All You Ever Wanted to Know About Dynamic Taint Analysis and Forward Symbolic Execution

(but might have been afraid to ask)

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Outline



Introduction

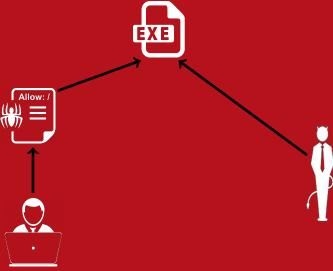
Dynamic Taint Analysis

Forward Symbolic Execution

Conclusions



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Input Analysis



There are two essential questions about the input analysis:

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Input Analysis



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- 1. Is the final value affected by user input?
 - Dynamic Taint Analysis!
 - Tracks information flow between sources and sinks

Input Analysis



There are two essential questions about the input analysis:

- 1. Is the final value affected by user input?
 - Dynamic Taint Analysis!
 - Tracks information flow between sources and sinks
- 2. What input will make execution reach this line of code?
 - Forward Symbolic Execution
 - Allows us to reason about the behavior of a program on many different inputs



goto y



```
Tainted
```

$$z := 42$$



```
Untainted = get_input()

Z := 42

y := x + z  Z is a "static"

goto y
```



$$x := get_input(x)$$
 $z := 42$
 $y := x + z$
 $goto y$
Is y taited?



```
x := get_input()
z := 42
                Js y taited?
goto v
            It depends on the
             selected policy
```

What's a policy?

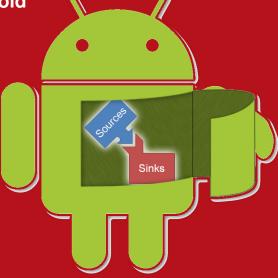


- ► A taint policy specifies three properties:
 - Taint Introduction
 - ▶ How is taint introduced into a system?
 - Taint Propagation
 - ▶ How does taint propagate into a system?
 - Taint Checking
 - Is the current operation secure?
- Undertainting vs Overtainting



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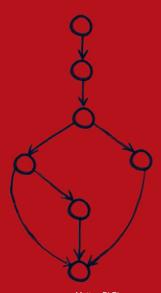




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Limitations

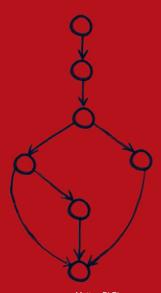




- Undertainting and overtainting are nearly unavoidable!
 - Time of detection vs Time of attack
- **▶** Sanitization problem
- Pure dynamic taint analysis considers data flows...

Limitations





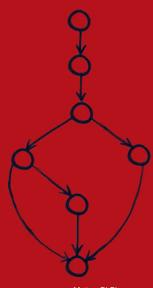
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- **▶** Sanitization problem

 $b := a \oplus a$

Pure dynamic taint analysis considers data flows...

Limitations





- Undertainting and overtainting are nearly unavoidable!
 - Time of detection vs Time of attack
- ► Sanitization problem
- Pure dynamic taint analysis considers data flows...
- ...but it ignores control-flows
 - What about different security policies for different I/O channels?
 - \rightarrow Static analysis

Forward Symbolic Execution



- We can reason about the behavior of a program using the logic...
- ... and it is conceptually a very simple process

```
x := 2 * get_input(src)
if x - 5 == 14 then goto 3 else goto 4
// line 3: catastrophic failure
// line 4: normal behaviour
```

Forward Symbolic Execution



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- ... and it is conceptually a very simple process

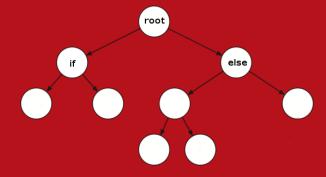
```
x := 2 * get_input(src)
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- get_input(src) now returns a symbol instead of a concrete value
- But now expressions cannot be fully evaluated to a concrete value

Path Selection and Performance



- Every conditional jump we must decide what path to follow first
 - But some path may never terminate
- Exponential blowup due to branches



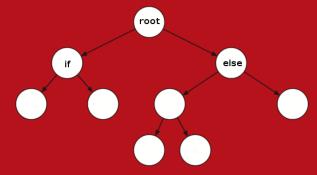
Path Selection and Performance



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while
$$(3^n + 4^n == 5^n) \{n++; \ldots\}$$

► Exponential blowup due to branches



Path Selection and Performance



- Every conditional jump we must decide what path to follow first
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while
$$(3^n + 4^n == 5^n) \{n++; \ldots\}$$

- Exponential blowup due to branches
- Solutions
 - Path Selection Heuristic
 - Concolic Testing
 - Depth-First or Random Search
 - More and faster hardware
 - Identify redundancies between formulas
 - Identify independent subformulas

Memory Address Problems



- What are we supposed to do if a referenced address is derived from user input?
 - LOAD, STORE \rightarrow Symbolic Memory Address
 - GOTO → Symbolic Jumps
- Solutions
 - Concolic testing
 - SMT (Satisfiability Modulo Theories) solvers
 - ▶ NP-Complete problem!
 - Static and alias analysis

A small comparison



Dynamic Taint Analysis

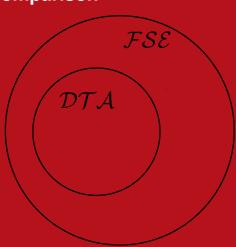
?

Forward Symbolic Execution

A small comparison



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Dynamic taint analysis analyzes only feasible paths

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Conclusions



- v Conceptually simple methods of analysis
- v There are a lot of possible use cases
 - Malware detection and analysis
 - Automatic testing
 - Automatic programs understanding

- x Usable with some care
- x The effectiveness depends on the application



Thank you for allowing me to taint your time!

Questions?



SimplL Grammar



```
stmt*
program
stmt \ s ::= var := exp \mid store(exp, exp)
                  goto exp assert exp
                  if exp then goto exp
                    else goto exp
            ::= load(exp) \mid exp \mid \Diamond_b exp \mid \Diamond_u exp
exp e
                  | var | get_input(src) | v
\Diamond_b
            ::= typical binary operators
\Diamond_u
            ::= typical unary operators
value v ::= 32-bit unsigned integer
```