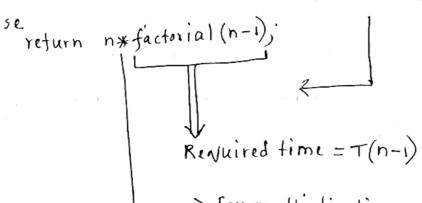
Time complexity of Recursive Algorithms

Total time = T(n)

else



> for multiplication, rewaited Time = b

Recurrence Relation for Time

$$T(n) = \begin{cases} a & n=1 \\ T(n-1)+b & n > 1 \end{cases}$$

solution:

$$\frac{301u\pi 10n :}{T(n) = T(n-1) + b} = T(n-2) + b + b = T(n-2) + 2b : :$$

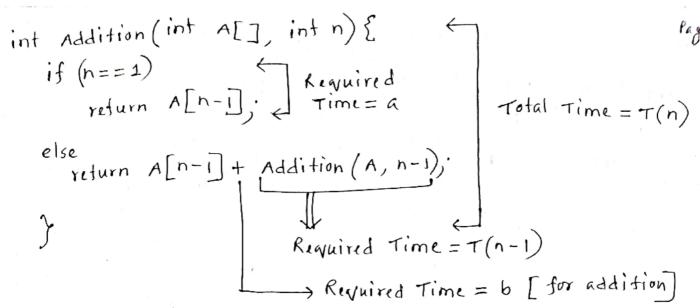
$$= T(n-2) + 2b : :$$

$$= T(n-2) + kb = T(n-1) + kb = T(n-1) + kb = T(n-1) + b(n-1) = bn + a - b = O(n)$$

$$\begin{bmatrix} \cdot \cdot \cdot \cdot \\ \tau(n-1) = \tau(n-2) + b \end{bmatrix}$$

Let
$$n-k=1$$

=) $k=n-1$
[$T(n)=a$ When $h=1$
=) $T(1)=a$]



Recurrence Relation for Time

$$T(n) = \begin{cases} a & n = 1 \\ T(n-1) + b & n > 1 \end{cases}$$

$$\frac{\text{Solution:}}{T(n) = T(n-1) + b}$$

$$=$$

$$=$$

$$\vdots$$

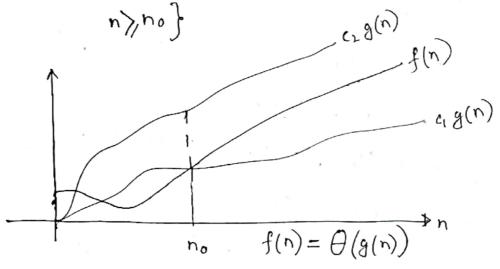
Different or

Different Notations

$$T(n) = \bigcap (n^{\gamma})$$

$$\Rightarrow f(n) = \Theta(g(n))$$
 where $g(n) = n^{\perp}$

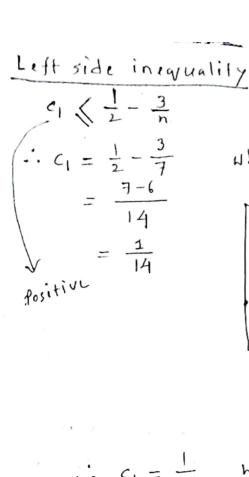
= {f(n): there exists positive constants c1, c2 and no such that cig(n) (f(n) (cig(n) for all



Prove that
$$f(n) = \frac{1}{2}n^{2} - 3n = \Theta(n^{2})$$

$$c_1g(n)$$
 $\langle f(n) \neq c_2g(n) \rangle$

$$\Rightarrow c_1 \left(\frac{1}{2} - \frac{3}{n} \right) \left(c_2 \right) \left[\text{Divide by n} \right]$$



Right side inequality fage-4
$$\frac{1}{2} - \frac{3}{n} \leqslant c_2 \qquad \text{(ositive)}$$

$$\therefore c_2 = \frac{1}{2} \quad \text{when } n > 1$$

$$\begin{array}{cccc}
n = 1 & \frac{1}{2} - 3 & & & & & \\
n = 2 & \frac{1}{2} - \frac{3}{2} & & & \\
n = 3 & \frac{1}{2} - \frac{3}{3} & & & \\
n = 4 & \frac{1}{2} - \frac{3}{4} & & & \\
n = 5 & \frac{1}{2} - \frac{3}{3} & & & & \\
n = 6 & \frac{1}{2} - \frac{3}{3} & & & & \\
n = 7 & \frac{1}{2} - \frac{3}{3} & & & & \\
n = 7 & \frac{1}{2} - \frac{3}{3} & & & & \\
\end{array}$$
Neg

$$\begin{array}{c}
n = 4 & \frac{1}{2} - \frac{3}{3} & & & & \\
n = 6 & \frac{1}{2} - \frac{3}{3} & & & & \\
n = 7 & \frac{1}{2} - \frac{3}{3} & & & & \\
\end{array}$$
Positive

When n>7

$$c_1 = \frac{1}{14} \quad \text{Whin } n > 7$$

$$c_2 = \frac{1}{2} \quad \text{Whin } n > 1$$

$$c_3 = \frac{1}{4} \quad \text{Whin } n > 7$$

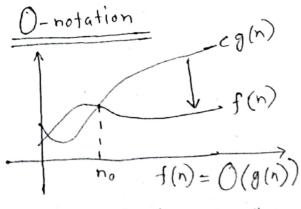
$$combindly \quad c_1 = \frac{1}{14}, \quad c_2 = \frac{1}{2} \quad \text{Whin } n > 7$$

$$\therefore c_1g(n) \leqslant f(n) \leqslant c_2g(n)$$

$$=) \frac{1}{14}g(n) \leqslant f(n) \leqslant \frac{1}{2}g(n) \quad \text{when } n \geqslant 7$$

: He can Write,
$$f(n) = \Theta(\vartheta(n))$$

=
$$\theta(n)$$
 proved



$$\frac{-\Omega - notation}{\int_{n_0}^{\infty} f(n)} f(n)$$

$$= \int_{n_0}^{\infty} f(n) = \int_{n_0}^{\infty} (g(n))$$

f(n) < cg(n) When n > no cg(n) < f(n) When n > no(Upper bound)

$$\frac{1}{f(n) = O(g(n))} \text{ when } \lim_{n \to \infty} \frac{f(n)}{g(n)} = 0$$

Here,
$$f(n) = 6n^{2}$$

 $g(n) = n^{3}$

$$\lim_{n \to \infty} \frac{f(n)}{g(n)} = \lim_{n \to \infty} \frac{6n^2}{n^3}$$

$$=\lim_{n\to\infty}\frac{6}{n}$$

$$f(n) = 6n^{2} = 0(n^{3})$$

other sorting techniques

Bubble Sort:

$$\frac{10.5 \times 47}{10.5 \times 47} \Rightarrow 5.10 \times 7 \Rightarrow 5.47 = 10$$

$$5 \begin{array}{c|c} 4 & 7 & 10 \Rightarrow 4 & 5 & 7 & 10 \Rightarrow 4 & 5 & 7 & 10 \end{array}$$

