by = d$, where $d = \gcd(a, b)$.

PROCEDURE

We can accomplish the above by the following:

- Suppose after running the Euclidean algorithm on a and b , we generate divisions with remainder

$$\begin{aligned}a &= bq_0 + r_1 && * \text{we leave out} \\b &= r_1 q_1 + r_2 && \text{the last step.} \\r_1 &= r_2 q_2 + r_3\end{aligned}$$

:

$$r_{n-2} = r_{n-1} q_{n-1} + r_n$$

- We now reverse the order of the equations, and isolate the remainder in each one:

$$r_n = r_{n-2} - r_{n-1} q_{n-1}$$

$$r_{n-1} = r_{n-3} - r_{n-2} q_{n-2}$$

:

$$r_2 = b - r_1 q_1$$

$$r_1 = a - b q_0.$$

- Then, we can "back-substitute" the below equations into the above one, which will eventually terminate when we have expressed $\gcd(a, b) = r_n = ax + by$ for some $x, y \in \mathbb{R}$, which is what we wanted to do.

EXAMPLE: $1009x + 33y = 1$

We can use the Extended Euclidean algorithm to find integer solutions to the equation $1009x + 33y = 1$.

First, we reverse and solve for each non-zero remainder the equations we obtained when applying the Euclidean algorithm:

$$\begin{aligned}1 &= 5 - 4 \cdot 1 && \textcircled{1} \\4 &= 14 - 5 \cdot 2 && \textcircled{2} \\5 &= 19 - 14 \cdot 1 && \textcircled{3} \\14 &= 33 - 19 \cdot 1 && \textcircled{4} \\19 &= 1009 - 23 \cdot 30 && \textcircled{5}.\end{aligned}$$

Then, observe if we substitute \textcircled{2} into \textcircled{1}, we get that

$$1 = 5 - (14 - 5 \cdot 2) \cdot 1 = 5 \cdot 3 - 14 \cdot 1.$$

If we substitute \textcircled{3} into this new equation, we subsequently get that

$$1 = 5 \cdot 3 - 14 \cdot 1 = (19 - 14 \cdot 1) \cdot 3 - 14 \cdot 1 = 19 \cdot 3 - 14 \cdot 4.$$

We can keep "back-substituting" like this, eventually arriving at the conclusion that

$$1 = 1009 \cdot 7 - 33 \cdot 214,$$

and so $1009x + 33y = 1$ has a solution $x=7$ and $y=-214$ over \mathbb{Z} .

LINEAR DIOPHANTINE EQUATIONS

Suppose R is an integral domain with a division algorithm.

Then, a linear Diophantine equation (in 2 variables) is any equation of the form $ax+by=c$, where $a, b, c \in R$ and x, y are the solutions to the equation.

IF A LINEAR DIOPHANTINE EQUATION HAS A SOLUTION, THEN $\gcd(a, b) | c$

Let R be an integral domain with a division algorithm, and let $a, b, c \in R$ be arbitrary. Then, if there exists a solution to the equation $ax+by=c$, with $x, y \in R$, then $\gcd(a, b) | c$.

Proof. Let $d = \gcd(a, b)$. Then $d | a$ and $d | b$, implying that $d | (ax+by) = c$. So $d | c$, and we are done. \blacksquare

IF $\gcd(a, b) | c$, $ax+by=c$ HAS A SOLUTION

Let R be an integral domain with a division algorithm, and $a, b, c \in R$ be arbitrary, with $a \neq 0$ and $b \neq 0$. Then, if $\gcd(a, b) | c$, the equation $ax+by=c$ has a solution with $x, y \in R$.

Proof. Let $d = \gcd(a, b)$. Then, the Euclidean algorithm tells us that $\exists x_0, y_0 \in R$ such that $ax_0 + by_0 = d$. Since $d | c$, it follows that $\exists k \in \mathbb{Z}$ such that $c = kd$, and so $c = kd = k(ax_0 + by_0) = a(kx_0) + b(ky_0)$. So that $ax+by=c$ has a solution $x = kx_0$ and $y = ky_0$. \blacksquare

$a | bc$ & $1 \sim \gcd(a, b) \Rightarrow a | c$

Suppose R is an integral domain with a division algorithm, and let $a, b, c \in R$ be arbitrary. Then, if both $a | bc$ and $1 \sim \gcd(a, b)$, necessarily $a | c$.

Proof. Since $a | bc$, there must exist a $k \in R$ such that $ak = bc$.

Moreover, since 1 is a gcd of a and b , we know there must exist $x, y \in R$ such that $ax+by=1$.

Thus, if we multiply both sides by c , we get $acx+bcy=c$.

But since $bc = ak$, so

$$acx+aky=c$$

$$\Rightarrow a(cx+ky)=c,$$

proving $a | c$, as required. \blacksquare

$d = \gcd(a, b) \Rightarrow a = da_0, b = db_0$

$\Rightarrow \gcd(a_0, b_0) \sim 1$

Let R be an integral domain with a division algorithm, and let $a, b \in R \setminus \{0\}$ be arbitrary.

Suppose $d = \gcd(a, b)$, so that $a = da_0$ and $b = db_0$ for some $a_0, b_0 \in R$.

Then $\gcd(a_0, b_0) \sim 1$ necessarily.

Proof. First, we know that $\exists x, y \in R$ such that $ax+by=d$.

$$\text{So } (da_0x) + (db_0y) = d$$

and cancelling out the $d \neq 0$ yields that $a_0x+b_0y=1$, which shows that $\gcd(a_0, b_0) \mid 1$, and so $\gcd(a_0, b_0) \sim 1$ (as $1 \mid \gcd(a_0, b_0)$ holds trivially.) \blacksquare

THE GENERAL SOLUTION OF A LINEAR DIOPHANTINE EQUATION

Let R be an integral domain with a division algorithm, and $a, b, c \in R$ be such that $a \neq 0$ & $b \neq 0$. Let $d = \gcd(a, b)$ be such that $d | c$, and write $a = da_0$ and $b = db_0$ for some $a_0, b_0 \in R$.

Then, the complete set of solutions to the equation $ax+by=c$, where $x, y \in R$,

is given by

$$(x, y) = (x_0 + kb_0, y_0 - ka_0),$$

where $k \in R$ is arbitrary, and (x_0, y_0) is a particular solution to $ax+by=c$.

Proof. By assumption, (x_0, y_0) is a solution to $ax+by=c$, implying that $ax_0+by_0=c$. —①

Let (x_1, y_1) be another solution to the equation, so that $ax_1+by_1=c$. —②

Subtracting ① from ② yields that

$$a(x_1-x_0) + b(y_1-y_0) = 0,$$

and making the substitutions $a = da_0$ and $b = db_0$ gives us

$$(da_0)(x_1-x_0) + (db_0)(y_1-y_0) = 0.$$

Rearranging this gets us that $a_0(x_1-x_0) = -b_0(y_1-y_0)$.

Hence $b_0 | a_0(x_1-x_0)$. But since $\gcd(a_0, b_0) \sim 1$, we can infer that $b_0 | (x_1-x_0)$, and so $x_1-x_0 = kb_0$ for some $k \in R$. (This implies $x_1 = x_0 + kb_0$.)

$$\text{So } a_0(kb_0) = -b_0(y_1-y_0),$$

and thus $ka_0 = -(y_1-y_0)$; then, if we solve for y_1 ,

$$\text{we get that } y_1 = y_0 - ka_0.$$

Thus if (x_1, y_1) is a solution to the equation, then

$$(x_1, y_1) = (x_0 + kb_0, y_0 - ka_0). \blacksquare$$

Conversely, note if $x_1 = x_0 + kb_0$ and $y_1 = y_0 - ka_0$, then

$$\begin{aligned} ax_1 + by_1 &= a(x_0 + kb_0) + b(y_0 - ka_0) \\ &= (ax_0 + by_0) + k(ab_0 - ba_0) \\ &= c + k((da_0)b_0 - (db_0)a_0) \\ &= c + 0 \\ &= c, \end{aligned}$$

and so (x_1, y_1) is a solution to the equation. \blacksquare

EXAMPLE: FINDING THE GENERAL SOLUTION

TO $1009x + 33y = 5$

We can use the formula given above to find all the solutions to $1009x + 33y = 5$ (in \mathbb{Z}).

First, by running the Euclidean algorithm on the eqn, we get that $\gcd(1009, 33) = 1$ and that

$$1009(7) + 33(-214) = 1.$$

Then, since $\gcd(1009, 33) \mid 5$, the above eqn has a solution.

Subsequently, multiplying both sides by 5 yields that

$$1009(35) + 33(-1070) = 5,$$

and so our equation has the particular solution $(x_0, y_0) = (35, -1070)$.

To get the general solution, note $a_0 = a = 1009$ & $b_0 = b = 33$, because the gcd is 1 in this case, and so the general solution of this eqn is

$$(x, y) = (35 + 33k, -1070 - 1009k),$$

where $k \in \mathbb{Z}$ is arbitrary.

MULTIPLICATIVE INVERSES IN $\mathbb{Z}/n\mathbb{Z}$

We can use linear Diophantine equations to calculate the multiplicative inverse of an element in $\mathbb{Z}/n\mathbb{Z}$, if it exists.

In particular, we can show $[a] \in \mathbb{Z}/n\mathbb{Z}$ has a multiplicative inverse if $ax+ny=1$ has a solution.

Proof. Observe $[a]$ has a multiplicative inverse if and only if there exists a $[x] \in \mathbb{Z}/n\mathbb{Z}$ such that $[a][x] = [1]$. This is equivalent to the assertion that $[ax] = [1]$, or that $ax \equiv 1 \pmod{n}$. In turn, this is the same as saying $n \mid (ax-1)$; ie $ny = 1-ax$ for some $y \in \mathbb{Z}$, or in other words whether $ax+ny=1$ has a solution. \square

Hence, this is only possible if $\gcd(a, n) = 1$.

Note if $n=p$, where p is prime, then $\gcd(a, p) = 1 \quad \forall [a] \in (\mathbb{Z}/p\mathbb{Z}) \setminus [0]$, and so $\mathbb{Z}/p\mathbb{Z}$ is a field.

POLYNOMIALS OVER A FIELD

Q1: Let F be a field. Then, we define the set of polynomials with coefficients in F , denoted by $F[x]$, by

$$F[x] = \{a_0 + a_1x + a_2x^2 + \dots + a_nx^n : n \geq 0, a_i \in F \forall i\}.$$

Q2: We can turn $F[x]$ into a commutative ring; given a $f = \sum_{i=0}^m a_i x^i$ and $g = \sum_{i=0}^n b_i x^i$, we can define an addition and multiplication by

$$f+g = \sum_{i=0}^{\max(m,n)} (a_i + b_i) x^i$$

and

$$fg = \sum_{i=0}^{m+n} c_i x^i, \quad \text{with } c_i = \sum_{j=0}^i a_j b_{i-j} \text{ for each } i.$$

DEGREE

Q: Let $F[x]$ be a set of polynomials with coefficients in some field F . Let $f \in F[x]$ be arbitrary. Then, the "degree" of f , denoted by $\deg(f)$, is the largest index i such that $a_i \neq 0$.

$$\deg(f+g) \leq \max(\deg(f), \deg(g))$$

Q: Let F be an arbitrary field, and $f, g \in F[x]$ be arbitrary. Then, if $\deg(fg) \neq 0$, then $\deg(fg) \leq \max(\deg(f), \deg(g))$.

Proof: By definition, $fg = \sum_{i=0}^{\max(m,n)} (a_i + b_i) x^i$.

So all coefficients of powers of x after $\max(m,n)$ are equal to 0, so if $fg \neq 0$, then $\deg(fg) \leq \max(m,n) = \max(\deg(f), \deg(g))$.

$$\deg(fg) = \deg(f) + \deg(g)$$

Q: Let F be an arbitrary field, and $f, g \in F[x]$ be arbitrary.

$$\text{Then } \deg(fg) = \deg(f) + \deg(g).$$

Proof: By defn, $fg = \sum_{i=0}^m c_i x^i$; hence all coefficients of powers of x after $m+n$ are equal to 0, implying $\deg(fg) \leq \deg(f) + \deg(g)$.

Next, observe that

$$c_{m+n} = \sum_{j=0}^m a_j b_{m+n-j} = a_0 b_{m+n} + a_1 b_{m+n-1} + \dots + a_m b_n + \dots + a_{m+n} b_0.$$

Then, since $a_j = 0 \forall j > m$, so all terms after $a_m b_n$ in the sum above are equal to 0.

Likewise, since $b_{m+n-j} = 0$ when $j < n$, all terms before $a_m b_n$ in the sum above are equal to 0.

Hence $c_{m+n} = a_m b_n$, and since $a_m \neq 0$ & $b_n \neq 0$, the fact that F is an integral domain tells us that $a_m b_n = c_{m+n} \neq 0$.

This is sufficient to prove that $\deg(fg) = m+n = \deg(f) + \deg(g)$. \square

$F[x]$ IS AN INTEGRAL DOMAIN

Q: Let F be an arbitrary field. We can show that $F[x]$ must be an integral domain.

Proof: Observe if $f, g \in F[x] \setminus \{0\}$, then $\deg(f)$ and $\deg(g)$ exist, and by the above $\deg(fg) = \deg(f) + \deg(g)$. So $fg \neq 0$; combined with the fact that $F[x]$ is a commutative ring, this is sufficient to show that $F[x]$ is an integral domain. \square

$F[x]$ HAS A DIVISION ALGORITHM

Q: Let F be an arbitrary field. Then, the integral domain $F[x]$ admits a division algorithm with divisor function $\deg(\cdot)$; ie for any polynomials $f, g \in F[x]$ with $g \neq 0$, there exist $q, r \in F[x]$ such that $f = gq+r$, and either $r=0$ or $\deg(r) < \deg(g)$.

Proof: Let $S = \{f - gg : g \in F[x]\}$.

If $0 \in S$, then there exists a $g \in F[x]$ such that $f = gg$, and we are done.

Otherwise, let $r \in S \setminus \{0\}$ be arbitrary.

By construction, $r = f - gg$ for some $g \in F[x]$, and so $f = gg+r$ as needed. *

Next, suppose $\deg(r) \geq \deg(g)$,

and let $r = a_n x^n + a_{n-1} x^{n-1} + \dots$ and $g = b_m x^m + b_{m-1} x^{m-1} + \dots$, where $b_m \neq 0$ since $g \neq 0$, and $n \geq m$.

Then, since $b_m \neq 0$ and F is a field, b_m^{-1} must exist.

$$\begin{aligned} \text{Let } r_1 &= r - a_n b_m^{-1} x^{n-m} g \\ &= (a_n x^n + a_{n-1} x^{n-1} + \dots) - (a_n x^n + a_n b_m^{-1} b_{m-1} x^{n-1} + \dots) \\ &= (a_{n-1} - a_n b_m^{-1} b_{m-1}) x^{n-1} + \dots. \end{aligned}$$

so that $\deg(r_1) < \deg(r)$.

However, $r_1 = r - a_n b_m^{-1} x^{n-m} g = (f - gg) - a_n b_m^{-1} x^{n-m} g = f - (g + a_n b_m^{-1} x^{n-m})g$, so $r_1 \in S$, contradicting our choice of r as an element of S of smallest degree.

Therefore $\deg(r) < \deg(g)$ after all, verifying that a decomposition $f = gg+r$ where $r=0$ or $\deg(r) < \deg(g)$ does exist.

EVALUATION

Q: Let F be an arbitrary field, and $f \in F[x]$ be arbitrary.

Then, the "evaluation of f " at c , where $c \in F$, is defined to be

$$f(c) = a_n c^n + a_{n-1} c^{n-1} \dots + a_0.$$

ROOT

Q: Let F be an arbitrary field, and $f \in F[x]$ be arbitrary.

Then a $c \in F$ is a "root of f " if $f(c) = 0$.

EVALUATION HOMOMORPHISM

Q: Let F be a field. Then, the "evaluation homomorphism" of F at some $c \in F$ is defined to be the mapping $\phi_c : F[x] \rightarrow F$ given by $\phi_c(f) = f(c) \quad \forall f \in F$.

Q2: Note that ϕ_c is a ring homomorphism.

THE FACTOR THEOREM

Let F be a field, and $c \in F$ be arbitrary. Then c is a root of f if and only if $(x-c) \mid f$ in $F[x]$; ie $\ker \phi_c = F[x](x-c)$.

Proof. First, assume $(x-c) \mid f$.

Then, by definition, there exists some $q \in F[x]$ such that $f = (x-c)q$.

Note $\phi_c(x-c) = c - c = 0$; thus,

$f(c) = \phi_c(f) = \phi_c(x-c)\phi_c(q) = 0 \cdot \phi_c(q) = 0$, showing that c is a root of f .

Conversely, suppose c is a root of f .

Applying division with remainder of f by $(x-c)$, we know there necessarily exists $q, r \in F[x]$ such that

$$f = (x-c)q + r,$$

where $r=0$ or $\deg(r) < \deg(x-c) = 1$.

Either way, this guarantees r is a constant polynomial, in the form $r=r_0$ for some $r_0 \in F$.

So, applying ϕ_c to the x^k yields that

$$\begin{aligned} 0 &= f(c) = \phi_c(f) = \phi_c((x-c)q + r) \\ &= \phi_c(q)\phi_c(x-c) + r_0 \\ &= 0 \cdot \phi_c(x-c) + r_0 \\ \therefore 0 &= r_0, \end{aligned}$$

Showing that $f = (x-c)q$, so $(x-c) \mid f$.

Thus $f(c)=0$ exactly when f is multiple of $(x-c)$,

implying that $\ker \phi_c = F[x](x-c)$. \blacksquare

f HAS AT MOST $\deg(f)$ DISTINCT ROOTS IN F (IF $f \neq 0$)

Q: Let F be a field. Then, any polynomial $f \in F[x] \setminus \{0\}$ has at most $\deg(f)$ distinct roots in F .

Proof. We prove this by induction.

First, if $\deg(f)=0$, then $f=f_0$, where $f_0 \in F \setminus \{0\}$. Hence $f_0 \neq 0 \forall c \in F$, showing f has no roots in F , establishing our first base case.

Similarly, if $\deg(f)=1$, then $f=ax+b$ for some $a, b \in F$, with $a \neq 0$.

Then note $f(a)=0$ iff $ac+b=0$, iff $ac=-b$, iff $c=-a^{-1}b$; so f has exactly one root in F , establishing our second base case.

Next, suppose $\deg(f)=n+1$ for some $n \geq 1$, $n \in \mathbb{Z}^+$, and assume all polynomials of degree $k < n+1$ have at most k roots in F .

If f has no roots in F the result is trivially true, so assume $c \in F$ is a root of f .

By the Factor Theorem, thus $f=(x-c)f_n$ for some $f_n \in F$.

Taking degrees of both sides yields that $\deg(f_n)=n$.

Then, note $\forall a \in F$, $f(a)=0$ if and only if $(a-c)f_n(a)=0$, which holds if and only if $a=c$ or $f_n(a)=0$.

But $f_n(a)=0$ has at most n distinct values for at most n distinct values of $a \in F$; hence, $f(a)=0$ for at most $n+1$ distinct values of a .

The claim follows by induction. \blacksquare

PRIMITIVE ROOTS MODULO p

$g \in G$, $\alpha(g)$ IS MAXIMAL $\Rightarrow h^{\alpha(g)} \forall h \in G$

Q: Let G be a finite abelian group, and suppose $g \in G$ is such that $\alpha(g)$ is maximal. Then, if $\alpha(g)=k$, then $h^k=e \forall h \in G$.

(Proof in A9)

F^* IS CYCLIC (F IS A FINITE FIELD)

Q: Let F be a finite field. Then the group F^* of units of F is cyclic.

Proof. Since F^* is a finite abelian group, we can choose an element $a \in F^*$ that has maximal order. Let this order be k .

Then by the above lemma, $a^k=1 \forall a \in F^*$.

In particular, the polynomial $x^k-1 \in F[x]$ has at least $|F^*|$ distinct roots, so $|F^*| \leq k$.

But we also know x^k-1 can only contain at most $|F^*|$ distinct roots, so $k \leq |F^*|$.

Thus $k=|F^*|$, and so there exists an element of F^* having order $|F^*|$.

Hence $|F^*|$ is cyclic, which we wanted to prove. \blacksquare

$(\mathbb{Z}/p\mathbb{Z})^*$ IS CYCLIC (p IS A PRIME)

Q: For any prime p , we can show that the group $(\mathbb{Z}/p\mathbb{Z})^*$ is cyclic.

Proof. Note for any prime p , the ring $\mathbb{Z}/p\mathbb{Z}$ is a finite field.

So, by the above theorem, necessarily we must get that $\mathbb{Z}/p\mathbb{Z}$ is cyclic. \blacksquare

CHINESE REMAINDER THEOREM

COPRIME / RELATIVELY PRIME

Let $a, b \in \mathbb{Z}$ be arbitrary.

Then, we say a and b are "coprime"

if $\gcd(a, b) = 1$.

PAIRWISE COPRIME

Let m_1, m_2, \dots, m_k be a list of integers.

We say this list is "pairwise coprime" if

every pair of distinct elements is coprime;

i.e. $\gcd(m_i, m_j) = 1$ if $i \neq j$.

CHINESE REMAINDER THEOREM FOR \mathbb{Z}

Let a_1, a_2, \dots, a_k denote arbitrary integers,

and let m_1, m_2, \dots, m_k be a list of

pairwise coprime integers.

Then, the system of congruences

$$\begin{cases} x \equiv a_1 \pmod{m_1} \\ x \equiv a_2 \pmod{m_2} \\ \vdots \\ x \equiv a_k \pmod{m_k} \end{cases}$$

has a solution $x \in \mathbb{Z}$, and this solution is also unique in modulo $m_1 m_2 \dots m_k$; i.e. if $y \in \mathbb{Z}$ is a solution to the system above, then $x \equiv y \pmod{m_1 m_2 \dots m_k}$.

Proof. First, assume there exists integers

b_1, b_2, \dots, b_k such that

$$\begin{cases} b_1 \equiv 1 \pmod{m_1} \\ b_2 \equiv 0 \pmod{m_2} \\ \vdots \\ b_k \equiv 0 \pmod{m_k} \end{cases}, \quad \begin{cases} b_1 \equiv 0 \pmod{m_1} \\ b_2 \equiv 1 \pmod{m_2} \\ \vdots \\ b_k \equiv 0 \pmod{m_k} \end{cases}, \quad \dots, \quad \begin{cases} b_1 \equiv 0 \pmod{m_1} \\ b_2 \equiv 0 \pmod{m_2} \\ \vdots \\ b_k \equiv 1 \pmod{m_k} \end{cases}.$$

Then, note that

$$\begin{aligned} \sum_{i=1}^k a_i b_i &= a_1 b_1 + a_2 b_2 + \dots + a_k b_k \\ &\equiv a_1(1) + a_2(0) + \dots + a_k(0) \pmod{m_1} \\ &\equiv a_1 \pmod{m_1}, \end{aligned}$$

so that $\sum_{i=1}^k a_i b_i \equiv a_1 \pmod{m_1}$.

Similar computations show $\sum_{i=1}^k a_i b_i \equiv a_j \pmod{m_j} \quad \forall j \in \{1, 2, \dots, k\}$,

so that $x = \sum_{i=1}^k a_i b_i$ is a solution to the system of equations.

We now show why b_1 exists.

Since $b_i \not\equiv 0 \pmod{m_i} \quad \forall i \in \{2, 3, \dots, k\}$, it follows

that $b_1 \mid m_j \quad \forall j \in \{2, 3, \dots, k\}$.

Let $M_1 = m_2 m_3 \dots m_k$.

Then if $b_1 = c_1 M_1$ for some $c_1 \in \mathbb{Z}$, then b_1 automatically satisfies $b_1 \equiv 0 \pmod{m_i}$ for $2 \leq i \leq k$.

Next, if b_1 satisfies the first congruence, we need to find a $c_1 \in \mathbb{Z}$ such that $c_1 M_1 \equiv 1 \pmod{m_1}$.

We know this can only happen if $[c_1] = [M_1]^{-1}$ in $\mathbb{Z}/m_1 \mathbb{Z}$,

which in turn can only occur iff $\gcd(M_1, m_1) = 1$.

But observe if $\gcd(M_1, m_1) \neq 1$, it implies M_1 & m_1 share a prime factor.

By defn, this factor would divide both m_1 and one of the other integers in $M_1 = m_2 m_3 \dots m_k$, implying this prime is a common factor of m_1 and some m_i for $i \neq 1$.

This is a contradiction, as $m_1, m_2, m_3, \dots, m_k$ is pairwise coprime by assumption!

Thus M_1 necessarily has a multiplicative inverse in $\mathbb{Z}/m_1 \mathbb{Z}$,

implying $\exists c_1 \in \mathbb{Z}$ such that $c_1 M_1 \equiv 1 \pmod{m_1}$.

So letting $b_1 = c_1 M_1$ gives us the integer we want.

A similar proof can be used to show why b_2, b_3, \dots, b_k exist as well, (we can set $b_j = c_j M_j$, where $c_j \in \mathbb{Z}$ and $M_j = \prod_{i \neq j} m_i$) and this is sufficient to show a solution exists. \square

Now, we can show why this solution is unique.

Suppose $x, y \in \mathbb{Z}$ both satisfy the system of congruences.

Then, in particular, $x \equiv y \pmod{m_i} \quad \forall i \in \{1, 2, \dots, k\}$,

implying $m_i \mid (x-y) \quad \forall i$.

So, by A9&4, since m_1, m_2, \dots, m_k are pairwise coprime,

it follows that $(m_1 m_2 \dots m_k) \mid (x-y)$,

implying $[x] = [y]$ in $\mathbb{Z}/m_1 m_2 \dots m_k \mathbb{Z}$. \square

EXAMPLE: $x \equiv 2 \pmod{3}$, $x \equiv 3 \pmod{5}$, $x \equiv 2 \pmod{7}$

We can use the Chinese Remainder Theorem to find the solutions $x \in \mathbb{Z}$ to the system

$$\begin{cases} x \equiv 2 \pmod{3} \\ x \equiv 3 \pmod{5} \\ x \equiv 2 \pmod{7}. \end{cases}$$

To start, let $M_1 = 5 \cdot 7 = 35$, $M_2 = 3 \cdot 7 = 21$ and $M_3 = 3 \cdot 5 = 15$, using the notation from the proof.

Then c_1, c_2, c_3 are determined by solving

$$\begin{aligned} \textcircled{1} \quad 35c_1 &\equiv 1 \pmod{3} \\ &\Rightarrow 35c_1 - 36c_1 \equiv 1 - 36 \pmod{3} \\ &\Rightarrow -c_1 \equiv 1 \pmod{3} \quad \because 36c_1 \equiv 0 \pmod{3} \\ &\therefore c_1 = 2; \\ \textcircled{2} \quad 21c_2 &\equiv 1 \pmod{5}; \text{ and} \\ &\Rightarrow 21c_2 - 20c_2 \equiv 1 - 20c_2 \pmod{5} \\ &\Rightarrow c_2 \equiv 1 \pmod{5} \quad \because 20c_2 \equiv 0 \pmod{5} \\ &\therefore c_2 = 1; \\ \textcircled{3} \quad 15c_3 &\equiv 1 \pmod{7}. \\ &\Rightarrow 15c_3 - 14c_3 \equiv 1 - 14c_3 \pmod{7} \\ &\Rightarrow c_3 \equiv 1 \pmod{7} \quad \because 14c_3 \equiv 0 \pmod{7} \\ &\therefore c_3 = 1. \end{aligned}$$

So, we get that $c_1 = -1$, $c_2 = 1$ and $c_3 = 1$.

Hence $b_1 = c_1 M_1 = -1(35) = -35$;

$b_2 = c_2 M_2 = 1(21) = 21$; and

$b_3 = c_3 M_3 = 1(15) = 15$,

so our solution to our system is

$$\begin{aligned} x &= a_1 b_1 + a_2 b_2 + a_3 b_3 \\ &= 2(-35) + 3(21) + 2(15) \\ &\Rightarrow x = 23. \end{aligned}$$

This solution is unique modulo $3 \cdot 5 \cdot 7 = 105$, so that if $y \in \mathbb{Z}$ is another integer solution to the system, then $y \equiv 23 \pmod{105}$.

ISOMORPHISM

Let $\phi: R \rightarrow S$ be a ring homomorphism.

Then, ϕ is an "isomorphism" if

ϕ is also a bijection.

Alternatively, ϕ is also an isomorphism

if there exists a ring homomorphism

$\psi: S \rightarrow R$ such that $\psi \circ \phi = \text{id}_R$ and $\phi \circ \psi = \text{id}_S$, where id_R and id_S are the identity maps on R and S respectively.

If an isomorphism exists between rings R and S , then we say R and S are "isomorphic", and write $R \cong S$.

Note the relation " \cong " on rings is an equivalence relation.

$(\mathbb{Z}/m_1\mathbb{Z} \cdots \mathbb{Z}/m_k\mathbb{Z}) \cong (\mathbb{Z}/m_1\mathbb{Z})(\mathbb{Z}/m_2\mathbb{Z}) \cdots (\mathbb{Z}/m_k\mathbb{Z})$

Suppose m_1, m_2, \dots, m_k are pairwise coprime positive integers. Then we must have that

$$(\mathbb{Z}/m_1\mathbb{Z} \cdots \mathbb{Z}/m_k\mathbb{Z}) \cong (\mathbb{Z}/m_1\mathbb{Z})(\mathbb{Z}/m_2\mathbb{Z}) \cdots (\mathbb{Z}/m_k\mathbb{Z}).$$

Proof. First, let $[x]_n = n\mathbb{Z} + x$ for any $n \in \mathbb{Z}$, $x \in \mathbb{Z}$. Then, let the homomorphism between rings

$\phi: (\mathbb{Z}/m_1\mathbb{Z} \cdots \mathbb{Z}/m_k\mathbb{Z}) \rightarrow (\mathbb{Z}/m_1\mathbb{Z})(\mathbb{Z}/m_2\mathbb{Z}) \cdots (\mathbb{Z}/m_k\mathbb{Z})$ be defined by

$$\phi([x]_{m_1 m_2 \cdots m_k}) = ([x]_{m_1}, [x]_{m_2}, \dots, [x]_{m_k}) \quad \forall [x]_{m_1 m_2 \cdots m_k} \in \mathbb{Z}/m_1\mathbb{Z} \cdots \mathbb{Z}/m_k\mathbb{Z}.$$

First, we show this map is well-defined.

If $[x]_{m_1 m_2 \cdots m_k} = [y]_{m_1 m_2 \cdots m_k}$, then $(m_1 m_2 \cdots m_k) | (x-y)$.

In particular, we have that $m_1 | (x-y)$, $m_2 | (x-y) \cdots$, $m_k | (x-y)$, so $[x]_{m_i} = [y]_{m_i}$ for each $1 \leq i \leq k$.

Hence $\phi([x]_{m_1 m_2 \cdots m_k}) = \phi([y]_{m_1 m_2 \cdots m_k})$, showing this map is well-defined.

Then, we can similarly check ϕ preserves addition, multiplication and the unity.

To show ϕ is injective, we show $\ker \phi = \{0\}$.

Suppose we have a coset $[x]_{m_1 m_2 \cdots m_k}$ such that

$$\phi([x]_{m_1 m_2 \cdots m_k}) = [0]_{m_1}, [0]_{m_2}, \dots, [0]_{m_k}.$$

Then x satisfies the system of congruences

$$\begin{cases} x \equiv 0 \pmod{m_1} \\ x \equiv 0 \pmod{m_2} \\ \vdots \\ x \equiv 0 \pmod{m_k}. \end{cases}$$

Clearly $x=0$ is a solution to the system, and by the uniqueness of solutions in the CRT we can deduce that $x \equiv 0 \pmod{m_1 m_2 \cdots m_k}$, so $[0]_{m_1 m_2 \cdots m_k} = [x]_{m_1 m_2 \cdots m_k}$.

So $\ker \phi$ is the zero ideal, proving that ϕ is injective.*

Next, to show ϕ is surjective, suppose we are given an arbitrary element

$$([a_1]_{m_1}, [a_2]_{m_2}, \dots, [a_k]_{m_k}) \in \prod_{i=1}^k (\mathbb{Z}/m_i\mathbb{Z}).$$

By CRT, $\exists x \in \mathbb{Z}$ such that

$$\begin{cases} x \equiv a_1 \pmod{m_1} \\ x \equiv a_2 \pmod{m_2} \\ \vdots \\ x \equiv a_k \pmod{m_k} \end{cases}$$

and since

$$\phi([x]_{m_1 m_2 \cdots m_k}) = ([a_1]_{m_1}, [a_2]_{m_2}, \dots, [a_k]_{m_k}),$$

it follows that ϕ is surjective.

So, since ϕ is both injective & surjective, it follows that ϕ is bijective, and so ϕ is an isomorphism

from $\mathbb{Z}/m_1\mathbb{Z} \cdots \mathbb{Z}/m_k\mathbb{Z}$ to $(\mathbb{Z}/m_1\mathbb{Z})(\mathbb{Z}/m_2\mathbb{Z}) \cdots (\mathbb{Z}/m_k\mathbb{Z})$. \blacksquare

FIELD OF FRACTIONS OF AN INTEGRAL DOMAIN

THE "FRACTION" EQUIVALENCE RELATION

Q: Let R be an integral domain. Let the relation \sim on $R \times (R \setminus \{0\})$ be such that $(a,b) \sim (c,d)$ if $ad = bc$. Then \sim is an equivalence relation.

Proof. First, since $ab = ba$, it follows that $(a,b) \sim (a,b)$, showing \sim is reflexive.

Then, if $(a,b) \sim (c,d)$, then necessarily $ad = bc$; thus $cb = da$, so that $(c,d) \sim (a,b)$.

Lastly, suppose we are given $(a,b), (c,d), (e,f) \in R \times (R \setminus \{0\})$ such that $(a,b) \sim (c,d)$ and $(c,d) \sim (e,f)$.

By defn, this implies $ad = bc$ and $cf = de$.

Multiplying the former eqn by f yields

$$\begin{aligned} adf &= bcf \\ \Rightarrow adf &= b(de) \\ \Rightarrow af &= be(de). \end{aligned}$$

Then, since R is an integral domain and $d \neq 0$, we can use the cancellation property to conclude that $af = be$, so that $(a,b) \sim (e,f)$, and thus that \sim is transitive. \blacksquare

Q: We denote the set of equivalence classes of \sim by $\mathbb{Q}(R)$.

Example: \mathbb{Q}

$[(a,b)]$ "stands for" $\frac{a}{b}$.

FIELD OF FRACTIONS

Q: Let R be an integral domain. Then, we can define an addition operation on $\mathbb{Q}(R)$ by

$$[(a,b)] + [(c,d)] = [(ac+bd, bd)]$$

and a multiplication operation by

$$[(a,b)] \times [(c,d)] = [(ac, bd)].$$

We can show that $\mathbb{Q}(R)$ is a field with respect to these two operations, called the "field of fractions" of R .

Proof. First, suppose $[(a_1, b_1)] = [(a_2, b_2)]$ and $[(c_1, d_1)] = [(c_2, d_2)]$.

Then, this implies $a_1b_2 = b_1a_2$ & $c_1d_2 = d_1c_2$.

Subsequently, note that

$$\begin{aligned} (a_1d_1 + b_1c_1)c_2d_2 &= (a_1b_2)d_1d_2 + (c_1d_2)b_1b_2 \\ &= (b_1a_2)d_1d_2 + (d_1c_2)b_1b_2 \\ &= (a_2d_2 + b_2c_2)c_1d_1, \end{aligned}$$

and thus $[(a_1d_1 + b_1c_1, b_1d_1)] = [(a_2d_2 + b_2c_2, b_2d_2)]$,

so $[(a_1, b_1)] + [(c_1, d_1)] = [(a_2, b_2)] + [(c_2, d_2)]$,

showing $+$ is well defined.

A similar proof shows \times is also well defined.

Next, we claim $[(0,1)] \in \mathbb{Q}(R)$ is the additive identity, where $b \in R \setminus \{0\}$ is arbitrary.

Indeed, note for any $[(c,d)] \in \mathbb{Q}(R)$, we have that

$$[(0,1)] + [(c,d)] = [(0 \cdot d + bc, bd)] = [(bc, bd)] = [(c,d)]$$

(the last equality holds since $bcd = bdc$)

A similar check shows that $[(c,d)] + [(0,1)] = [(c,d)]$, so that

$[(0,1)]$ is the additive identity.

A similar proof shows $[(1,1)]$ is the multiplicative identity of $\mathbb{Q}(R)$.

Subsequently, we claim for any $[(a,b)] \in \mathbb{Q}(R)$ with $[(a,b)] \neq [(0,0)]$ for any $c \in R \setminus \{0\}$, $[(cb,a)] = [(a,b)]^{-1}$.

First, note $a \neq 0$, since $[(a,b)] \neq [(0,0)]$ implies $ab \neq 0$, and so $a \neq 0$ as $c \neq 0$.

Thus $[(cb,a)] \in \mathbb{Q}(R)$, and

$$[(cb,a)] \cdot [(a,b)] = [(cba, ab)] = [(1,1)],$$

with a similar proof showing that $[(a,b)] \cdot [(cb,a)] = [(1,1)]$, so that $[(a,b)] = [(cb,a)]^{-1}$.

We can use similar lines of reasoning to show

$\mathbb{Q}(R)$ obeys the other conditions of a

field. \blacksquare

$$R \cong \left\{ \frac{c}{d} : c \in R \right\}$$

Q: Let R be an integral domain. Then R is isomorphic to the subring $\left\{ \frac{c}{d} : c \in R \right\} \subseteq \mathbb{Q}(R)$.

Proof. We first verify $R_0 = \left\{ \frac{c}{1} : c \in R \right\}$ is a subring of $\mathbb{Q}(R)$.

Clearly since $\frac{1}{1} \in R_0$, so that R_0 has the unity of $\mathbb{Q}(R)$.

Then, given $\frac{a}{1}, \frac{b}{1} \in R_0$, notice that

$$\frac{a}{1} - \frac{b}{1} = \frac{a-1-b-1}{1 \cdot 1} = \frac{a-b}{1} \in R_0,$$

and

$$\frac{a}{1} \cdot \frac{b}{1} = \frac{ab}{1 \cdot 1} = \frac{ab}{1} \in R_0,$$

so by the Subring Test R_0 is a subring of R .

Subsequently, let the function $\sigma: R \rightarrow R_0$ by $\sigma(r) = \frac{r}{1} \quad \forall r \in R$.

We claim σ is a ring homomorphism.

Indeed, for any $r, s \in R$, we have that

$$\textcircled{1} \quad \sigma(r+s) = \frac{r+s}{1} = \frac{r}{1} + \frac{s}{1} = \sigma(r) + \sigma(s);$$

$$\textcircled{2} \quad \sigma(rs) = \frac{rs}{1} = \frac{r}{1} \cdot \frac{s}{1} = \frac{r}{1} \cdot \sigma(s) \text{; and}$$

$$\textcircled{3} \quad \sigma(1) = \frac{1}{1},$$

so that σ is a ring homomorphism.

Moreover, σ is clearly surjective, and the fact that $\ker \sigma = \{0\}$ (since $\sigma(r) = \frac{0}{1}$ implies $r = 0$) tells us σ is also injective.

Thus σ is bijective, and hence must also be an isomorphism, proving the claim. \blacksquare

EVERY ELEMENT OF R HAS AN INVERSE IN $\mathbb{Q}(R)$

Q: Let R be an integral domain. Then for every element $a \in R$, there exists an $a^{-1} \in \mathbb{Q}(R)$ such that $a \cdot a^{-1} = 1$.

Proof. Note since $\mathbb{Q}(R)$ is a field, every non-zero element of R_0 has an inverse in $\mathbb{Q}(R)$.

But since $R_0 \cong R$, the claim follows from here.

EVERY ELEMENT OF $\mathbb{Q}(R)$ CAN BE WRITTEN AS ab^{-1} ; $a \in R$, $b \in R$

Q: Let R be an integral domain. Then for any element $q \in \mathbb{Q}(R)$, we can show there must exist some $a \in R$, $b \in R$ such that $q = ab^{-1}$.

Proof. This follows from the fact that

$$\frac{a}{b} = \frac{a}{1} \cdot \frac{1}{b} = \left(\frac{a}{1}\right) \cdot \left(\frac{1}{b}\right)^{-1} = ab^{-1}$$

(and that $\frac{a}{b} \in \mathbb{Q}(R)$). \blacksquare

LOCALISATION

MULTIPLICATIVE SET

Let R be an integral domain.

Then, a subset $S \subseteq R$, where $S \neq \emptyset$, is called a "multiplicative set" if $1 \in S$ and S is closed under multiplication; ie if $a \in S$ and $b \in S$, then $ab \in S$.

LOCALISATION

Let R be an integral domain, and

S be a multiplicative set in R .

Then the "localisation of R at S ", denoted as $S^{-1}R$, is defined to be the set of equivalence classes of ordered pairs in $R \times S$, with the equivalence relation being that $(a,b) \sim (c,d)$ if $ad = bc$ in R .

The addition and multiplication operations are the same in $S^{-1}R$ as they are in $(Q(R))$.

Also, R is isomorphic to a similar subring of $S^{-1}R$ (ie with denominator 1) and every element of S has a multiplicative inverse in $S^{-1}R$.

COMPLEX NUMBERS

The set of complex numbers, or \mathbb{C} , is given to be the set of ordered pairs $(a,b) \in \mathbb{R} \times \mathbb{R}$, where we represent each ordered pair (a,b) by $a+bi$, where $i^2 = -1$.

Moreover, we define an addition on \mathbb{C} by

$$(a+bi) + (c+di) = (a+c) + (b+d)i$$

and a multiplication on \mathbb{C} by

$$(a+bi)(c+di) = (ac-bd) + (ad+bc)i.$$

\mathbb{C} IS A FIELD

With respect to the two operations defined above, we can show \mathbb{C} is a field.

Proof: Since $+$ agrees with the ordinary additive structure on $\mathbb{R} \times \mathbb{R}$, we can deduce \mathbb{C} is an abelian group under addition.

Then, note that

$$\begin{aligned} (a+bi)(c+di)(e+fi) &= ((ac-bd) + (ad+bc)i)(e+fi) \\ &= ((ac-bd)e - (ad+bc)f) + ((ad+bc)e + (ac-bd)f) \\ &= (ace - bde - ade - bcf) + (acf - bdf + ade + bce)i \end{aligned}$$

and

$$\begin{aligned} (a+bi)((c+di)(e+fi)) &= (a+bi)((ce-df) + (de+cf)i) \\ &= (a(ce-df) - b(de+cf)) + (b(ce-df) + a(de+cf))i \\ &= (ace - adf - bde - bcf) + (bce - baf + ade + acf)i \\ &= ((ac-bd) + (ad+bc)i)(e+fi), \end{aligned}$$

showing multiplication in \mathbb{C} is associative.

Moreover, $1 \in \mathbb{C}$ is a multiplicative identity, and note that

$$\begin{aligned} (a+bi)(c+di) &= (ac-bd) + (bc+ad)i \\ &= (ca-db) + (cb+da)i \\ &= (c+di)(a+bi), \end{aligned}$$

showing multiplication is commutative in \mathbb{C} .

Next, observe that for any $a+bi \in \mathbb{C}$, we have

$$\begin{aligned} (a+bi)\left(\frac{a}{a^2+b^2} - \frac{b}{a^2+b^2}i\right) &= \left(\frac{a^2}{a^2+b^2} - \left(\frac{-1^2}{a^2+b^2}\right)\right) + \left(\frac{ab}{a^2+b^2} - \frac{ab}{a^2+b^2}\right)i \\ &= 1 + 0i \\ &= 1, \end{aligned}$$

so that every $a+bi \in \mathbb{C}$ has a multiplicative inverse

$$\frac{a}{a^2+b^2} - \frac{b}{a^2+b^2}i \text{ in } \mathbb{C}.$$

We could also show the distributive laws hold in \mathbb{C} ; this would be sufficient to show \mathbb{C} is a field. \blacksquare

TERMINOLOGY

STANDARD FORM

Let $z \in \mathbb{C}$. Then, the form $z = a+bi$, where $a, b \in \mathbb{R}$, is called the "standard form" for z .

REAL PART

Let $z \in \mathbb{C}$, and suppose $z = a+bi$, where $a, b \in \mathbb{R}$. Then we say $a \in \mathbb{R}$ is the "real part" of z , denoted by $\operatorname{Re}(z)$.

IMAGINARY PART

Let $z \in \mathbb{C}$, and suppose $z = a+bi$, where $a, b \in \mathbb{R}$. Then, we say $b \in \mathbb{R}$ is the "imaginary part" of z , denoted by $\operatorname{Im}(z)$.

COMPLEX CONJUGATE

Let $z \in \mathbb{C}$, and suppose $z = a+bi$, where $a, b \in \mathbb{R}$. Then the "complex conjugate" of z , denoted by \bar{z} , is defined to be the complex number

$$\bar{z} = a - bi.$$

MODULUS

Let $z \in \mathbb{C}$, and suppose $z = a+bi$, where $a, b \in \mathbb{R}$.

Then, the "modulus" of z , denoted by $|z|$, is defined to be

$$|z| = \sqrt{a^2 + b^2}.$$

PROPERTIES OF COMPLEX NUMBERS

$$\overline{z+w} = \overline{z} + \overline{w}$$

💡 Let $z, w \in \mathbb{C}$ be arbitrary.

Then $\overline{z+w} = \overline{z} + \overline{w}$.

Proof. Write $z = a+bi$ and $w = c+di$, where $a, b, c, d \in \mathbb{R}$.

Then $z+w = (a+c) + (b+d)i$.

Thus $\overline{z+w} = (a+c) - (b+d)i$

$$= a+c - bi - di$$

$$= (a-bi) + (c-di)$$

$$= \overline{z} + \overline{w}, \text{ as needed. } \blacksquare$$

$$\overline{zw} = \overline{z} \cdot \overline{w}$$

💡 Let $z, w \in \mathbb{C}$ be arbitrary.

Then $\overline{zw} = \overline{z} \cdot \overline{w}$.

Proof. Again, write $z = a+bi$ and $w = c+di$, where $a, b, c, d \in \mathbb{R}$.

Then $zw = (ac-bd) + (ad+bc)i$,

thus,

$$\overline{zw} = (ac-bd) - (ad+bc)i.$$

$$\text{But } \overline{z} \cdot \overline{w} = (a-bi)(c-di)$$

$$= (ac - (-b)(-d)) + (-bc-ad)i$$

$$= (ac+bd) - (ad+bc)i$$

$$= \overline{zw},$$

completing the proof. \blacksquare

$$\overline{z^{-1}} = \overline{\overline{z}}^{-1}$$

💡 Let $z \in \mathbb{C}$. Then $(\overline{\frac{1}{z}}) = \frac{1}{\overline{z}}$.

Proof. Write $z = a+bi$, where $a, b \in \mathbb{R}$.

Then $\frac{1}{z} = \frac{a}{a^2+b^2} - \frac{b}{a^2+b^2}i$, implying that

$$(\overline{\frac{1}{z}}) = \frac{a}{a^2+b^2} + \frac{b}{a^2+b^2}i.$$

But notice that

$$\begin{aligned} \frac{1}{\overline{z}} &= \frac{1}{a-bi} = \frac{1}{a-bi} \cdot \frac{a+bi}{a+bi} \\ &= \frac{a+bi}{a^2+b^2} \\ &= \frac{a}{a^2+b^2} + \frac{b}{a^2+b^2}i \\ &= (\overline{\frac{1}{z}}), \end{aligned}$$

completing the proof. \blacksquare

$$\overline{\overline{z}} = z$$

💡 Let $z \in \mathbb{C}$. Then $(\overline{\overline{z}}) = z$.

(Trivial proof.)

$$z + \overline{z} = 2\operatorname{Re}(z)$$

💡 Let $z \in \mathbb{C}$.

Then $z + \overline{z} = 2\operatorname{Re}(z)$.

(Trivial proof.)

$$z - \overline{z} = 2i \operatorname{Im}(z)$$

💡 Let $z \in \mathbb{C}$.

Then $z - \overline{z} = 2i \operatorname{Im}(z)$.

(Trivial proof.)

$$z \cdot \overline{z} = |z|^2$$

💡 Let $z \in \mathbb{C}$.

Then $z \cdot \overline{z} = |z|^2$.

Proof. Let $z = a+bi$, where $a, b \in \mathbb{R}$.

Then $z \cdot \overline{z} = (a+bi)(a-bi)$

$$= a^2 + b^2$$

$$= (\sqrt{a^2+b^2})^2$$

$$= |z|^2, \text{ as needed. } \blacksquare$$

$$\operatorname{Re}(z) \leq |\operatorname{Re}(z)| \leq |z|$$

💡 Let $z \in \mathbb{C}$.

Then $|\operatorname{Re}(z)| \leq |\operatorname{Re}(z)| \leq |z|$.

Proof. Let $z = a+bi$, where $a, b \in \mathbb{R}$. Then clearly $|a| \leq |z|$, establishing the first inequality.

Next, note $|\operatorname{Re}(z)|^2 = a^2$

$$\leq a^2 + b^2 \leq b^2 \leq 0$$

$$= |z|^2,$$

so that $|\operatorname{Re}(z)| \leq |z|$, establishing the second inequality. \blacksquare

$$\operatorname{Im}(z) \leq |\operatorname{Im}(z)| \leq |z|$$

💡 Let $z \in \mathbb{C}$.

Then $|\operatorname{Im}(z)| \leq |\operatorname{Im}(z)| \leq |z|$.

(Similar proof to the "Re" section.)

$$|zw| = |z||w|$$

💡 Let $z, w \in \mathbb{C}$.

Then $|zw| = |z||w|$.

Proof. Write $z = a+bi$ and $w = c+di$, for some $a, b, c, d \in \mathbb{R}$.

Then $zw = (a+bi)(c+di)$

$$\therefore zw = (ac-bd) + (ad+bc)i,$$

so that

$$\begin{aligned} |zw|^2 &= (ac-bd)^2 + (ad+bc)^2 \\ &= a^2c^2 - 2acbd + b^2d^2 + b^2c^2 + 2adbc + a^2d^2 \\ &= a^2c^2 + b^2d^2 + b^2c^2 + a^2d^2 \\ &= (a^2 + b^2)(c^2 + d^2) \\ &= |z|^2 |w|^2, \end{aligned}$$

implying that $|zw| = |z||w|$, as required. \blacksquare

$$|\overline{z}| = |z|$$

💡 Let $z \in \mathbb{C}$.

Then $|\overline{z}| = |z|$.

(Proof is trivial.)

TRIANGLE INEQUALITY FOR \mathbb{C}

💡 Let $z, w \in \mathbb{C}$ be arbitrary.

Then necessarily $|z+w| \leq |z| + |w|$.

Proof. First, note that

$$\begin{aligned} |z+w|^2 &= (z+w)(\overline{z+w}) \\ &= (z+w)(\overline{z} + \overline{w}) \\ &= z\overline{z} + w\overline{z} + \overline{z}\overline{w} + w\overline{w} \\ &= |z|^2 + |w|^2 + (\overline{z}\overline{w} + \overline{(z+w)}) \\ &= |z|^2 + |w|^2 + 2\operatorname{Re}(\overline{z}\overline{w}). \end{aligned}$$

Then, since $\operatorname{Re}(\overline{z}\overline{w}) \leq |\overline{z}\overline{w}| = |z||w| = |z||w| = |z||w|$, it follows that

$$|z+w|^2 \leq |z|^2 + |w|^2 + 2|z||w|$$

$$\therefore |z+w|^2 \leq (|z| + |w|)^2,$$

and taking the square root of both sides yields

$$|z+w| \leq |z| + |w|,$$

which we wanted to show. \blacksquare

POLAR FORM

$$e^{i\theta}$$

Let the angle $\theta \in \mathbb{R}$ (in radians) be arbitrary.
Then, we define the complex exponential function at θ , denoted by $e^{i\theta}$, by

$$e^{i\theta} = \cos \theta + i \sin \theta.$$

Note that regardless of θ ,

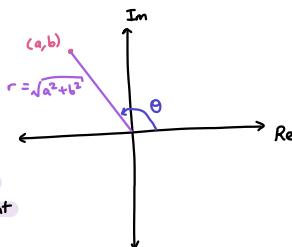
$$|e^{i\theta}| = \sqrt{\cos^2 \theta + \sin^2 \theta} = \sqrt{1} = 1.$$

POLAR FORM

Let $z \in \mathbb{C}$ be arbitrary.
Then the "polar" form of z is given by

$$z = re^{i\theta},$$

where $r = |z|$ and θ is the angle from the positive x -axis to the line segment from (a, b) to the origin.



ARGUMENT

Let $z \in \mathbb{C}$ be arbitrary, and write $z = re^{i\theta}$, with $r, \theta \in \mathbb{R}$. Then θ is called the "argument" of z .

MULTIPLICATION IN POLAR FORM

Suppose $z_1, z_2 \in \mathbb{C}$ are given in polar form, with $z_1 = r_1 e^{i\theta_1}$ and $z_2 = r_2 e^{i\theta_2}$, for some $r_1, r_2 \in \mathbb{R}^+$ and $\theta_1, \theta_2 \in \mathbb{R}$. Then $z_1 z_2 = (r_1 r_2) e^{i(\theta_1 + \theta_2)}$.

Proof. By definition, $z_1 = r_1 e^{i\theta_1} = r_1 (\cos \theta_1 + i \sin \theta_1)$

$$\text{and } z_2 = r_2 e^{i\theta_2} = r_2 (\cos \theta_2 + i \sin \theta_2).$$

Thus

$$\begin{aligned} z_1 z_2 &= (r_1 (\cos \theta_1 + i \sin \theta_1)) (r_2 (\cos \theta_2 + i \sin \theta_2)) \\ &= r_1 r_2 (\cos \theta_1 + i \sin \theta_1)(\cos \theta_2 + i \sin \theta_2) \\ &= r_1 r_2 [(\cos \theta_1 \cos \theta_2 - \sin \theta_1 \sin \theta_2) + i(\sin \theta_1 \cos \theta_2 + \cos \theta_1 \sin \theta_2)] \\ &= r_1 r_2 [(\cos(\theta_1 + \theta_2)) + i(\sin(\theta_1 + \theta_2))] \\ \therefore z_1 z_2 &= r_1 r_2 e^{i(\theta_1 + \theta_2)}, \text{ which we wanted to show.} \quad \square \end{aligned}$$

DE MOIREE'S THEOREM

Let $z \in \mathbb{C} \setminus \{0\}$ be arbitrary, and suppose $z = re^{i\theta}$, where $r \in \mathbb{R}^+$ and $\theta \in \mathbb{R}$.

Then, for any $n \in \mathbb{Z}$, we have that

$$z^n = r^n e^{i(n\theta)}$$

Proof. First, if $n=0$, then $z^0 = 1 = r^0 e^{i(0\theta)} = 1$, proving the base case.

Then, assume the result is true for some $n \in \mathbb{N}$.

This implies $z^n = r^n e^{i(n\theta)}$.

$$\begin{aligned} \text{So, } z^{n+1} &= z^n \cdot z = (r^n e^{i(n\theta)}) (r e^{i\theta}) \\ &= r^{n+1} e^{i(n\theta + \theta)} \\ \therefore z^{n+1} &= r^{n+1} e^{i((n+1)\theta)} \end{aligned}$$

proving the claim holds for $n+1$ as well.

It follows by induction that the claim is true $\forall n \in \mathbb{N}$.

Moreover, observe for any $n \in \mathbb{Z}$,

$$(r^{-n} e^{i(-n\theta)}) (r^n e^{i(n\theta)}) = 1.$$

But since $r^{-n} e^{i(-n\theta)} = (re^{i\theta})^{-n}$, it follows that $(r^{-n} e^{i(-n\theta)}) = (re^{i\theta})^{-n}$, which exists by uniqueness of inverses.

Therefore the claim is true $\forall n \in \mathbb{Z}$, and we are done. \square

THE FUNDAMENTAL THEOREM OF ALGEBRA

The Fundamental Theorem of Algebra states that any non-constant polynomial $f \in \mathbb{C}[x]$ has a root in \mathbb{C} ; ie there always exists a $c \in \mathbb{C}$ such that $f(c)=0$.

ALGEBRAICALLY CLOSED

Let F be a field. Then, we say F is "algebraically closed" if every non-constant polynomial $f \in F[x]$ has a root in F .

Alternatively, F is algebraically closed if and only if every non-constant polynomial $f \in F[x]$ can be factored as a product of linear polynomials

$$f = c(x-a_1)(x-a_2) \dots (x-a_n),$$

where $c, a_1, \dots, a_n \in F$ and $n = \deg(f)$.

Proof. First, suppose F is algebraically closed.

Then, if $n=1$, $f=ax+b$ for some $a, b \in F$ with $a \neq 0$, so that $f=a(x+a^{-1}b)$ gives the factorization of f in the required form.

Now, assume the claim is true for all polynomials of degree n in $F[x]$, where $n \geq 1$.

let $f \in F[x]$ has degree $n+1$.

Then, since F is algebraically closed, f must have a root in F .

let this root be a^{nti} .

So, by the Factor Theorem, $f=g(x-a^{nti})$, where $g \in F[x]$.

In particular, $\deg(g)=n$, so that by the induction hypothesis, we have that

$$g = c(x-a_1)(x-a_2) \dots (x-a_n)$$

so that

$$f = g(x-a^{nti}) = c(x-a_1)(x-a_2) \dots (x-a_n)(x-a^{nti}),$$

proving the claim is true for $n+1$ also.

By induction, it follows that the claim is true $\forall n \in \mathbb{N}$.

Conversely, assume F is a field for which every non-constant polynomial in $F[x]$ factors as a product of linear polynomials.

This implies for any $f \in F[x]$ which is non-constant, we can write

$$f = c(x-a_1)(x-a_2) \dots (x-a_n),$$

where $n = \deg(f) \geq 1$.

Indeed, as a_1, a_2, \dots, a_n are roots of f , it follows that F is algebraically closed, proving the backward argument. \square

SOLVING EQUATIONS IN \mathbb{C}

SOLVING $z^n = a$ IN \mathbb{C}

Let $a \in \mathbb{C}$ and $n \in \mathbb{Z}^+$ be arbitrary.

Write a in polar form, so that $a = r e^{i\theta}$

for some $r \in \mathbb{R}^+$ and $\theta \in \mathbb{R}$.

Then there are exactly n distinct solutions

to the equation $z^n = a$, given by

$$z = \sqrt[n]{r} \cdot e^{i(\frac{\theta + 2k\pi}{n})}$$

where $k \in \{0, 1, \dots, n-1\}$.

Proof: First, assume z is a solution to the eqn (which we knew exists due to the fundamental theorem of algebra),

and write $z = s e^{i\phi}$, where $s > 0$ and $\phi \in \mathbb{R}$.

This implies that

$$z^n = (s e^{i\phi})^n = s^n e^{i n \phi} = a = r e^{i\theta}.$$

Then, since $|s^n e^{i n \phi}| = |r e^{i\theta}|$ it follows that $r = s^n$,

so that $s = \sqrt[n]{r}$.

Additionally, $n\phi = \theta + 2k\pi$ for some $k \in \mathbb{Z}$,

so that $\phi = \frac{\theta}{n} + \frac{2k\pi}{n}$.

But note that only taking $k=0, 1, \dots, n-1$ yields distinct values for the argument of z , so that there are exactly n possible n th roots.

Finally, we can verify these are indeed roots of the equation:

$$(\sqrt[n]{r} \cdot e^{i(\frac{\theta}{n} + \frac{2k\pi}{n})})^n = r e^{i(\theta + 2k\pi)} = r e^{i\theta} = a. \quad \square$$

EXAMPLE: $z^4 = (1+i)$

We can use the method described above to find the solutions in \mathbb{C} of the equation $z^4 = (1+i)$.

First, note $r = \sqrt{1^2 + 1^2} = \sqrt{2}$, and that

$$\theta = \tan^{-1}(\frac{1}{1}) = \frac{\pi}{4}, \text{ so that}$$

$$1+i = \sqrt{2} e^{i\frac{\pi}{4}}.$$

It follows that

$$z = (\sqrt{2})^{\frac{1}{4}} e^{i(\frac{\pi}{4} + \frac{2k\pi}{4})}, \quad k \in \{0, 1, 2, 3\}$$

$$\Rightarrow z = 2^{\frac{1}{4}} e^{i(\frac{\pi}{16} + \frac{k\pi}{2})}, \quad k \in \{0, 1, 2, 3\}. \quad \square$$

QUADRATIC FORMULA FOR A FIELD

Let F be a field with char $F \neq 2$, and suppose

we are given $a, b, c \in F$ with $a \neq 0$.

Suppose further there exists a $y \in F$ such that $y^2 = b^2 - 4ac$.

Then the quadratic equation

$$ax^2 + bx + c = 0$$

has solutions given exactly by

$$x = (-b \pm y)(2a)^{-1}.$$

Proof. First, notice that

$$a((-b \pm y)(2a)^{-1})^2 + b(-b \pm y)(2a)^{-1} + c = (4a)^{-1}(a^2 \mp 2by + y^2) + (2a)^{-1}(-b \pm by) + c \\ = (4a)^{-1}(b^2 \mp 2by + b^2 - 4ac - 2b^2 \pm 2by + 4ac) \\ = 0,$$

so that both vals of x given above are solutions to the quadratic equation. *

Conversely, suppose that $x \in F$ satisfies $ax^2 + bx + c = 0$.

Then, $a(x^2 + a^{-1}bx) + c = 0$

$$\Rightarrow a(x^2 + a^{-1}bx + (2a)^{-2}b^2) - (4a)^{-1}b^2 + c = 0$$

$$a(x + (2a)^{-1}b)^2 + (4a)^{-1}(4ac - b^2) = 0$$

$$\Rightarrow a(x + (2a)^{-1}b)^2 = (b^2 - 4ac)(4a)^{-1}$$

$$\therefore (x + (2a)^{-1}b)^2 = (b^2 - 4ac)(2a)^{-2}.$$

From here, note the two square roots of the RHS are

$y \cdot (2a)^{-1}$ and $-y \cdot (2a)^{-1}$, since $y^2 = b^2 - 4ac$.

Thus $x + (2a)^{-1}b = \pm y \cdot (2a)^{-1}$

and so $x = -(2a)^{-1}b \pm y \cdot (2a)^{-1}$,

$$\therefore x = (-b \pm y)(2a)^{-1},$$

proving the theorem. \square

IRREDUCIBLE POLYNOMIALS

Let F be a field, and $f \in F[x]$ a non-constant polynomial.

Then we say that f is "reducible" if f admits a proper factorisation $f = gh$, where $g, h \in F[x]$ and $\deg(g), \deg(h) > 1$.

Otherwise, we say f is "irreducible".

In other words, f is irreducible if whenever we have a factorisation $f = gh$ with $g, h \in F[x]$, necessarily either g or h must be constant.

$\deg f \geq 2$, f IS IRREDUCIBLE

$\Rightarrow f$ HAS NO ROOTS IN F

Let F be an arbitrary field, and $f \in F[x]$ be a polynomial with $\deg(f) \geq 2$.

Suppose f is irreducible.

Then f has no roots in F .

Proof. Suppose f has a root $c \in F$.

Then, by the Factor Theorem, this implies we can write $f = (x-c)h$ for some $h \in F[x]$.

But since $\deg(x-c) = 1$ and $\deg(h) = \deg(f)-1 \geq 1$, this is a contradiction to our assumption that f is irreducible.

Thus f cannot have any roots in F , proving the claim. \square

$\deg(f) = 2$ or 3 , f HAS NO ROOTS IN F

$\Rightarrow f$ IS IRREDUCIBLE

Let F be an arbitrary field, and $f \in F[x]$ be a polynomial such that $\deg(f)=2$ or $\deg(f)=3$.

Suppose f has no roots in F .

Then f is irreducible.

Proof. Suppose f is reducible.

Then we can write $f = gh$, where $g, h \in F[x]$ are non-constant polynomials.

It follows that $\deg(f) = \deg(g) + \deg(h)$, and so (since $\deg(f) = 2$ or $\deg(f) = 3$) it forces $\deg(g) = 1$ or $\deg(h) = 1$.

Either way, this means f has a linear factor, and this linear factor necessarily has a root in F .

In turn, this implies f has a root in F , which is a contradiction.

Thus f is irreducible after all, which we wanted to prove. \square

f IS IRREDUCIBLE $\Leftrightarrow \deg(f) = 1$

(f IS A NON-CONSTANT POLYNOMIAL IN AN ALGEBRAICALLY CLOSED FIELD)

Let F be an algebraically closed field. Then a non-constant polynomial $f \in F[x]$ is irreducible if and only if $\deg(f) = 1$.

Proof. Clearly, a linear polynomial over any field is irreducible.

Conversely, if $f \in F[x]$ is irreducible and $\deg(f) > 1$, then f has no roots in F .

But this is impossible since F is algebraically closed, so that there are no irreducible polynomials of degree larger than 1 in $F[x]$. \square

IRREDUCIBLE POLYNOMIALS IN $\mathbb{R}[x]$

$$f(c) = 0 \Rightarrow f(\bar{c}) = 0$$

Suppose $f \in \mathbb{R}[x]$ is such that $f(c) = 0$, where $c \in \mathbb{C}$.

Then necessarily $f(\bar{c}) = 0$ also.

Proof. First, write

$$f = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$$

where $a_0, a_1, \dots, a_n \in \mathbb{R}$ and $a_n \neq 0$.

Then, since $f(c) = 0$, it follows that

$$0 = a_n c^n + a_{n-1} c^{n-1} + \dots + a_1 c + a_0$$

Taking conjugates of both sides yields that

$$\bar{0} = \bar{0} = \overline{a_n c^n + a_{n-1} c^{n-1} + \dots + a_1 c + a_0}$$

$$= \overline{a_n} \bar{c}^n + \overline{a_{n-1}} \bar{c}^{n-1} + \dots + \overline{a_1} \bar{c} + \bar{a}_0$$

$$= a_n \bar{c}^n + a_{n-1} \bar{c}^{n-1} + \dots + a_1 \bar{c} + a_0,$$

so that $f(\bar{c}) = 0$, completing the proof. \square

f IS IRREDUCIBLE IN $\mathbb{R}[x] \Leftrightarrow \deg(f) = 1$

OR $\deg(f) = 2$ AND f HAS NO REAL ROOTS

Let $f \in \mathbb{R}[x]$ be a non-constant polynomial.

Then f is irreducible in $\mathbb{R}[x]$ if and only if $\deg(f) = 1$, or $\deg(f) = 2$ and f has no real roots.

Proof. First, if $\deg(f) = 1$ or $\deg(f) = 2$ and f has no real roots, then the words in the previous sections tell us that f is irreducible.

Conversely, suppose $f \in \mathbb{R}[x]$ is an irreducible, non-constant polynomial.

Since $R \subseteq \mathbb{C}$, it follows that $f \in \mathbb{C}[x]$ also.

Then, by the Fundamental Theorem of Algebra, f has a complex root $c \in \mathbb{C}$.

If $c \in \mathbb{R}$, then f has a real root, and the irreducibility of f forces $\deg(f) = 1$.

Otherwise, $c \notin \mathbb{R}$, so that $\bar{c} \neq c$ and so c is also a root of f .

Then, applying the Factor Theorem, we get that

$$f = (x-c)(x-\bar{c})h \text{ for some other polynomial } h \in \mathbb{C}[x].$$

Note that

$$g = (x-c)(x-\bar{c}) = x^2 - (c+\bar{c})x + c\bar{c} = x^2 - (2\operatorname{Re} c)x + |c|^2 \in \mathbb{R}[x].$$

Thus, if we carry out division with remainder of f by g in $\mathbb{R}[x]$, we get that $f = gh+r$,

where $h, r \in \mathbb{R}[x]$ and $\deg(r) < 2$.

But if we do the same division with remainder over $\mathbb{C}[x]$, we know we get that $f = gh+r$.

Since the remainders are unique, it follows that $r=0$, so that $h=h'$.

Consequently, $f = gh$ in $\mathbb{R}[x]$, where $\deg(g) = 2$.

By the irreducibility of f , h must be a constant polynomial.

This implies $f = kg$, where k is constant, so that

f is also of degree 2 with no real roots. \square

IRREDUCIBLE POLYNOMIALS IN $\mathbb{Q}[x]$

RATIONAL ROOTS THEOREM

Let $f \in \mathbb{Q}[x]$ be a non-constant polynomial, and suppose $r \in \mathbb{Q}$ is a root of f . Suppose further that $f = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$, where $a_0, a_1, \dots, a_n \in \mathbb{Z}$ and $a_n \neq 0$.

Then, if $r = \frac{p}{q}$, where $p, q \in \mathbb{Z}$ and $\gcd(p, q) = 1$, we must have that $q | a_n$ and $p | a_0$ in \mathbb{Z} .

Proof. Since $f(\frac{p}{q}) = 0$, it follows that

$$a_n \left(\frac{p}{q}\right)^n + a_{n-1} \left(\frac{p}{q}\right)^{n-1} + \dots + a_1 \left(\frac{p}{q}\right) + a_0 = 0.$$

Multiplying both sides by q^n and re-arranging gives us that

$$\begin{aligned} a_0 q^n &= - (a_n p^n + a_{n-1} p^{n-1} q + \dots + a_1 p q^{n-1}) \\ &= -p (a_n p^{n-1} + a_{n-1} p^{n-2} q + \dots + a_1 q^{n-1}), \end{aligned}$$

showing that $p | a_0 q^n$.

Then since $\gcd(p, q^n) = 1$, it follows that $p | a_0$.

Similarly, if we isolate for $a_n p^n$ instead of $a_0 q^n$, we deduce that $q | a_n p^n$, and since $\gcd(q, p^n) = 1$, we get that $q | a_n$, completing the proof. \square

SHOWING NUMBERS ARE IRRATIONAL

We can use the Rational Roots Theorem to evaluate whether a given $r \in \mathbb{R}$ is irrational.

Example: we can show that $\sqrt{2} + \sqrt{3} \in \mathbb{R} \setminus \mathbb{Q}$.

First, let $\alpha = \sqrt{2} + \sqrt{3}$.

Then note that

$$\begin{aligned} \alpha^2 &= 2 + 2\sqrt{6} + 3 \\ &= 5 + 2\sqrt{6} \end{aligned}$$

so that

$$\alpha^2 - 5 = 2\sqrt{6}$$

and hence

$$\begin{aligned} (\alpha^2 - 5)^2 &= (2\sqrt{6})^2 \\ &= 24 \end{aligned}$$

$$\Rightarrow \alpha^4 - 10\alpha^2 + 25 = 24$$

and so α is a solution to the polynomial

$$f = x^4 - 10x^2 + 1.$$

Then, by the Rational Roots Theorem, the only candidate rational roots of f are 1 and -1 ; but since $f(1) = f(-1) = -8 \neq 0$, this implies that α , being a root of f , cannot be rational!