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Module Name : Algorithms & Data Structures Using Java.
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# DAY-01:
# Introduction to data structures?
Q. Why there is a need of data structures?
- if we want to store marks of 100 students
int m1, m2, m3, m4, ...., m100; //400 bytes if
sizeof(int) = 4 bytes
int marks [ 100 ]; //400 bytes - if sizeof(int) = 4 bytes
- we want to store rollno, marks & name
rollno
        : int
marks
        : float
        : char [] / String / String
name
struct student
    int rollno;
    char name[ 32 ];
    float marks;
};
struct student s1;
class Employee
    //data members
    int empid;
    String empName;
    float salary;
    //member functions/methods
};
Employee e1;
Employee e2;
=> to learn data structures is not learn any programming
language, it is a programming concept i.e. it is nothing
```

but to learn algorithms, and algorithms learned in data structures can be implemented by using any programming language.

```
# algorithm to do sum of array elements => end user
(human being)
step-1: initially take sum as 0.
step-2: traverse an array and add each array element
into sum sequentially
from first element max till last element.
step-3: return final sum.
# pseudocode => programmer user
Algorithm ArraySum(A, n)//whereas A is an array of size
"n"
{
    sum = 0;
    for ( index = 1; index \leq n; index++) {
        sum += A[ index ];
    return sum;
}
# program => compiler => machine
int arraySum(int [] arr, int size){
    int sum = 0;
    for( int index = 0 ; index < size ; index++ ) {</pre>
        sum += arr[ index ];
    return sum;
}
Bank Manager => Algorithm => Project Manager => Software
Architect
=> Pseudocode => Developers => Program => Machine
e.g.
Problem : "Searching" => to search / to find an element
(can be referred as a
key element) into a collection/list of elements.
```

```
1. Linear Search
2. Binary Search
Problem : "Sorting" => to arrange data elements in a
collection / list of elements either in an ascending
order or in a descending order.
1. Selection Sort
2. Bubble Sort
3. Insertion Sort
4. Merge Sort
5. Quick Sort
etc....
- when one problem has many solutions we need to go for
an efficient solution.
City-1:
City-2:
multiple paths exists => optimum path
distance, condition, traffic situation ....
- to traverse an array => to visit each array element
sequentially from first
element max till last element.
- there are two types of algorithms:

    iterative approach (non-recursive)

2. recursive approach
- recursion
- recursive function
- recursive function call
- tail-recursive function
- non-tail recursive function
Class Employee
    int empid;
    String name;
    float salary;
};
```

```
- object el is an instance of Employee class
Employee e1(1, "sachin", 1111.11);
Employee e2(2, "nilesh", 2222.22);
Employee e3(3, "amit", 3333.33);
# Space Complexity:
for size of an array = 5 \Rightarrow index = 0 to 5 \Rightarrow only 1 mem
copy of index = 1 unit
for size of an array = 10 \Rightarrow index = 0 to 10 \Rightarrow only 1
mem copy of index = 1 unit
for size of an array = n \Rightarrow index = 0 to n \Rightarrow only 1 mem
copy of index = 1 unit
for any input size array we require only 1 memory copy of
index var =>
simple var
+ sum:
for size of an array = 5 \Rightarrow sum \Rightarrow only 1 mem copy of sum
= 1 unit
for size of an array = 10 => sum => only 1 mem copy of
sum = 1 unit
for any input size array we require only 1 memory copy of
sum var => simple var
+ n = input size of an array -> instance characteristics
of an algo
for size of an array = 5 \Rightarrow if n = 5 \Rightarrow 5 memory copies
required to store 5 ele's => 5 units
for size of an array = 10 \Rightarrow if n = 10 \Rightarrow 10 \text{ memory}
copies required to store 10 ele's => 10 units
```

for size of an array =  $20 \Rightarrow if n = 20 \Rightarrow 20$  memory copies required to store  $20 \text{ ele's} \Rightarrow 20$  units

for size of an array =  $100 \Rightarrow if n = 100 \Rightarrow 100 memory$  copies required to store  $100 ele's \Rightarrow 100 units$ 

size

- for any input size array no. of instructions inside an algo remains fixed i.e. space required for instructions i.e. code space for any size array will going to remains fixed or constant.

```
int sum( int n1, int n2 )//n1 & n2 are formal params
{
   int res;//local var
   res = n1 + n2;
   return res;
}
```

- When any function completes its excution control goes back into its calling function as an addr of next instruction to be executed in its calling function gets stored into FAR of that function as a "return addr".

FAR contains:

local vars

formal params

return addr => addr of next instruction to be executed in
its calling function

old frame pointer => an addr of its prev stack frame/FAR.
etc...

```
# Linear Search:
for ( index = 1; index \leq n; index++) {
    if( key == arr[ index ] )
        return true;
}
return false;
- In Linear Search best case "if key found at very first
position"
for size of an array = 10, no. of comparisons = 1
for size of an array = 20, no. of comparisons = 1
for size of an array = 50, no. of comparisons = 1
for any input size array => no. of comparisons = 1
Running Time \Rightarrow O(1)
- In Linear Search worst case "if either key found at
last first position or key does not exists"
for size of an array = 10, no. of comparisons = 10
for size of an array = 20, no. of comparisons = 20
for size of an array = 50, no. of comparisons = 50
for size of an array = n, no. of comparisons = n
Running Time \Rightarrow O(n)
Lab Work => Implement Linear Search => by non-rec as well
rec way
```