GRAPH THEORY

Presented by

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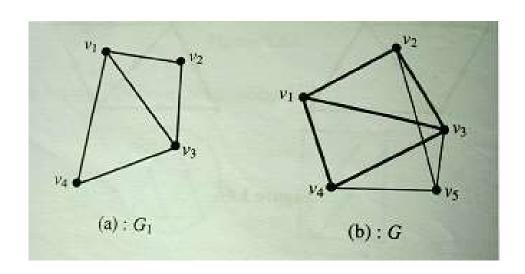
Topics

- Sub-Graphs
- Connected and Disconnected Graphs
- Eulerian Graph, Hamiltonian Graphs
- Trees: Spanning and minimal spanning tree: Kruskal's and Prim's Algorithms
- Shortest Path: Dijkstra's algorithms

Subgraphs

Let G(V,E) and $G_1(V_1,E_1)$ be two graphs. The graph G_1 is said to be **SUBGRAPH** of G, if the following conditions hold

- 1) All the vertices and all the edges of G₁ are in G.
- 2) Each edge of G_1 has the same end vertices in G as in G_1 In general, a subgraph is a graph which is a part of another graph.



Subgraphs

Properties:

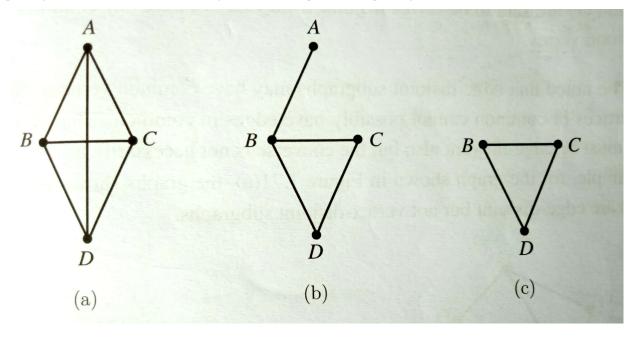
- 1) Any graph isomorphic to a subgraph of a graph G is also referred to as a subgraph of G.
- 2) Every graph is a subgraph of itself.
- 3) Every simple graph of n vertices is a subgraph of complete graph K_n
- 4) If G1 is a subgraph of a graph G2 and G2 is a subgraph of a graph G, then G1 is a subgraph of G.
- 5) A single vertex in a graph G is a subgraph of G.
- 6) A single edge in a graph G, together with its end vertices, is a subgraph of G.

Spanning Subgraph

Let G(V,E) be graph. If there is a subgraph $G_1(V_1,E_1)$ of G such that **V1=V**, then G1 is called a **SPANNING SUBGRAPH** of G.

In other words, a subgraph G1 of a graph G is a spanning subgraph of G whenever G1 contains all vertices of G.

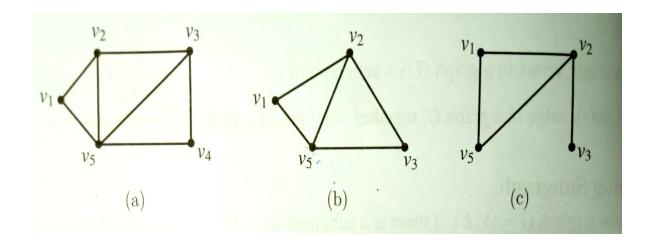
Note: Every graph is its own spanning subgraph.



Induced Subgraph

Let G(V,E) be graph. suppose there is a subgraph $G_1(V_1,E_1)$ of G such that every edge $\{A,B\}$ of G, where $A,B \in V1$ is an edge of G1 also. Then G1 is called an **INDUCED SUBGRAPH** of G (induced by V1) and is denoted <V1>.

It follows that a subgraph G1 = (V1, E1) of a graph G=(V,E) is not an induced subgraph of G if for some A,B \in V1, there is an edge $\{A,B\}$ which is in G but not in G1.



Edge-disjoint and Vertex-disjoint Subgraphs

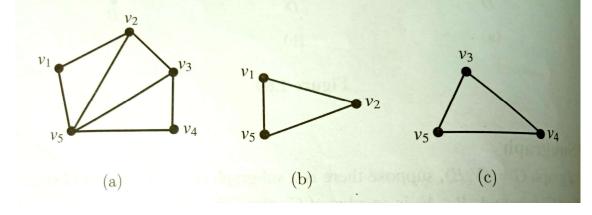
Let G(V,E) be graph and G1 and G2 be two subgraphs of G. Then:

- 1) G1 and G2 are said to be **EDGE-DISJOINT** if they do not have any edge in common.
- G1 and G2 are said to be VERTEX-DISJOINT if they do not have any common edge and any common vertex.

Note: 1) Edge-disjoint subgraphs may have common vertices.

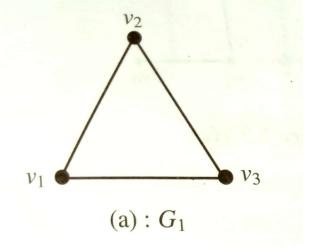
2) Vertex disjoint subgraphs must be edge-disjoint also but converse

need not be true.



Given a graph G1, can there exist a graph G2 such that G1 is a subgraph of G2 but not a spanning subgraph of G2 and yet G1 and G2 have the

same size?

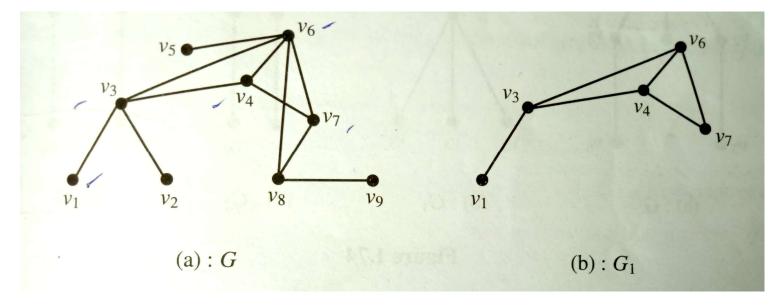


Consider the graph G shown in figure

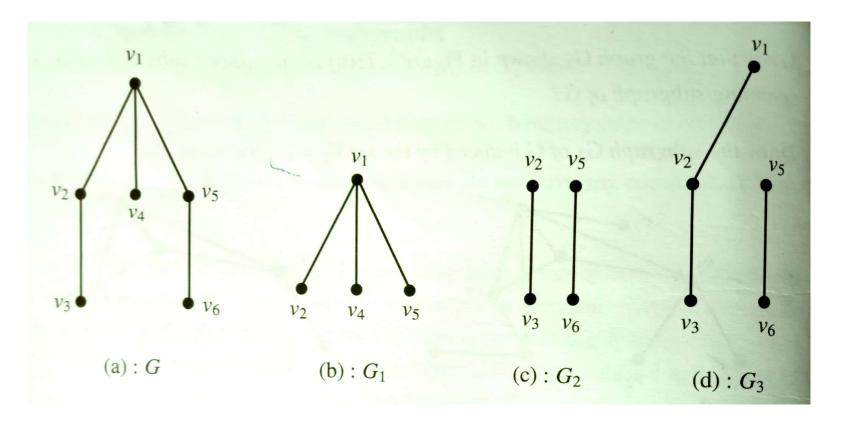
1) Verify that the graph G1 is an induced subgraph of G. Is this a spanning subgraph of G?

2) Draw the subgraph G2 of G induced by the set V2={v3, v4, v6,

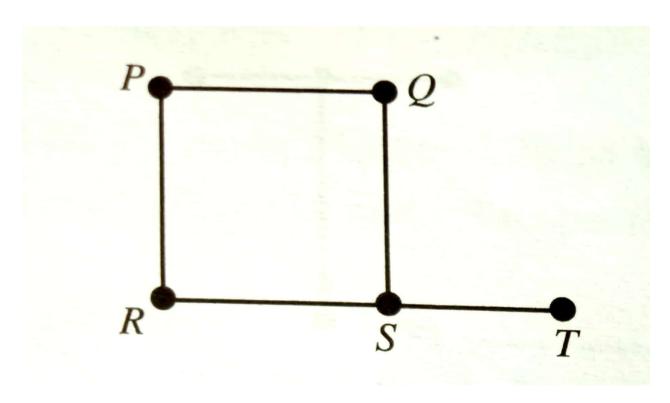
v8,v9}.



Consider the graph G shown in figure. Verify that the graphs G1, G2 and G3 are induced subgraphs of G.



Fog the given graph, find two edge-disjoint subgraphs and two vertex-disjoint subgraphs.



- 1) Walk
- 2) Trail
- 3) Circuit
- 4) Path
- 5) Cycle

WALK: Let G be a graph having at least one edge. In G, consider a finite, alternating sequence of vertices and edges of the form

$$v_i e_j v_{i+1} e_{j+1} v_{i+2},...., e_k v_m$$

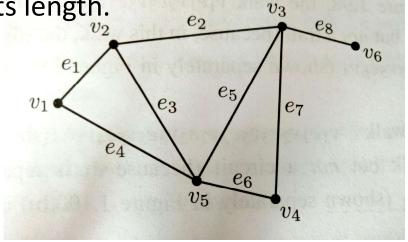
which begins and ends with vertices and which is such that each edge in the sequence is incident on the vertices preceding and following it in the sequence. Such a sequence is called a WALK in G.

Note: In a walk, a vertex or an edge or both can appear more than

once.

The number of edges present in a walk is called its LENGTH. e_7 V5

Write different walks from graph and find its length. v_2



The vertex with which a walk begins is called the **initial vertex (or origin)** of the walk and the vertex with which a walk ends is called the **FINAL VERTEX (OR THE TERMINUS)** of the walk.

The initial and final vertex of a walk are together called its terminal vertices.

Note: Terminal vertices need not be distinct.

Nonterminal vertices of a walk are called its **INTERNAL** vertices.

Notation: A walk having u as the initial vertex and v as the final vertex is called a walk from u to v, or briefly a u-v walk.

OPEN AND CLOSED WALK

A walk that begins and ends at the same vertex is called a **CLOSED WALK**.

A walk which is not closed is called an **OPEN WALK.**

TRAIL AND CIRCUIT

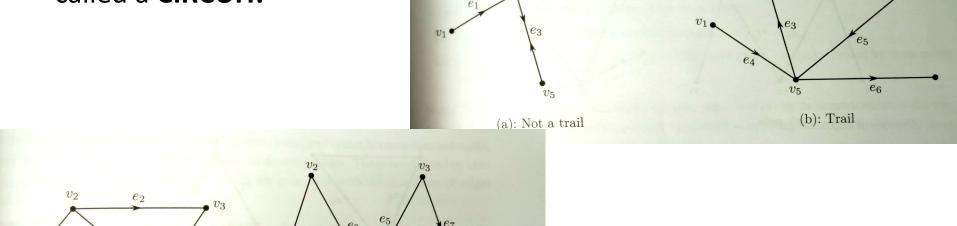
If in an open walk no edge appears more than once, then the walk is called a **TRAIL.**

If in a closed walk no edge appears more than once, then the walk is

called a CIRCUIT.

 e_4

(a): Not a circuit



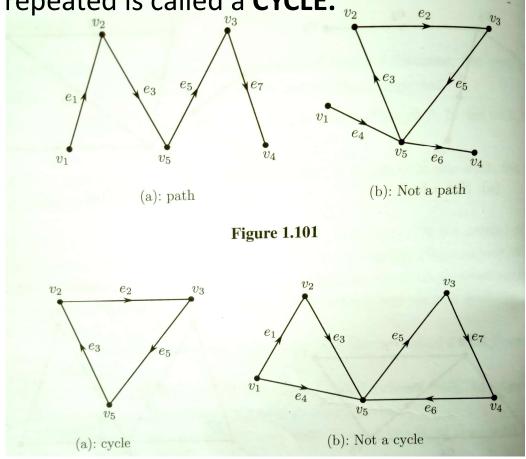
(b): circuit

PATH AND CYCLE

A trail in which no vertex appears more than once is called a PATH.

A circuit in which the terminal vertex does not appear as an internal vertex

and no internal vertex is repeated is called a CYCLE.



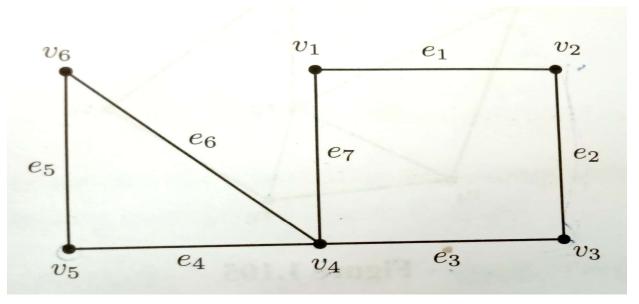
PROPERTIES

- A walk can be open or closed. In a walk (closed or open), a vertex and/or an edge can appear more than once.
- A trail (circuit) is an open (closed) walk in which a vertex can appear more than once but an edge cannot appear more than once.
- A path (Cycle) is an open (closed) walk in which neither a vertex nor an edge can appear more than once.
- Every path is a trail but a trail need not be a path
- Every cycle is a circuit but a circuit need not be a cycle.
- If a cycle contains only one edge, it has to be a loop.
- Two parallel edges (when they occur form a cycle.
- In a simple graph, a cycle must have at least three edges.
- A Cycle formed by three edges is called a TRIANGLE.

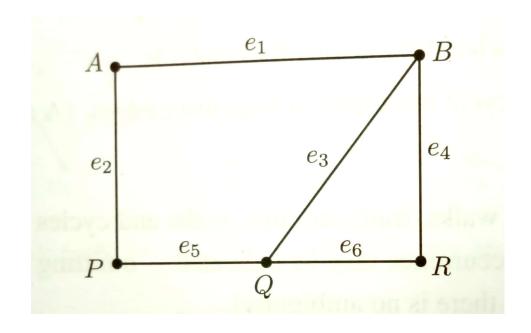
For the graph shown in figure, indicate the nature of the following walks.

(i) v1e1v2e2v3e2v2 (ii) v4e7v1e1v2e2v3e3v4e4v5 (iii) v1e1v2e2v3e3v3e7v1

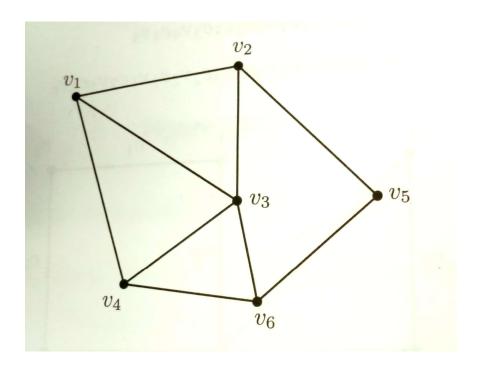
(iv) v1e1v2e2v3e3v4e4v5 (v) v6e5v5e4v4e3v3e2v2e1v1e7v4e6v6



From the graph find all paths from vertex A to vertex R. Also indicate their lengths?

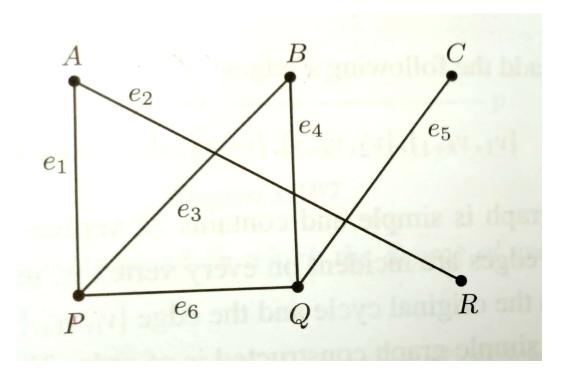


Determine the number of different paths of length 2 in the graph shown below?



If G is a simple graph in which every vertex has degree at least k, prove that G contains a path of length at least k.

Find all the cycles present in the graph



Prove the following:

- (i) A path with n vertices is of length n-1
- (ii) If a cycle has n vertices it has n edges
- (iii) The degree of every vertex in a cycle is two.

