

Assignment #3

Nikhil Unni

1. Write out the sentences G and the sentence D for the following machine (not pictured) with input $n = 2$.

We have the same “background information”:

(1)

$$\forall u \forall v \forall w (((S_{uv} \wedge S_{uw}) \implies v = w) \wedge ((S_{vu} \wedge S_{wu}) \implies v = w))$$

(2)

$$\forall u \forall v (S_{uv} \implies u < v) \wedge \forall u \forall v \forall w ((u < v \wedge v < w) \implies u < w)$$

(3)

$$\forall u \forall v (u < v \implies u \neq v)$$

(4)

$$\forall u \forall v \forall w (((S_{muv} \wedge S_{mw}) \implies v = w) \wedge ((S_{mvu} \wedge S_{mw}) \implies v = w))$$

(5)

$$\forall u \forall v \forall w (((S_{muv} \implies u < v) \text{ if } m \neq 0$$

(6)

$$\forall u \forall v \forall w (((S_{muv} \implies u \neq v) \text{ if } m \neq 0$$

(7)

$$\forall u \forall v \forall w ((S_{mwu} \wedge S_{uv}) \implies S_{kwv} \text{ if } k = m + 1$$

(8)

$$\forall u \forall v \forall w ((S_{kwv} \wedge S_{uv}) \implies S_{mwu} \text{ if } m = k - 1$$

(9)

$$p \neq q \text{ if } p \neq q$$

(10)

$$\forall v (S_{mv} \implies v = k) \text{ where } k = m + 1$$

(11)

$$\forall u (S_{uk} \implies u = m) \text{ where } m = k - 1$$

Now, the initial state of our machine with input $n = 2$ can be described as:

(12)

$$Q_1 0 \wedge @00 \wedge M00 \wedge M01 \wedge \forall x ((x \neq 0 \wedge x \neq 1) \implies \neg M0x)$$

And, finally, the encoding of our instructions:

(13)

$$\forall t \forall x ((Q_1 t \wedge @tx \wedge M_1 tx) \implies \exists u (Stu \wedge (\exists y (Sxy \wedge @uy)) \wedge Q_1 u \wedge \\ \forall y (M_0 ty \implies M_0 uy \wedge M_1 ty \implies M_0 uy)))$$

(14)

$$\forall t \forall x ((Q_1 t \wedge @tx \wedge M_0 tx) \implies \exists u (Stu \wedge (\exists y (S_{-1} xy \wedge @uy)) \wedge Q_2 u \wedge \\ \forall y ((M_0 ty \implies M_0 uy \wedge M_1 ty \implies M_0 uy))))$$

$$\begin{aligned}
(15) \quad & \forall t \forall x ((Q_2 t \wedge @tx \wedge M_1 tx) \implies \exists u (Stu \wedge (@ux \wedge M_0 ux) \wedge Q_j u \wedge \forall y ((y \neq x \wedge M_1 ty) \\
& \implies M_1 uy) \wedge \forall y ((y \neq x \wedge M_0 ty) \implies M_0 uy)))
\end{aligned}$$

Now for the parts of D:

$$(16) \quad \exists t \exists x (Q_2 t \wedge @tx \wedge M_0 tx)$$

So then G consists of the set of equations (1) to (15), and D is just equation (16).

2. *Show that*

$$\forall w \forall v (Twv \iff \exists y (Rwy \wedge Syv))$$

and

$$\forall u \forall v \forall y ((Suv \wedge Syv) \implies u = y)$$

together imply

$$\forall u \forall v \forall w ((Twv \wedge Suv) \implies Rwu)$$

Assume that for some u, v, w , Twv and Suv . Then, from the first formula, we know that there exists some y such that Rwy and Syv . But now we have Suv and Syv , so from the second formula, we know $y = u$. So then we know that since Rwy is true, Rwu is true as well.

3. *Show that*

$$\forall x (\neg Ax \implies \neg \exists t (Bt \wedge Rtx))$$

and

$$\neg \exists x (Cx \wedge Ax)$$

together imply

$$\neg \exists t \exists x (Bt \wedge Cx \wedge Rtx)$$

We can rephrase our assumptions as:

$$\forall x (\neg Ax \implies \forall t (\neg Bt \vee \neg Rtx))$$

$$\forall x (\neg Cx \vee \neg Ax)$$

and the final implication as:

$$\forall t \forall x (\neg Bt \vee \neg Cx \vee \neg Rtx)$$

We know that, for some x , either $\neg Cx$ or $\neg A$. If we assume $\neg Cx$, then $\forall t \forall x (\neg Bt \vee \neg Cx \vee \neg Rtx)$ is trivially true, since $\neg Cx$ is true. So let's assume that $\neg A$ is true. Then, we know that $\forall t (\neg Bt \vee \neg Rtx)$. This makes $\forall t \forall x (\neg Bt \vee \neg Cx \vee \neg Rtx)$ true as well, since $\forall x (\forall t (\neg Bt \vee \neg Rtx))$ is true.

4. *Verify that $G \models D$ for $n = 2$ where G and D are as defined in exercise 1. (You can appeal to facts about S and $<$ that have been established in the book, i.e. equations (1) to (11).)*

Let's start with the initial state of the machine, given by (12). From this, we know that at time 0, we're at Q_1 , and the only marked squares are at $x = 0$ and $x = 1$. So then from equation (13), we know that this together implies that there is a successor time (which can be expressed as $t = 1$, thanks to our background formulae), and a successor square ($x = 1$), such that at time $t = 1$, we're at square $x = 1$, and all of the squares' markings are unchanged. From equation (13) again, we move another square to the right, meaning we're at $x = 2$ and $t = 2$, and all squares remain unchanged. Now, from (12), we know that $\neg M02$, and since we know that the markings haven't changed from $t = 0$ to $t = 2$, we know that $\neg M22$.

So at this point, we have:

$$Q_1 2 \wedge @22 \wedge M20 \wedge M21 \wedge \forall x (x \neq 0 \wedge x \neq 1 \implies \neg M2x)$$

Then, from (14), we move to the left by 1, move to state Q_2 , and all the squares remain unchanged. This means we have:

$$Q_23 \wedge @31 \wedge M30 \wedge M31 \wedge \forall x(x \neq 0 \wedge x \neq 1 \implies \neg M3x)$$

From equation (15), we know this all implies that there's a successor time $u = 4 = t + 1$ such that @41 with a blank square 1: $\neg M41$, with all other squares besides x remaining unchanged. This leaves us at:

$$Q_24 \wedge @41 \wedge M40 \wedge \neg M41 \wedge \forall x(x \neq 0 \wedge x \neq 1 \implies \neg M4x)$$

Since we have $Q_24 \wedge @41 \wedge M40 \wedge \neg M41$, this makes statment (16) :

$$\exists t \exists x (Q_2t \wedge @tx \wedge M_0tx)$$

a true statement, with $t = 4, x = 1$ (since M_0tx is equivalent to $\neg Mtx$).