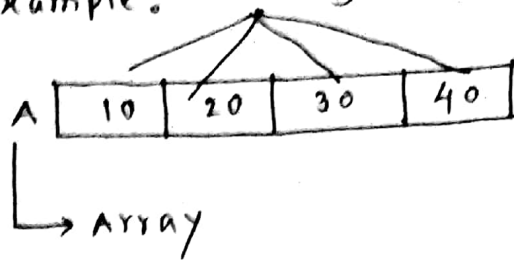


## Data structures:

A particular way of organizing data in a computer so that it can be used effectively.

For example: integer data



We can store a list of items having the same data-type using the array data structure.

## Different Data structures

- Linear: arrays, lists
- Tree: Binary, heaps, space partitioning, etc
- stack
- Queue
- Graph
- Files
- Records

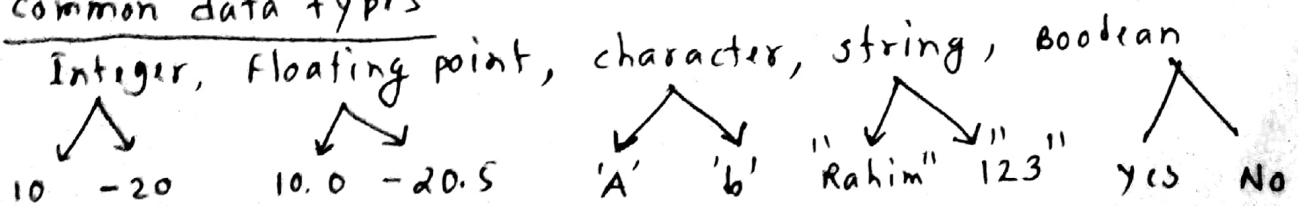
## Uses of Data structures

- Operating system
- Compiler design
- Artificial Intelligence
- Graphics

## Fundamental data structures

- Array
- Linked list

## Common data types



## What is algorithms?

step by step problem solution technique

## Algorithm Representation

- Pseudocode
- flowchart
- Program

# Array Data structure

Page-2

## Sorting:

A

10	20	30
----	----	----

0 1 2

Ascending order

A

30	20	10
----	----	----

0 1 2

Descending order

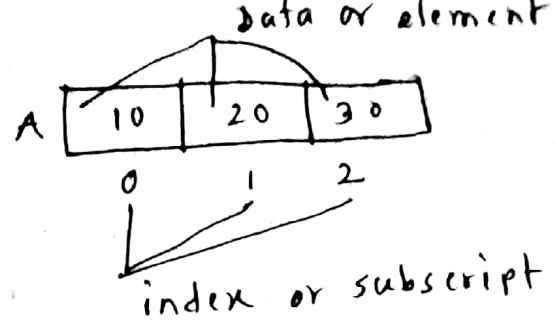
Name

Abu	Babu	Robi
-----	------	------

0 1 2

Alphabetical order

A  
B  
R



$A[0] = 10$

$A[1] = 20$

$A[2] = 30$

A

9	8	6	3	3	1
---	---	---	---	---	---

0 1 2 3 4 5

Non-increasing order

A

1	3	3	6	8	9
---	---	---	---	---	---

0 1 2 3 4 5

Non-decreasing order

## Different types of sorting Techniques

- Bubble Sort
- Selection Sort
- Merge Sort \*
- Insertion Sort \*
- Quick Sort \*
- Heap Sort \*

- Counting Sort \*
- Radix Sort \*
- Bucket Sort

Linear  
time  
sorting

Different cases of  
sorting: (sort in ascending  
order)

Best case: A

10	20	30
----	----	----

0 1 2

Worst case: A

30	20	10
----	----	----

0 1 2

Average case: A

10	30	20
----	----	----

0 1 2

# Counting Sort

Sort the Array A in ascending order using counting sort

A	4	2	3	2	4	1
	1	2	3	4	5	6

Ans:

Input Data: A 

4	2	3	2	4	1
1	2	3	4	5	6

Find Max number:  $k=4$

Initialize Mapping Array: C 

0	0	0	0	0	0
0	1	2	3	4=k	

count number stored in Array A: C 

0	1	2	1	2	
0	1	2	3	4	

find Proper Mapping Array: C 

0	1	3	4	6	
0	1	2	3	4	

$j=6$   $A[6]=1$   $C[1]=1$   $B[1]=1$ 

1					
1	2	3	4	5	6

 $C[1]=0$  C 

0	0	3	4	6	
0	1	2	3	4	

$j=5$   $A[5]=4$   $C[4]=6$   $B[6]=4$ 

1				4	
1	2	3	4	5	6

 $C[4]=5$  C 

0	0	3	4	5	
0	1	2	3	4	

$j=4$   $A[4]=2$   $C[2]=3$   $B[3]=2$ 

1	2			4	
1	2	3	4	5	6

 $C[2]=2$  C 

0	0	2	4	5	
0	1	2	3	4	

$j=3$   $A[3]=3$   $C[3]=4$   $B[4]=3$ 

1	2	3		4	
1	2	3	4	5	6

 $C[3]=3$  C 

0	0	2	3	5	
0	1	2	3	4	

$j=2$   $A[2]=2$   $C[2]=2$   $B[2]=2$ 

1	2	2	3		4
1	2	3	4	5	6

 $C[2]=1$  C 

0	0	1	3	5	
0	1	2	3	4	

$j=1$   $A[1]=4$   $C[4]=5$   $B[5]=4$ 

1	2	2	3	4	4
1	2	3	4	5	6

 $C[4]=4$  C 

0	0	1	3	4	
0	1	2	3	4	

## Algorithm

countingsort(A, B, k) { output

for  $i = 0$  to  $k$

$C[i] = 0$

for  $j = 1$  to  $n$

$C[A[j]] = C[A[j]] + 1$

for  $i = 2$  to  $k$

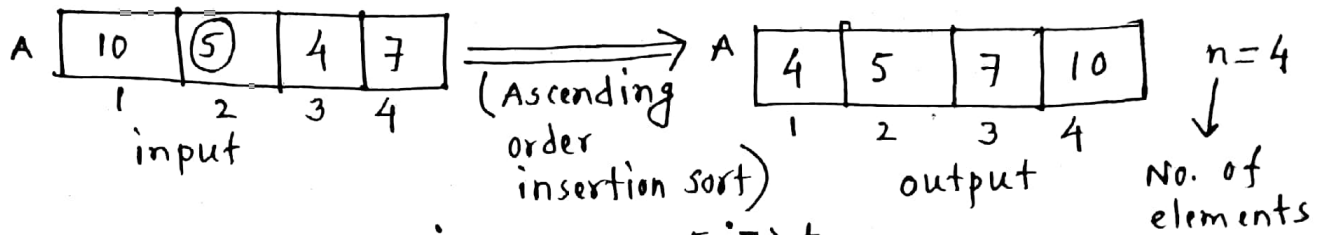
$C[i] = C[i] + C[i-1]$

for  $j = n$  down to  $1$   
 $B[C[A[j]]] = A[j]$   
 $C[A[j]] = C[A[j]] - 1$

# Insertion Sort

Page-4

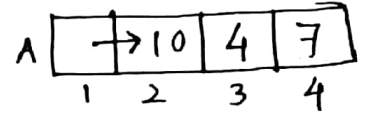
Show the mechanism of insertion sort for the following data



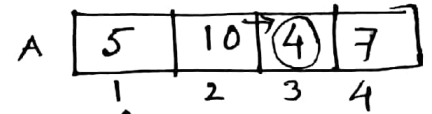
Ans:

$$j=2 \quad t=A[2] \quad i=j-1 \\ =5 \quad =2-1 \\ =1$$

$$i \geq 1 \text{ AND } A[i] > t \\ 1 \geq 1 \text{ AND } A[1] > 5 \\ \Rightarrow \text{True AND } 10 > 5 = \text{True}$$

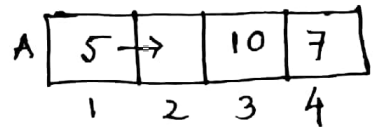


$$i=0 \quad 0 \geq 1 \text{ AND } - \\ \Rightarrow \text{False AND } - \\ i+1 = 0+1 = 1$$

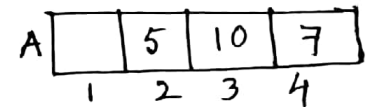


$$j=3 \quad t=A[3] \quad i=j-1 \\ =4 \quad =3-1 \\ =2$$

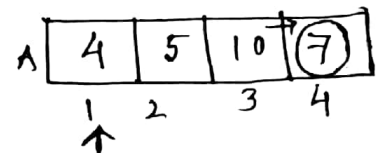
$$2 \geq 1 \text{ AND } A[2] > 4 \\ \Rightarrow \text{True AND } 10 > 4 = \text{True}$$



$$i=1 \quad 1 \geq 1 \text{ AND } A[1] > 4 \\ \Rightarrow \text{True AND } 5 > 4 = \text{True}$$

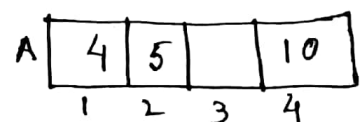


$$i=0 \quad 0 \geq 1 \text{ AND } - = \text{False} \\ \text{False} \quad i+1 = 0+1 = 1$$

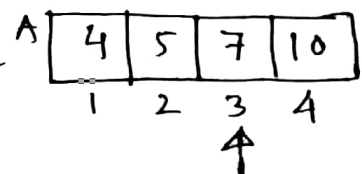


$$j=4 \quad t=A[4] \quad i=j-1 \\ \downarrow n \quad =7 \quad =4-1 \\ =3$$

$$3 \geq 1 \text{ AND } A[3] > 7 \\ \Rightarrow \text{True AND } 10 > 7 = \text{True}$$



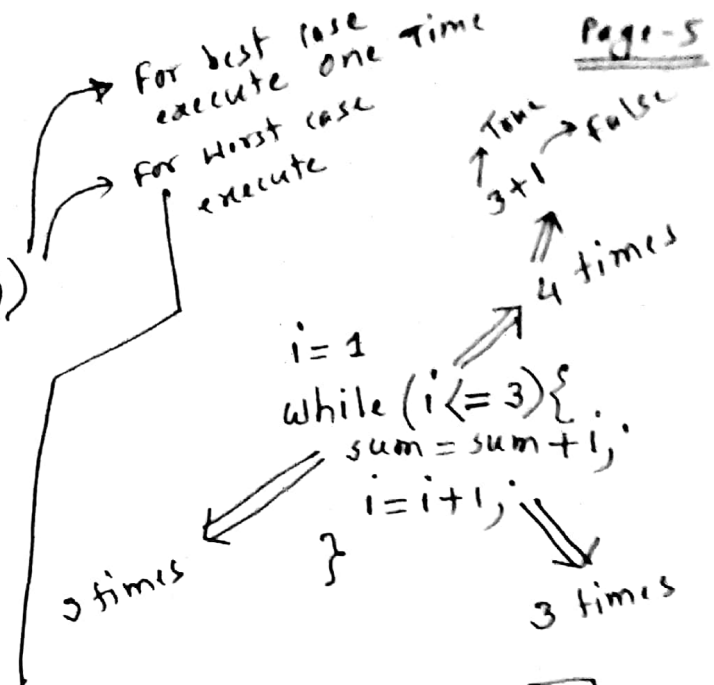
$$i=2 \quad 2 \geq 1 \text{ AND } A[2] > 7 \\ \Rightarrow \text{True AND } 5 > 7 = \text{False} \\ i+1 = 2+1 = 3$$



## Algorithm

```

L1: for j = 2 to n
L2:   t = A[j]
L3:   i = j-1
L4:   while ((i >= 1) AND (A[i] > t))
L5:     A[i+1] = A[i]
L6:     i = i-1
      end while
L7:   A[i+1] = t
      end for
  
```



## Time Complexity Analysis

### Best case Analysis:

A	10	20	30
	1	2	3

1
2
3
4

4-1+1 = 4

Line Number	Iteration	Time/Iteration	Time
L1	$(n-2+1)+1$ True False	$c_1$	$c_1 n$
L2	$n-1$	$c_2$	$c_2(n-1)$
L3	$n-1$	$c_3$	$c_3(n-1)$
L4	$n-1$	$c_4$	$c_4(n-1)$
L5	X	X	X
L6	X	X	X
L7	$n-1$	$c_7$	$c_7(n-1)$

$$\begin{aligned}
 \text{Total running time, } T(n) &= c_1 n + c_2(n-1) + c_3(n-1) + c_4(n-1) + c_7(n-1) \\
 &= (c_1 + c_2 + c_3 + c_4 + c_7) n + (-c_2 - c_3 - c_4 - c_7) \\
 &= An + B = O(n)
 \end{aligned}$$



# Worst Case Analysis

Page-6

A	30	20	10
	1	2	3

$j=2$   $t=20$   $i=1$  True } 2 times  
 $i=0$  False }  $\downarrow j$

$j=3$   $t=10$   $i=2$  True } 3 times  
 $i=1$  True }  $\downarrow j$   
 $i=0$  False }  $j$  times

Line Number	Iteration	Time/Iteration	Time
L1	$(n-2+1) + \underbrace{1}_{\text{True False}} = n$	$c_1$	$c_1 n$
L2	$n-1$	$c_2$	$c_2 (n-1)$
L3	$n-1$	$c_3$	$c_3 (n-1)$
L4	$\sum_{j=2}^n j$	$c_4$	$c_4 \sum_{j=2}^n j$
L5	$\sum_{j=2}^n (j-1)$	$c_5$	$c_5 \sum_{j=2}^n (j-1)$
L6	$\sum_{j=2}^n (j-1)$	$c_6$	$c_6 \sum_{j=2}^n (j-1)$
L7	$n-1$	$c_7$	$c_7 (n-1)$

Total running time,  $T(n) = c_1 n + c_2 (n-1) + c_3 (n-1) + c_4 \sum_{j=2}^n j$   
 $+ c_5 \sum_{j=2}^n (j-1) + c_6 \sum_{j=2}^n (j-1) + c_7 (n-1)$

$$\begin{pmatrix} n \\ 2 \end{pmatrix} \sqrt{\frac{n}{2}} \quad \begin{pmatrix} n \\ 2 \end{pmatrix} \sqrt{\frac{n}{2}}$$

$$\begin{matrix} n \\ + \\ n \\ + \\ \dots \\ + \\ 1+2+\dots+n \end{matrix} \quad \begin{matrix} n \\ + \\ n-1 \\ + \\ n-2 \\ + \\ \dots \\ + \\ 1 \end{matrix}$$

$$= c_1 n + c_2 (n-1) + c_3 (n-1) + c_4 (2+3+\dots+n)$$

$$+ c_5 (1+2+\dots+n-1) + c_6 (1+2+3+\dots+n-1)$$

$$+ c_7 (n-1)$$

$$= c_1 n + c_2 (n-1) + c_3 (n-1) + c_4 (1+2+3+\dots+n-1)$$

$$+ c_5 \frac{n(n-1)}{2} + c_6 \frac{n(n-1)}{2} + c_7 (n-1)$$

$$= c_1 n + c_2(n-1) + c_3(n-1) + c_4 \left\{ \frac{n(n+1)}{2} - 1 \right\} + c_5 \frac{n(n-1)}{2} + c_6 \frac{n(n-1)}{2} + c_7(n-1)$$

$$= c_1 n + c_2 n - c_2 + c_3 n - c_3 + \frac{c_4}{2} n^2 + \frac{c_4}{2} n - c_4 + \frac{c_5}{2} n^2 - \frac{c_5}{2} n + \frac{c_6}{2} n^2 - \frac{c_6}{2} n + c_7 n - c_7$$

$$= \left( \frac{c_4}{2} + \frac{c_5}{2} + \frac{c_6}{2} \right) n^2 + \left( c_1 + c_2 + c_3 + \frac{c_4}{2} - \frac{c_5}{2} - \frac{c_6}{2} + c_7 \right) n + (-c_2 - c_3 - c_4 - c_7)$$

$$= An^2 + Bn + C$$

$$= O(n^2)$$

### Average case Analysis

$$\text{condition of while loop} = \frac{\text{best} + \text{Worst}}{2} = \frac{1+j}{2}$$

Line Number	Iteration	Time/Iteration	Time
L1	n	c <sub>1</sub>	c <sub>1</sub> n
L2	n-1	c <sub>2</sub>	c <sub>2</sub> (n-1)
L3	n-1	c <sub>3</sub>	c <sub>3</sub> (n-1)
L4	$\sum_{j=2}^n \frac{1+j}{2}$	c <sub>4</sub>	$c_4 \sum_{j=2}^n \frac{1+j}{2}$
L5	$\sum_{j=2}^n \left\{ \frac{1+j}{2} - 1 \right\}$	c <sub>5</sub>	$c_5 \sum_{j=2}^n \frac{j-1}{2}$
L6	$\sum_{j=2}^n \left\{ \frac{1+j}{2} - 1 \right\}$	c <sub>6</sub>	$c_6 \sum_{j=2}^n \frac{j-1}{2}$
L7	<del>Σ</del> (n-1)	c <sub>7</sub>	c <sub>7</sub> (n-1)

$$\text{Total running time, } T(n) = c_1 n + c_2(n-1) + \dots + c_7(n-1)$$

$$= O(n^2)$$