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# 03. SPECIAL LINKED STRUCTURES

‣ *Introduction*

‣ *Linked Lists*

‣ *Doubly Linked Lists*

‣ *Circular Linked Lists*

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# 03. SPECIAL LINKED STRUCTURES

‣ *Introduction*

‣ *Singly Linked Lists (revisit)*

‣ *Doubly Linked Lists*

‣ *Circular Linked Lists*

# Introduction

Array.  A collection of contiguous data elements whose order is given by the physical memory location.

Linked List.  A collection of data elements whose order is NOT given by the physical memory location.

Arrays and Linked Lists are the most fundamental memory organizations that are used to create more data structure abstractions such as Stacks and Queues.

# Sequential and Linked Data Structures

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# SPECIAL LINKED STRUCTURES

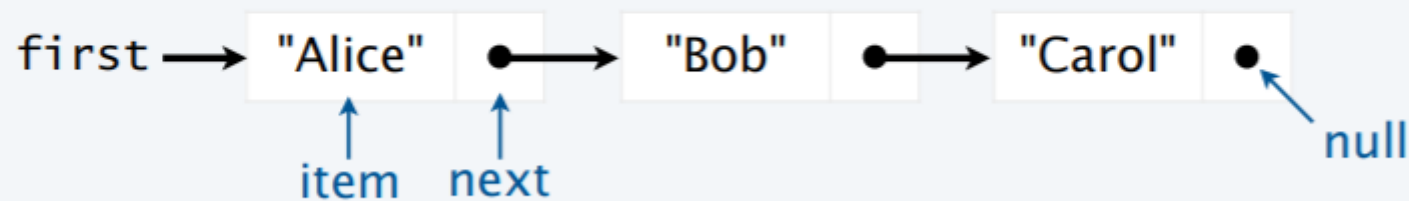# Simplest singly-linked data structure – Linked List

## Linked list

- A recursive data structure.
- Def. A *linked list* is null or a reference to a *node*.
- Def. A *node* is a data type that contains a reference to a node.
- Unwind recursion: A linked list is a sequence of nodes.

```
private class Node
{
    private String item;
    private Node next;
}
```
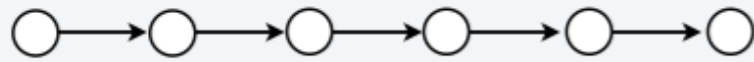
## Representation

- Use a private nested class Node to implement the node abstraction.
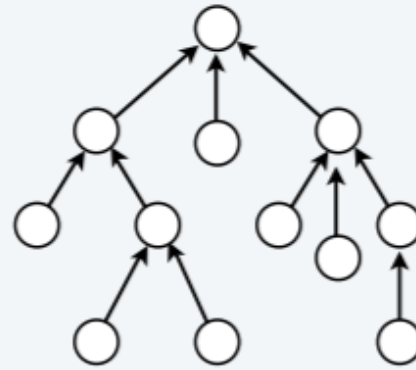- For simplicity, start with nodes having two values: a String and a Node.
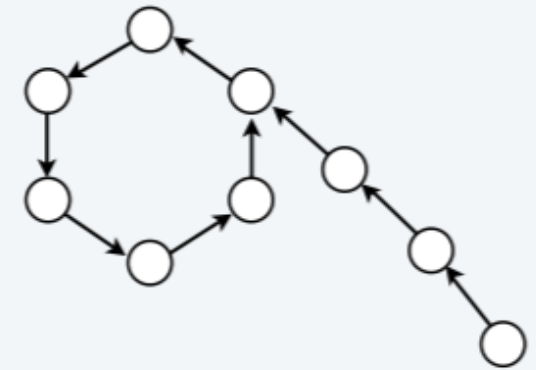
**A linked list**

first ⟶ "Alice" ● ⟶ "Bob" ● ⟶ "Carol" ●⟶ null

    item  next

# Linked Data Structures

**Linked list (this lecture)**

**Tree**

**Rho**

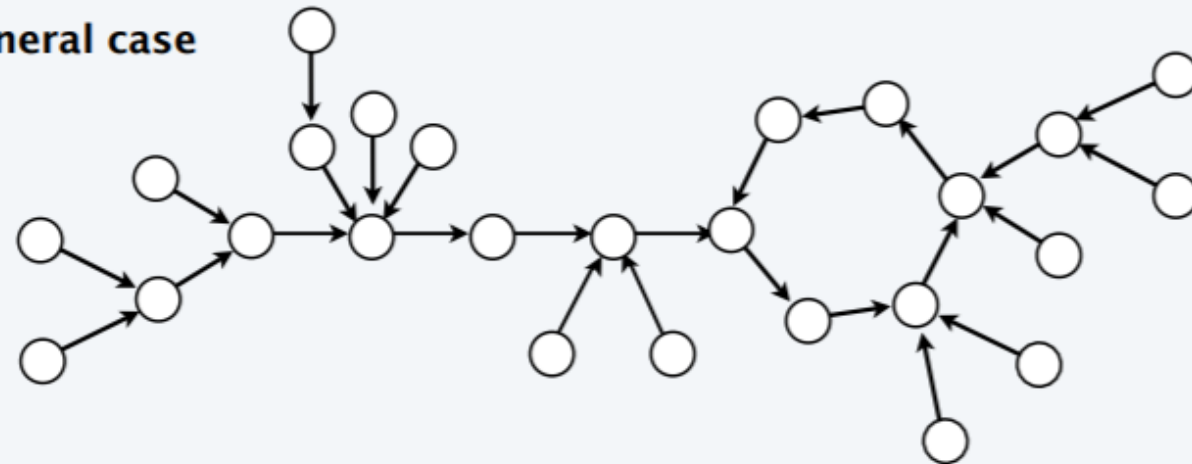**Circular list (TSP)**
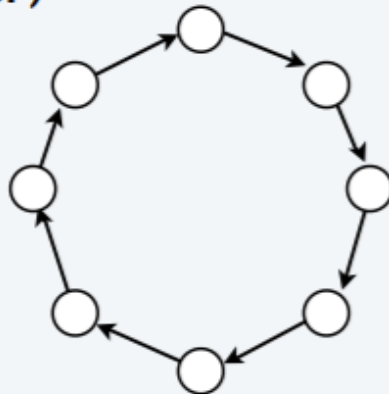
**General case**

Multiply linked structures: many more possibilities!

From the point of view of a particular object, all of these structures look the same.

# Building a Singly Linked List

| addr | value |
|------|-------|
| C0 | "Carol" |
| C1 | null |
| C2 | |
| C3 | |
| C4 | "Alice" |
| C5 | CA |
| C6 | |
| C7 | |
| C8 | |
| C9 | |
| CA | "Bob" |
| CB | C0 |

third  C0

second  CA

first  C4

first        second        third

"Alice"  →  "Bob"  →  "Carol"  ●

null

first.next = second     second.next = third
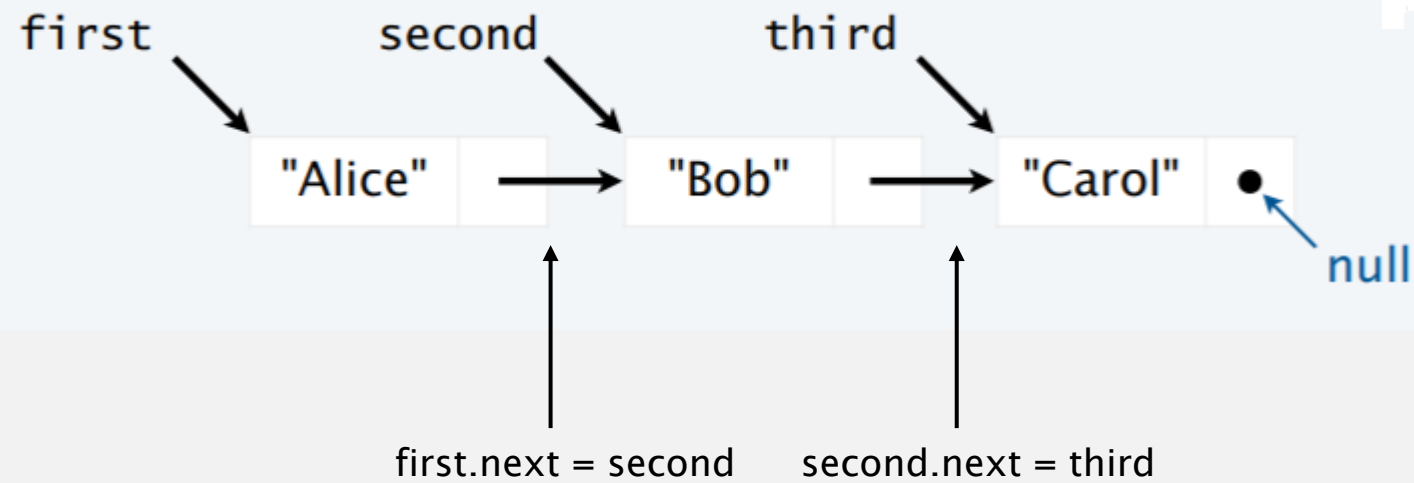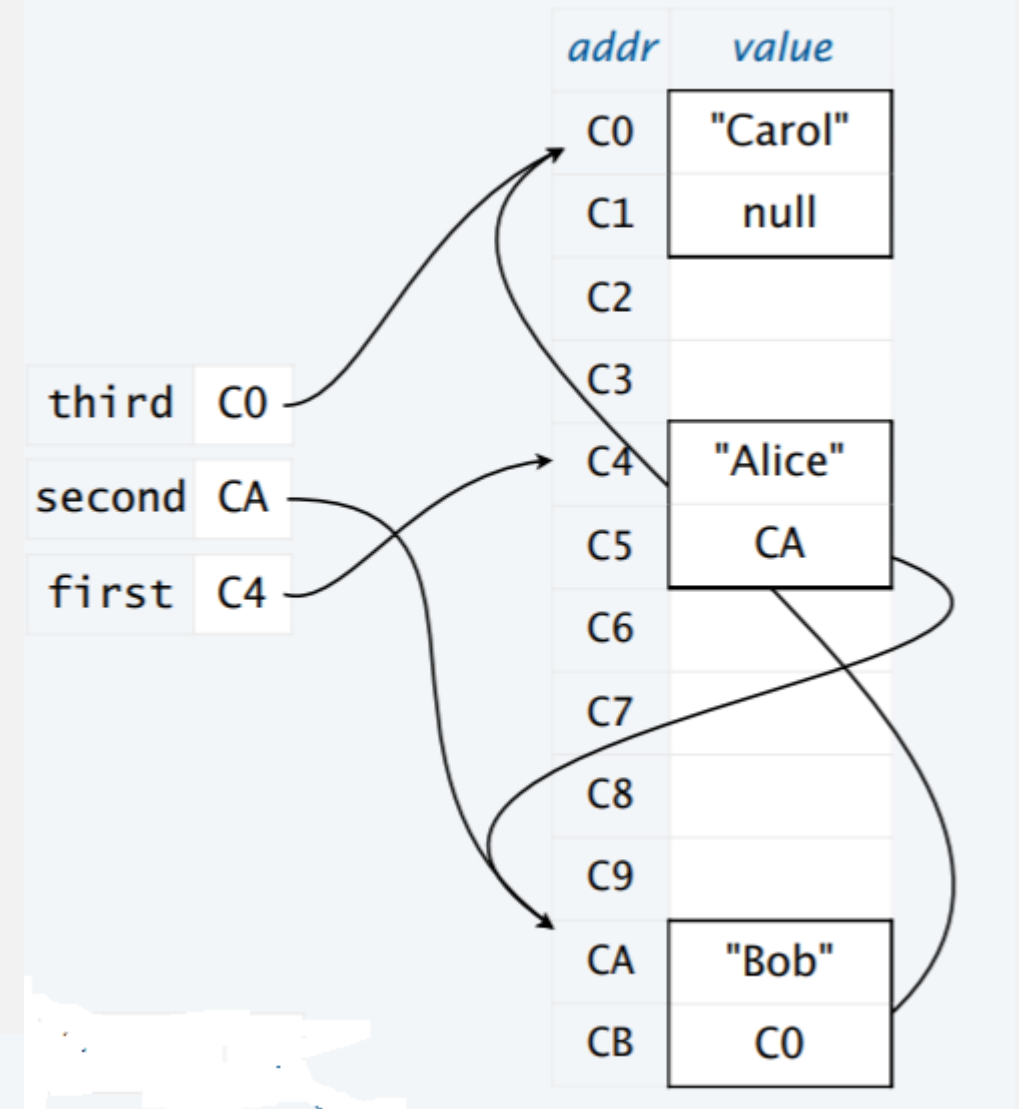
# List Processing Code

Standard operations for processing data structured as a singly-linked list

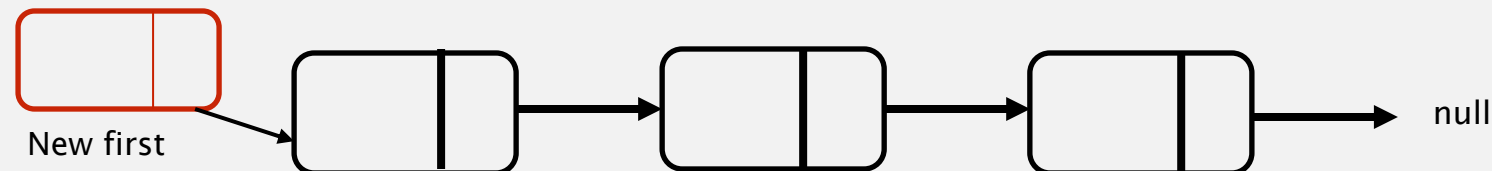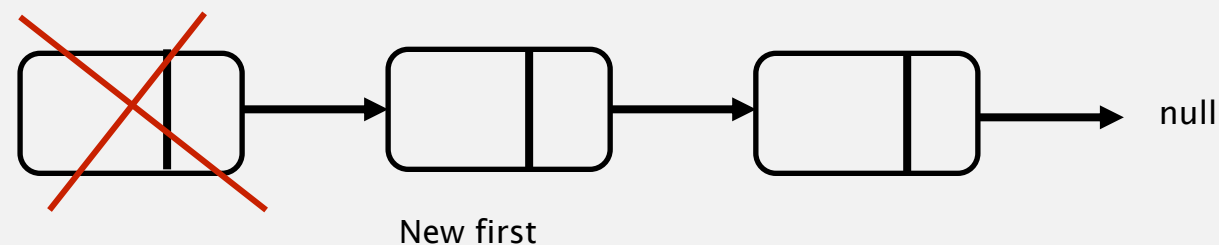- Add a node at the beginning of the list



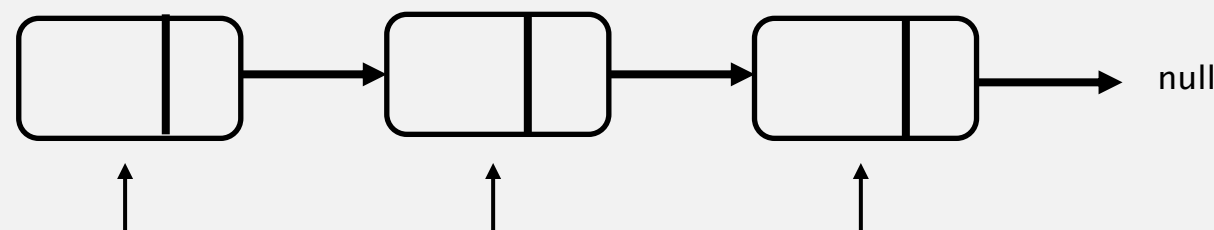New first

null

- Remove and return the node at the beginning



New first

null

- Add a node at the end (requires a reference to the last node)



last

New last

null
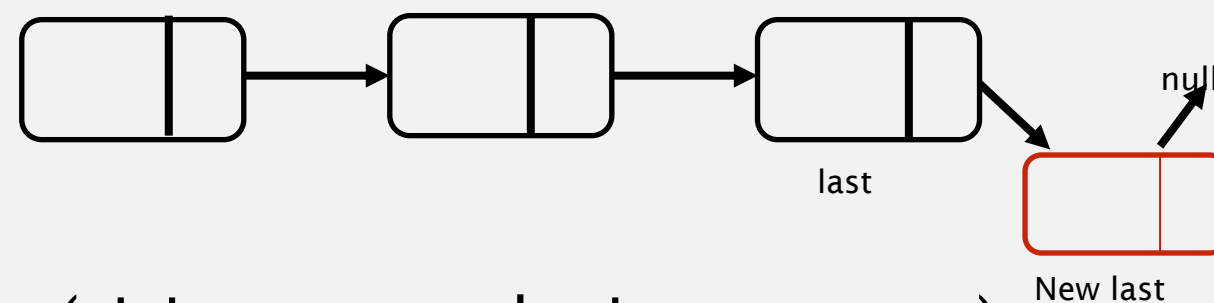
- Traverse the list (visit every node, in sequence).



null

# List Processing code – Remove and Return First Item

Goal. Remove and return the first item in a linked list `first`.

first → "Alice" •→ "Bob" •→ "Carol" •

```
item = first.item;
```

item
"Alice"

first → "Alice" •→ "Bob" •→ "Carol" •

```
first = first.next;
```

item
"Alice"

first → "Alice" •→ "Bob" •→ "Carol" •

available for garbage collection

```
return item;
```

item
"Alice"

first → "Bob" •→ "Carol" •
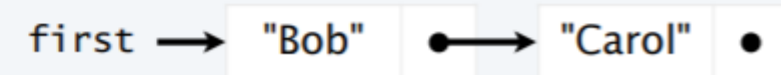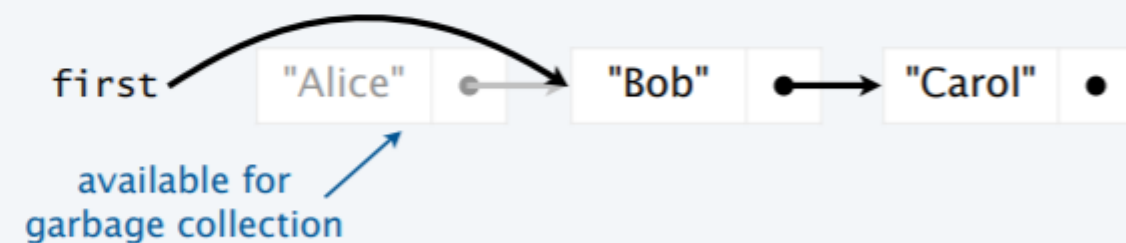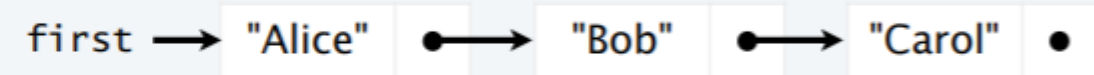
# List Processing Code – Add a new node at the beginning

Goal. Add `item` to a linked list `first`.

```
item
"Dave"
```

first → "Alice" ●→ "Bob" ●→ "Carol" ●

```
Node second = first;
```

second
first → "Alice" ●→ "Bob" ●→ "Carol" ●

```
first = new Node();
```

second
first → ☐ ● "Alice" ●→ "Bob" ●→ "Carol" ●

```
first.item = item;
first.next = second;
```

second
first → "Dave" ●→ "Alice" ●→ "Bob" ●→ "Carol" ●
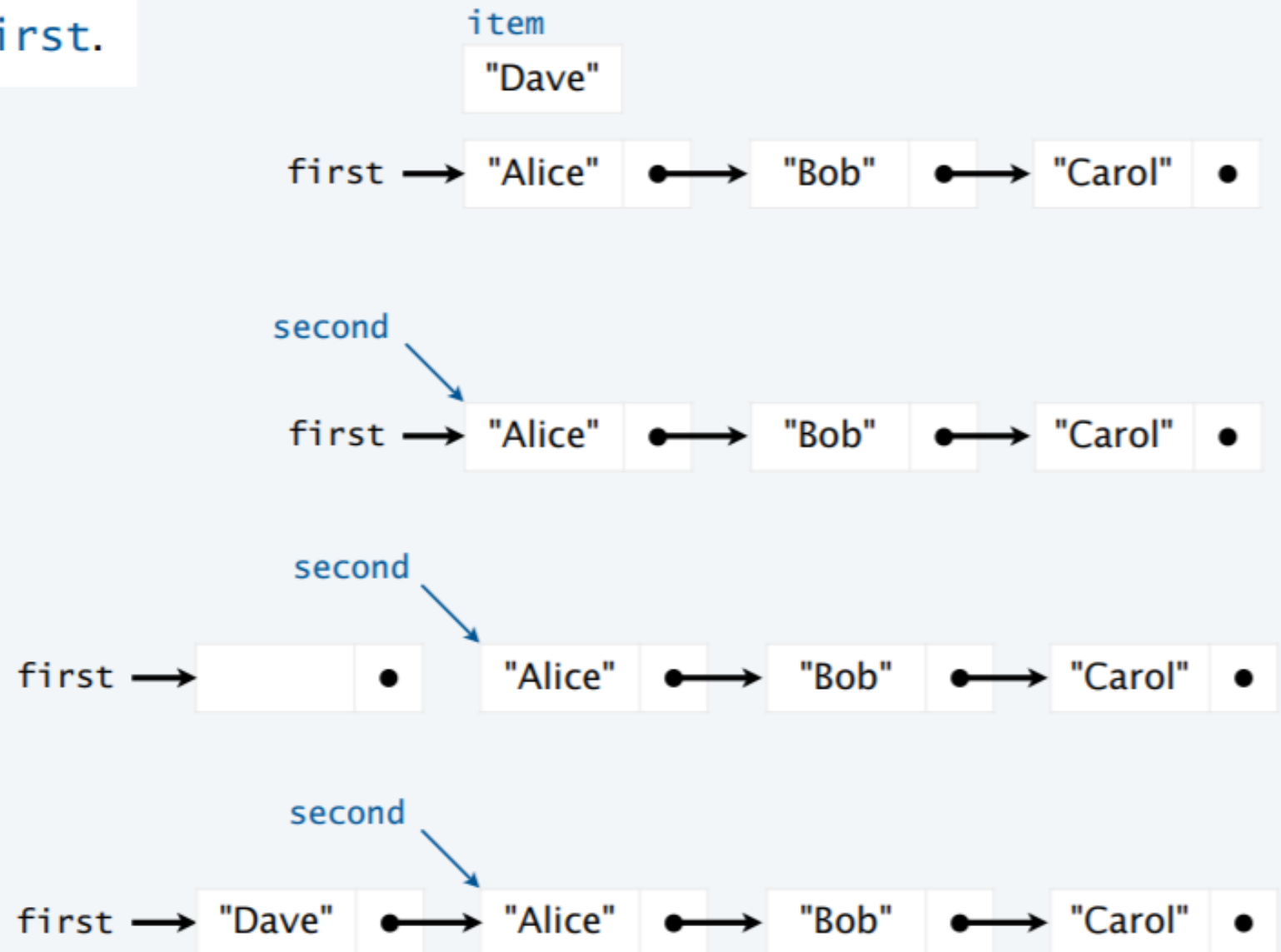
# List processing code – Traverse a List

Goal. Visit every node on a linked list `first`.
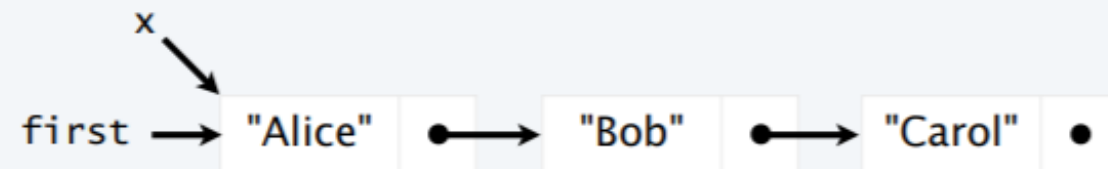
```
Node x = first;
while (x != null)
{
    StdOut.println(x.item);
    x = x.next;
}
```

x

first →  "Alice" ●——→ "Bob" ●——→ "Carol" ●

StdOut
```
Alice
Bob
Carol
```

# Pop quiz1 on Linked Lists

Q. What is the effect of the following code (not-so-easy question)?

```
...
Node list = null;
while (!StdIn.isEmpty())
{
    Node old = list;
    list = new Node();
    list.item = StdIn.readString();
    list.next = old;
}
for (Node t = list; t != null; t = t.next)
    StdOut.println(t.item);
...
```

# Pop quiz 2 on Linked Lists

Q. What is the effect of the following code (not-so-easy question)?

```
...
Node list = new Node();
list.item = StdIn.readString();
Node last = list;
while (!StdIn.isEmpty())
{
    last.next = new Node();
    last = last.next;
    last.item = StdIn.readString();
}
...
```
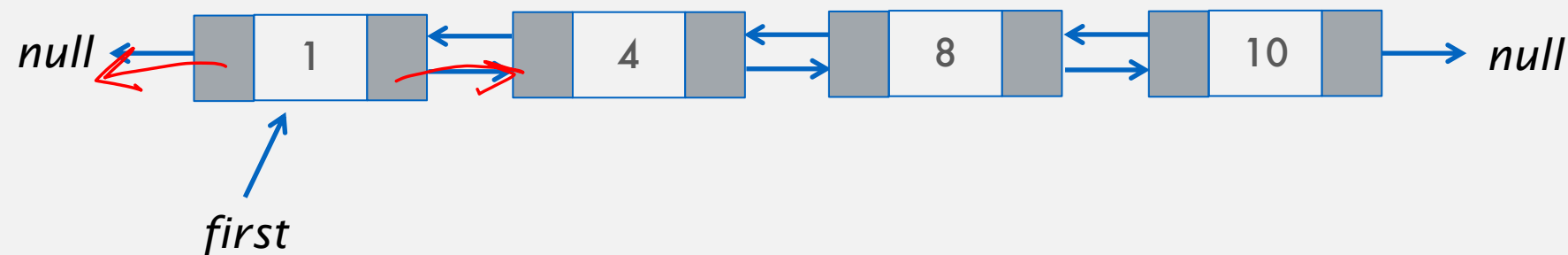
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# SPECIAL LINKED STRUCTURES

# Doubly Linked Lists

A linked list where every node refers to its previous and next nodes



*first*

Each node has three parts:
- a data part
- a link to the previous node
- a link to the next node

link to previous        link to next

data

```
class DoublyLinkedList <T> {
    Node first;
    int n;
}
```

```
class Node <T> {
    T data;
    Node prev;
    Node next;
}
```

## Insert after a target

*target*

Create the new node 8 and insert it after 4

first ——→ | 1 | | 4 | | 10 |

ptr

first ——→ | 1 | | 4 | | 10 |

8

n
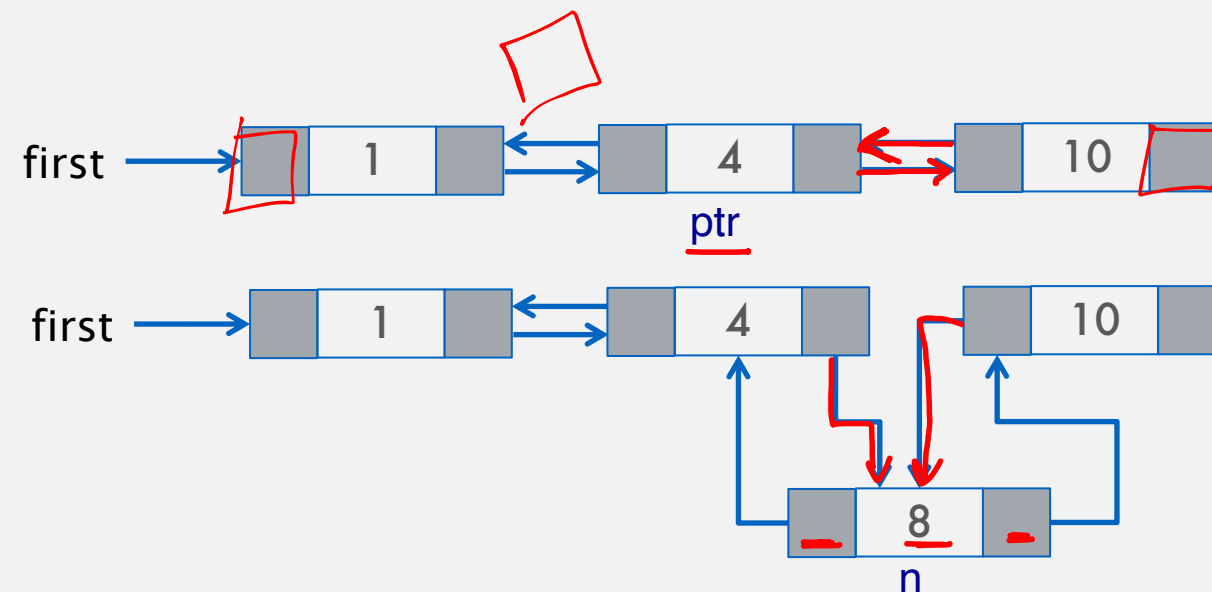
```
Node n = new Node(8);
n.data = 8;
n.next = ptr.next;
ptr.next = n;
n.next.prev = n;
```

$O(1)$

n.prev = PTR;

PTR

F → 1 ⇄ 4 ⇄ 10

8

n

Running time? ~~O(1)~~

B  $O(1)$

W  $O(n)$
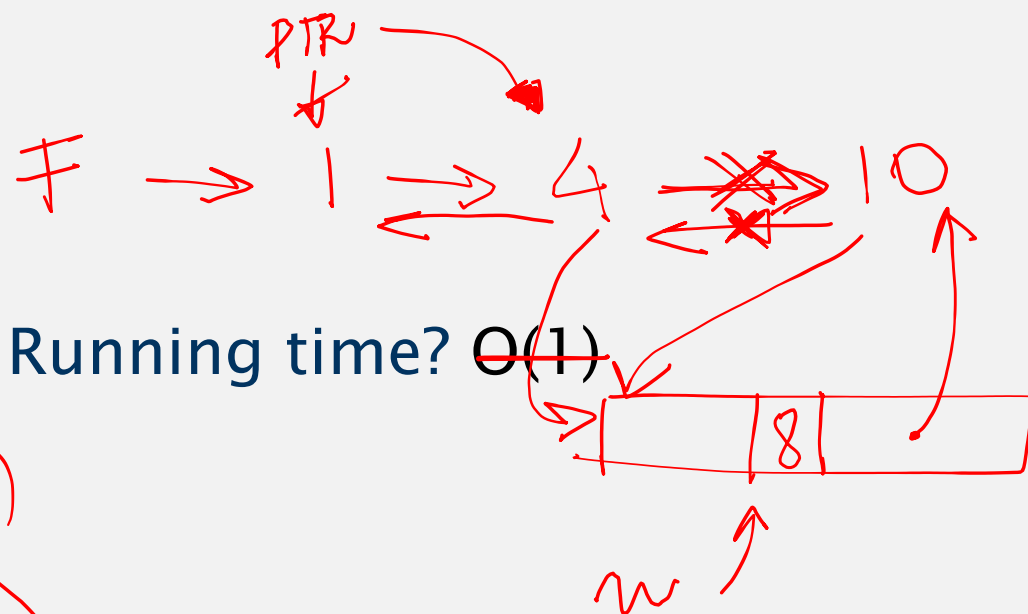
## Remove a target

Delete node containing 4
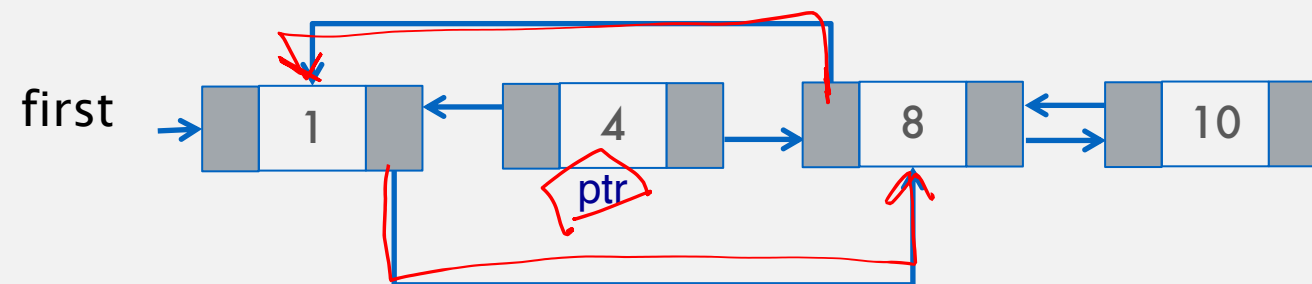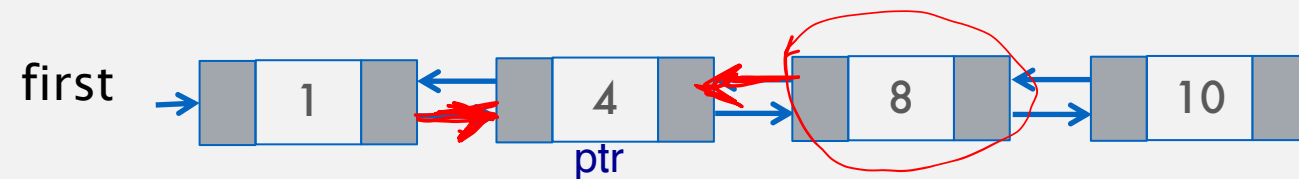


```
ptr.prev.next = ptr.next;
ptr.next.prev = ptr.prev;
```

Running time?
- Best case if target is the first node: O(1)
- Worst case if target is the last node: O(n)

# Applications of Doubly Linked Lists

- A doubly linked might be suitable for many applications where it is important to be able to navigate the linked structure in any direction (forward or backward)

- An application is to maintain the history of web pages visited where both backward and forward navigation of the history is possible.

- In some applications it can be used to maintain most and least recently used memory cache.

- One can design a music player with a next and previous buttons to maintain the history of the playlist, both forward and backward.

## Disadvantages

- Each node requires extra 8 bytes of memory for previous reference
- It is harder to implement methods as operations such as inserting a node to the middle requires 4 statements

## Advantages

- More versatile than singly linked list as one can traverse in both directions.
- When implementing a dynamic application, reference to mostRecentlyAccessednode can be saved in order to move forward or backward.
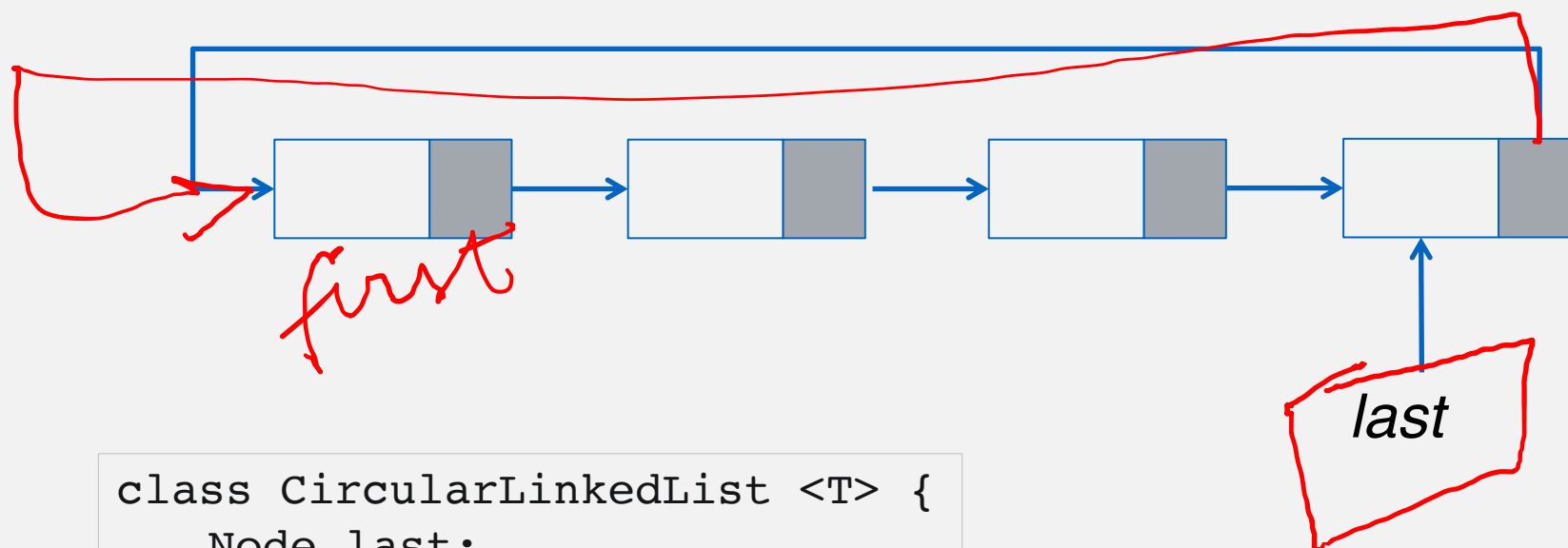
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# SPECIAL LINKED STRUCTURES

A circular link list is a special form of a singly linked list where all nodes are in a continuous cycle. There is no use of null in a circular link list to indicate the end of the list.



*first*

*last*

By keeping a pointer to the last entry we have access to the *first* and *last* entry in constant time.
last: last
first: last.next

```
class CircularLinkedList <T> {
    Node last;
    int n;
}
```

Circular Linked lists have the property that the last node refers back to the first node:  last.next is the reference to the first node

# Advantages of Circular Linked Lists

A circular linked list can be traversed beginning at any node in the list.

A circular linked list can be used to implement a data structure such as a queue where two references front, and back are needed. However, we note that back.next = front

A circular linked list can also be singly or doubly.
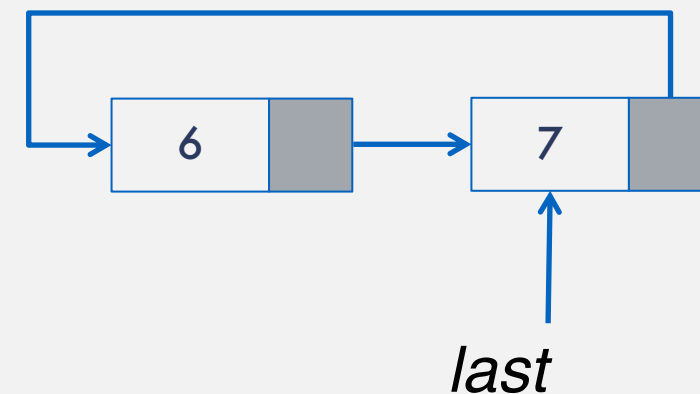
**Add to front**: add an element to the front of the list

**Two cases:**

1. List is empty

2. List is not empty
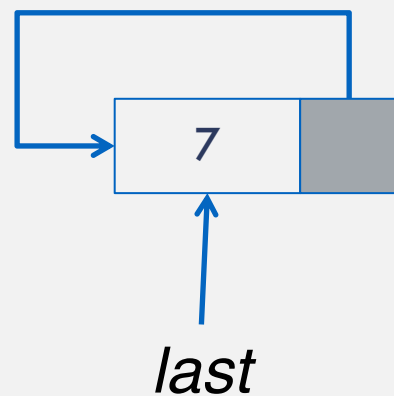
*last*

*last*

```
Node n = new Node();
n.data = 7;
n.next = n;
last = n;
```

```
Node n = new Node();
n.data = 6;
n.next = last.next;
last.next = n;
```
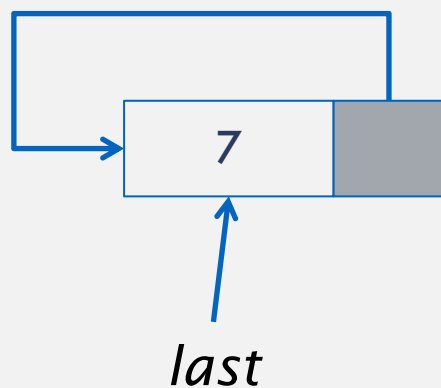
Running time: O(1)

# Implementation of a circular Linked List

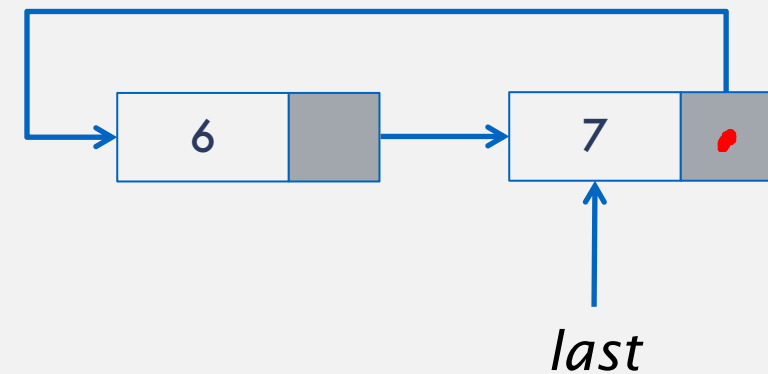**Remove front**: deletes the first element of the list

**Three cases:**

1. List is empty
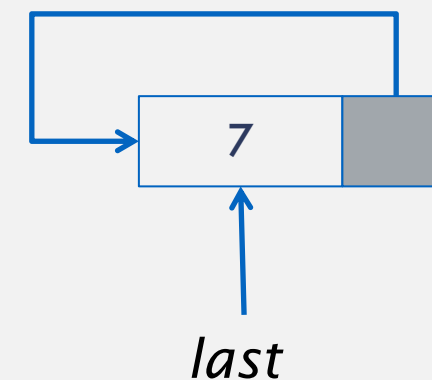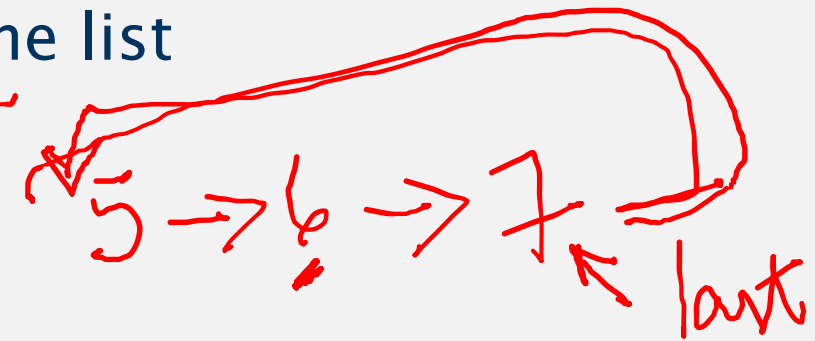2. One element

3. More than one element



```
last = null;
```

Running time: O(1)

```
last.next = last.next.next;
```
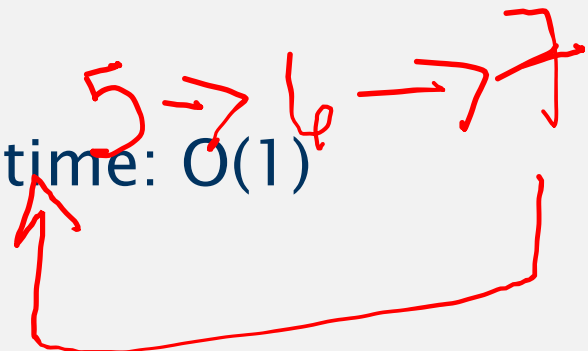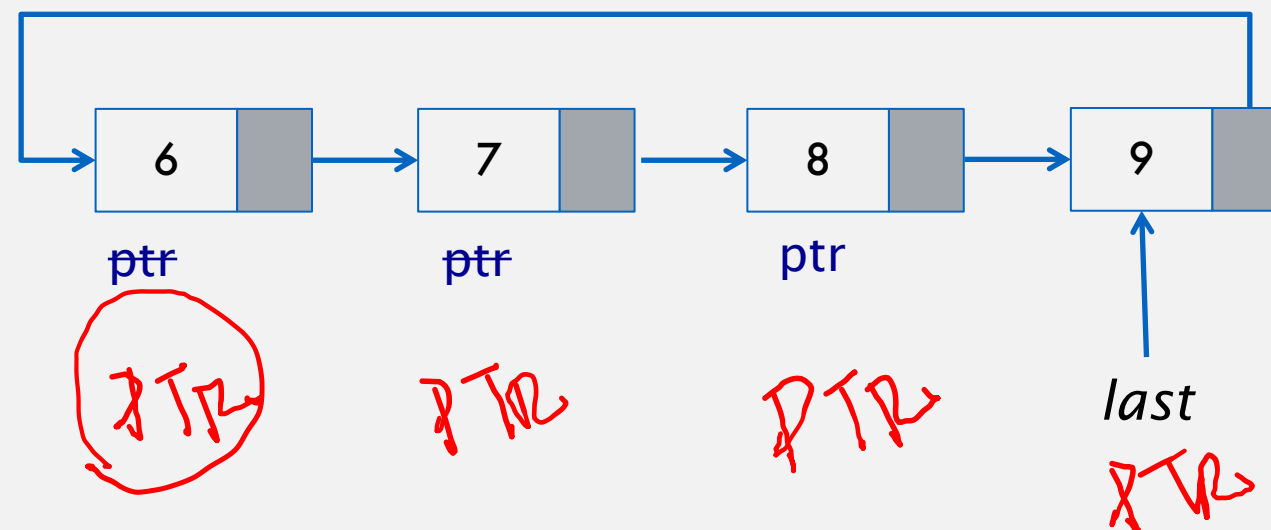
# Implementation of a circular Linked List

To search in a CLL: say target is 8

1. Start a pointer at the front
2. Advance pointer until target is found or the beginning of the list is reached again.



```
Node ptr = last.next;
do{
    if (ptr.data == target) {
    break;
    }
    ptr = ptr.next;
} while (ptr != last.next);
```
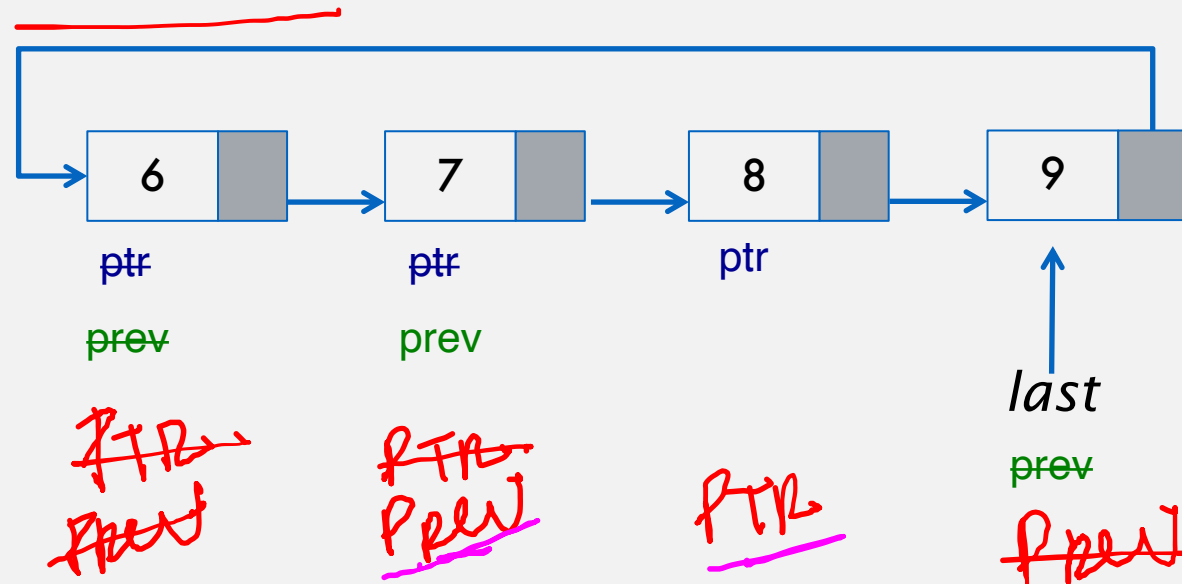
What is the running time?
Worst: O(n)
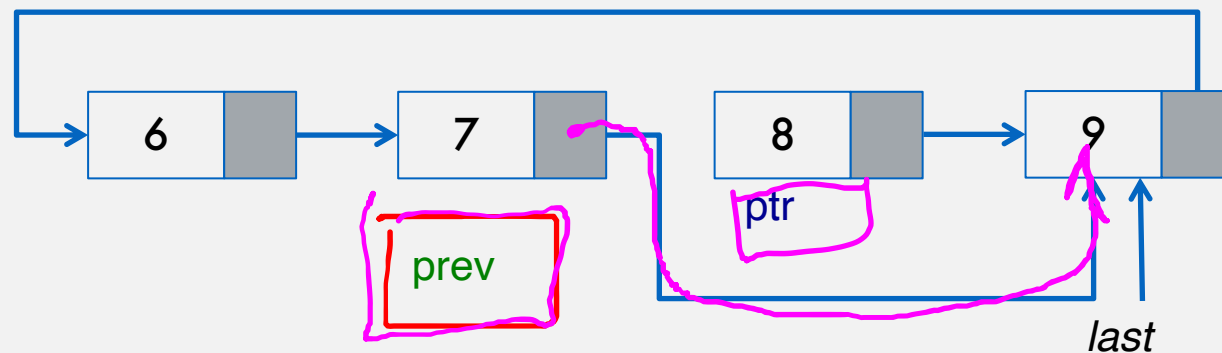Best: O(1)

# Implementation of a circular Linked List

## Delete target: removes the node with the target value

### 1. Find target



6 | 7 | 8 | 9

ptr  ptr  ptr

prev  prev  last  prev

```
Node ptr = last.next;
Node prev = last;
do{
    if (ptr.data == target) {
    break;
    }
    prev = ptr;
    ptr = ptr.next;
} while (ptr != last.next);
```

### 2. Delete target



6 | 7 | 8 | 9

prev  ptr  last

```
prev.next = ptr.next;
```

What is the running time?
  Worst if removing the last node: O(n)
  Best if removing the first node: O(1)

Three cases:
  1.  target not found (ptr == last.next)
  2.  target is last in the list (ptr == last)
  3.  target is found and is not the last in the list

- A circular linked list can be used to implement a queue, where only one reference is needed to implement enqueue or deque operations

- A circular linked list can be used to allocate a shared resource to multiple applications where each application is given some fixed amount of time and the operations can be continued until all applications complete their tasks

- A circular linked list can be used in a multi player game, where each player is given a chance in sequence, and first player is given a chance again in the second round once the last player is done in the first round.

## Disadvantages

- Need to maintain a reference to the last node at all times
- It may be bit trickier to implement some operations such as addLast or addFirst

## Advantages

- This is a great data structure to implement an application such as a multi-player game where players can be maintained in a circular linked list.
- Circular doubly linked list are used for implementation of practical data structures. (later)

# SPECIAL LINKED STRUCTURES

- ‣ *Introduction*
- ‣ *Linked Lists*
- ‣ *Doubly Linked Lists*
- ‣ *Circular Linked Lists*

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Credits: Some slides Adopted and modified from Sedgewick and Wayne

Last updated on 1/11/21 12:10 PM