



CIS 263 Introduction to Data Structures and Algorithms

Dynamic Programming



Dynamic Programming

Recursion:

```
int fib(int n)
{
    if (n <= 1)
        return n;
    return fib(n-1) + fib(n-2);
}
```

Exponential time



Dynamic Programming

Recursion:

```
int fib(int n)
{
    if (n <= 1)
        return n;
    return fib(n-1) + fib(n-2);
}
```

Exponential time

Dynamic Programming :

```
f[0] = 0;
f[1] = 1;

for (i = 2; i <= n; i++)
{
    f[i] = f[i-1] + f[i-2];
}

return f[n];
```

Linear time



Recursion

-1	-1	-1	-1	-1	-1
0	1	2	3	4	5

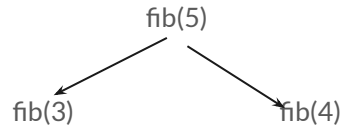
fib(5)

$$\text{fib}(n) = \begin{cases} 0 & \text{if } n = 0 \\ 1 & \text{if } n = 1 \\ \text{fib}(n-2) + \text{fib}(n-1) & \text{if } n > 2 \end{cases}$$

```
def fib(n):  
    if n <= 1:  
        return n  
    else:  
        return fib(n-2) + fib(n-1)
```

Recursion

-1	-1	-1	-1	-1	-1
0	1	2	3	4	5

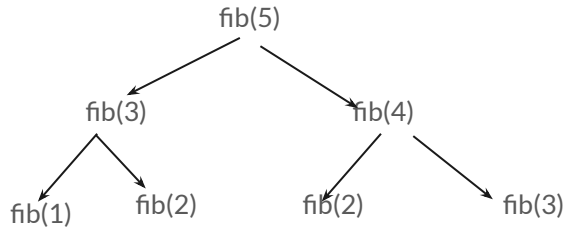


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Recursion

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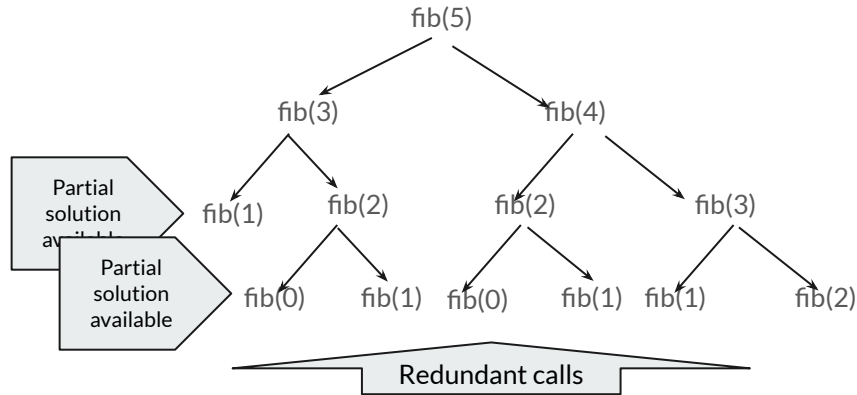
Partial
solution
available

$$\text{fib}(n) = \begin{cases} 0 & \text{if } n = 0 \\ 1 & \text{if } n = 1 \\ \text{fib}(n-2) + \text{fib}(n-1) & \text{if } n > 2 \end{cases}$$

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def fib(n):  
    if n <= 1:  
        return n  
    else:  
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Recursion

-1	-1	-1	-1	-1	-1
0	1	2	3	4	5

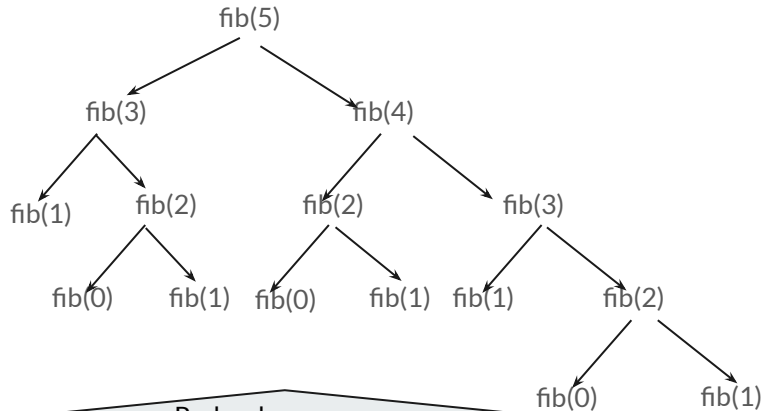


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def fib(n):  
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Recursion

-1	-1	-1	-1	-1	-1
0	1	2	3	4	5



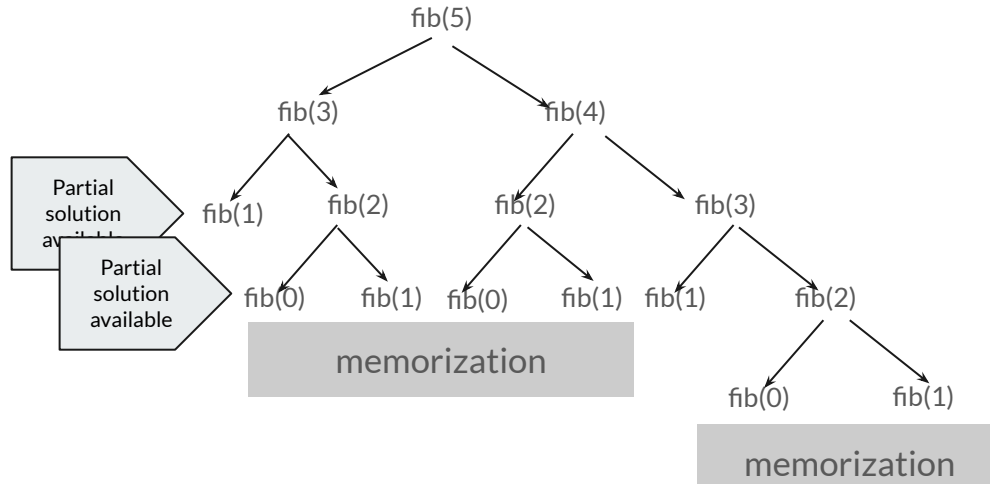
Redundancy grows
exponentially

$$\text{fib}(n) = \begin{cases} 0 & \text{if } n = 0 \\ 1 & \text{if } n = 1 \\ \text{fib}(n-2) + \text{fib}(n-1) & \text{if } n > 2 \end{cases}$$

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def fib(n):  
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Recursion

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def fib(n):  
    if n <= 1:  
        return n  
    else:  
        return fib(n-2) + fib(n-1)
```



Dynamic programming

-1	-1	-1	-1	-1	-1
0	1	2	3	4	5

```
def fib(n):
```

```
    f = [0]*n
```

```
    f[0] = 0
```

```
    f[1] = 1
```

```
    for i in range(2, n):
```

```
        f[i] = f[i-1] + f[i-2]
```

```
    return f
```



Longest Common Subsequence

```
str1 = "ABCDE"  
str2 = "ACE"
```



Longest Common Subsequence

str1 = "ABCDE"
str2 = "ACE"

A	B	C	D	E
---	---	---	---	---

A
C
E



Longest Common Subsequence

str1 = "ABCDE"
str2 = "ACE"

	0	A	B	C	D	E
0						
A						
C						
E						



Longest Common Subsequence

str1 = "ABCDE"
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	0	A	B	C	D	E
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Longest Common Subsequence

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	0	A	B	C	D	E
0	0	0	0	0	0	0
A						
C						
E						



Longest Common Subsequence

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	0	A	B	C	D	E
0	0	0	0	0	0	0
A	0					
C	0					
E	0					

Longest Common Subsequence

str1 = "ABCDE"
str2 = "ACE"

	0	A	B	C	D	E
0	0	0	0	0	0	0
A	0					
C	0					
E	0					

match

```
If str1[i-1] == str2[j-1]:  
    lcd[i][j] = 1 + lcd[i-1][j-1]
```

no-match

```
else:  
    lcd[i][j] = max(  
        lcd[i-1][j],  
        lcd[i][j-1]  
    )
```

Longest Common Subsequence

str1 = "ABCDE"
str2 = "ACE"

	0	A	B	C	D	E
0	0	0	0	0	0	0
A	0	1				
C	0					
E	0					

match

```
If str1[i-1] == str2[j-1]:  
    lcd[i][j] = 1 + lcd[i-1][j-1]
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Longest Common Subsequence

str1 = "ABCDE"
str2 = "ACE"

	0	A	B	C	D	E
0	0	0	0	0	0	0
A	0	1	1	1	1	1
C	0					
E	0					

match

```
If str1[i-1] == str2[j-1]:  
    lcd[i][j] = 1 + lcd[i-1][j-1]
```

no-match

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else:  
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Longest Common Subsequence

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str2 = "ACE"

	0	A	B	C	D	E
0	0	0	0	0	0	0
A	0	1	1	1	1	1
C	0	1	1	2	2	2
E	0	1	1	2	2	3

match

```
If str1[i-1] == str2[j-1]:  
    lcd[i][j] = 1 + lcd[i-1][j-1]
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else:  
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        lcd[i-1][j],  
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Longest Common Subsequence

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E	0	1	1	2	2	3

Recurrence relations

Longest Common Subsequence

str1 = "ABCDE"
str2 = "ACE"

	0	A	B	C	D	E
0	0	0	0	0	0	0
A	0	1	1	1	1	1
C	0	1	1	2	2	2
E	0	1	1	2	2	3

Recurrence relations



Longest Common Subsequence

str1 = "ABCDE"
str2 = "ACE"

	0	A	B	C	D	E
0	0	0	0	0	0	0
A	0	1	1	1	1	1
C	0	1	1	2	2	2
E	0	1	1	2	2	3

Backtracking

Longest Common Subsequence

str1 = "ABCDE"
str2 = "ACE"

	0	A	B	C	D	E
0	0	0	0	0	0	0
A	0	1	1	1	1	1
C	0	1	1	2	2	2
E	0	1	1	2	2	3

E

Backtracking

Longest Common Subsequence

str1 = "ABCDE"
str2 = "ACE"

	0	A	B	C	D	E
0	0	0	0	0	0	0
A	0	1	1	1	1	1
C	0	1	1	2	2	2
E	0	1	1	2	2	3

A

C

E

Backtracking



Minimum Edit Distance

str1 = "ABCDE"

str2 = "ACE"



Minimum Edit Distance

str1 = "ABCDE"
str2 = "ACE"

	0	A	B	C	D	E
0						
A						
C						
E						



What should be the initial values?



Minimum Edit Distance

str1 = "ABCDE"
str2 = "ACE"

	0	A	B	C	D	E
0	0	1	2	3	4	5
A	1					
C	2					
E	3					



Minimum Edit Distance

str1 = "ABCDE"
str2 = "ACE"

	0	A	B	C	D	E
0	0	1	2	3	4	5
A	1	0	1	2	3	4
C	2	1	1	1	2	3
E	3	2	2	2	2	2

```
If str1[i-1] == str2[j-1]:  
    med[i][j] = med[i-1][j-1]
```

```
else:  
    med[i][j] = 1 + min (  
        med[i-1][j],  
        med[i][j-1]  
        med[i-1][j-1]  
    )
```

Minimum Edit Distance

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str2 = "ACE"

	0	A	B	C	D	E
0	0	1	2	3	4	5
A	1	0	1	2	3	4
C	2	1	1	1	2	3
E	3	2	2	2	2	2

AB → ACE

A == A

add "C": ACB

replace "B" with E: AC~~B~~E

replace "B" with "C": A~~B~~C

add "E": A~~B~~CE



Other applications

- Matrix chain multiplication
- Sequence alignment
- Tower of Hanoi puzzle
- Markov decision process, Viterbi Algorithm
- Dynamic Time warping
- Floyd's shortest path algorithm