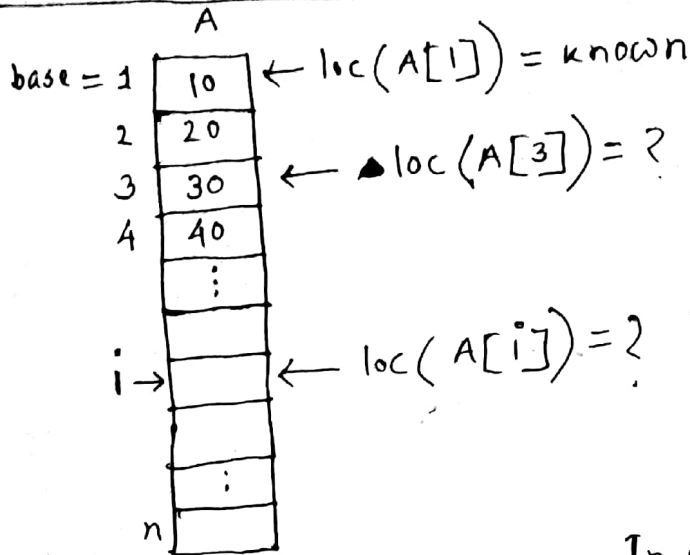


Memory Allocation

One dimensional array:



Each integer data = 4 bytes
= width of int
= w

$$\begin{aligned}\text{loc}(A[3]) &= \text{loc}(A[1]) + 4 + 4 \\ &= \text{loc}(A[1]) + 4 * 2 \\ &= \text{loc}(A[1]) + 4 * (3-1)\end{aligned}$$

In general,

$$\begin{aligned}\text{loc}(A[i]) &= \text{loc}(A[1]) + w * (i-1) \\ &= \text{loc}(A[\text{base}]) + w * (i - \text{base})\end{aligned}$$

$$\text{loc}(A[i]) = \text{loc}(A[\text{base}]) + w * (i - \text{base})$$

Example

double A[100];
if $\text{loc}(A[10]) = (5CDF)_{16}$, find $\text{loc}(A[60])$.

Ans:

For double, w = 8 bytes

$$\text{loc}(A[10]) = (5CDF)_{16}$$

$$i = 60$$

$$\text{base} = 10$$

From Formula,

$$\begin{aligned}\text{loc}(A[i]) &= \text{loc}(A[\text{base}]) + w * (i - \text{base}) \\ \Rightarrow \text{loc}(A[60]) &= \text{loc}(A[10]) + 8 * (60 - 10) \\ &= (5CDF)_{16} + 8 * 50 \\ &= (5CDF)_{16} + (400)_{10} \\ &= (5CDF)_{16} + (190)_{16} = (5E6F)_{16}\end{aligned}$$

$$\begin{array}{r} 16 \overline{) 400} \\ 16 \overline{) 25-0} \\ 16 \overline{) 1-9} \\ \hline 0-1 \end{array}$$
$$\begin{array}{r} 5CDF \\ + 190 \\ \hline 5E6F \end{array}$$

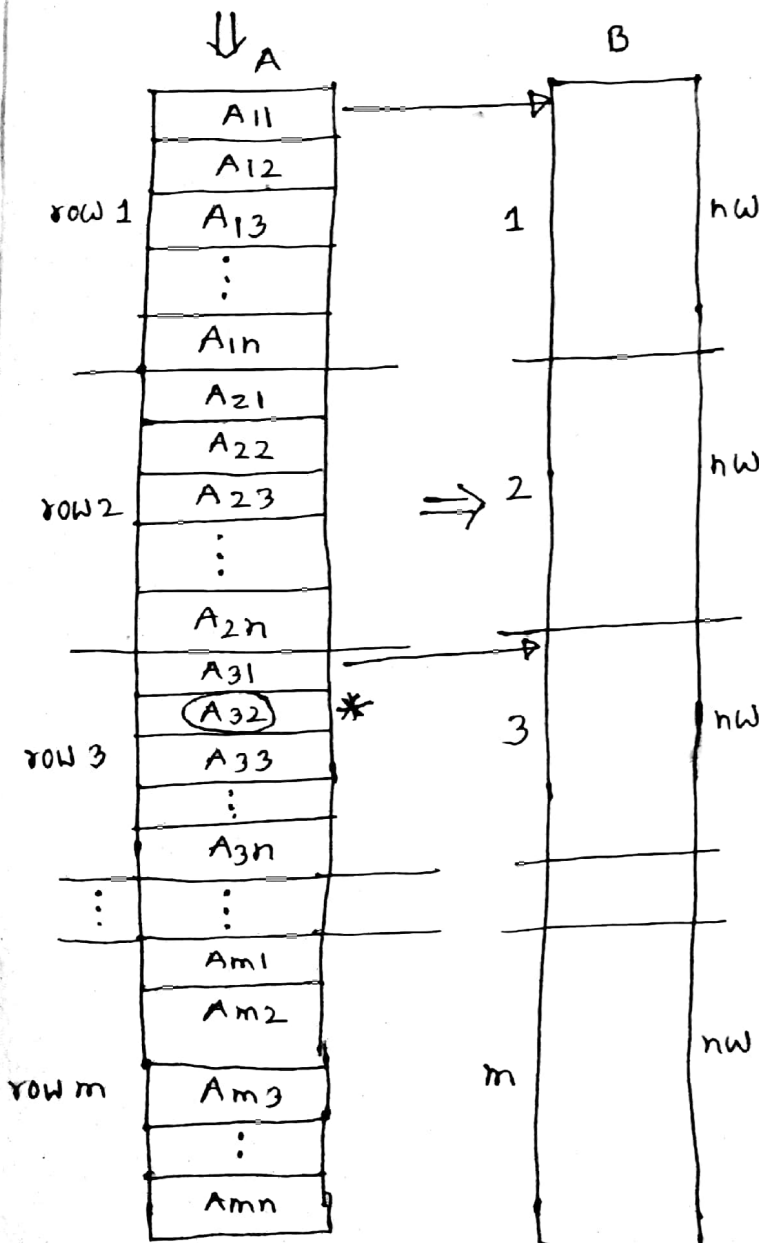
Two dimensional Array:

Page-2

Row major Allocation:

row 1 →	A_{11}	A_{12}	A_{13}	...	A_{1n}
row 2 →	A_{21}	A_{22}	A_{23}	...	A_{2n}
row 3 →	A_{31}	A_{32}	A_{33}	...	A_{3n}
⋮	⋮	⋮	⋮	⋮	⋮
row m →	A_{m1}	A_{m2}	A_{m3}	...	A_{mn}

One data width = w



$$\text{loc}(B[i]) = \text{loc}(B[1]) + nw(i-1)$$

$$\Rightarrow \text{loc}(B[3]) = \text{loc}(B[1]) + nw(3-1)$$

$$\Rightarrow \text{loc}(A_{31}) = \text{loc}(A_{11}) + nw(3-1)$$

Again,

$$\text{loc}(A_{32}) = \text{loc}(A_{31}) + w$$
$$= \text{loc}(A_{31}) + w(2-1)$$

$$= \text{loc}(A_{11}) + nw(3-1) + w(2-1)$$

$$= \text{loc}(A_{11}) + w[n(3-1) + (2-1)]$$

In general,

$$\text{loc}(A_{ij}) = \text{loc}(A_{11}) + w[n(i-1) + (j-1)]$$

for row-wise allocation

$$\text{loc}(A_{ij}) = \text{loc}(A_{11}) + w[n(i-1) + (j-1)]$$

for column-wise allocation

$$\text{loc}(A_{ij}) = \text{loc}(A_{11}) + w[m(j-1) + (i-1)]$$

Examplefloat $A[10][20];$ if $\text{loc}(A[0][0]) = (5CDF)_{16}$, Find $\text{loc}(A[6][15])$.

Assume row-major allocation

Ans:For float, $w = 4$ bytes

$$\text{loc}(A[0][0]) = (5CDF)_{16}$$

$$i = 6$$

$$j = 15$$

$$m = 10$$

$$n = 20$$

From formula,

$$\text{loc}(A[i][j]) = \text{loc}(A[0][0]) + w[n(i-0) + (j-0)]$$

$$= (5CDF)_{16} + 4[20(6-0) + (15-0)]$$

$$= (5CDF)_{16} + 4[20*6 + 15]$$

$$= (5CDF)_{16} + 4[120 + 15]$$

$$= (5CDF)_{16} + 4*135$$

$$= (5CDF)_{16} + (540)_{16}$$

$$= (5CDF)_{16} + (21C)_{16}$$

$$= (5EFB)_{16}$$

$$\begin{array}{r} 16 \overline{) 540} \\ 16 \overline{) 33 - (12)C} \\ 16 \overline{) 2 - 1} \\ \hline 0 - 2 \end{array}$$

$$\begin{array}{r} 5 C D F \\ 2 1 C \\ \hline 5 E F B \end{array}$$

Memory Allocation for structure :

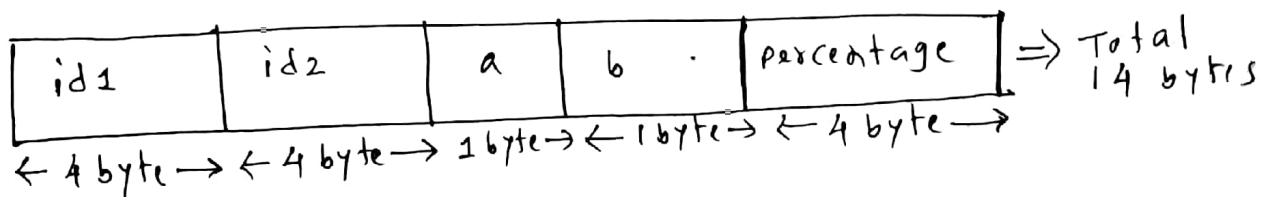
Page - 4

```
struct student {  
    int id1;  
    int id2;  
    char a;  
    char b;  
    float percentage;  
};
```

Find Memory allocation for this structure. Assume that 32 bit computer.

Ans:

From structure



Since 32 bit computer

Word size = 32 bit

= 4 × 8 bit

= 4 × 1 byte = 4 byte

First CPU cycle = id1 (4 byte)

second CPU cycle = id2 (4 byte)

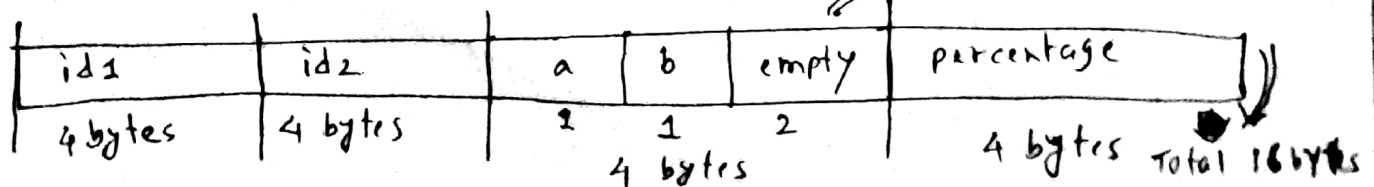
3rd CPU cycle = a, b, empty (4 byte)

↓ ↓ ↓
1 1 2 bytes

4th CPU cycle = percentage (4 byte)

↓
4 byte

Actual structure memory allocation → structure padding



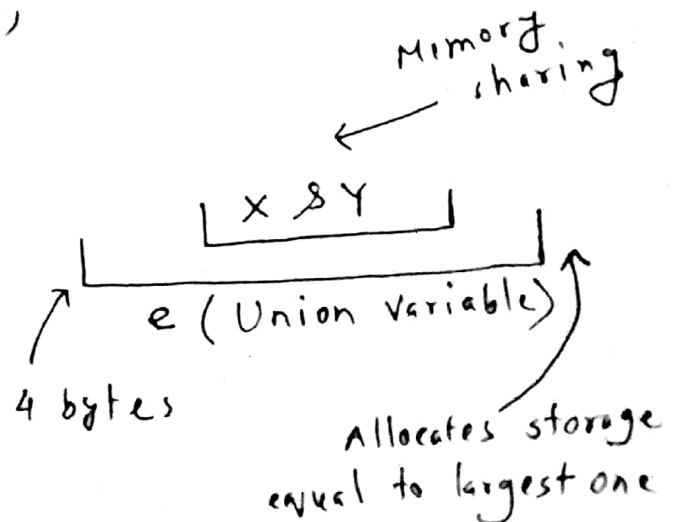
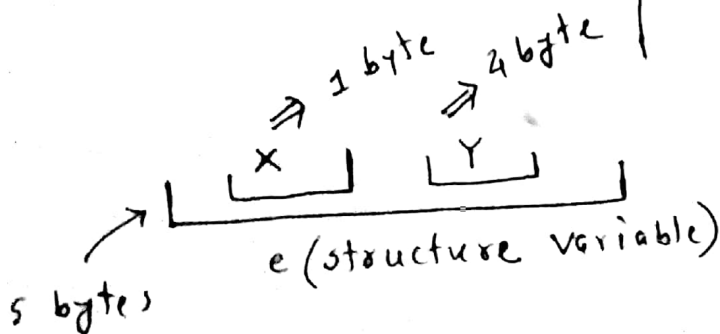
Data types	Memory allocation in C 32 bit computer		
	From Address	To Address	Total bytes
int id1	675376768	675376771	4
int id2	675376772	675376775	4
char a	675376776	675376776	1
char b	675376777	675376777	1
structure Padding (675376778 to 675376779) empty			2
float percentage	675376780	675376783	4

→ 1 CPU cycle
→ 1 CPU cycle
→ 1 CPU cycle
→ 1 CPU cycle

Difference between Memory Allocation of structure and union:

```
struct Emp
{
    char x; // 1 byte
    float y; // 4 byte
};
```

```
union Emp
{
    char x;
    float y;
};
```



Linked list:

Address	data	next
1001	E	1003
1002		
1003	A	NULL
1004		
1005	C	1007
1006		
1007	N	1009
1008		
1009	D	1001
1010		

