Software Synchronization

Race Conditions

- A race condition is a situation where several processes or threads attempt to manipulate the same data, and the outcome of the execution depends on the particular order the access takes place.
- The section of code where a process or thread accesses a shared resource is referred to as the critical section.
- To prevent race conditions, we use synchronization. This ensures that no more than one process is able to execute it's critical section at a time.

Syncronization Solution Requirements

- Any **solution** to **synchronization** should satisfy the following requirements:
 - 1. The solution should have **mutual exclusion (mutex)**: Only **one process** can execute it's **critical section** at any given time.
 - 2. The solution should have **progress**: When **no process** is currently in a **critical section**, **any process** that **requests entry** into the **critical section** must be **permitted without delay**.
 - 3. The solution must prevent starvation (bounded wait): There is an upper bound on the number of times a process enters the critical section while another is waiting.
 - Such a solution will prevent race conditions.

Synchronization Solutions

- A simple way to achieve synchronization is the use interrupts; when interrupts are disabled, context switches will not happen.
 - This is not a good solution, because user processes generally cannot disable interrupts.
 - This also does not work on a multi-core system.
- Another way to achieve synchronization is by using a variable as a flag to determine if any other processes are executing their critical section. When a process wants to execute their critical section and another process has already locked the resource, the process will use a conditional loop until the lock is removed.
 - This acheives mutual exclusion, but wastes a lot of processor time.
 - This also does not prevent starvation.
- Modern computers use hardware solutions to synchronize processes.

Hardware Syncronization