

# CSCI235 Database Systems

## Database Design Quality

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# Database Design Quality

## Outline

Why not ONE BIG TABLE !?

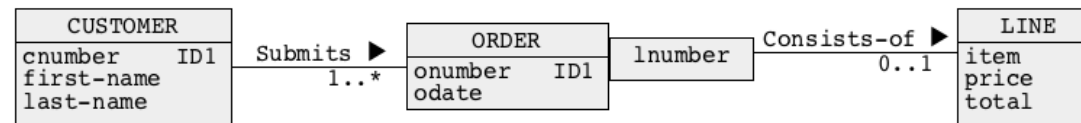
Where is a problem ?

Functional dependency

# Why not ONE BIG TABLE !?

Let us consider the following database domain:

- A **customer** is described by a unique **customer number**, **first**, and **last name**
- **Customers** submit **orders**. An **order** is described by a **unique order number** and **order date**
- **Orders** consist of **lines**. A **line** contains information about a **name of ordered item**, **price per single item**, and **total number of ordered items**



Logical design provides the following relational schemas:

CUSTOMER(cnumber, first-name, last-name)

**PRIMARY KEY** = (cnumber)

ORDERS(onumber, odate, cnumber) **PRIMARY KEY** = (onumber)

**FOREIGN KEY** = (cnumber) **REFERENCES** CUSTOMER(cnumber)

LINE(onumber, lnumber, item, price total)

**PRIMARY KEY** = (onumber, lnumber)

**FOREIGN KEY** = (onumber) **REFERENCES** ORDERS(onumber)

Relational schemas

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# Why not ONE BIG TABLE !?

Why not one relational schema ?

Big relational schema

CUSTOMER(cnumber, first-name, last-name, onumber, odate, cnumber, onumber, lnumber, item, price total)  
**PRIMARY KEY** = (cnumber, onumber, lnumber)

Insertion of information about one customer who submitted 2 orders such that each order consists several lines reveals a problem !

Big relational table

cnumber	fname	lname	onumber	odate	lnumber	item	price	total
7	James	Bond	7	2017-01-01	1	bolt	23.04	5
7	James	Bond	7	2017-01-01	2	screw	29.01	3
7	James	Bond	7	2017-01-01	3	nut	4.55	2
7	James	Bond	8	2018-01-01	1	bolt	23.04	1
7	James	Bond	8	2018-01-01	2	screw	23.04	1
7	James	Bond	8	2018-01-01	3	nut	23.04	2
7	James	Bond	8	2018-01-01	4	lock	23.04	1

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# Why not ONE BIG TABLE !?

Big relational table

cnumber	fname	lname	onumber	odate	lnumber	item	price	total
7	James	Bond	7	2017-01-01	1	bolt	23.04	5
7	James	Bond	7	2017-01-01	2	screw	29.01	3
7	James	Bond	7	2017-01-01	3	nut	4.55	2
7	James	Bond	8	2018-01-01	1	bolt	23.04	1
7	James	Bond	8	2018-01-01	2	screw	23.04	1
7	James	Bond	8	2018-01-01	3	nut	23.04	2
7	James	Bond	8	2018-01-01	4	lock	23.04	1

A **number**, **first name**, and **last name** of a customer is repeated as many times as the total number of different items purchased in all orders and ...

... and **order number** is repeated together with **order date** as many times as the total number of different items purchased in an order

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# Why not ONE BIG TABLE !?

A multitable design does not have such a problem:

CUSTOMER(cnumber, first-name, last-name)  
PRIMARY KEY = (cnumber)

CUSTOMER schema

cnumber	fname	lname
7	James	Bond

CUSTOMER table

ORDERS(onumber, odate, cnumber) PRIMARY KEY = (onumber)  
FOREIGN KEY = (cnumber) REFERENCES CUSTOMER(cnumber)

ORDERS schema

onumber	odate	cnumber
7	2017-01-01	7
8	2018-01-01	7

ORDERS table

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# Why not ONE BIG TABLE !?

A multitable design does not have such a problem:

```
LINE(onumber, lnumber, item, price total)
PRIMARY KEY = (onumber, lnumber)
FOREIGN KEY = (onumber) REFERENCES ORDERS(onumber)
```

LINE schema

onumber	lnumber	item	price	total
7	1	bolt	23.04	5
7	2	screw	29.01	3
7	3	nut	4.55	2
8	1	bolt	23.04	1
8	2	screw	23.04	1
8	3	nut	23.04	2
8	4	lock	23.04	1

LINE table

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## Outline

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# Where is a problem ?

Why do we get redundancies in an incorrectly designed relational table ?

TABLE_NAME			
COLUMN_1	COLUMN_2	...	COLUMN_N
Green	Red	...	Blue
Green	Red	...	Orange
Green	Red	...	Red
Blue	Yellow	...	Yellow
Blue	Yellow	...	Magenta
Orange	Red	...	Yellow
Orange	Red	...	Green

Data dependencies:

- If **COLUMN\_1** is green then **COLUMN\_2** is red
- If **COLUMN\_1** is blue then **COLUMN\_2** is yellow
- If **COLUMN\_1** is orange then **COLUMN\_2** is red

For any colour x if **COLUMN\_1** is x then **COLUMN\_2** is y

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# Where is a problem ?

Data dependencies can be represented as a separate relational table ...

TABLE\_1

COLUMN_1	COLUMN_2

... and **COLUMN\_2** can be removed from the original table

TABLE\_2

COLUMN_1	...	COLUMN_N
	...	
	...	
	...	
	...	
	...	
	...	
	...	

# Where is a problem ?

Do data dependencies exist in BIG TABLE ?

Big relational table									
cnumber	fname	lname	onumber	odate	lnumber	item	price	total	
7	James	Bond	7	2017-01-01	1	bolt	23.04	5	
7	James	Bond	7	2017-01-01	2	screw	29.01	3	
7	James	Bond	7	2017-01-01	3	nut	4.55	2	
7	James	Bond	8	2018-01-01	1	bolt	23.04	1	
7	James	Bond	8	2018-01-01	2	screw	23.04	1	
7	James	Bond	8	2018-01-01	3	nut	23.04	2	
7	James	Bond	8	2018-01-01	4	lock	23.04	1	

Data dependencies:

- If **cnumber** = 7 then **fname** = James
- If **cnumber** = 7 then **lname** = Bond

For any customer number x if **cnumber** = x then **fname** = y and **lname** = z

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# Where is a problem ?

Do data dependencies exist in BIG TABLE ?

Big relational table									
cnumber	fname	lname	onumber	odate	lnumber	item	price	total	
	7	James	Bond	7	2017-01-01	1	bolt	23.04	5
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	7	James	Bond	8	2018-01-01	2	screw	23.04	1
	7	James	Bond	8	2018-01-01	3	nut	23.04	2
	7	James	Bond	8	2018-01-01	4	lock	23.04	1

Data dependencies:

- If **onumber** = 7 then **odate** = 2017-01-01
- If **onumber** = 8 then **odate** = 2018-01-01

For any order number x if **onumber** = x then **odate** = y

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# Functional dependency

What does it mean: if a value in column A is x then a value in column B is always y ?

It means that every value x in a column A is associated with only one value y in a column B

For example, every customer number in a column **cnumber** is associated with only one first name in a column **fname**, i.e. a customer has only one first name

For example, every customer number in a column **cnumber** is associated with only one last name in a column **lname** i.e. a customer has only one last name

For example, every order number in a column **onumber** is associated with only one order date in a column **odate** i.e. an order has only one date

# Functional dependency

Such data dependency does not hold for **item name** and **order number** because an **item name** in a column **item** can be associated with many **order numbers** in a column **onumber** and the opposite ...

... an **order number** in a column **onumber** can be associated with many **item names** in a column **item**

# Functional dependency

If every value in a column **A** is associated with only one value in a column **B** then it means that the columns **A** and **B** represent a function **f** that maps the values in a column **A** into the values in a column **B**

$$f : \text{domain}(A) \rightarrow \text{domain}(B)$$

If every value in a column **cnumber** is associated with only one value in a column **fname** then the columns **cnumber** and **fname** represent a function

$$f : \text{domain}(\text{cnumber}) \rightarrow \text{domain}(\text{fname})$$

If every value in a column **cnumber** is associated with only one value in a column **lname** then the columns **cnumber** and **lname** represent a function

$$f : \text{domain}(\text{cnumber}) \rightarrow \text{domain}(\text{lname})$$



# Functional dependency

If every value in a column **onumber** is associated with only one value in a column **odate** then the columns **onumber** and **odate** represent a function

$$f : \text{domain}(\text{onumber}) \rightarrow \text{domain}(\text{odate})$$

If the columns **A** and **B** in a relational table **R** represent a function

$$f : \text{domain}(A) \rightarrow \text{domain}(B)$$

then in the future it will be denoted by

$$A \rightarrow B$$

and we shall say that a **functional dependency**  $A \rightarrow B$  is valid in a **relational table R** or that **A functionally determines B**

# Functional dependency

Therefore, the following functional dependencies are valid in a big table

**CUSTOMER:**

$\text{cnumber} \rightarrow \text{fname}$

$\text{cnumber} \rightarrow \text{lname}$

$\text{onumber} \rightarrow \text{odate}$

$\text{onumber} \rightarrow \text{cnumber}$

$\text{onumber} \rightarrow \text{fname}, \text{lname}$

$\text{onumber}, \text{lnumber} \rightarrow \text{item}$

$\text{onumber}, \text{lnumber} \rightarrow \text{price}, \text{total}$

... and the others

# Functional dependency

**Functional dependency** is a special kind of so called **data dependency** which is a reflection of the real world **consistency constraint**

**Functional dependencies** can be used to describe the **semantics** (meaning) of data

**Functional dependencies** can be used to determine whether a relational schema (header of relational table) is constructed in a correct way

**Functional dependencies** can be used to design a database, however such approach is used very rarely

# References

T. Connolly, C. Begg, Database Systems, A Practical Approach to Design, Implementation, and Management, Chapter 14.1 The Purpose of Normalization, Chapter 14.2 How Normalization Supports Database Design, 14.3 Data Redundancies and Update Anomalies, Pearson Education Ltd, 2015