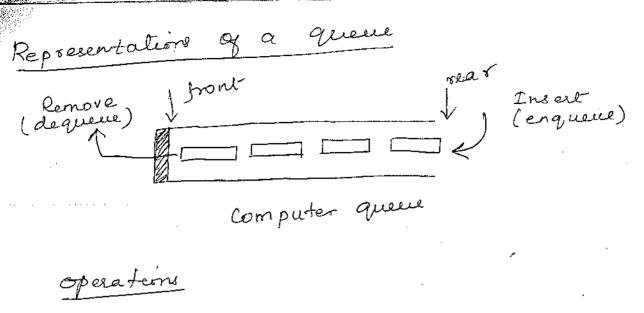
Queues

Definition

A queue is a linear list in which data can only be inserted at one end, called the That, and deleted from the other end, called the front. These restrictions ensure that the data are processed through the queel in the order in which they are received. In other. words, a queue is a first in - first out (FIFO) Structure

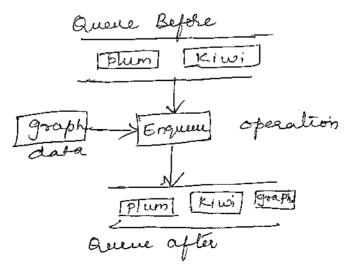
- A line of people waiting for the bus at a bus station is a queue
- A list of calls put on hold to be answered by a telephone operator is queue
- I list of waiting jobs to be processed by a computer is a queen.



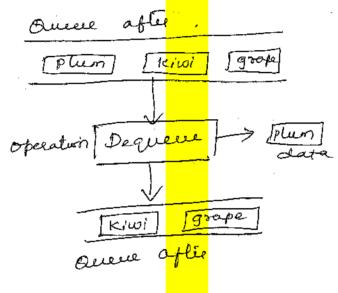
Operations
(1) Enqueue (2)
(2) Dequeue ()

Enqueue (Frqueue inserts an element at the real of the queue)

The queue insert operation is known as enqueue pflix the data have been inserted enqueue pflix the data have been inserted in to the queue, the new element becomes the into the queue, the queue, the queue another element in the queue, the queue is in an overflow state.



The queue delete operation is known as dequeue. The data at the front of the queue are returned to the user and removed from the quere. If there are no dava in the queue when a dequeue is attempted, the queue is in an underflow stale.



Applications of Quene

D'uner a resource is shared among multiple

eg: CPU Schooleeling, Diek Schooleeling

@ what data is transferred asynchronously Codata not necessarily received at some rate as sent) between two processes. eg Io Buffers, Pipes, file I/o etc.

```
Implementation of queue using array
```

```
His lude LStolio. h)
Hetapine que near [
# define Max 10
int quevears [Max];
 int reas = -1
 int front: -1;
 Void main()
     int choice;
      while (1)
       pointf ("1. Instert \n");
       painty ("2. Delete \n");
        paint ("3. Display \n");
        Paint ("4. Quit \n").
        Paint ("Enter your choice:");
         Scarf ("1.d", & choice).
         Switch (choice)
                      : insert();
                        break;
                        delete (), 1
                         break;
```

```
case 3: display();
                    break;
            Case 4: exit()
            default: prist ["woong choice ["]-
insert ()
   int item:
    if ( rear == Max -1)
     Priorf ("Queue overflow |n");
    else
      if ( front == - 1)
       prestf ("Insert the element in queue: ");
      efeort = 0;
      scarf ("1.d", 2 item);
       Alax = 9002+1;
       quevears [rear] = iten;
```

```
delet()
    if (front == -1) | front > near)
                  " Queue underflow \ n'');
         Refuso;
    else
       paintf ("Element deleted from queue is: 1/4 | ""),
                                           queueass [fent]);
         front = front + 1;
display ()
    if [ front == -1)
         printf (" Dueu is enty \n");
     else
       prints (" Ouere is:\n");
         for (e= front; if= rear; i++)
         printf ("/.d" que vearr [i]);
printf ("/");
printf ("\n");
```

Disadvantage of orglinary Queue 0 1 2 3 4 30 40 50 The above setuation arises when 5 elements day 10,20,30,40 and 50 are inserted and then defeted first too eleme 10 and 20. Now, if up to insect an item we get the message " Queue Overflow". In the above situation, near insertion ie denied even if space is available of the front end. This is because in our functions Defore insetting an element, we test who there lear is equal to Queve-Size-1. If so, we say . Onene is full and can't viscet. This is a disadvantage. This disadvantage Car be overcome using 2 methods. * Shift left method * Useng circular representation of

Note that each time an item is debted, all the elements towards right are moved to left by one position and are moved to left by one position and reac is decremented by 1. But, shifting the sear is costly in terms of computer time if deta is costly in terms of computer time if the data being stored is very large. The data being stored is very large.

The data being stored is recommended.

d

Circular Quene

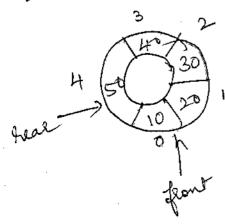
In circular queue, the elements of a given queue can be stored efficiently in an array so as to "wrap around" so that end of the queue is followed by the queue. If followed by the

To warp around the end of the queue the queue to follow the feant of queue the feelowing operation is performed.

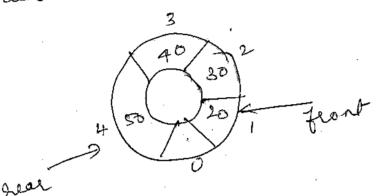
The following operation is performed.

hear (sear +!) 1/. Queve-Size

Assure Queue Size is 5 and 5 elements. are inserted into queue . The circular representation of queue as shown as follow



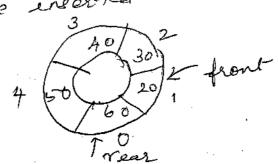
In the above queue, etem to is the streng deleterin, the item first element. So desired deleterin, the item to has to be deleted. This is achieved by to has to be deleted this is achieved by excrementing front by I so that front excrementing front by I so that front excrement the index of the second element.



Now if we want to enect an item to we have to increment rear by 1. In the value of rear is 4. The value of rear is 4. The wall will the obove figure, the value of rear will the walve increment hear by 1, it value will be 5. But, we have assumed that que we is circular. So, after incrementing by 1, it is cultared. So, after incrementing by 1, it is cultared. Should be a instead 5. This is achieved should be a instead 5. This is achieved.

Rear = (rear +1) 1. Queue_Size

After executing the above statement. The value of rear well be 0 so that item be comed at 0th position



n earD we increment front by , each time an item is deleted using the statement

•

/ front = (front +1) °/. Queue_sixe]

```
# is clude < s+die. h>
# include < Conio. h>
# define max 5
 void insert();
 void del();
 void display ();
 int Cq[max], sear = -1, front = -1
  void main()
      int choice;
     chrocres;
     do
          Printf ("1. Insect "\n");
          Peintf ("2. delete |n");
          Paint ("3. display \n");
          print ("4. ereit |n");
          Prints ("Enter your choice");
          Scarf ("1-d" & choice);
```

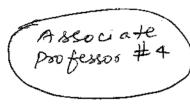
```
Switch (choice)
               insect(); break;
               del (); break;
     case 2:
              display (); break;
     case 3:
              exit(0);
3 while (D)
void insert ()
     int ele;
     if ( (seas +1) 1. max = = front)
        Peint (" Ouere full");
         if ( front = = -1)
          front = 0;
          Printf ("Enter the element | "");
          Scarf ("/. d", & ele).
          real = (real+1) / Max;
         CP[rear] = ele;
```

```
void del ()
      if ( front = = -1)
       Printy (" Queue Emply");
      prints ("Deleted Element "1.d", cg[front]);
       if (front = 2 sear)
           front =
         front = (front +1) 1. max;
void display ()
     if ("front = = -1)
       prints (" auere End Pty ");
      painté ("Content og the queue");
      forli = front; i != real; i = (i+1) 1. max)
       g print ("/d", cq [i]);
         Printy ("7. d'; cq l'i]);
```

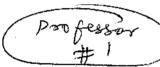
Priority Queue

A Special form of queue from which items are removed according to their designated priority and not the order is which they extered.

eg



Serior Professor #3



Assistant propessions)

tems entered the queue is Sequential Order but will be removed in the order #3, #1, #4, #21

29

Network souting - give priority to packets

voits strickest quality- of-Service requirements

In a priority queue each item (element)
las an associated priority of data

The main action that must be very efficient in a priority queue is finding the item with highest priority (usually the one with with highest priority (usually the one with the maximum of minimum key value)

or more elements -> lace element has a priority or value

- certife the FIFO queue, the order of deletion from the a priority queue (eg. who gets sorted nont) is determined by the element priority of
- Elements are deleted by increasing or decreasing order of priority Rather than order in which they arrived in a queue
- operations performed on priority queues * Find an element * Insert a new element
 - & Delete an element ete

Two kends of (Nin, Max) priority queues

L> In a Mis Priority queue, find | delete operation finde I delete the clement with minimum

Ly In a Max priority queue, find/delete find & / delete the element with maximum Priority

Peiosity Queue is a Collection of A priority queue is a Collection of elements, each one having an assigned to one-and-the same.

Insertion and defetion are 2 basic operations to be performed for this type of operations queue as well. However, deletion, operation is different from normal queue.

we classify a types of priority quants based upon the way in which the elements are deleted.

- (a) Ascerding priority Queue (min priority

 Queue) element with minimum priority

 or value is deleted first.
- (b) Descending priority Queue (Max priority
 queue) element with maximum priority
 or value is deleted first.

Priority Queue operations defined for a Priority Queue 2 operations defined for a

- 1. experience: enter an element, à cirto the queue crosspective of its priority.
- 2. depqueue: defeté an element from the queux cohose priority és the highest.

An efficient inplementation for the priority queue is to use a heap.

The below discussion is based on implementing priority Queue using array and Store the queue elements in an ensorted mannel.

Priority Queue Insertion - enfqueue.

The priority queue clesign, especially insertion, is same as the traditional queue one small different, however, is that we one small different, however, is that we shall not use a pointire front and many but manage with only one pointie, india

index = - I Initial Condition

10/25/15/17/11/21

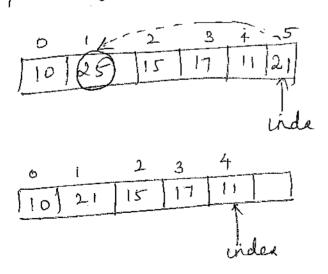
After insecting 10,25,15,17,11,21

we shall assume that the elemental values are the priorities. The task of rear position is taken by the variable index. when is taken by the variable index. when ender leader N-1 it signifies that queue ender leader N-1 it signifies that queue

Priority Queue Deletion - depqueue

ve Shall assume a decerding Priosity queue for our discussion. This means, the largest element must be removed. Seise the elements are not stored in any Particular Order, we must find the largest element and hemoire it.

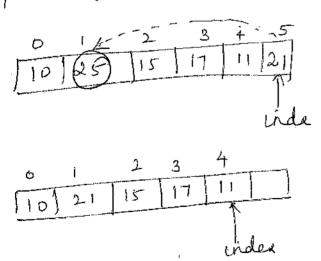
To defeti the element, first find the largest element and Sove it in temporary variable, k. copy the last element to occupy the position of the largest element. In our example, the largest element 25 is at position 1. That 25 and 21 and Swapped. Then docrement the index by one and refer the largest element (ie) the element with the highest paid sity.



Priority Queue Deletion - depqueue

De Shall assume a decerding Priosity queue for our discussion. This means, the largest element must be removed. Seise the elements are not stored in any Particular Order, we must find the largest element. and hemove it.

To delete the element, first find the largest element and save it in temporary variable, k. copy the last element to occupy the position of the largest element. In our example, the largest element 25 is at position 1. That 25 and 21 are swapped. Then document the index by one and referr the largest element. (ie) the element with the highest priority.



Every variable has a memory location and every memory localin has its address defined which can be accessed using ampensard(x) operator, which denotes an address in memory.

Consider the following example, which will prent the address of the variables defined

include LStdis. h>

void main ()

int- vari;

char var2 [10];

Printf ("Address of var variable: 1.x | n", & voe 1); Print ("Address q Vall variable: 1/x | n", 4 Vall);

Definition

A pointer is a variable whose value is the address of another variable. (ie) direct address of the memory location.

The general form of a pointer Variable declaration is

I type * var-name;

int *ip; double *dp; float *fp;

The actual data type of the value of all pointers, whether integer, float, character or otherwise, is the same, a long hexadecimal neember that represents a memory address The only difference between pointees of different data types is the data type of the variable or constant that the pointer points to

The following from mogram is the example q how to define a pointer variable, assign the address of a variable to a pointer and access the value at the address available is me pointer variable.

include Lstdio'h>

Void main()

{
 cre var = 20;
 int *ip';
 ip = e vae;

Printf("Address of a vae: '/x | n', & vae);

Printf("Address & to rad is ip vaeiable: '/x | n', ip).

Printf("Value of tip vaeiable: '/d | n', tip);

Printf("Value of tip vaeiable: '/d | n', tip);

NULL Pointers us a grood practice

It is always a grood practice

to asseg a NULL value to a pointer variable

to asseg you do not have exact address to be

is case you do not have exact address to be

assigned. This is done at the fine of variable

assigned. A pointer that is assigned NULL

declaration. A pointer

is called a new pointer.

The NULL Pointer is a Constant with a value of zero defined in Several standard libraries.

include < Stoler.h>

Void main()

sit * plx = NULL;

printf ("The value of ptx is: "/.x | n", ptx).

On most of the OS, programs are not Permitted to access memory at address a because that memory is reversed by the OS. However, the memory address a has special significance. it slights that the pointer is not intended to point to an accessible memory location. But by convention, if a pointer contains the null But by convention, if a pointer contains the null (zero) value, it assumed to point to nothing.

Pointer Asithmetic

that can be used on painters tt, --, + and -

Let us consider that ple is an integer pointer which points to the address wood Assuming 32 bit integers, let us possoon the following arithmetic operation on the pointer

ple++)

Point the location 1004 because each time ptr is point the location 1004 because each time ptr is in cremented, it will present to the next integer location. Which is 4 by the next to the current location. This operation will move the pointer to next memory location without compacting actual value at the memory location. If ptr points to a character whose address is 1000, then above operation will point to the location tool because next character point to the location tool because next character will be available (00).

in cleide (Stdio h)

const. int MAX = 3;

void main ()

int var [j = {10,100,200};

int i, * pti;

pti = var;

for (i=0; i < MAX; i+1)

{

Printf ("Address of var ['/d]='/.x\n", i, pti);

Printf ("value of var ['/d]='/.d|n", i, * pti);

Pti ++;

}

Array of pointers

There roay be a situation when we want to maintain an array, which can store pointers to an int prochar or any other data type available. Following is the declaration of an array of pointers to an integer.

[int *pta[3];

This declares ptr as an array of 3 integer pointers. Thus, each element in ptr, now holds a pointer to an int value. Following examples makes use of three cirtegers, which will be stored in an array of pointers as follows.

is clede (Stelio h)
void main ()

cit Var [] = \$10,100,200 };

for (1=0) (< 3 ; (++)

of platei) = & vartil;

Ĵ

for (i=0; i<3; (++)

f print ("value of var [1,d]=1,d\n",i", + Pta [i]).

```
Pointers with Strings
   # is clude (Stolio. A)
   mais ()
     char name EJ= "Pointers";
      char + pti;
         pto = 2 nameloj;
        while (+pta!="\0")
           printy ("1.c", + ptr);
            PJa++;
Pointers as Fanction Arguments
# in chede (Stoliv . h)
void swaf(int *, int *);
 void main ()
     int not, no2.
     no1=5, no2=10;
     swap (2 no1, 2 no2)
     Printf ("no1= 1.d, no2=1.d", no1, no2);
```

1

Vocal swap (int *a, int *b) int temp; temp = *a; *a= *b; *b=lemp; Frenction Returning painters Function can return address as they return integers, float or character type down. Relieve type neist explicitly be declased using Poen-ters # isclude (stdie. h) int feine(); Void main () int * b; b = fune (); Printf("b=1/x/n *b=7/d", b, *b); int & func () Septeent j=50

, selun(2);

Dynamic rosemony allocateurs

The process of allocating memory during program executein is called dynamic memory allocation.

Library. function There are 4 under "Stallib.h" for dynamic memory allocation. They are

- 1. malloc ()
- 2. calloc()
- 3. realloc()
- 4. free ()

The name malloc stands for malloct)
"memory allocation". The function malloct) reserves a block of memory of specified Size and retern a pointer of type void which can be easted into pointer of any form

Syntax

pla = (cast-type *) malluc (byte-size)

Here, pli is pointer of cost-type. The malloc() feine tein reteirne a pointer to an area of memory with Size of byte size. If the Space is insenfficient, allocation fails and returns NULL pointer.

in eleder (Stdio. h) # include & Streng. h) # isclude (Stellib.h) void main () char * mem_alloc; menalloc = malloc (20* Sixely(char)); if (mem_alloc == NULL) Print (" could not able to allocate requested memory"); f stropy (men-alloc, "kayawizhy"); pain of ("Dynamically allocated memory ement;"

7.5", mem-alloc);

free (mem-alloc);

The name called stands for "contiguous allocation". The only difference between mallocis and callocis is that mallocis allocates Length block of memory whereas (allocis) allocates block of memory each of Same nultiple block of memory each of Same size and sele all bytes to xero:

Syntas

pte = (cast-type +) calloc (n, element- fize); /

This statement will allocate contiguous space in memory for an array of netements.

eg

Pth = (float *) calloc (25, Size of (float));

This statement allocates contiguous space in memory for an array of 25 elements each of sixe of float (ie) 4 bytes.

Dynamically allocated memory with either calloc() or malloc() does not get return (alloc() or malloc() does not get return on its our . The programmer must use free() explicitly to release space

Thee (pts);

the space is memory
pointer by pte to be deallocated

```
C program to find Seen of n elements.
entered by user. To feeform this program,
altocat memory dynamically.
# include & Stollo. h>
# include < Stalib. h>
  void main ()
       int n, i, + pla, sum = 0
         Priotif ("Enter number of elements: ");
         Scarf["1.d", 20)",
         pta = (int *) · callo c (n, cize of (int));
          4 (ptr == NULL)
             print ("Error! memory not allocated").
           q exit(0);
           paintf ("Enter elements of array:");
           for(i=0', i<n', i++)
                Scanf (" -1.d", pta+i);
             3. Sum: Sum + + (pte+1);
          Print[" Sum = 1/d", Sum);
           free (pts);
```

to

If the previously allocated memory is in sufficient of more than sufficient. Then, you can change memory sixe previously allocated using realloc().

Syntax

pti = realloc (pti, newsize);

Here, ple is reallocated with size of new cize

#include < stdio. h)
#include < stdib. h>

void main ()

Ş

int +pte, i, n1, n2;

Printf ("Enla to size of array");

Scanf (+ 1, 4", 2 n 1);

pti = (cit *) malloc (n1 * size of (int));

Printf (" Adds of previously allocated minory");

tor (i=0; [< n1; (++)

Peistf ("Y. U1, Plati);

plint ("Enter new size of array");

Scanf ("1.d" 4n2);

pla = realloc (Pla, n2); for (i=0; i< A2; i++);

Pain # (" 1/4", pla+1);

1