CIS 263 Introduction to Data Structures and Algorithms

Graph Algorithms

Linear Data Structures

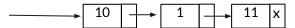
Array



Linear Data Structures

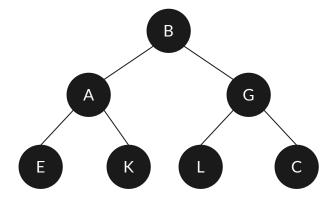
- Array
- Linked List





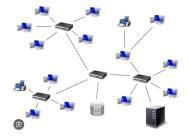
Hierarchical relationship

- Tree
 - Parent child relationship



Graphs

- Pairwise relationship
- Examples
 - Computer networks



Graphs

- Pairwise relationship
- Examples
 - Computer networks
 - Social networks

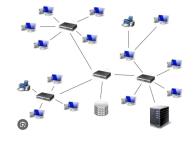


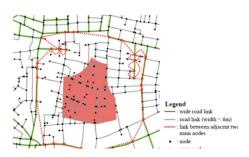


Graphs

- Pairwise relationship
- Examples
 - Computer networks
 - Social networks
 - Road networks

0

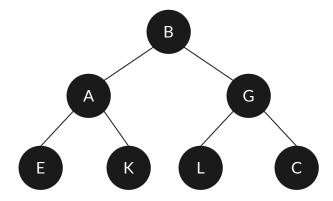






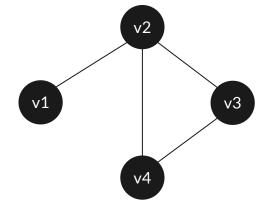
A Tree is a special instance of Graph

- Tree
 - Parent child relationship



Graph (Formal Definition)

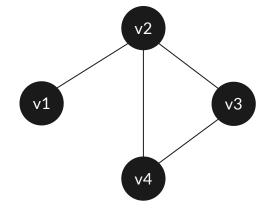
- Pairwise relationship
- A Graph, G, is an ordered pair of vertices, V, and edges, E.



Graph (Formal Definition)

- Pairwise relationship
- A Graph, G, is an ordered pair of vertices, V, and edges, E.

Undirected Graph

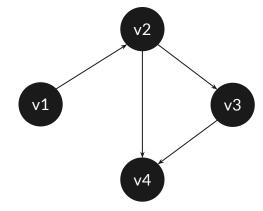


Directed Graph

- Pairwise relationship
- A Graph, G, is an ordered pair of vertices, V, and edges, E.

$$V = \{v1, v2, v3, v4\}$$

$$E = \{(v1,v2), (\{v2,v3\}, (v3,v4), (v2,v4)\}$$



Graph (Data Structure) - Undirected

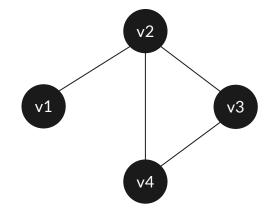
G = (V, E)

$$V = \{v1, v2, v3, v4\}$$

$$E = \{\{v1,v2\}, \{v2,v3\}, \{v3,v4\}, \{v2,v4\}\}$$

Using pointer

- non-static node definitions



Graph (Data Structure) - Undirected

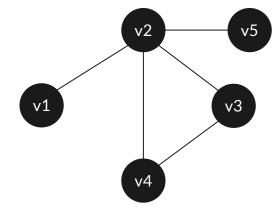
G = (V, E)

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Using pointer

- non-static node definitions



Representation of an undirected graph

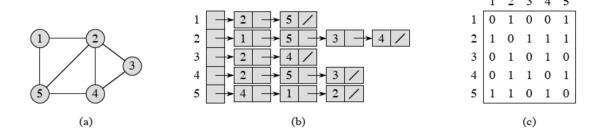


Figure 22.1 Two representations of an undirected graph. (a) An undirected graph G with 5 vertices and 7 edges. (b) An adjacency-list representation of G. (c) The adjacency-matrix representation of G.

Ref: Introduction to Algorithms, by Thomas Cormen

Representation of an undirected graph

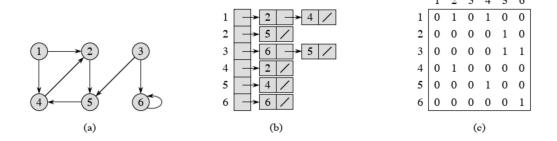
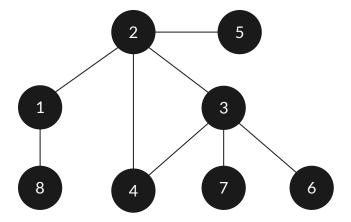


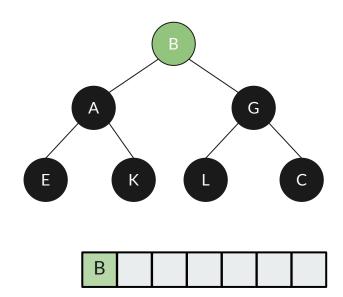
Figure 22.2 Two representations of a directed graph. (a) A directed graph G with 6 vertices and 8 edges. (b) An adjacency-list representation of G. (c) The adjacency-matrix representation of G.

Ref: Introduction to Algorithms, by Thomas Cormen

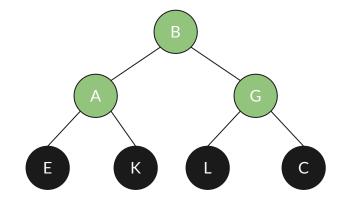
- Breadth First (level order)
- Depth First (pre order in case of BT)



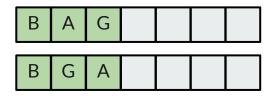
- Breadth First (level order)
- 1. Start with any node, say (B)
- 2. Traverse one level at a time

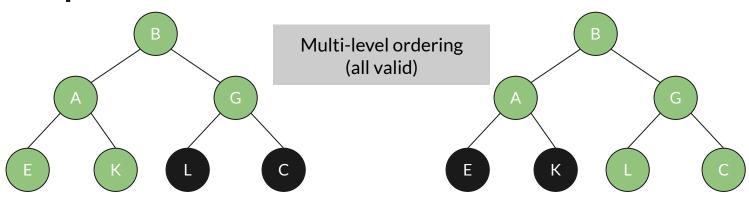


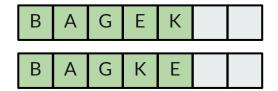
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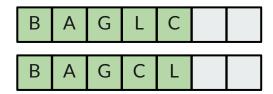


Both valid; we will track the first instance

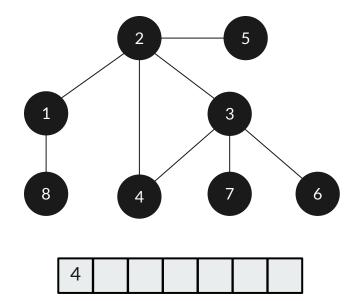




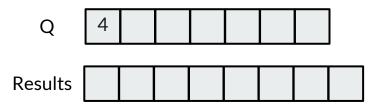


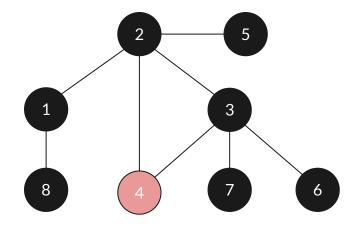


- Breadth First (level order)
- 1. Start with any node, say (4)
- 2. Traverse one level at a time
- 3. To generalize, we will use an additional Queue

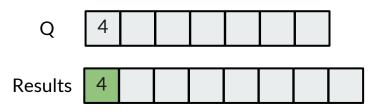


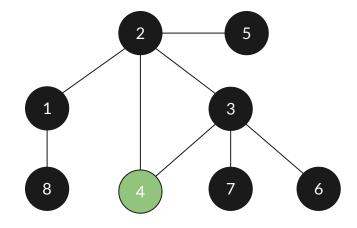
- Breadth First
- 1. Start with any node, say (4)
- 2. Traverse one level at a time
- To generalize, we will use an additional Queue,
 Q
 - a. Append data/key in Q



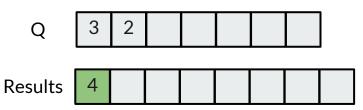


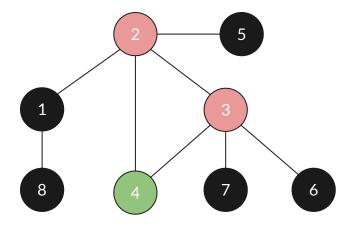
- Breadth First
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 - b. Process data/key



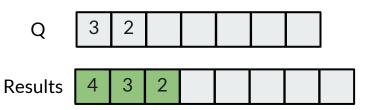


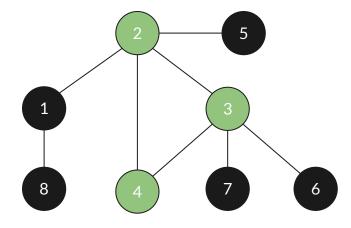
- Breadth First
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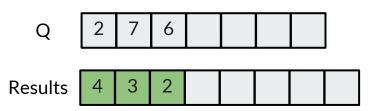


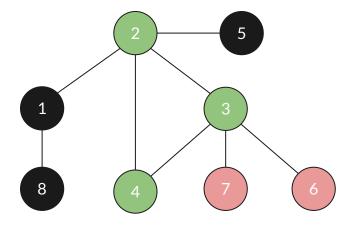
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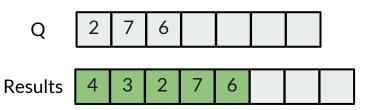


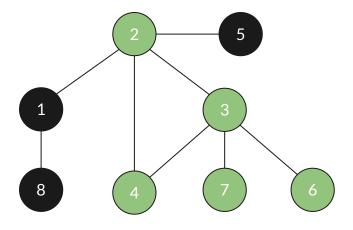
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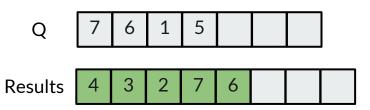


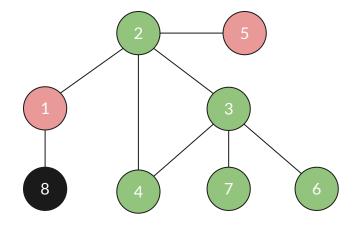
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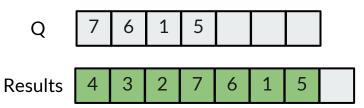


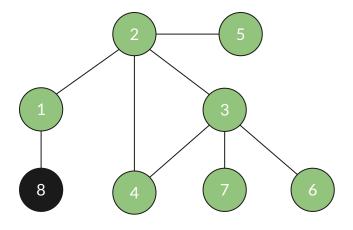
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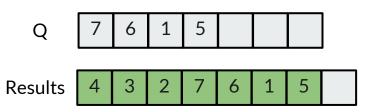


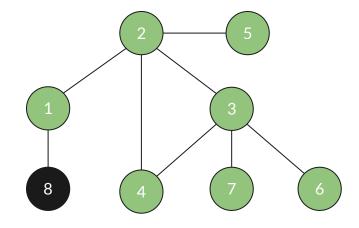
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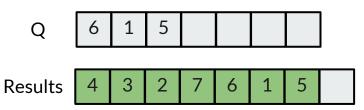
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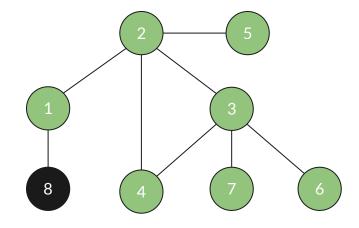




7's neighbor 3 is already visited ...

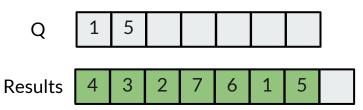
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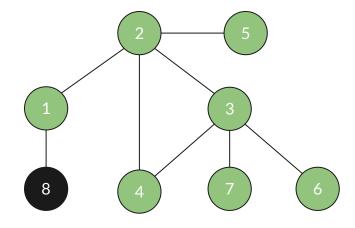




6's neighbor 3 is already visited ...

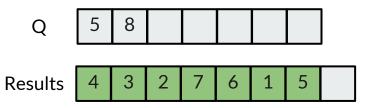
- Breadth First
- 1. Start with any node, say (4)
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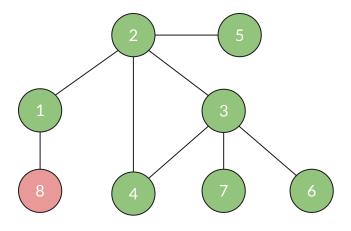




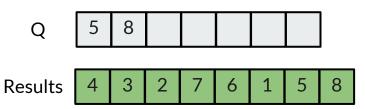
1's neighbor 2 is already visited; 8 not

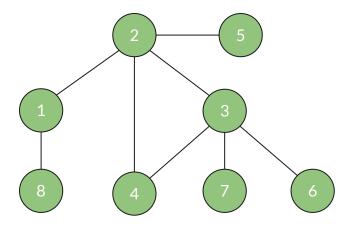
- Breadth First
- 1. Start with any node, say (4)
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- To generalize, we will use an additional Queue,
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 - b. Process data/key



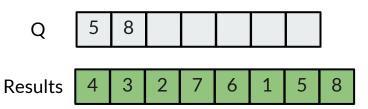


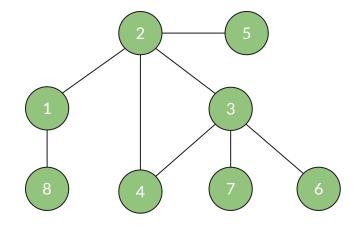
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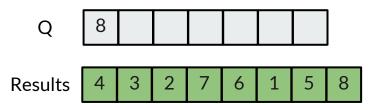
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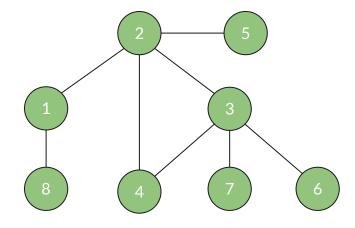




5's neighbor 2 is already visited ...

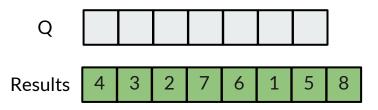
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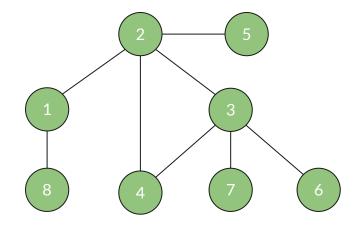




8's neighbor 1 is already visited ...

- Breadth First
- 1. Start with any node, say (4)
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 - b. Process data/key





Q is empty; we finish



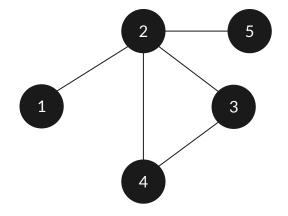
Graph (Data Structure) - Undirected

G = (V, E)

$$V = \{v1, v2, v3, v4\}$$

 $E = \{\{v1,v2\}, \{v2,v3\}, \{v3,v4\}, \{v2,v4\}\}$

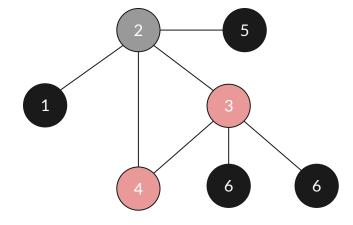
		1	2	3	4	5
	1		1			
	2			1	1	1
	3				1	
	4					
	4					



Using 2D Array

- n x n Matrix
- Sparse Matrix

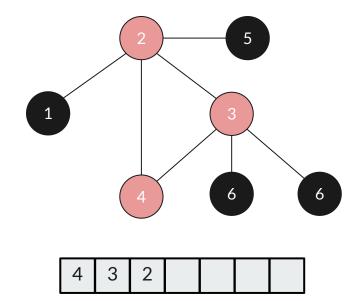
- Breadth First
- 1. Start with any node, say (4)
- (a) Perform node operation, say (print (key)), if not already done
 - (b) Perform Exploration (1-st level neighbors), any order
 - (c) Iterate until all node visited



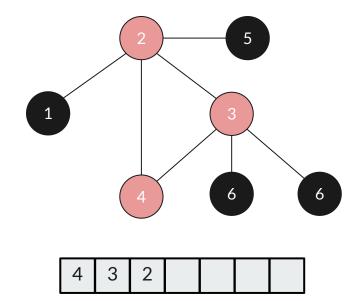
Results 4 3

Q 4 3

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QA