Module 2-6

Database Design

Objectives

- Database Design Exercise
- Database Definition Language
- Database Control Language

Constraints - Review

Constraints are rules imposed on the table, upon creation, that limits the ability to change the data.

- NOT NULL: A value must be specified
- PRIMARY KEY: Define that certain column/columns are part of the key
 - A primary key value cannot be NULL.
- FOREIGN KEY: Defines a foreign key based on a primary key from a different table
- CHECK: Only certain values can be inserted or updated
- e.g. CONSTRAINT CHK_duration CHECK (duration >= 20)

Transactions

A large number of SQL statements can be rolled into a single transaction.

The following syntax is observed:

START TRANSACTION; -- or BEGIN TRANSACTION;

// Lots of SQL statements.

COMMIT TRANSACTION; -- or COMMIT;

Your INSERT or UPDATE SQL statements will only commit (permanently save in the database) if all the SQL statements in the transaction end successfully.

Transaction Syntax

START TRANSACTION

Do the UPDATE/INSERT/DELETE statements

COMMIT (ends the transaction and saves the changes)

OR

ROLLBACK (ends the transaction without saving the changes)

Transactions can be used to safely test a statement that changes the database during development/testing.

Database Design Exercise

Gallery Customer History Form Customer Name Jackson, Elizabeth Phone (206) 284-6783 123 - 4th Avenue Fonthill, ON L3J 4S4 Purchases Made Title Purchase Date Sales Price Artist 03 - Carol Channing Laugh with Teeth 09/17/2000 7000.00 South toward Emerald Sea 15 - Dennis Frings 05/11/2000 1800.00 03 - Carol Channing At the Movies 02/14/2002 5550.00 15 - Dennis Frings South toward Emerald Sea 07/15/2003 2200.00

The Gill Art Gallery wishes to maintain data on their customers, artists and paintings. They may have several paintings by each artist in the gallery at one time. Paintings may be bought and sold several times. In other words, the gallery may sell a painting, then buy it back at a later date and sell it to another customer.

Normal Forms

Before a single CREATE statement is run, the tables and their relationships need to be well thought out.

Normalization is the process of organizing a database to reduce data redundancy and improve data integrity.

We normalize data to:

- 1. Avoid duplicate data
- 2. Fix anomalies
- 3. Simplify search queries.

Normal Forms: Before normalization

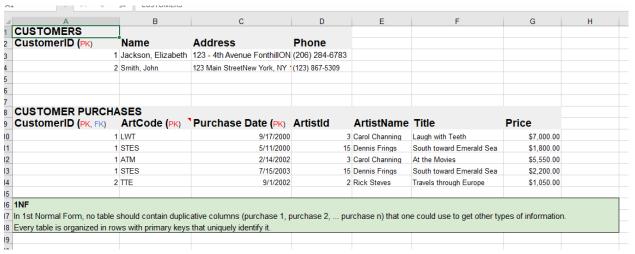
| Gallery Customer H | listory Form | | |
|-------------------------------------------------------------------------|--------------------------|---------------|-------------|
| Customer Nar | me | | |
| Jackson, Eliza 123 – 4 th Aver Fonthill, ON L3J 4S4 | ' ' | 284-6783 | |
| Purchases Made | | | |
| Artist | Title | Purchase Date | Sales Price |
| 03 - Carol Channing | Laugh with Teeth | 09/17/2000 | 7000.00 |
| 15 - Dennis Frings | South toward Emerald Sea | 05/11/2000 | 1800.00 |
| 03 - Carol Channing | At the Movies | 02/14/2002 | 5550.00 |
| | South toward Emerald Sea | 07/15/2003 | 2200.00 |

The Gill Art Gallery wishes to maintain data on their customers, artists and paintings. They may have several paintings by each artist in the gallery at one time. Paintings may be bought and sold several times. In other words, the gallery may sell a painting, then buy it back at a later date and sell it to another customer.

| 1 | А | В | С | D | Е |
|---|--------------------|----------------------------|------------------|--------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|---|
| 1 | Name | Address | Phone | Purchases(ArtistId, ArtistName, ArtTitle, PurchaseDate, Price) | |
| 2 | Jackson, Elizabeth | 123 - 4th Avenue FonthillO | 1 (206) 284-6783 | 03, Carol Channing, Laugh with Teeth, 9/17/2000, \$7000.00 15, Dennis Frings, South toward Emerald Sea, 5/11/2000, \$1800.00 03, Carol Channing, At the Movies, 2/14/2002, \$5550.00 15, Dennis Frings, South toward Emerald Sea, 7/15/2003, \$2200.00 | |
| 3 | Smith, John | 123 Main StreetNew York, | (123) 867-5309 | 02, Rick Steves, Travels through Europe, 9/01/2002, \$1050.00 | |
| 4 | | | | | |
| 5 | | | | | |
| 6 | | | | | |
| 7 | | | | | |
| 8 | | | | | |

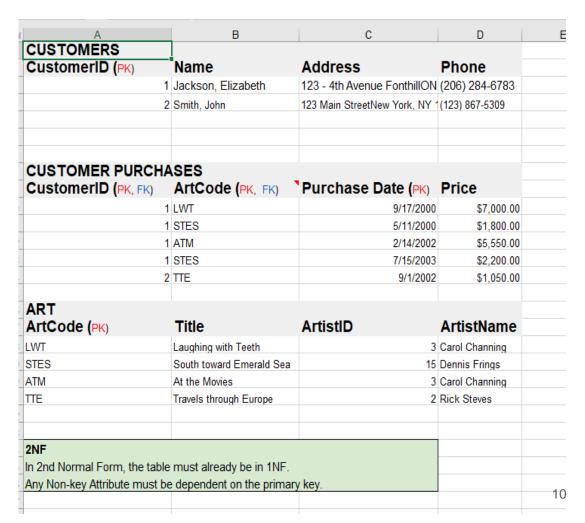
Normal Forms: 1NF

- No table should contain duplicative columns that one could use to get other types of information.
- Every table should be organized in rows with primary keys that uniquely identify it.



Normal Forms: 2NF

- Must be in 1NF
- Any non-key attribute must be dependent on the primary key



2nd Normal Form – example2

| ¶ Supplier# | 🕴 Part# | City | Quantity |
|-------------|---------|------------|----------|
| S1 | P1 | London | 1000 |
| S1 | P2 | London | 1500 |
| S1 | P3 | London | 3400 |
| S 1 | P4 | London | 2100 |
| S2 | P2 | Paris | 3400 |
| S2 | P3 | Paris | 1000 |
| S4 | P1 | Nuku alofa | 5 |
| S4 | P4 | Nuku alofa | 7 |

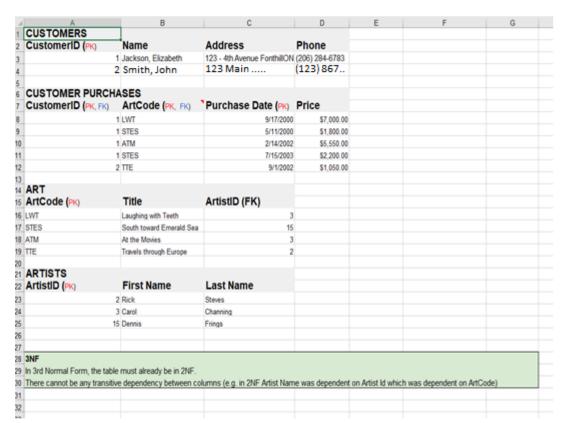
Table in Second Normal Form

| Supplier# | City |
|------------|------------|
| S 1 | London |
| S2 | Paris |
| S4 | Nuku alofa |

| Supplier# | Part# | Quantity |
|------------|-------|----------|
| S 1 | P1 | 1000 |
| S1 | P2 | 1500 |
| S1 | P3 | 3400 |
| S1 | P4 | 2100 |
| S2 | P2 | 3400 |
| S2 | P3 | 1000 |
| S4 | P1 | 5 |
| S4 | P4 | 7 |

Normal Forms: 3NF

- Must be in 2NF
- No transitive functional dependency – if column A is dependent on column B and column B is dependent on C, then column A is dependent on C



Third Normal Form

A table is in the third normal form if it is the second normal form and there are no non-key columns dependant on other non-key columns that could not act as the primary key.

Normal Forms: 3NF

There are several levels above 3NF of "normal form" compliance, but generally the third normal form is good enough for 99% of all situations.

An informal intuitive definition of 3NF is as follows:

There are no fields in a table that are not directly determined by the values of the primary key.

Therefore, all fields in a table should be directly related to (determined by) the primary key of that table.

Normal Forms: 3NF Example

Suppose we have the following table:

| InvoiceNumber (PK) | InvoiceDate | Inventory ID | Inventory Description |
|--------------------|-------------|--------------|-----------------------|
| 1000 | 10/1/2019 | 45 | Hammer |
| 1001 | 10/3/2019 | 28 | Nails |
| 1002 | 10/3/2019 | 17 | Screwdriver |
| 1003 | 10/4/2019 | 45 | Hammer |

Some questions to consider:

- Is an invoice date/inventory ID directly related to an invoiceNumber?
- Is an inventory description directly related to an invoiceNumber?

Yes

No

Normal Forms: 3NF Example

Suppose we need a Spanish version of this database, and we need to value to show *Martillo* instead of Hammer. This would entail an UPDATE statement that targets 2 rows.

| InvoiceNumber (PK) | InvoiceDate | Inventory ID | Inventory Description |
|--------------------|-------------|--------------|-----------------------|
| 1000 | 10/1/2019 | 45 | Martillo |
| 1001 | 10/3/2019 | 28 | Nails |
| 1002 | 10/3/2019 | 17 | Screwdriver |
| 1003 | 10/4/2019 | 45 | Martillo |

Normal Forms: 3NF Example

In this situation, we could have split up the data into 2 tables, thus we end up with a less risky query, affecting only 1 row:

| InvoiceNumber (PK) | InvoiceDate | Inventory ID |
|-----------------------|-------------|--------------|
| 1000 | 10/1/2019 | 45 |
| 1001 | 10/3/2019 | 28 |
| 1002 | 10/3/2019 | 17 |
| 1003 | 10/4/2019 | 45 |

| Inventory ID (pk) | Description | |
|-------------------|-------------|--|
| 28 | Nails | |
| 17 | Screwdriver | |
| 45 | Martillo | |

The Relational Attestment to 3NF

The Database Designer's Swearing In Ceremony

Raise your left mouse button finger and repeat after me:

"I solemnly swear that all attributes will be related to, The KEY, The WHOLE Key, And NOTHING BUT the Key.

So help me Codd"

Your finger may now be rested



Many to Many relationships

Generally speaking, when there are 2 entities for which there is a "many to many" relationship, we will end up with 3 tables when considering 3NF as part of our design.



Many to Many relationships Example

Consider the MovieDB example:

- An actor can be a cast member of several movies.
 - A movie can have several actors.

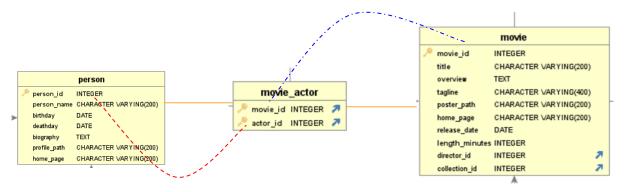
This is a "many to many" relationship.

Many to Many Relationships: The Problem

| Actor | | | | Film | |
|------------------------------------------|-----------------|--------------------|-----|---------|----------------------------------|
| Actor_Id | Name | | | Film_ld | Title |
| 1 | Clint Westworld | | | 1 | The Ok, Horrible, and Ridiculous |
| 2 | Adam Aimy | | | 2 | Two Donkeys for Brother Samuel |
| | | | | 3 | Departure |
| | | | 12 | 4 | Find Me If You Haven't |
| | | 4 Dat | A ! | | |
| Actor | | Micaleu. | | Film | |
| Actor_Id | Film_Id | Na Admiticated Dat | | Film_ld | Title |
| 1 | | Clint Westworld | | 1 | The Ok, Horrible, and Ridiculous |
| 1 | 2 | Clint Westworld | | 2 | Two Donkeys for Brother Samuel |
| Mpound of | 3 | Clint Westworld | | 3 | Departure |
| endent or 2 | 3 | Adam Aimy | | 4 | Find Me If You Haven't |
| Mnound PK 2 endent on her Table 19 | 4 | Adam Aimy | | | |

Many to Many relationships Example

We use a bridge table to break the Many to Many relationship so that the tables can be in 3rd normal form.

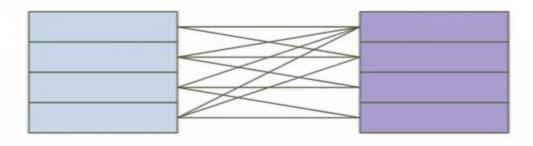


We have defined two foreign keys in the movie_actor table, the primary keys of each of the other two tables.

Many to Many Relationship

Many records in one table relate to

Many records in another table



Student and Class Tables

student

| ID | Name | |
|----|--------|--|
| 1 | John | |
| 2 | Mark | |
| 3 | Sarah | |
| 4 | Claire | |

class

| ID | Name |
|----|-------|
| 2 | DB01 |
| 5 | PH01 |
| 7 | WEB01 |
| 8 | WEB02 |

Student and Class Tables

student

| ID | Name | Class ID |
|----|--------|----------|
| 1 | John | 2, 5 |
| 2 | Mark | 5, 7 |
| 3 | Sarah | 8 |
| 4 | Claire | 2, 5, 8 |
| | | |

class

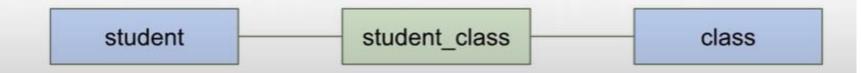
| Student ID |
|------------|
| 1, 4 |
| 1, 2, 4 |
| 2 |
| 3, 4 |
| |

Joining Table

Instead of this:

student class

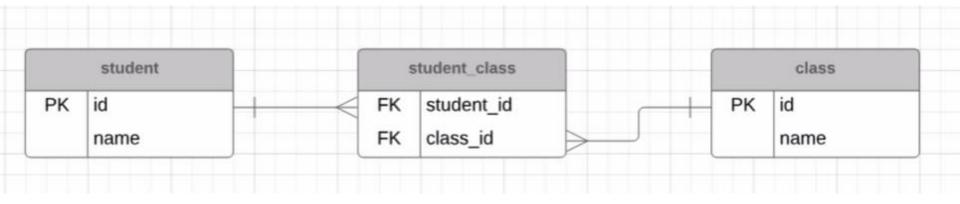
We have this:



The Joining Table



Captures every instance of student and class



| ID | Name |
|----|--------|
| 1 | John |
| 2 | Mark |
| 3 | Sarah |
| 4 | Claire |

| Student ID | Class ID |
|------------|----------|
| 1 | 2 |
| 1 | 5 |
| 2 | 5 |
| 2 | 7 |
| 3 | 8 |
| 4 | 2 |
| 4 | 5 |
| 4 | 8 |

| ID | Name | |
|----|-------|--|
| 2 | DB01 | |
| 5 | PH01 | |
| 7 | WEB01 | |
| 8 | WEB02 | |

DML vs DDL vs TCL

The SQL statements we have seen so far fall into a number of different categories:

- Data Manipulation Language (DML): SELECT, INSERT, UPDATE, DELETE
- Data Definition Language (DDL): CREATE, ALTER, DROP
- Transaction Control Language (TCL): BEGIN, TRANSACTION, COMMIT
- Data Control Language (DCL): GRANT, REVOKE

The focus of this lecture will be DDL statements with appropriate constraints.

Creating Tables Example

We are now ready to evaluate the syntax for table creation and alteration. This is the Create table syntax for all 3 of the previous tables:

```
CREATE TABLE person (
    person_id serial NOT NULL,
    person_name varchar(200) NOT NULL,
    birthday DATE NULL,
    CONSTRAINT pk_person PRIMARY KEY
(person_id)
);
```

In movie_actor are actor_id and movie_id foreign keys yet?

```
No!
```

```
CREATE TABLE movie_actor (
    actor_id integer NOT NULL,
    movie_id integer NOT NULL,
    CONSTRAINT pk_movie_actor PRIMARY KEY (actor_id, movie_id)
);
```

```
CREATE TABLE movie (
movie_id int NOT NULL DEFAULT nextval('movie_serial'),
title varchar(200) NOT NULL,
overview text NULL,
tagline varchar(400) NULL,
poster_path varchar(200) NULL,
home_page varchar(200) NULL,
release_date date NULL,
length_minutes int NOT NULL,
director_id int NULL,
collection_id int NULL,
CONSTRAINT pk_movie PRIMARY KEY (movie_id)
);
```

Creating Tables Example

We finish by specifying that actor_id and film_id are actually foreign keys. The DBMS does not assume this just because it has the same name, we must use the ALTER command:

ALTER TABLE movie_actor
ADD FOREIGN KEY(movie_id)
REFERENCES movie(movie_id);
ALTER TABLE movie_actor
ADD FOREIGN KEY(actor_id)

REFERENCES person(person id);

CREATE/DROP syntax

```
CREATE DATABASE database_name;
DROP DATABASE database_name;
```

```
CREATE TABLE table_name
(
    column_name1 data_type(size),
    column_name2 data_type(size) NOT NULL,
    column_name3 data_type(size),
    CONSTRAINT pk_column_1 PRIMARY KEY (column_name1),
    CONSTRAINT fk_column_2 FOREIGN KEY (column_name2)
        REFERENCES table_name(column_1)
);
```

ALTER syntax

```
ALTER TABLE table_name

ADD CONSTRAINT pk_constraint_name

PRIMARY KEY (column_name(s));
```

```
ALTER TABLE table_name

ADD CONSTRAINT fk_constraint_name

FOREIGN KEY (column_name) REFERENCES

table(column_name);
```

```
ALTER TABLE table_name

ADD CONSTRAINT chk_constraint_name

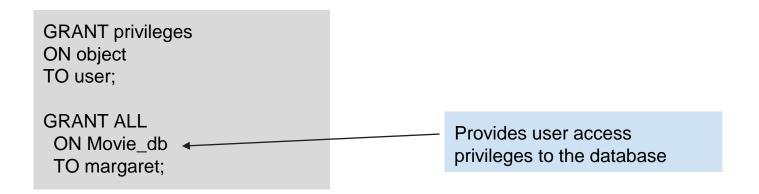
CHECK (column_name = 'value' OR

column_name IN (values));
```

DCL commands deal with the permissions, rights and other controls for the database system.

- CREATE USER Allows the creation of a user to the database
 - Users have permission to log in to the database by default
- CREATE ROLE Allows the creation of a role to the database
 - Roles do not have access to log in to the database (but can be granted this)
- GRANT allow a role or user access privileges to a database or table
- ALTER ROLE allows a role to be modified
- REVOKE remove access privileges to a database or table

DCL commands deal with the permissions, rights and other controls for the database system. Examples are GRANT and REVOKE



DCL commands deal with the permissions, rights and other controls fo the database system. Examples are GRANT and REVOKE

REVOKE privileges
ON object
FROM user;

REVOKE INSERT
ON Movie_db
FROM margaret;

Removes user access
privileges to the database

To change attributes of a role, you use the following form of ALTER ROLE statement:

