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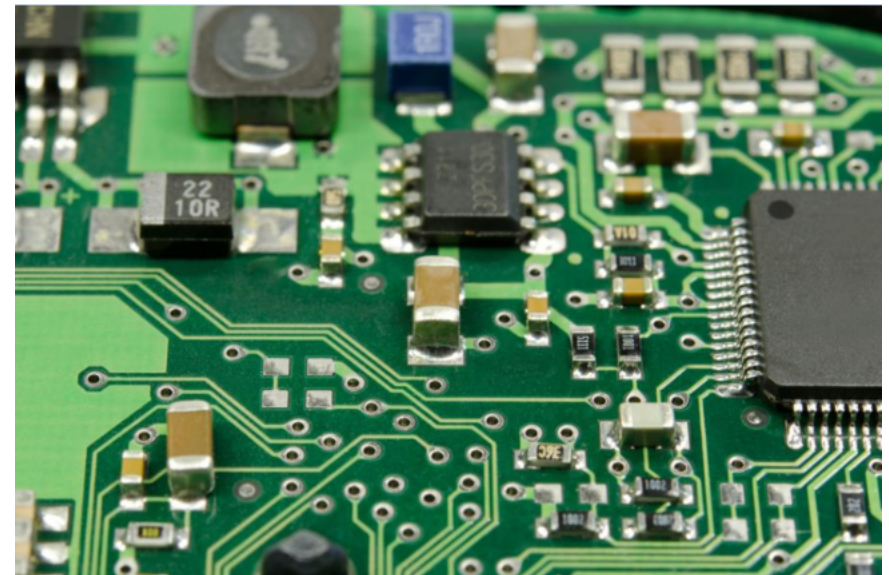
Overview

- Boolean Expressions
- Representing Boolean Functions
- Logic Gates
- Karnaugh Maps



Why is this important?

- We study Boolean algebra as a foundation for designing and analyzing digital systems!
- Boolean Algebra: a useful mathematical system for specifying and transforming logic functions.



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Boolean Expressions

Complement, Sum, Product

- Correspond to logical NOT, OR, and AND.
- We will denote the two logic values as $0 \equiv F$ and $1 \equiv T$, instead of **False** and **True**.
 - Using numbers encourages algebraic thinking.
- New, more algebraic-looking notation for the most common Boolean operators:

$$\bar{x} \equiv \neg x$$

$$x \cdot y \equiv x \wedge y$$

$$x + y \equiv x \vee y$$

Precedence order \rightarrow



- Boolean values:
 - True
 - False
- Operation
 - Conjunction (AND)
 - $A \cdot B$; $A \wedge B$
 - Disjunction (OR)
 - $A + B$; $A \vee B$
 - Negation (NOT)
 - \bar{A} ; $\neg A$; A'

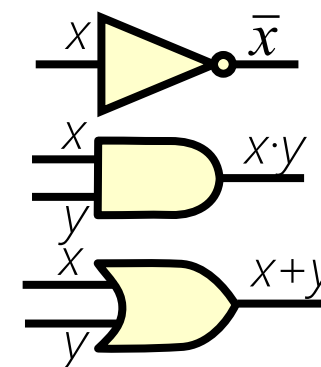
- Truth tables

A	B	$A \cdot B$
0	0	0
0	1	0
1	0	0
1	1	1

A	B	$A + B$
0	0	0
0	1	1
1	0	1
1	1	1

A	A'
0	1
1	0

- Logic gates



Boolean Functions

- Let $B = \{0, 1\}$, the set of Boolean values.
- For all $n \in \mathbb{Z}^+$, any function $f: B^n \rightarrow B$ is called a *Boolean function of degree n* .
- There are 2^{2^n} (wow!) distinct Boolean functions of degree n .
 - B/c $\exists 2^n$ rows in truth table, w. 0 or 1 in each.

<u>Degree</u>	<u>How many</u>	<u>Degree</u>	<u>How many</u>
0	2	4	65,536
1	4	5	4,294,967,296
2	16	6	18,446,744,073,709,551,616.
3	256		

Boolean Expressions

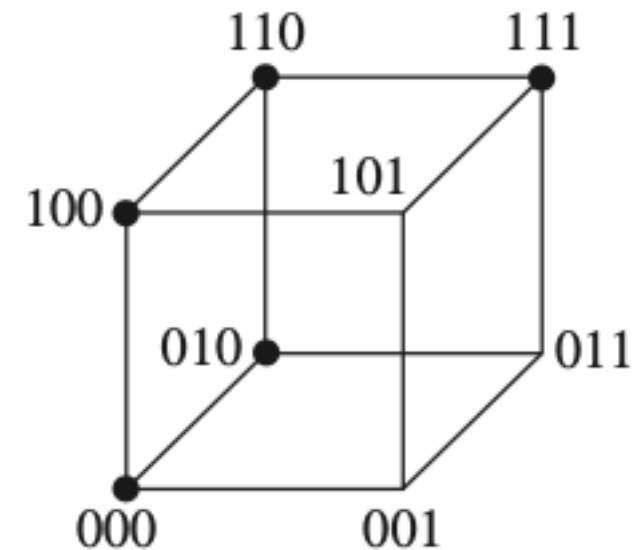
- Let x_1, \dots, x_n be n different Boolean variables.
 - n may be as large as desired.
- A *Boolean expression* (recursive definition) is a string of one of the following forms:
 - Base cases: 0 , 1 , x_1 , ..., or x_n .
 - Recursive cases: $\overline{E_1}$, $(E_1 E_2)$, or $(E_1 + E_2)$, where E_1 and E_2 are Boolean expressions.
- A Boolean expression represents a Boolean function.
 - Furthermore, *every* Boolean function (of a given degree) can be represented by a Boolean expression.

Boolean equivalents

- Two Boolean expressions e_1 and e_2 that represent the exact *same* function f are called *equivalent*. We write $e_1 \Leftrightarrow e_2$, or just $e_1 = e_2$.
 - Implicitly, the two expressions have the same value for *all* values of the free variables appearing in e_1 and e_2 .
 - E.g., Boolean expressions xy , $xy+0$, $xy \cdot 1$

Hypercube Representation

- A Boolean function of degree n can be represented by an n -cube (hypercube) with the corresponding function value at each vertex.





Boolean Identities

TABLE 5 Boolean Identities.

<i>Identity</i>	<i>Name</i>
$\overline{\overline{x}} = x$	Law of the double complement
$x + x = x$ $x \cdot x = x$	Idempotent laws
$x + 0 = x$ $x \cdot 1 = x$	Identity laws
$x + 1 = 1$ $x \cdot 0 = 0$	Domination laws
$x + y = y + x$ $xy = yx$	Commutative laws

$x + (y + z) = (x + y) + z$ $x(yz) = (xy)z$	Associative laws
$x + yz = (x + y)(x + z)$ $x(y + z) = xy + xz$	Distributive laws
$\overline{(xy)} = \overline{x} + \overline{y}$ $\overline{(x + y)} = \overline{x} \overline{y}$	De Morgan's laws
$x + xy = x$ $x(x + y) = x$	Absorption laws
$x + \overline{x} = 1$	Unit property
$x\overline{x} = 0$	Zero property



Example 10

- Prove the absorption law $x(x+y) = x$ using other identities of Boolean algebra.


$x(x + y) = (x + 0)(x + y)$	Identity law for the Boolean sum
$= x + 0 \cdot y$	Distributive law of the Boolean sum over the Boolean product
$= x + y \cdot 0$	Commutative law for the Boolean product
$= x + 0$	Domination law for the Boolean product
$= x$	Identity law for the Boolean sum.



Duality

- The *dual* e^d of a Boolean expression e representing function f is obtained by exchanging $+$ with \cdot , and 0 with 1 in e .
 - The function represented by e^d is denoted f^d .

Find the duals of $x(y + 0)$ and $\bar{x} \cdot 1 + (\bar{y} + z)$.

Solution: Interchanging \cdot signs and $+$ signs and interchanging 0s and 1s in these expressions produces their duals. The duals are $x + (y \cdot 1)$ and $(\bar{x} + 0)(\bar{y}z)$, respectively. 



The Abstract Definition of a Boolean Algebra

A *Boolean algebra* is a set B with two binary operations \vee and \wedge , elements 0 and 1, and a unary operation $\bar{}$ such that these properties hold for all x , y , and z in B :

$$\left. \begin{array}{l} x \vee 0 = x \\ x \wedge 1 = x \end{array} \right\} \quad \text{Identity laws}$$

$$\left. \begin{array}{l} x \vee \bar{x} = 1 \\ x \wedge \bar{x} = 0 \end{array} \right\} \quad \text{Complement laws}$$

$$\left. \begin{array}{l} (x \vee y) \vee z = x \vee (y \vee z) \\ (x \wedge y) \wedge z = x \wedge (y \wedge z) \end{array} \right\} \quad \text{Associative laws}$$

$$\left. \begin{array}{l} x \vee y = y \vee x \\ x \wedge y = y \wedge x \end{array} \right\} \quad \text{Commutative laws}$$

$$\left. \begin{array}{l} x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z) \\ x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z) \end{array} \right\} \quad \text{Distributive laws}$$



Tentukan ekspresi Boolean dari tabel kebenaran berikut ini dan sederhanakan!

x	y	z	$F(x, y, z)$
1	1	1	1
1	1	0	1
1	0	1	0
1	0	0	1
0	1	1	0
0	1	0	1
0	0	1	0
0	0	0	1

$$F(x, y, z) = xyz + xy\bar{z} + x\bar{y}\bar{z} + \bar{x}y\bar{z} + \bar{x}\bar{y}\bar{z}$$

Dengan identitas diatas fungsi tersebut dapat kita sederhanakan menjadi

$$= xyz + \bar{z}(xy + x\bar{y} + \bar{x}y + \bar{x}\bar{y})$$

$$= xyz + \bar{z}(x + \bar{x})(y + \bar{y})$$

$$= xyz + \bar{z}$$

Dengan *absorption law*

$$= xyz + (\bar{z} + \bar{z}xy)$$

$$= xyz + xy\bar{z} + \bar{z}$$

$$= xy(z + \bar{z}) + \bar{z}$$

$$= xy + \bar{z}$$



Tentukan ekspresi Boolean dari tabel kebenaran berikut ini dan sederhanakan!

x	y	z	$F(x,y,z)$
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1



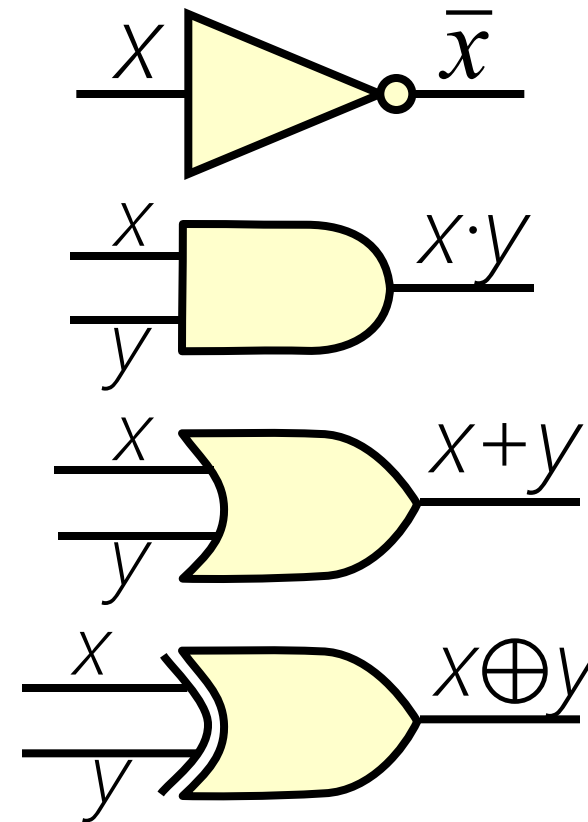
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Logic Gates



Logic Gate Symbols

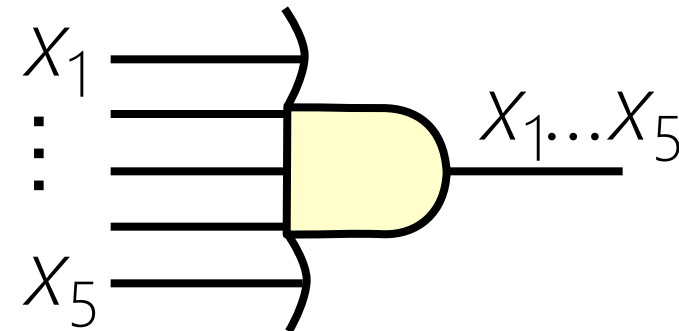
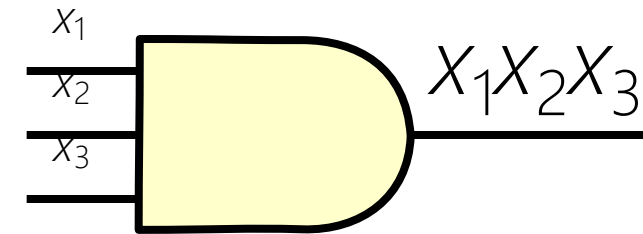
- Inverter (logical NOT, Boolean complement).
- AND gate (Boolean product).
- OR gate (Boolean sum).
- XOR gate (exclusive-OR, sum mod 2).





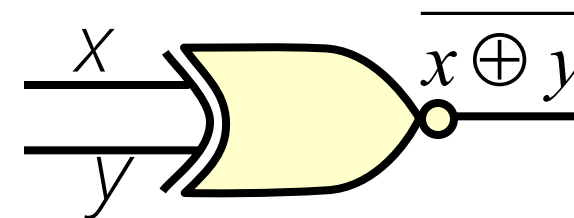
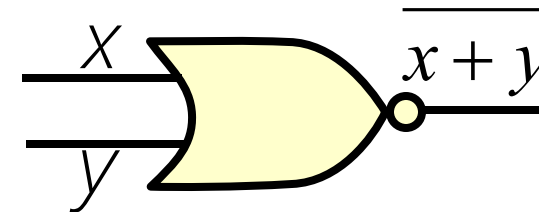
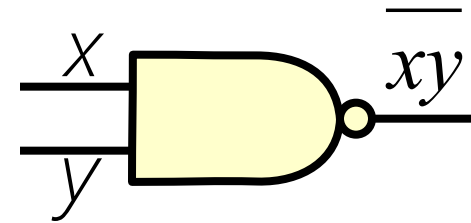
Multi-input AND, OR, XOR

- Can extend these gates to arbitrarily many inputs.
- Two commonly seen drawing styles:
 - Note that the second style keeps the gate icon relatively small.



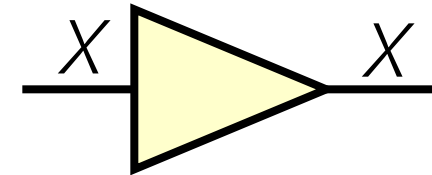
NAND, NOR, XNOR

- Just like the earlier icons, but with a small circle on the gate's output.
 - Denotes that output is complemented.
- The circles can also be placed on inputs.
 - Means, input is complemented before being used.



Buffer

- What about an inverter symbol *without* a circle?
- This is called a *buffer*. It is the identity function.
- It serves no logical purpose, but...
- It represents an explicit delay in the circuit.
 - This is sometimes useful for timing purposes.
- All gates, when physically implemented, incur a non-zero delay between when their inputs are seen and when their outputs are ready.





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- Gates: The basic elements of electronic circuits
- Combinational circuits: the circuits whose output depends only on the input and not on the current state of the circuit (no memory).



Karnaugh Maps

- K-map is a graphical method for finding terms to combine for Boolean functions involving a relatively small number of variables.

$$wx\bar{y}\bar{z} + w\bar{x}yz + w\bar{x}y\bar{z} + w\bar{x}\bar{y}z + \bar{w}x\bar{y}\bar{z} + \bar{w}x\bar{y}z + \bar{w}x\bar{y}\bar{z}$$

$$x\bar{y} + \bar{x}y + \bar{x}\bar{y}$$

	y	\bar{y}
x		1
\bar{x}	1	1

$$\bar{x} + \bar{y}$$

$$xy\bar{z} + x\bar{y}\bar{z} + \bar{x}yz + \bar{x}\bar{y}\bar{z}$$

	yz	$y\bar{z}$	$\bar{y}\bar{z}$	$\bar{y}z$
x		1	1	
\bar{x}	1		1	

	yz	$y\bar{z}$	$\bar{y}\bar{z}$	$\bar{y}z$
wx			1	
$w\bar{x}$	1	1	1	
$\bar{w}\bar{x}$		1	1	
$\bar{w}x$			1	

Don't Care Conditions

- When we care only about the output for some combinations of input values, because other combinations of input values are not possible or never occur.
 - Those combinations that never occur can be arbitrarily chosen.