Lecture 01 Logic and Mathematical Proofs

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Mathematics Review Course, Summer 2023 University of Minnesota August 7th, 2023

Course Preview

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THIS COURSE

- ▶ Review of graduate-level mathematics necessary for the 1st year sequence.
- ▶ Aimed at PhD-track. MS-track is encouraged.
- ► This sets the foundation (Not exhaustive).
- ▶ By the end you should feel confident tackling a variety of math situations in a short period.
- ▶ Syllabus on Github repo. Repo is the most up-to-date place for course content.
- ► This course is **optional**.

PREVIEW OF COURSE

- 1. Logic, Proofs, Sets, & Topology
- 2. Uni-variate Calculus & Multi-variate Calculus
- 3. Linear Algebra
- 4. Functions & Optimization
- 5. Probability & Statistics
- 6. Dynamic Programming

ABOUT THE INSTRUCTOR



Ryan McWay

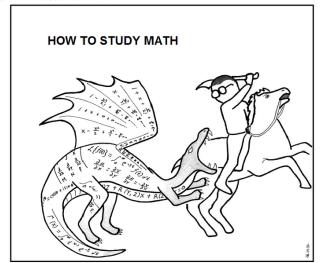
- ► Current: 2nd Year APEC PhD student
- ▶ Background: $SLU \rightarrow USF \rightarrow UMich \rightarrow UMN$
- ► Research: Development, Behavior, Urban, Environment

DAILY ICEBREAKER

- ► Attendance via prompt:
 - ► Name
 - ► Hometown
 - ► Program and track
 - ► Research interests
 - ▶ Daily Icebreaker: Imagine you are a professional baseball player or wrestler. What is you walk up (intro) song?



FIGHT WITH MATH...



Don't just read it; fight it!

Logic

Topic: Logic

- ► General background
 - ▶ Logic is at the heart of reasoning and arguments.
 - ► Expressed in words and formalized through math, this is a foundation of theoretical arguments.
 - ▶ Deduce information correctly. Not deducing correct information.
- ▶ Why do economists' care?
 - ► Foundation for theory
 - ► Criteria to evaluate arguments
- ► Application in this career
 - ► Creating logical arguments
 - ► How you think about research
 - Evaluating theory and conclusions from empirical evidence

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- 1. Logical Statements
- 2. Necessary Conditions
- 3. Sufficient Conditions

0. TERMINOLOGY

\forall	For al

- √ For an
- : Such that
- Given
- · Therefore
- · · Because
- ∧ And
- ∨ Or
- √ Negation
- Equivalent or identical.
- ⇒ Implies, then, or sufficient
- ⇔ If and only if, or necessary and sufficient
- $\subset \quad \text{Strict subset} \quad$
- ⊆ Subset
- ∈ In, or an element of the following set
- \square End of Proof. QED (quod erat demonstrandum \sim it has been demonstrated).

0. TERMINOLOGY

Aa ALPHA [a]

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Bβ BETA [b]

Γγ GAMMA [g] γάμμα $\Delta\delta$ DELTA [d]

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Math Review 2023: Logic & Proofs

Aug. 7th, 2023

1. LOGICAL STATEMENTS

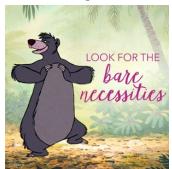
- ▶ Logical Statement: Use a set of facts to infer/assume a new fact.
 - ► Hypothesis (If): Premise with set of facts
 - ► Conclusion (Then): New set of facts inferred if hypothesis is true.
 - e.g., **If** I study throughout the course, **then** I earn a higher grade.
- ► Family of statements:
 - ightharpoonup Tautologies: Statement is always true (1 = 1)
 - ightharpoonup Contradictions: Statement is always false (2=3)
 - ightharpoonup Statement: $A \implies B$
 - ightharpoonup Contrapositive: $\neg B \implies \neg A$
 - ightharpoonup Converse: $B \implies A$
 - ▶ Inverse: $\neg A \implies \neg B$

1. LOGICAL STATEMENTS

- Axiom: Statements assumed to be true.
 - ightharpoonup e.g., a = b, $b = c \implies a = c$
- ► Theorem: A statement proven to be true.
- ▶ Corollary: A theorem that follows from another theorem.
- ▶ Lemma: A minor theorem used to prove another theorem.

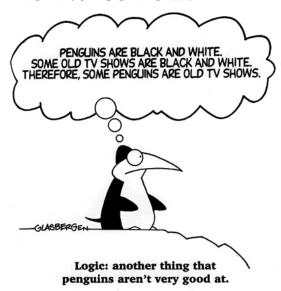
2. Necessary Condition

- \triangleright A is necessary for B
 - ▶ If B is true, A must be true: $B \implies A$.
 - ▶ If A is not true, B is not true: $\neg A \implies \neg B$
- \triangleright A is needed to make the argument.



3. SUFFICIENT CONDITION

- \triangleright A is sufficient for B
 - ▶ If A is true, B must be true: $A \implies B$
 - ▶ If B is not true, A is not either: $\neg B \implies \neg A$
- ightharpoonup A allows you to state B, but not necessary to make argument.



4. Necessary and Sufficient (If and Only If \sim 1ff)

- ightharpoonup If A is sufficient for B, B is necessary for A.
- ▶ If $A \implies B$ and $B \implies A$, then $A \iff B$ (iff)
 - ightharpoonup A is necessary and sufficient for B.
 - \triangleright A and B are equivalent statements.
 - \triangleright A is true iff B is true: A iff B

DEMONSTRATION: NECESSARY AND SUFFICIENT

Ouestion:

Is this statement true: "If I open the door, I used the key."

DEMONSTRATION: NECESSARY AND SUFFICIENT

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Answer:

Logic: Open Door (A) \Longrightarrow Used Key (B) Necessary: You need a key (B) to open the door (A). $B \Longrightarrow A$. Sufficient: If you do not have the key $(\neg B)$, then there is no way to open the door $(\neg A)$. So $A \Longrightarrow B$.

PRACTICE: NECESSARY AND SUFFICIENT CONDITIONS

1. Scoring more touchdowns than your opponent in American football means you won the game.

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Answer: Show Work

Sufficient but not necessary.

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- 1. Scoring more touchdowns than your opponent in American football means you won the game.
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Answer: \ Show Work

Necessary but not sufficient.

- 1. Scoring more touchdowns than your opponent in American football means you won the game.
- 2. Obtaining a learner's permit will lead to earning a driver's license.
- 3. All even positive whole numbers must be divisible by two.

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Answer: Show Work

Both necessary and sufficient.

Topic: Proofs

- ► General background
 - ▶ Method for proving or disproving a logical statement
- ▶ Why do economists' care?
 - ▶ Determine which theories are incorporated into economic theory
- ► Application in this career
 - ► Theory papers and well-developed theory sections of empirical papers.
 - ▶ Often in appendix sections to prove statements articulated as part of an argument in a paper.

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- 1. Truth Tables
- 2. Existence & Uniqueness
- 3. Direct Proofs
- 4. Proof by Contradiction
- 5. Proof by Induction
- 6. Proof by Contrapositive

ASSUMPTIONS ARE THE CORE OF PROOFS...





Proofs



Smbc-comics.com

- ▶ Shows how the truth/falsity of a compound statement depends on the truth/falsity of the simple statements from which it's constructed.
- ► Statements = {Known to be true, known to be false, truth unknown }
- ▶ Truth table for $(P \rightarrow Q)$:

P	Q	$P \iff Q$
Τ	Т	T
Τ	\mathbf{F}	F
F	Т	F
F	\mathbf{F}	${ m T}$

DEMONSTRATION: TRUTH TABLE

Question:

Construct a truth table for $(P \to Q) \lor (Q \to P)$

P	Q	$P \rightarrow Q$	$Q \rightarrow P$	$(P \to Q) \lor (Q \to P)$
T	Т	T	T	T
T	F	F	T	T
F	Т	Τ	F	T
F	F	Т	Т	Т

Proofs

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Answer:

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T	F	F	${ m T}$	${ m T}$
F	T	T	F	${ m T}$
\mathbf{F}	F	T	${ m T}$	T

Proofs

2. Existence and Uniqueness

Existence: Can an answer to the truth of a statement (viz., proof) be found.

- Sometimes you can prove no answer can exist.
- ▶ Uniqueness: An assertion that there is exactly one statement that is true for that family of statements.
 - \triangleright x is the family P(x).
- ▶ Ideally, you want a statement which exists and is unique.
 - \triangleright Ex. If "x + 2 = 3, then x = 1" is a statement that exists and is unique.

APPLICATION: EQUILIBRIUM EXISTS

Existence of Equilibrium II Theorem:

Suppose that each consumer's preferences are continuous, strongly monotonic, and convex. Suppose also that, for each consumer $i, \omega_i \gg 0$. Then there exists a Walrasian equilibrium (p^*, x^*) for ε .

3. Proof by Deduction (Direct Proof)

- ightharpoonup Show $A \implies B$
- ▶ Deductive reasoning: Use a set of premises that lead to a conclusion.
- \triangleright Sometimes we need to strengthen A but adding assumptions (e.g., weak assumptions are preferred).

DEMONSTRATION: DIRECT PROOF

Question:

Let m be an even integer and p be any integer. Then $m \times p$ is an even integer.

Answer

Proof.

m is an even integer so \exists an integer q such that $m=2\times q$ by the definition of an even integer. Therefore, we can make the statement:

$$m \times p = (2 \times q) \times p = 2 \times (q \times p)$$

So, $m \times p$ is an even integer.

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4. Proof by Contradiction

- $\blacktriangleright A \implies B \equiv \neg A \text{ and } \neg B \equiv \neg B \implies \neg A.$
- ▶ E.g., If the conclusion is not untrue, then the premise must be untrue.

DEMONSTRATION: PROOF BY CONTRADICTION

Ouestion:

Walras' Law: $\forall x \in x(p, w)$ that maximizes consumer utility, then $x \times p = w$.

Answer.

Proof

Suppose $\exists x \in x(p, w) : x \times p < w \ (\neg B)$, then there must be another $y \in x(p, w)$ that is affordable and $y \succ x$ by the assumption of "local non-satiation". Therefore, since y exists and is affordable, then x does not maximize utility $(\neg A)$.

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5. Proof by Induction



Professor Schmidt demonstrates the concept of proof by induction.

- ► Inductive reasoning: Drawing conclusions by reasoning a series of specific examples generalizes.
- ▶ Often used by indexing through integers.

DEMONSTRATION: PROOF BY INDUCTION

Question:

$$P(n): 1+2+3+\cdots+n = \frac{n(n+1)}{2}$$

Answer:

Proof

Note that P(1) is true because $1 = \frac{1 \times 2}{2}$. Assume P(n) is true for $k \in n$ integers: $1 + 2 + \cdots + k = \frac{k(k+1)}{2}$. Add (k+1) to both sides.

$$1 + 2 + \dots + k + (k+1) = \frac{k(k+1)}{2} + (k+1) = \frac{(k+1)(k+2)}{2}$$

This is P(k+1), implying that P(k) is true for all P(n).

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This is P(k+1), implying that P(k) is true for all P(n).

6. PROOF BY CONTRAPOSITIVE

▶ Double Negation:

$$\neg B \implies \neg A \equiv \neg \neg A \implies \neg \neg B \equiv A \implies B.$$

 \triangleright Convenient when there is a universal quantifier (\forall) present included by the contrapositive.

DEMONSTRATION: PROOF BY CONTRAPOSITIVE

Question:

Suppose $x \in \mathbb{Z}$. If 7x + 9 is even, then x is odd.

Answer:

Proof

Suppose x is **not** odd (i.e., even) implying x = 2a for some integer a. Then,

$$7x + 9 = 7(2a) + 9 = 14a + (8+1) = 2(7a+4) + 1 = 2b + 1$$

if b = 7a + 4. Consequently, 2b + 1 is odd for all b. Therefore 07x + 9 is **not** even.

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Math Review 2023: Logic & Proofs

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Answer: Show Work

 x^2 is odd by definition of an odd number.

- 1. If x is odd, then x^2 is odd.
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Answer: Show Work

By contrapositive, $x^2 = 6x + 5$ is odd and therefore not even.

- 1. If x is odd, then x^2 is odd.
- 2. Suppose $x \in \mathbb{Z}$. If $x^2 = 6x + 5$ is even, then x is odd.
- 3. There are infinitely many prime numbers.

PRACTICE: PROOFS

- 1. If x is odd, then x^2 is odd.
- 2. Suppose $x \in \mathbb{Z}$. If $x^2 = 6x + 5$ is even, then x is odd.
- 3. There are infinitely many prime numbers.

Answer: Show Work

Proof by contradiction.

Review

Course Preview

REVIEW OF LOGIC

- 1. Logical Statement: Necessary to provide clarity to your statements
- 2. Necessary vs. Sufficient Conditions: Aiming to make statements that are both (iff)

- 1. Truth Tables
- 2. Four methods to prove a statement:
 - ▶ Direct proof
 - ▶ Proof by contradiction
 - ▶ Proof by induction
 - ► Proof by Contrapositive

- ▶ Readings on Sets & Topology before Lecture 02:
 - ► B&S Appendix A
 - ► Hammack Ch 4. & 10
- ► Assignment:
 - ► Problem Set 01 (PS01)
 - ► Solution set will be available following end of Lecture 02
- ► Struggling?
 - 1. Read the 'Encouraged Reading'
 - 2. Review 'Supplementary material'
 - 3. Reach out directly

N & S CONDITIONS QUESTION 1 ANSWER:



- ▶ Not Necessary: If you won the game (B), you may have scored other points but less touchdowns $(\neg A)$.
- ▶ Sufficient: If you score more touchdowns (A) (and therefore more overall points), then you will win the game (B).

N & S CONDITIONS QUESTION 2 ANSWER:



- ▶ Necessary: A learner's permit (A) is required before you can get a drivers license (B).
- ▶ Not Sufficient: Not all learners $(\neg A)$ successfully earn their drivers license $(\neg B)$.

N & S CONDITIONS QUESTION 3 ANSWER:



- ▶ Necessary: To be divisible by 2 (B), you must be a even whole number (A).
- ▶ Sufficient: If you are an even whole number (A), you will have no remainder if divided by two (B).

PROOFS QUESTION 1 ANSWER:

◆ QUESTION

Proof.

Suppose x is odd. Then x = 2a + 1 for some $a \in \mathbb{Z}$, by definition an odd number. Thus $x^2 = (2a + 1)^2 = 4a^2 + 4a + 1$. This $2(2a^2 + 2a) + 1$. So $x^2 = 2b + 1$ for an integer b. Therefore, x^2 is odd, by definition of an odd number.

PROOFS QUESTION 2 ANSWER:

◆ QUESTION

Proof.

Suppose *x* is **not** odd. Thus *x* is even, so x = 2a for some integer *a*. So $x^2 - 6x + 5 = (2a)^2 - 6(2a) + 5 = 2(2a^2 - 6a + 2) + 1$. Then $x^2 - 6x + 5 = 2b + 1$ for $b = 2a^2 - 6a + 2$. Consequently, $x^2 - 6x + 5$ is odd, and therefore not even.

PROOFS QUESTION 3 ANSWER:



Proof.

Suppose there are only finite prime numbers. Then they can be listed as p_1, p_2, \ldots, p_n . Then p_n is the final and largest prime number. Consider a number $a = (p_1 \cdot p_2 \cdots p_n) + 1$. a has at least one prime divisor (e.g., p_k in the list). So there is some integer c such that $(p_1 \cdot p_2 \cdots p_{k-1}p_kp_{k+1} \cdots p_n) + 1 = c \cdot p_k$. Divide both sides by p_k . Now we have $\frac{1}{p_k} = c - (p_1 \cdot p_2 \cdots p_{k-1}p_{k+1} \cdots p_n)$. The expression on the right is an integer (i.e., which prime is a part of) **but** the left is not an integer. This is a contradiction. Therefore, there must be no finite range of prime numbers. \square